



## Counting Operations

On this worksheet, we will only count 'basic units': assignments statements (A) and comparisons (C).

### 1) Count the number of basic operations for the worst-case:

```

a ← 3 * n           1A
count ← 0           1A
while a > n do      2n(1C) + 1C
    a ← a - 1       2n(1A)
    count ← count + 1  2n(1A)
end

```

$T(n) = 2 + 2n(3) + 1 = 6n + 3$

### 2) Count the number of basic operations for the worst-case:

```

s ← 0               1A
for k ← 1 to n do   1A + n(1C + 1A) + 1C
    s ← s + (k * k)  n(1A)
end

```

$T(n) = 2 + n(3) + 1 = 3n + 3$

### 3) Count the number of basic operations for the worst-case:

Algorithm find( $A, n, key$ ):

**Input:** An array  $A$  storing  $n \geq 1$  integers; an integer  $key$  to search for

**Output:** The index of  $key$  in  $A$

```

k ← 0               1A
while k < n do      n(1C + 1C + 1A) + 1C
    if A[k] = key do
        return k    T(n) = 1 + n(3) + 1 + 1 = 3n + 3
    end
    k ← k + 1
end
return "not found"  1A

```

### 4) Count the number of basic operations for the worst-case:

```

for k ← 0 to n - 1 do  1A + n(1C + 1A) + 1C
    for j ← 0 to n - 1 do  n(1A + n(1C + 1A + 1C + 3A) + 1C)
        if A[k] ≤ A[j] then
            swap(A, j, k)
        end
    end
end

```

$T(n) = 2 + 2n + 2n + 6n^2 = 6n^2 + 4n + 2$

### 5) Count the number of basic operations for the worst-case:

```

s ← 0               1A
for i ← 1 to n do   1A + n(1C + 1A) + 1C
    for j ← 1 to i do  n(1A + ??(1C + 1A + 2A) + 1C) we can re-write as:
        s ← s + i
        s ← s + i
    end
end

```

$T(n) = 2 + n(4) + \frac{n(n+1)}{2}(4) + 1$

$T(n) = 3 + 4n + 2n^2 + 2n = 2n^2 + 6n + 3$

How many times are the 4 operations in X executed?

when  $i$  is 1: once, when  $i$  is 2: twice, ...

$1 + 2 + 3 + 4 + 5 + \dots + n-2 + n-1 + n$  times.

If we add  $(1+n)$ ,  $(2+n-1)$ ,  $(3+n-2)$  we have  $n/2$  pairs of  $(n+1)$ .

So the 4 operations are executed  $\frac{n(n+1)}{2}$  times