

# Graduate AI

## Lecture 19:

## Game Theory I

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# NORMAL-FORM GAME

- A game in normal form consists of:
  - Set of players  $N = \{1, \dots, n\}$
  - Strategy set  $S$
  - For each  $i \in N$ , utility function  $u_i: S^n \rightarrow \mathbb{R}$ : if each  $j \in N$  plays the strategy  $s_j \in S$ , the utility of player  $i$  is  $u_i(s_1, \dots, s_n)$



# THE PRISONER'S DILEMMA

- Two men are charged with a crime
- They are told that:
  - If one rats out and the other does not, the rat will be freed, other jailed for nine years
  - If both rat out, both will be jailed for six years
- They also know that if neither rats out, both will be jailed for one year



# THE PRISONER'S DILEMMA

	Cooperate	Defect
Cooperate	-1,-1	-9,0
Defect	0,-9	-6,-6

What would you do?



# ON TV



<http://youtu.be/S0qjK3TWZE8>

# THE PROFESSOR'S DILEMMA

		Class	
		Listen	Sleep
Professor	Make effort	$10^6, 10^6$	-10, 0
	Slack off	0, -10	0, 0

Dominant strategies?



# NASH EQUILIBRIUM

- Each player's strategy is a **best response** to strategies of others
- Formally, a **Nash equilibrium** is a vector of strategies  $s = (s_1 \dots, s_n) \in S^n$  such that for all  $i \in N, s'_i \in S$ ,  
 $u_i(s) \geq u_i(s_1, \dots, s_{i-1}, s'_i, s_{i+1}, \dots, s_n)$



# THE PROFESSOR'S DILEMMA

		Class	
		Listen	Sleep
Professor	Make effort	$10^6, 10^6$	-10, 0
	Slack off	0, -10	0, 0

Nash equilibria?



# ROCK-PAPER-SCISSORS

	R	P	S
R	0,0	-1,1	1,-1
P	1,-1	0,0	-1,1
S	-1,1	1,-1	0,0

Nash equilibria?

# MIXED STRATEGIES

- A **mixed strategy** is a probability distribution over (pure) strategies
- The mixed strategy of player  $i \in N$  is  $x_i$ , where

$$x_i(s_i) = \Pr[i \text{ plays } s_i]$$

- The utility of player  $i \in N$  is

$$u_i(x_1, \dots, x_n) = \sum_{(s_1, \dots, s_n) \in S^n} u_i(s_1, \dots, s_n) \cdot \prod_{j=1}^n x_j(s_j)$$



# EXERCISE: MIXED NE

- **Exercise:** player 1 plays  $\left(\frac{1}{2}, \frac{1}{2}, 0\right)$ , player 2 plays  $\left(0, \frac{1}{2}, \frac{1}{2}\right)$ . What is  $u_1$ ?
- **Exercise:** Both players play  $\left(\frac{1}{3}, \frac{1}{3}, \frac{1}{3}\right)$ . What is  $u_1$ ?

	R	P	S
R	0,0	-1,1	1,-1
P	1,-1	0,0	-1,1
S	-1,1	1,-1	0,0

# EXERCISE: MIXED NE

- Poll 1: Which is a NE?

1.  $\left(\left(\frac{1}{2}, \frac{1}{2}, 0\right), \left(\frac{1}{2}, \frac{1}{2}, 0\right)\right)$

2.  $\left(\left(\frac{1}{2}, \frac{1}{2}, 0\right), \left(\frac{1}{2}, 0, \frac{1}{2}\right)\right)$

3.  $\left(\left(\frac{1}{3}, \frac{1}{3}, \frac{1}{3}\right), \left(\frac{1}{3}, \frac{1}{3}, \frac{1}{3}\right)\right)$

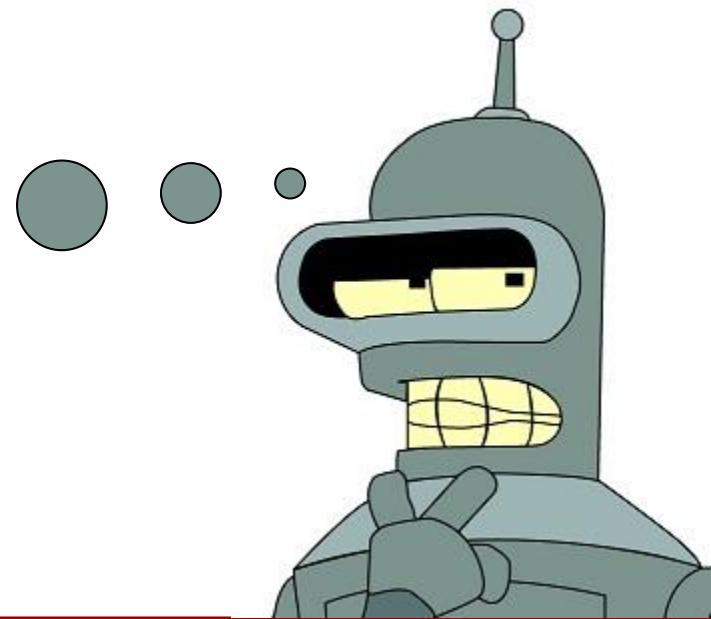
4.  $\left(\left(\frac{1}{3}, \frac{2}{3}, 0\right), \left(\frac{2}{3}, 0, \frac{1}{3}\right)\right)$

	R	P	S
R	0,0	-1,1	1,-1
P	1,-1	0,0	-1,1
S	-1,1	1,-1	0,0

# NASH'S THEOREM

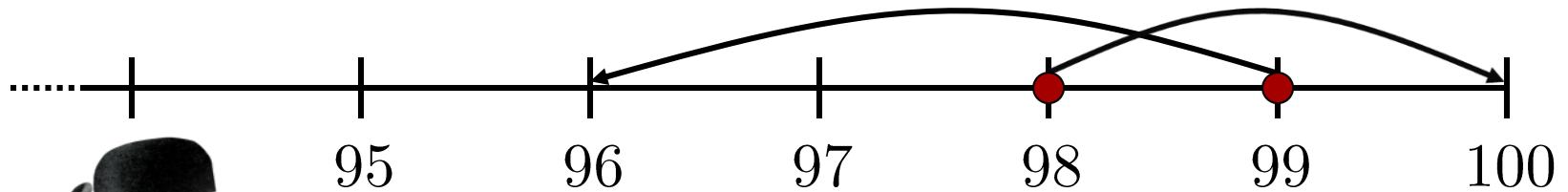
- Theorem [Nash, 1950]: In any (finite) game there exists at least one (possibly mixed) Nash equilibrium

What about  
computing a Nash  
equilibrium?



# DOES NE MAKE SENSE?

- Two players, strategies are  $\{2, \dots, 100\}$
- If both choose the same number, that is what they get
- If one chooses  $s$ , the other  $t$ , and  $s < t$ , the former player gets  $s + 2$ , and the latter gets  $s - 2$
- Poll 2: What would you choose?



# CORRELATED EQUILIBRIUM

- Let  $N = \{1,2\}$  for simplicity
- A mediator chooses a pair of strategies  $(s_1, s_2)$  according to a distribution  $p$  over  $S^2$
- Reveals  $s_1$  to player 1 and  $s_2$  to player 2
- When player 1 gets  $s_1 \in S$ , he knows that the distribution over strategies of 2 is

$$\Pr[s_2|s_1] = \frac{\Pr[s_1 \wedge s_2]}{\Pr[s_1]} = \frac{p(s_1, s_2)}{\sum_{s'_2 \in S} p(s_1, s'_2)}$$

# CORRELATED EQUILIBRIUM

- Player 1 is best responding if for all  $s'_1 \in S$ 
$$\sum_{s_2 \in S} \Pr[s_2 | s_1] u_1(s_1, s_2) \geq \sum_{s_2 \in S} \Pr[s_2 | s_1] u_1(s'_1, s_2)$$
- Equivalently,
$$\sum_{s_2 \in S} p(s_1, s_2) u_1(s_1, s_2) \geq \sum_{s_2 \in S} p(s_1, s_2) u_1(s'_1, s_2)$$
- $p$  is a **correlated equilibrium (CE)** if both players are best responding



# GAME OF CHICKEN



<http://youtu.be/u7hZ9jKrwvo>

# GAME OF CHICKEN

- Social welfare is the sum of utilities
- Pure NE: (C,D) and (D,C), social welfare = 5
- Mixed NE: both  $(1/2, 1/2)$ , social welfare = 4
- Optimal social welfare = 6

	Dare	Chicken
Dare	0,0	4,1
Chicken	1,4	3,3



# GAME OF CHICKEN

- Correlated equilibrium:

- $(D,D): 0$

- $(D,C): \frac{1}{3}$

- $(C,D): \frac{1}{3}$

- $(C,C): \frac{1}{3}$

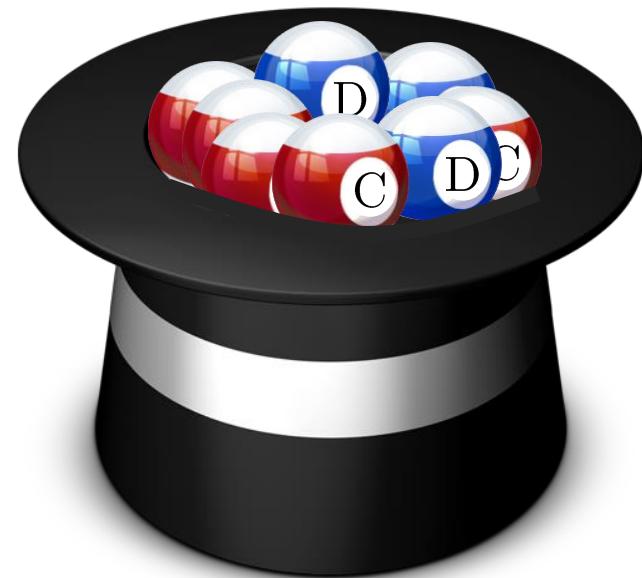
	Dare	Chicken
Dare	0,0	4,1
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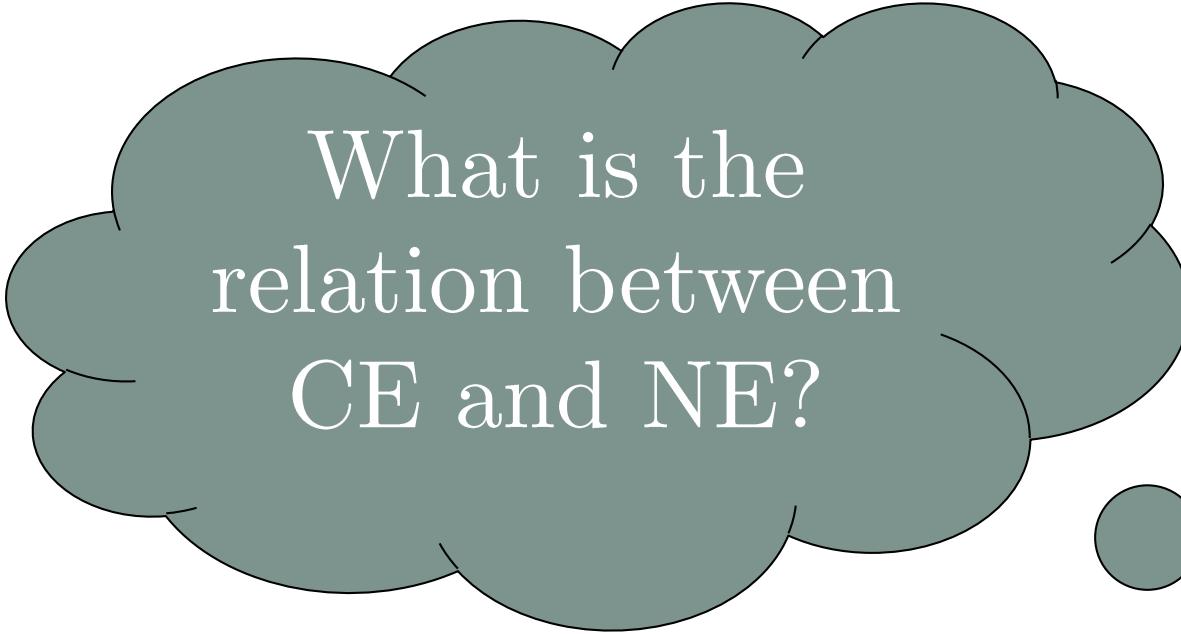
- Social welfare of CE =  $\frac{16}{3}$



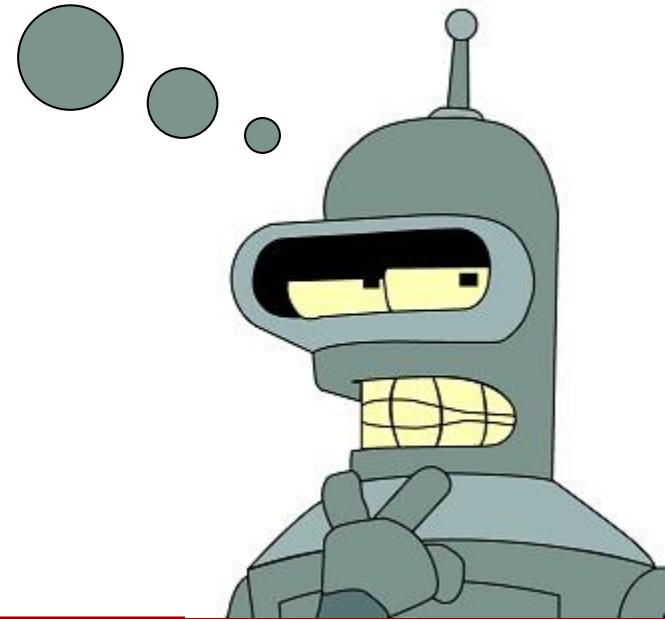
# IMPLEMENTATION OF CE

- Instead of a mediator, use a hat!
- Balls in hat are labeled with “chicken” or “dare”, each blindfolded player takes a ball
- **Poll 3:** Which balls implement the distribution of slide 19?
  1. 1 chicken, 1 dare
  2. 2 chicken, 1 dare
  3. 2 chicken, 2 dare
  4. 3 chicken, 2 dare





What is the  
relation between  
CE and NE?



# CE AS LP

- Can compute CE via linear programming in polynomial time!

find  $\forall s_1, s_2 \in S, p(s_1, s_2)$

s.t.  $\forall s_1, s'_1 \in S, \sum_{s_2 \in S} p(s_1, s_2)u_1(s_1, s_2) \geq \sum_{s_2 \in S} p(s_1, s_2)u_1(s'_1, s_2)$

$\forall s_2, s'_2 \in S, \sum_{s_1 \in S} p(s_1, s_2)u_2(s_1, s_2) \geq \sum_{s_1 \in S} p(s_1, s_2)u_2(s_1, s'_2)$

$$\sum_{s_1, s_2 \in S} p(s_1, s_2) = 1$$

$\forall s_1, s_2 \in S, p(s_1, s_2) \in [0, 1]$



# SUMMARY

- Terminology:
  - Normal-form game
  - Nash equilibrium
  - Mixed strategies
  - Correlated equilibrium
- Algorithms:
  - LP for CE

