



COMP9020

Foundations of Computer Science
Term 3, 2024

Lecture 1-2: Introduction, Number Theory

Outline

Course introduction

- Who are we?
- Why are we here?
- How will you be assessed?

Number Theory

- Number Theory in Computer Science
- Numbers and Numerical Operations
- Divisibility
- Greatest Common Divisor and Least Common Multiple
- Euclidean Algorithm
- Modular Arithmetic
- Euclidean Algorithm (again)

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Lectures

Lecturers: Jiaojiao Jiang (LiC), Paul Hunter
Times: Thursday 11-1pm and Friday 11-1pm

Admin

Name: Hao Ren
Course email: cs9020@cse.unsw.edu.au

Tutorials

Tutors: Different tutors each session
Times: Check the detailed [timetable](#) on WebCMS

Links

Course webpages:

- [WebCMS](#)
- [Moodle](#)

Lectures:

- Recordings available on echo360 (through [Moodle](#))

Other points of contact:

- [Course forums \(Ed Forum\)](#)
- Email: `cs9020@cse.unsw.edu.au`

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What is this course about?

Computer Science is about exploring the ability, and limitation, of computers to solve problems. It covers:

- **What** are computers capable of solving?
- **How** can we get computers to solve problems?
- **Why** do these approaches work?

This course aims to increase your level of mathematical maturity to assist with the fundamental problem of **finding, formulating, and proving** properties of programs.

Key skills you will learn:

- Working with abstract concepts
- Giving logical (and rigorous) justifications
- Formulating problems so they can be solved computationally

Course Structure

The actual content is taken from a list of topics that constitute the basis of the tool box of every serious practitioner of computing:

- number theory week 1
- set theory week 2
- relation week 3
- function and boolean week 4
- propositional and sequence & Induction week 5
- **mid-term test** (no lectures) week 6
- recursion, counting week 7
- probability and statistics week 8
- graph week 9
- algorithm analysis & formal languages week 10

Course Material

Textbooks:

- KA Ross and CR Wright: [Discrete Mathematics](#)
- E Lehman, FT Leighton, A Meyer:
[Mathematics for Computer Science](#)

Alternatives:

- K Rosen: Discrete Mathematics and its Applications

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Assessment Summary

- 1 online quizzes (weeks 1, 2, 3, 4, 5, 7, 8, 9) — max. marks 20
- 2 mid-term test — max. marks 20
- 3 final exam — max. marks 60

Take Notice

*To pass the course, your overall score must be **50 or higher** and your mark for the final exam must be **24 or higher**.*

The weekly quiz:

- becomes available after the Thursday lecture each week
- is due **Friday, 23:59** in the following week

Late policy and Special Consideration

All assessments are submitted through the course website

Lateness policy

- Quizzes: Late submissions not accepted
- Exams: Late submissions not accepted

If you cannot meet a deadline through illness or misadventure you need to apply for [Special Consideration](#).

Credits

COMP9020 credit for material goes to:

- Michael Thielscher
- Paul Hunter
- Katie Clinch
- Sebastian Sequoiah-Grayson
- more...

Pre-course polls



Pre-course questionnaire



Pre-course poll

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Reading Material

If you'd like to read more about the topics covered in this lecture, check out the following chapters of the recommended textbooks:

	[LLM]	[RW]
Week 1 Number Theory	Ch. 8	Ch. 1, 3

- **[RW]** is KA Ross and CR Wright: [Discrete Mathematics](#)
- **[LLM]** is Lehman, Leighton, Meyer: [Mathematics for Computer Science](#)

Number Theory in Computer Science

In this course, we are interested in **discrete mathematics**. This is the theory of e.g. the integers.

Continuous mathematics instead considers number systems with no “gaps”, e.g. the real numbers.

Applications of **discrete number theory** include:

- Cryptography/Security (primes, divisibility)
- Large integer calculations (modular arithmetic)
- Date and time calculations (modular arithmetic)
- Solving optimization problems (integer linear programming)
- Interesting examples for future topics in this course

Question

What is something that is easy to do with real numbers but hard to do with integers?

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Notation for numbers

Definition

- Natural numbers $\mathbb{N} = \{0, 1, 2, \dots\}$
- Integers $\mathbb{Z} = \{\dots, -1, 0, 1, 2, \dots\}$
- Positive integers $\mathbb{N}_{>0} = \mathbb{Z}_{>0} = \{1, 2, \dots\}$
- Rational numbers (fractions) $\mathbb{Q} = \left\{ \frac{m}{n} : m, n \in \mathbb{Z}, n \neq 0 \right\}$
- Real numbers (decimal or binary expansions) \mathbb{R}
 $r = a_1 a_2 \dots a_k . b_1 b_2 \dots$

In \mathbb{N} and \mathbb{Z} different symbols denote different numbers.

$$1 \neq 2 \neq 3$$

In \mathbb{Q} and \mathbb{R} the standard representation is not necessarily unique.

$$\frac{1}{2} = \frac{2}{4} = \frac{3}{6}$$

Floor and ceiling

Definition

$\lfloor \cdot \rfloor : \mathbb{R} \longrightarrow \mathbb{Z}$ — **floor** of x , the greatest integer $\leq x$

$\lceil \cdot \rceil : \mathbb{R} \longrightarrow \mathbb{Z}$ — **ceiling** of x , the least integer $\geq x$

Example

$$\lfloor \pi \rfloor = 3 = \lceil e \rceil \quad \pi, e \in \mathbb{R}; \lfloor \pi \rfloor, \lceil e \rceil \in \mathbb{Z}$$

Floor and ceiling

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Example

$$\lfloor \pi \rfloor = 3 = \lceil e \rceil \quad \pi, e \in \mathbb{R}; \lfloor \pi \rfloor, \lceil e \rceil \in \mathbb{Z}$$

Simple properties

- $\lfloor -x \rfloor = -\lceil x \rceil$, hence $\lceil x \rceil = -\lfloor -x \rfloor$
- For all $t \in \mathbb{Z}$:
 - $\lfloor x + t \rfloor = \lfloor x \rfloor + t$ and
 - $\lceil x + t \rceil = \lceil x \rceil + t$

Fact

Let $k, m, n \in \mathbb{Z}$ such that $k > 0$ and $m \geq n$. The number of multiples of k between n and m (inclusive) is

$$\left\lfloor \frac{m}{k} \right\rfloor - \left\lfloor \frac{n-1}{k} \right\rfloor$$

Absolute value

Definition

$$|x| = \begin{cases} x & , \text{ if } x \geq 0 \\ -x & , \text{ if } x < 0 \end{cases}$$

Example

$$|3| = |-3| = 3 \quad 3, -3 \in \mathbb{Z}; \quad |3|, |-3| \in \mathbb{N}$$

Exercises

Exercises

RW: 1.1.4

(b) $2 \lfloor 0.6 \rfloor - \lfloor 1.2 \rfloor =$

$2 \lceil 0.6 \rceil - \lceil 1.2 \rceil =$

(d) $\lceil \sqrt{3} \rceil - \lfloor \sqrt{3} \rfloor =$

RW: 1.1.19

(a) Give x, y such that $\lfloor x \rfloor + \lfloor y \rfloor < \lfloor x + y \rfloor$:

20T2: Q1 (a)

(i) True or false for all $x \in \mathbb{R}$:

$$\lceil \lceil x \rceil \rceil = \lceil x \rceil$$

Exercises

Exercises

RW: 1.1.4

(b) $2 \lfloor 0.6 \rfloor - \lfloor 1.2 \rfloor = -1$

$$2 \lceil 0.6 \rceil - \lceil 1.2 \rceil = 0$$

(d) $\lceil \sqrt{3} \rceil - \lfloor \sqrt{3} \rfloor = 1$

RW: 1.1.19

- (a) Give x, y such that $\lfloor x \rfloor + \lfloor y \rfloor < \lfloor x + y \rfloor$:
 $x = y = 0.9$

20T2: Q1 (a)

- (i) True or false for all $x \in \mathbb{R}$:
 $\lceil |x| \rceil = |\lceil x \rceil|$ — false (e.g. $x = -1.5$)

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Divisibility

Definition

For $m, n \in \mathbb{Z}$, we say m **divides** n if $n = k \cdot m$ for some $k \in \mathbb{Z}$.

We denote this by $m|n$

Also stated as: ' n is **divisible** by m ', ' m is a **divisor** of n ', ' n is a **multiple** of m '

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$m \nmid n$ is the negation of $m|n$.

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In other words, $m \nmid n$ means ' m **does not divide** n '

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Take Notice

Notion of divisibility applies to all integers — positive, negative and zero.

Exercises

Exercises

True or False for all $n \in \mathbb{Z}$:

- $1|n$
- $-1|n$
- $0|n$
- $n|0$

RW: 1.2.2

- (a) $n|1$
- (b) $n|n$
- (c) $n|n^2$

Exercises

Exercises

True or False for all $n \in \mathbb{Z}$:

- $1|n$ — true
- $-1|n$ — true
- $0|n$ — false (only when $n = 0$)
- $n|0$ — true

RW: 1.2.2

- (a) $n|1$ — false (only when $n = \pm 1$)
- (b) $n|n$ — true
- (c) $n|n^2$ — true

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gcd and lcm

Definition

Let $m, n \in \mathbb{Z}$.

- The **greatest common divisor** of m and n , $\gcd(m, n)$, is the largest **positive** $d \in \mathbb{Z}$ such that $d|m$ and $d|n$.
- The **least common multiple** of m and n , $\text{lcm}(m, n)$, is the smallest **positive** $k \in \mathbb{Z}$ such that $m|k$ and $n|k$.
- Exception: $\gcd(0, 0) = \text{lcm}(0, n) = \text{lcm}(m, 0) = 0$.

Example

$$\begin{aligned}\gcd(-4, 6) &= \gcd(4, -6) = \gcd(-4, -6) = \gcd(4, 6) = 2 \\ \text{lcm}(-5, -5) &= \dots = 5\end{aligned}$$

gcd and lcm

Take Notice

$\gcd(m, n)$ and $\text{lcm}(m, n)$ are always taken as *non-negative* even if m or n is negative.

Fact

$$\gcd(m, n) \cdot \text{lcm}(m, n) = |m| \cdot |n|$$

Primes and relatively prime

Definition

- A number $n > 1$ is **prime** if it is only divisible by ± 1 and $\pm n$.
- m and n are **relatively prime** if $\gcd(m, n) = 1$

Examples

- 2, 3, 5, 7, 11, 13, 17, 19 are all the primes less than 20.
- 4 and 9 are relatively prime; 9 and 14 are relatively prime.

Exercises

Exercises

RW: 1.2.7(b) $\gcd(0, n) \stackrel{?}{=}$

RW: 1.2.12 Can two even integers be relatively prime?

RW: 1.2.9 Let m, n be positive integers.

(a) What can you say about m and n if $\text{lcm}(m, n) = m \cdot n$?

(b) What if $\text{lcm}(m, n) = n$?

Exercises

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RW: 1.2.7(b) $\gcd(0, n) \stackrel{?}{=} |n|$

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(a) What can you say about m and n if $\text{lcm}(m, n) = m \cdot n$?

They must be relatively prime since always $\text{lcm}(m, n) = \frac{mn}{\gcd(m, n)}$

(b) What if $\text{lcm}(m, n) = n$?

Exercises

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m must be a divisor of n

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Euclid's gcd Algorithm

Question. How do we compute the greatest common divisor $\gcd(m, n)$? Especially when the numbers m, n are large?

Answer. [Euclid's algorithm](#) gives a way of doing this by repeatedly replacing m and n with smaller numbers. This method is over 2000 years old!

$$\gcd(m, n) = \begin{cases} m & \text{if } m = n \\ \gcd(m - n, n) & \text{if } m > n \\ \gcd(m, n - m) & \text{if } m < n \end{cases}$$

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Example

$$\gcd(45, 27) =$$

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Example

$$\begin{aligned} \gcd(45, 27) &= \gcd(18, 27) \\ &= \gcd(18, 9) \\ &= \gcd(9, 9) \\ &= 9 \end{aligned}$$

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Example

$$\gcd(108, 8) =$$

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Example

$$\begin{aligned} \gcd(108, 8) &= \gcd(100, 8) \\ &= \gcd(92, 8) \\ &= \cdots = \gcd(8, 4) \\ &= \gcd(4, 4) \\ &= 4 \end{aligned}$$

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Fact

For $m > 0, n > 0$ the algorithm always terminates.

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For $m > 0, n > 0$ the algorithm always terminates.

Fact

For $m, n \in \mathbb{Z}$, if $m > n$ then $\gcd(m, n) = \gcd(m - n, n)$

Euclid's gcd Algorithm

Fact

For $m, n \in \mathbb{Z}$, if $m > n$ then $\gcd(m, n) = \gcd(m - n, n)$

Proof.



Euclid's gcd Algorithm

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For $m, n \in \mathbb{Z}$, if $m > n$ then $\gcd(m, n) = \gcd(m - n, n)$

Proof.

We first show that for all $d \in \mathbb{Z}$, $(d|m \text{ and } d|n)$ if, and only if, $(d|m - n \text{ and } d|n)$:



Euclid's gcd Algorithm

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For $m, n \in \mathbb{Z}$, if $m > n$ then $\gcd(m, n) = \gcd(m - n, n)$

Proof.

We first show that for all $d \in \mathbb{Z}$, $(d|m \text{ and } d|n)$ if, and only if, $(d|m - n \text{ and } d|n)$:

“ \Rightarrow ”: if $d|m$ and $d|n$ then $m = a \cdot d$ and $n = b \cdot d$, for some $a, b \in \mathbb{Z}$,
so $m - n = (a - b) \cdot d$,
hence $d|m - n$



Euclid's gcd Algorithm

Fact

For $m, n \in \mathbb{Z}$, if $m > n$ then $\gcd(m, n) = \gcd(m - n, n)$

Proof.

We first show that for all $d \in \mathbb{Z}$, ($d|m$ and $d|n$) if, and only if, ($d|m - n$ and $d|n$):

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so $m - n = (a - b) \cdot d$,
hence $d|m - n$

“ \Leftarrow ”: if $d|m - n$ and $d|n$ then $m - n = a \cdot d$ and $n = b \cdot d$, for some $a, b \in \mathbb{Z}$,
so $m = (m - n) + n = (a + b) \cdot d$,
hence $d|m$



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For $m, n \in \mathbb{Z}$, if $m > n$ then $\gcd(m, n) = \gcd(m - n, n)$

Proof.

We first show that for all $d \in \mathbb{Z}$, $(d|m \text{ and } d|n)$ if, and only if, $(d|m - n \text{ and } d|n)$:

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“ \Leftarrow ”: if $d|m - n$ and $d|n$ then $m - n = a \cdot d$ and $n = b \cdot d$, for some $a, b \in \mathbb{Z}$,

so $m = (m - n) + n = (a + b) \cdot d$,
hence $d|m$

Therefore, any common divisor of m and n is a common divisor of $m - n$ and n , and vice versa.



Euclid's gcd Algorithm

Fact

For $m, n \in \mathbb{Z}$, if $m > n$ then $\gcd(m, n) = \gcd(m - n, n)$

Proof.

We first show that for all $d \in \mathbb{Z}$, $(d|m \text{ and } d|n)$ if, and only if, $(d|m - n \text{ and } d|n)$:

“ \Rightarrow ”: if $d|m$ and $d|n$ then $m = a \cdot d$ and $n = b \cdot d$, for some $a, b \in \mathbb{Z}$,
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“ \Leftarrow ”: if $d|m - n$ and $d|n$ then $m - n = a \cdot d$ and $n = b \cdot d$, for some $a, b \in \mathbb{Z}$,
so $m = (m - n) + n = (a + b) \cdot d$,
hence $d|m$

Therefore, any common divisor of m and n is a common divisor of $m - n$ and n , and vice versa.

Therefore, the greatest common divisor of m and n is the greatest common divisor of $m - n$ and n . □

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Euclid's division lemma

Fact

For $m \in \mathbb{Z}$, $n \in \mathbb{Z}_{>0}$ there exists $q, r \in \mathbb{Z}$ with $0 \leq r < n$ such that

$$m = q \cdot n + r$$

Euclid's division lemma

Fact

For $m \in \mathbb{Z}$, $n \in \mathbb{Z}_{>0}$ there exists $q, r \in \mathbb{Z}$ with $0 \leq r < n$ such that

$$m = q \cdot n + r$$

Observe:

- $q = \lfloor \frac{m}{n} \rfloor$

Euclid's division lemma

Fact

For $m \in \mathbb{Z}$, $n \in \mathbb{Z}_{>0}$ there exists $q, r \in \mathbb{Z}$ with $0 \leq r < n$ such that

$$m = q \cdot n + r$$

Observe:

- $q = \lfloor \frac{m}{n} \rfloor$
- $r = m - q \cdot n$

mod and div

Definition

Let $m, p \in \mathbb{Z}$, $n \in \mathbb{Z}_{>0}$.

- $m \text{ div } n = \lfloor \frac{m}{n} \rfloor$
- $m \% n = m - (m \text{ div } n) \cdot n$
- $m =_{(n)} p$ if $n | (m - p)$

mod and div

Definition

Let $m, p \in \mathbb{Z}$, $n \in \mathbb{Z}_{>0}$.

- $m \text{ div } n = \lfloor \frac{m}{n} \rfloor$
- $m \% n = m - (m \text{ div } n) \cdot n$
- $m =_{(n)} p$ if $n \mid (m - p)$

Important!

$m =_{(n)} p$ is **not standard**. More commonly written as

$$m = p \pmod{n}$$

Fact

- $0 \leq (m \% n) < n.$

Fact

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- $m =_{(n)} p$ if, and only if, $(m \% n) = (p \% n)$.
- $m =_{(n)} (m \% n)$
- If $m =_{(n)} m'$ and $p =_{(n)} p'$ then:
 - $m + p =_{(n)} m' + p'$ and
 - $m \cdot p =_{(n)} m' \cdot p'$.

Exercises

Exercises

- $42 \text{ div } 9 \stackrel{?}{=}$
- $42 \% 9 \stackrel{?}{=}$
- $(-42) \text{ div } 9 \stackrel{?}{=}$
- $(-42) \% 9 \stackrel{?}{=}$
- *True or False:*
 $(a + b) \% n = (a \% n) + (b \% n)?$

Exercises

Exercises

- $42 \text{ div } 9 \stackrel{?}{=}$ 4
- $42 \% 9 \stackrel{?}{=}$
- $(-42) \text{ div } 9 \stackrel{?}{=}$
- $(-42) \% 9 \stackrel{?}{=}$
- *True or False:*
 $(a + b) \% n = (a \% n) + (b \% n)?$

Exercises

Exercises

- $42 \text{ div } 9 \stackrel{?}{=} 4$

- $42 \% 9 \stackrel{?}{=} 6$

- $(-42) \text{ div } 9 \stackrel{?}{=}$

- $(-42) \% 9 \stackrel{?}{=}$

- *True or False:*

$$(a + b) \% n = (a \% n) + (b \% n)?$$

Exercises

Exercises

- $42 \text{ div } 9 \stackrel{?}{=} 4$
- $42 \% 9 \stackrel{?}{=} 6$
- $(-42) \text{ div } 9 \stackrel{?}{=} -5$
- $(-42) \% 9 \stackrel{?}{=}$
- *True or False:*
 $(a + b) \% n = (a \% n) + (b \% n)?$

Exercises

Exercises

- $42 \operatorname{div} 9 \stackrel{?}{=} 4$
- $42 \% 9 \stackrel{?}{=} 6$
- $(-42) \operatorname{div} 9 \stackrel{?}{=} -5$
- $(-42) \% 9 \stackrel{?}{=} 3$
- *True or False:*
 $(a + b) \% n = (a \% n) + (b \% n)?$

Exercises

Exercises

- $42 \operatorname{div} 9 \stackrel{?}{=} 4$
- $42 \% 9 \stackrel{?}{=} 6$
- $(-42) \operatorname{div} 9 \stackrel{?}{=} -5$
- $(-42) \% 9 \stackrel{?}{=} 3$
- *True or False:*
 $(a + b) \% n = (a \% n) + (b \% n)?$
False (take $a = b = 1, n = 2$)

Exercises

Exercises

- $10^3 \% 7 \stackrel{?}{=}$
- $10^6 \% 7 \stackrel{?}{=}$
- $10^{2021} \% 7 \stackrel{?}{=}$
- What is the last digit of 7^{2023} ?

Exercises

Exercises

- $10^3 \% 7 \stackrel{?}{=}$ 6
- $10^6 \% 7 \stackrel{?}{=}$
- $10^{2021} \% 7 \stackrel{?}{=}$
- What is the last digit of 7^{2023} ?

Exercises

Exercises

- $10^3 \% 7 \stackrel{?}{=}$ 6
- $10^6 \% 7 \stackrel{?}{=}$ 1
- $10^{2021} \% 7 \stackrel{?}{=}$
- What is the last digit of 7^{2023} ?

Exercises

Exercises

- $10^3 \% 7 \stackrel{?}{=}$ 6
- $10^6 \% 7 \stackrel{?}{=}$ 1
- $10^{2021} \% 7 \stackrel{?}{=}$ 5
- What is the last digit of 7^{2023} ?

Exercises

Exercises

- $10^3 \% 7 \stackrel{?}{=}$ 6
- $10^6 \% 7 \stackrel{?}{=}$ 1
- $10^{2021} \% 7 \stackrel{?}{=}$ 5
- What is the last digit of 7^{2023} ? 3

Exercises

Exercises

RW: 3.5.20

- (a) Show that the 4 digit number $n = abcd$ is divisible by 2 if and only if the last digit d is divisible by 2.
- (b) Show that the 4 digit number $n = abcd$ is divisible by 5 if and only if the last digit d is divisible by 5.

RW: 3.5.19

- (a) Show that the 4 digit number $n = abcd$ is divisible by 9 if and only if the digit sum $a + b + c + d$ is divisible by 9.

Outline

Course introduction

- Who are we?
- Why are we here?
- How will you be assessed?

Number Theory

- Number Theory in Computer Science
- Numbers and Numerical Operations
- Divisibility
- Greatest Common Divisor and Least Common Multiple
- Euclidean Algorithm
- Modular Arithmetic
- Euclidean Algorithm (again)

Faster Euclidean gcd Algorithm

$$\text{gcd}(m, n) = \begin{cases} m & \text{if } m = n \text{ or } n = 0 \\ n & \text{if } m = 0 \\ \text{gcd}(m \% n, n) & \text{if } m > n > 0 \\ \text{gcd}(m, n \% m) & \text{if } 0 < m < n \end{cases}$$

Fact

For $m, n \in \mathbb{Z}$, if $m > n$ then $\text{gcd}(m, n) = \text{gcd}(m \% n, n)$

Proof.

Let $k = m \text{ div } n$. Then $m \% n = m - k \cdot n$.

Faster Euclidean gcd Algorithm

Example

$$\text{gcd}(108, 8) =$$

Faster Euclidean gcd Algorithm

Example

$$\gcd(108, 8) = \gcd(4, 8)$$

Faster Euclidean gcd Algorithm

Example

$$\begin{aligned}\gcd(108, 8) &= \gcd(4, 8) \\ &= \gcd(4, 0)\end{aligned}$$

Faster Euclidean gcd Algorithm

Example

$$\begin{aligned}\gcd(108, 8) &= \gcd(4, 8) \\ &= \gcd(4, 0) \\ &= 4\end{aligned}$$