



* Assignment No. 7 *

- Title:- program for constructing telephone book database of a client. make use of a hash table implementation to quickly lookup client telephone no.

* Objective :-

- 1) Basic concept.
- 2) Concept of recursion of insertion.

* problem Statement.

consider telephone book database n clients.
make use of hash table implementation to quickly look up telephone number.

* Outcomes:

Input :-

- 1) Enter name:-
- 2) Enter phone number.

Output:-

Insert element into table Search element from key delete element of a key.

• Hardware Requirement :-

8GB RAM, 500GB and Intel Processor.

• Software Requirement :-

gedit s/w terminal, gcc/g++ compiler.



* Theory:-

Hash table are an efficient implemented of a key array data structure. a struct sometime known as an associative array.

* Hashing function:-

used to We are designing a container which will be ~~wired~~ holds some number of items of given set K . In this context we all the element of the set K key the generated key in the array is given by function is called a hash function.

• method of hash function:-

- 1) Division method
- 2) mid-square method.
- 3) Folding method.
- 4) Digit analysis.
- 5) length dependent method.
- 6) Algebraic coding.
- 7) multiplicative hashing.

* Collision Resolution:-

collision resolution in the main program in hashing. If the element to be inserted is mapping to the same locate where element is inserted. then we have a collision is it must be resolved. There are several strategies for collision resolution. The most commonly used are

collision resolution strategies.



Dispersive channing
2) open addressing

- a) Linear problem
- b) Quadratic probing.
- c) double hashing.

Consider an Example.:

Sr. No.	Name	mobile no.
1.	Shreyas	9298182696.
2.	Deepika Nimbalkar	9826688210.
3.	Harshali Kale	9928342612
4.	Neha Landge	9712118018.
5.	Teju Salunkhe	9421349868
6.	Oneha Kale	984442816.

In above table it is a telephone breaks list. of peoples. Now by using linear probing we can create a hash tables by using absolute of list name letter of name lie.

$$H(83) \% 10 = 3$$

$$H(72) \% 10 = 2$$

$$H(69) \% 10 = 9$$

$$H(78) \% 10 = 8 \text{ collision,}$$

$$H(84) \% 10 = 4$$

$$H(83) \% 10 = 3 \text{ collision,}$$



	Name	mod. no.
0		
1	Harshali kale	
2	Harsh.	
3	Shrey a padve.	9928342618
4	Teju salunke.	72991826188
5	Sneha kale.	8087125421
6		
7		
8	Deepika Nimbalkar	8018762182
9	Neha landge.	9712118018

Basic Operation:

following are the primary operation of a hash table.

- 1) Search.: search on element in a hash table.
- 2) Insert an element in a hash table.
- 3) delete an element from a hash table.

Search Operation..:

Whenever an element is to be compared the hash table code of the key passed, and locate the element using hash table code is index in the array. use it. the element is not found at the computed hash code

Example.:

struct node

{

int data;

int key;

};



struct node * Search (int key).

{

int hashindex = hashcode (key).

while (hashArray[hashIndex] != key)

{

if (hashArray[hashIndex] → key == key)

return hashArray[hashIndex];

hashindex : hashindex + 1; size;

}

return NULL;

}

• Insert operation:

Whenever an element is to be insert compute the hash table code of key passed code as an index in the array. Use linear probing for empty location. If an element is found at the computed hash code.

Ex:

void insert (int key, int data).

{

struct node * item (struct node k) malloc
(size of (struct node));

item → data = data;

item → key = key;

int hashIndex = hashcode (key);

hashArray[hashIndex] → key != -1;

{

++hashIndex;

hashIndex % size;

}



g. $\text{hashArray}[\text{hashArray}] = \text{item};$

• Delete Operation:

Whenever an element is to be deleted compute the hash code of the key deleted compute the hash code and locate using that use linear probing to get the element has if an element is not found at the computed item there to keep performance of the hashtable.

• Conclusion:

Hence, we have studied and implemented the telephone book directory using hashing.