

## \* Assignment No.: 3 \*

\* Title:- program for constructing a postorder Traversal.

\* Objective:-

To study and implement the postorder Traversal

\* problem Statements:-

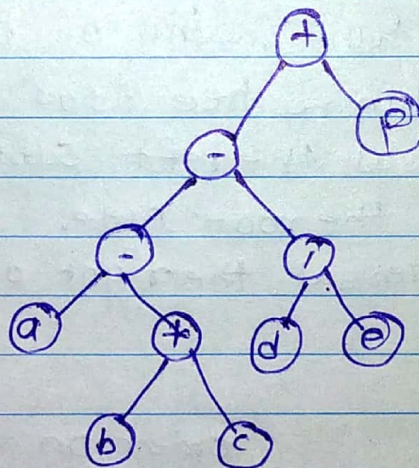
for a given Expressions the postorder Traversal preorders and inorder it using  $a-b*c-d/ef$ .

\* Outcome:-

◦ Input :-

Enter the valid expression  
 $a-b*c-d/ef$

output



\* Theory:-

Binary tree traversal (DFS) :- Most of the operation requires traversing a tree in a particular order. Traversing a tree and exactly once.

Since, a binary tree is defined in a recursive manner. tree traversal could be defined recursively.





for Example :- to traverse a tree, one may visit the root first, then the left subtree and finally traverse the right subtree. If we impose the restriction that left subtree and finally traverse the right subtree.

- 1) Visit the root, traverse, left subtree, traverse right subtree.
- 2) Traverse left subtree, visit the root, traverse subtree.
- 3) Traverse left subtree, traverse right subtree, visit the root.

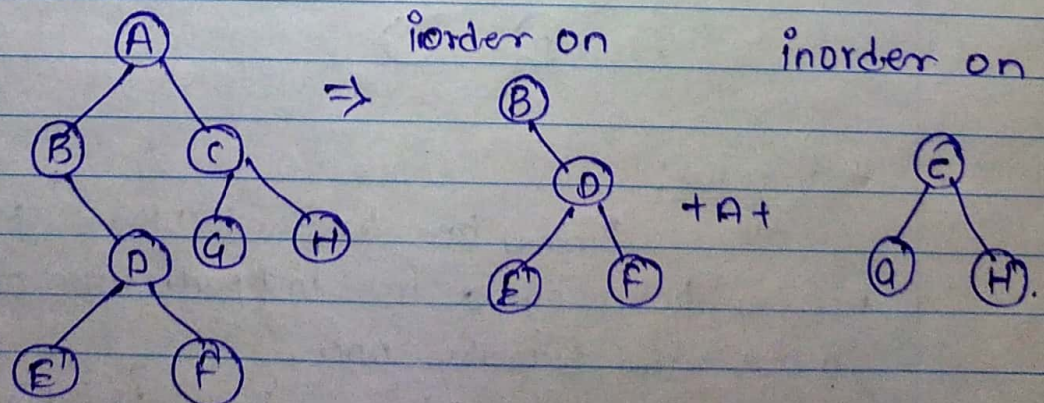
These three techniques of traversal are known as postorder, inorder, and preorder traversal of a binary tree.

### 1) Inorder traversal:-

The functioning of inorder traversal of a non-empty, binary tree is as follows:-

- i) firstly, traversal the left subtree in order.
- ii) Next, visit the root node.
- iii) At last, traversal the right of tree shown below

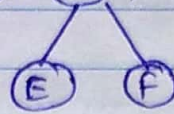
\*inorder on :-







$$= (B + \text{inorder on } (D)) + A + (\text{inorder on } (G)) + C +$$



$$(\text{inorder on } (H))$$

$$= (B + (\text{inorder on } (E)) + D + (\text{inorder on } (F))) + A + (G(H))$$

$$= BEDFAGCH.$$

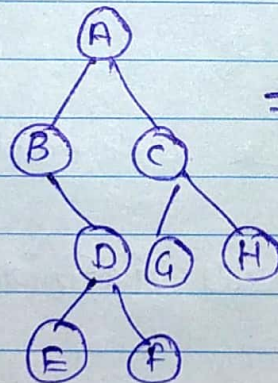
\* postorder traversal :-

The functioning of postorder traversal of a non-empty binary tree is as follows:-

- i) firstly, traversal the left subtree in postorder.
- ii) Next, traverse the right subtree in postorder.
- iii) At last, visit the root node.

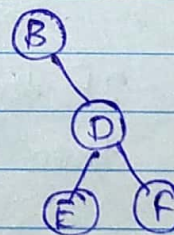
Stepwise postorder traversal of tree shown below.

postorder on :-



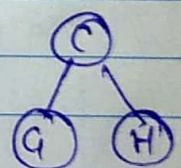
postorder

$\Rightarrow$

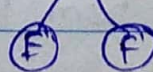


+

postorder on



$$\Rightarrow ((\text{postorder on } (B) + B) + ((\text{postorder on } (G)) + (\text{postorder on } (H) + C)) + A.$$



$$\Rightarrow EFDBGHC A.$$





### Algorithm for non-recursive inorder traversal:-

Algorithm for non-recursive inorder traversal of binary tree too works similar to non-recursive preorder traversal, a node is visited before it is pushed into the stack. Immediately after, it is popped from the stack. In non-recursive inorder traversal a node is visited before it is pushed into the stack. Immediately after, it is popped from the stack.

Step 1:- Start traversing from root (say T), traverse left and continue traversing left. All the traversed are pushed into the stack (say S).

```
while (T != NULL)
```

```
{
```

```
    S.push(T);
```

```
    T = T → left;
```

```
}
```

Step 2:-

If the stack S is empty

then

traversal is finished

else

{

Visited the right subtree of the node popped from the stack.

```
T = S.pop(), Visit()
```

```
T = T → right();
```

```
while (T != NULL)
```





5

S.push(T);  
T = T → left();

2.

2.

Step 3: Go to Step 3.

\* Algorithm for non-recursive postorder traversal.

Non-Recursive postorder traversal works in slightly different way. In postorder traversal, element is visited after the right subtree is traversed.

Thus, the address of the node should be preserved in stack until it has been printed. Non-Recursive algorithm encounters a node three times:-

- 1) while going to the left.
- 2) During backtracking to traverse the right subtree.
- 3) while returning from the right.

An additional field 'flag' in stack is used to differentiate between the two situations.

• while going to left, address of the node along with flag=0 is pushed onto stack with flag to 1

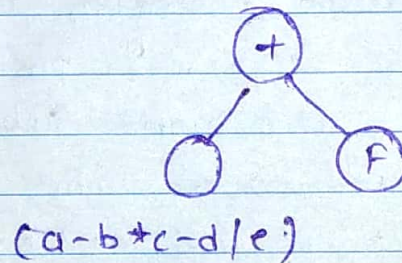
• During backtracking to traverse the right subtree, an element is popped from the stack and if flag field is found to be 0, it is pushed back with flag to 1.



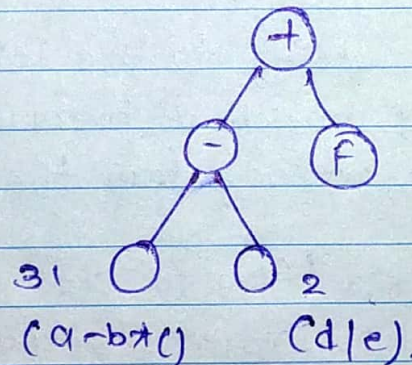
When we return from the right subtree and popped element from the stack, the flag field will be 1. This is the time when we visit the node. For given expression e.g.  $a-b*c-d/e+f$  construct inorder sequence & traverse it using sequence & traverse it using postorder traversal.

The operator which evaluate the last make it as root node. Divide the Expression.

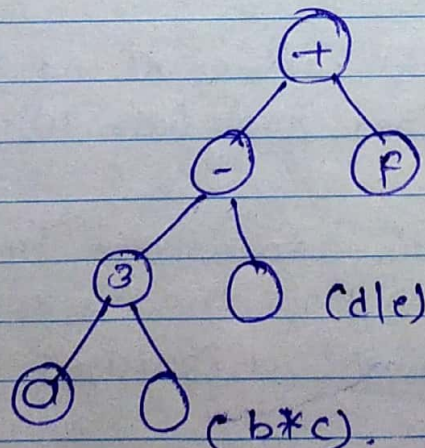
Step 1



Step 2:-



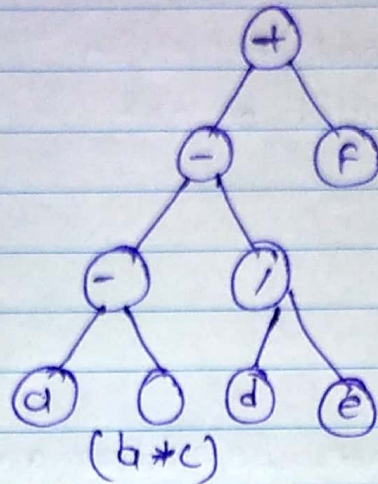
Step 3:-



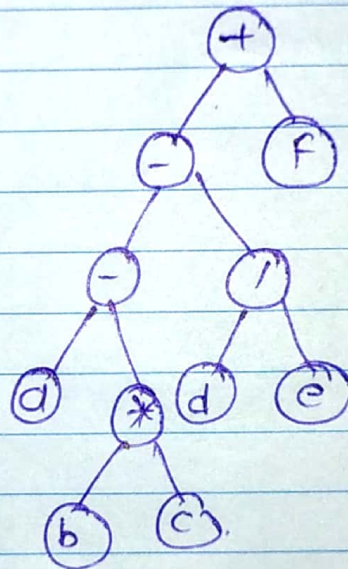




Step :- 4



\* Step 5:-



postorder traversal :-  $a b c * - d e / - f +$

Test case :-

Test Id.

Test Scenario.

Test step.

Tc-1.

i) check wheather  
tree is empty  
or not

check tree  
is  $(T == NULL)$   
then it is empty





T.C - 2

ii) Check wheather  
external expression  
is valid or not

check tree.  
expression check  
contain the valid  
Operand with  
valid operator.

\* Conclusion:-

Hence, we studied and implementation  
the program for construct inorder Sequence Epression  
Using postorder traversal.