$$\lim_{n \to \infty} \frac{f(n)}{g(n)} = c : f(n) \in \Theta(g(n))$$

# Big-Oh, Big-Omega, Big-Theta, Little-Oh, and Little-Omega

## 1) Show that $n \log n \in O(n^2)$

 $n \log n \le n^2$  for all  $n \ge 1$ 

$$f(n) \in O(n^2) \text{ with } c, n_0 = 1$$

 $\lim_{n\to\infty}\frac{n\log n}{n^2}=\lim_{n\to\infty}\frac{\log n}{n}=\frac{\infty}{\infty}.$  Use L'Hospital's rule: In the case  $\lim_{n\to\infty}\frac{f(n)}{g(n)}=\frac{\infty}{\infty}$ , we have:  $\lim_{n\to\infty}\frac{f(n)}{g(n)}=\lim_{n\to\infty}\frac{f'(n)}{g'(n)}$ 

So, with the derivatives: 
$$\lim_{n \to \infty} \frac{\left(\frac{1}{n \ln 2}\right)}{(1)} = \lim_{n \to \infty} \frac{1}{n \ln 2} = 0$$

So,  $n \log n \in o(n^2)$ 

This means  $n \log n \notin \Omega(n^2)$ , and so  $n \log n \notin \Theta(n^2)$ 

### **2)** Show that $10^{-10}n^2 \in \Theta(n^2)$

Let  $c = 10^{-10}$  and  $n_0 = 0$ .  $T(n) \le cn^2$  for all  $n \ge n_0$ .  $\therefore 10^{-10}n^2 \in O(n^2)$  $T(n) \ge cn^2$  for all  $n \ge n_0$ .  $\therefore 10^{-10}n^2 \in \Omega(n^2)$ 

Using limits:  

$$\lim_{n \to \infty} \frac{10^{-10} n^2}{n^2} = \lim_{n \to \infty} \frac{10^{-10}}{1} = 10^{-10}$$

$$\therefore 10^{-10} n^2 \in \Theta(n^2)$$

Since  $10^{-10}n^2 \in O(n^2)$  and  $10^{-10}n^2 \in \Omega(n^2)$ ,  $10^{-10}n^2 \in O(n^2)$ 

#### 3) Show that $n \log n \in \Omega(n)$

 $n \log n \ge n$  for all  $n \ge 2$ 

$$\therefore n \log n \in \varOmega(n) \text{ with } c = 1, n_0 = 2$$

$$\lim_{n \to \infty} \frac{n \log n}{n} = \lim_{n \to \infty} \frac{\log n}{1} = \infty$$

 $n \log n \in \omega(n^2)$ , so  $n \log n \notin O(n^2)$ , and so  $n \log n \notin O(n^2)$ 

### 4) If d(n) is $\Omega(f(n))$ and e(n) is $\Omega(f(n))$ then is it always true that d(n) + e(n) is $\Omega(f(n))$ ? If so, prove it. If not, provide a counterexample.

 $d(n) \in \Omega(f(n))$  means there exists positive constants c and  $n_0$  such that:

 $d(n) \ge cf(n)$  for all  $n \ge n_0$ 

Similarly,  $e(n) \in \Omega(f(n))$  means there exists positive constants c' and  $n'_0$  such that:  $e(n) \ge c' f(n)$  for all  $n \ge n'_0$ 

This means that  $d(n) + e(n) \ge cf(n) + c'f(n)$  for all  $n \ge \max\{n_0, n'_0\}$ = (c + c')(f(n)) so our new constant is c + c'

Therefore, we have satisfied the definition for Big-Omega.

## **5)** Show that $n^3 + 4 \in o(n^4)$

$$\lim_{n \to \infty} \frac{T(n)}{f(n)} = \lim_{n \to \infty} \frac{n^3 + 4}{n^4} = \lim_{n \to \infty} \frac{n^3}{n^4} + \lim_{n \to \infty} \frac{4}{n^4} = 0 + 0 = 0$$

$$\therefore n^3 + 4 \in o(n^4)$$

## 6) Show that $n^3 + 4 \in \omega(n^2)$

$$\lim_{n \to \infty} \frac{T(n)}{f(n)} = \lim_{n \to \infty} \frac{n^3 + 4}{n^2} = \lim_{n \to \infty} \frac{n^3}{n^2} + \lim_{n \to \infty} \frac{4}{n^2} = \lim_{n \to \infty} n + \lim_{n \to \infty} \frac{4}{n^2} = \infty + 0 = \infty$$

$$\therefore n^3 + 4 \in \omega(n^2)$$

### 7) Find a tight bound on $T(n) = \log n!$

1) Big-Oh:  $T(n) = \log(n * (n-1) * (n-2) * ... * 2 * 1)$  $= \log n + \log n - 1 + \log n - 2 + \dots + \log 2 + \log 1$  $\leq \log n + \log n + \log n + \dots + \log n + \log n$ 

 $T(n) \in O(n \log n)$  with c, n = 1

2) Big-Omega:  $T(n) = \log n + \log n - 1 + \log n - 2 + \dots + \log 2 + \log 1$  $\geq \log \frac{n}{2} + \log \frac{n}{2} + \log \frac{n}{2} + \dots + 0 + 0$  $= n/2(\log n/2)$  $= n/2(\log n - \log 2)$ 

> $= n/2(\log n - 1)$  $\geq n/4 \log n$  for all  $n \geq 4$

 $T(n) \in O(n \log n)$  with c = 1/4 and  $n_0 = 4$ 

Since  $T(n) \in O(n \log n)$  and  $T(n) \in O(n \log n)$ ,  $T(n) \in O(n \log n)$