

EECS 280

Programming and Introductory Data Structures

Recursion, Function Pointers and Testing

Review: recursion

- Tail recursion and iteration are the same
 - We proved this last time with a *constructive* proof
 - Should you ever use tail recursion?

Pro: Compiler handles update dependencies

Pro: Sometimes tail recursion (base/rec. cases) is more intuitive

Con: Sometimes iteration is more intuitive

Con: Can be tricky to ensure tail call optimization

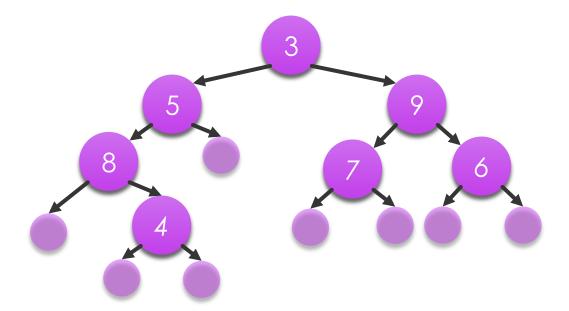
- "General" recursion is expensive
 - Stack frames sit and store pending work
 - Why would you ever use this??
 - Sometimes you need it

Review: recursive data structures

• List (last time)

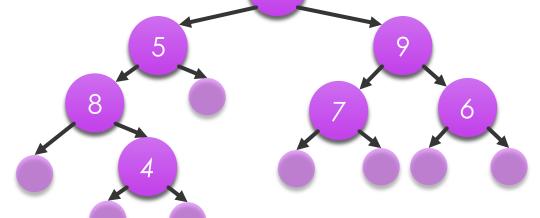


• Tree (today)



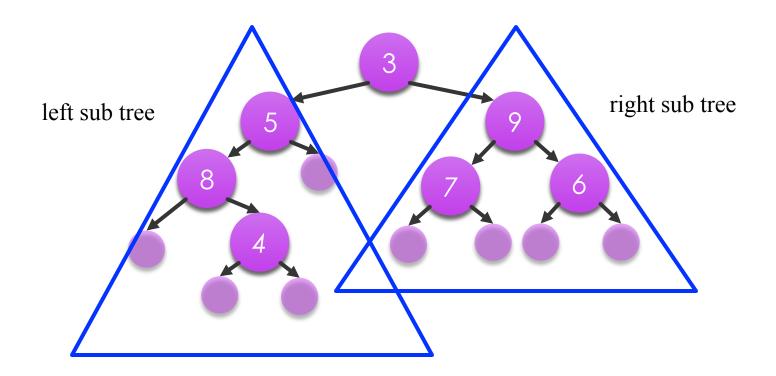
Binary tree definition

• Tree

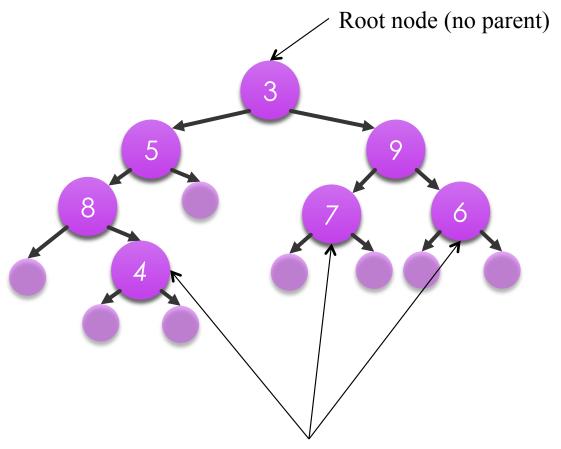


- A binary tree is:
 - The empty tree, or
 - An integer element, plus two children, called the left subtree and the right subtree, both binary trees
- Recursively defined data structure

Tree terminology



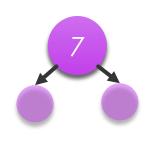
Tree terminology

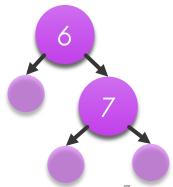


Leaf node: both children are empty tree

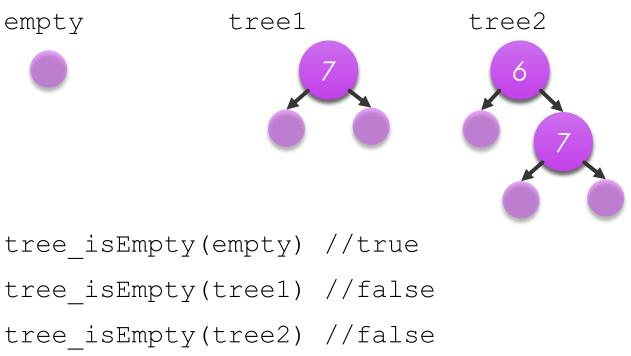
Tree type provided in project 2

- In addition to list_t, the type tree_t is provided in project 2
- tree t empty = tree make();

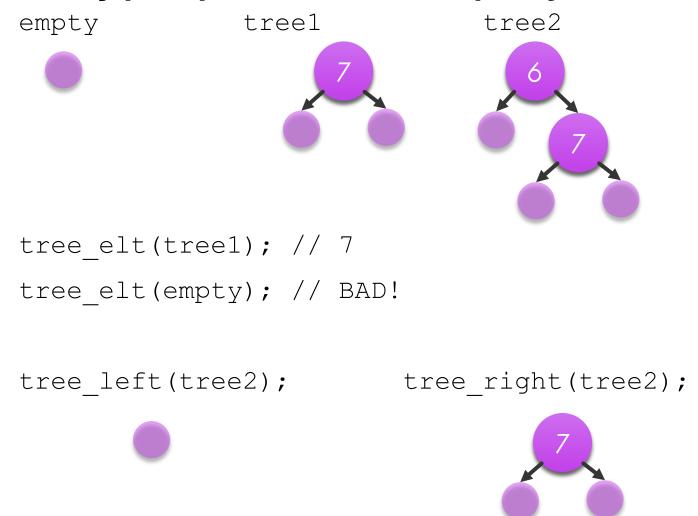




Tree type provided in project 2



Tree type provided in project 2



Tree properties

- Trees are *immutable*
 - There is no tree setElement
- Think of lists as values
 - To "modify" a tree, you just have to "compute" a new one with desired changes
 - This means putting the three parts of the tree back together using tree_make()...
 - element, left subtree, right subtree

Exercise: tree_depth()

```
tree2
     empty
                       tree1
tree depth (empty) //0
                  tree depth(tree1) //1
                                       tree depth(tree2) //2
//EFFECTS: Returns the depth of tree
int tree depth (tree t tree) {
    //use general recursion
   //you can use the max(a, b) function from
    //last time
```

Hint

```
• Base case
if (tree isEmpty(tree)) {
  return 0;
• Recursive case
depth = 1 + max(L, R)
                      Left
                                         Right
```

```
L = depth(tree left(tree)) R = depth(tree right(tree))
```

Solution: tree_depth()

```
int max(int x, int y) { return x > y ? x : y; }
//EFFECTS: Returns the depth of tree.
int tree depth (tree t tree) {
  if (tree isEmpty(tree)) return 0; // BASE CASE
  // RECURSIVE CASE
  int L = tree depth(tree left(tree));
  int R = tree depth(tree right(tree));
  return 1 + \max(L, R);
```

Exercise: tree_depth() tail

```
//EFFECTS: Returns the depth of tree
int tree_depth (tree_t tree) {
    // Use tail recursion OR iteration
}
```

Solution: tree_depth() tail

```
//EFFECTS: Returns the depth of tree
int tree_depth (tree_t tree) {
    // Use tail recursion OR iteration
}
```

• You can't do it!!!!

• On to function pointers

Recall tail recursive list_max()

```
int max(int x, int y) { return x > y ? x : y; }
int list max h(list t list, int so far) {
  if (list isEmpty(list)) return so far;
  return list max h(
    list rest(list),
   max(list first(list), so_far)
  );
int list max (list t list) {
  assert(!list isEmpty(list);
  return list max h(list, list first(list));
    Exercise: how would you change this
    code to implement list min()?
```

Solution: list_min()

```
int max(int x, int y) { return x > y ? x : y; }
int min(int x, int y) { return x < y ? x : y; }
int list min h(list t list, int so far) {
  if (list isEmpty(list)) return so_far;
  return list min h(
    list rest(list),
max min(list first(list), so far)
  );
int list min (list t list) {
  assert(!list isEmpty(list);
  return list min h(list, list first(list));
```

- We've copy pasted our code for list_max() to write list min()
- Let's make one new function that can do the work of both

```
int list extreme(list t list, int so far);
```

• Now we need to pass the behavior of min and max as an input to list extreme

```
int max(int x, int y) { return x > y ? x : y; }
int min(int x, int y) { return x < y ? x : y; }
int list_extreme(list_t list, int so_far, /* min or max */) {
    //...
}</pre>
```

• How can we make a function behave like a variable so we can make it an input to list extreme()?

- Function pointers are variables
- Function pointers are like a second name for a function
- Function pointers let us decide between two (or more) functions at run time

```
int max(int x, int y) { return x > y ? x : y; }
int min(int x, int y) { return x < y ? x : y; }
int main() {
  int (*comp) (int, int);
  comp = max;
  cout << comp(-42, 42); //42
}</pre>
```

```
int max(int x, int y) { return x > y ? x : y; }
int min(int x, int y) { return x < y ? x : y; }
int main() {
  int (*comp)(int,int);
  comp = max;
  cout << comp(-42, 42); //42
}</pre>
```

• What does this mean?

```
int max(int x, int y) { return x > y ? x : y; }
int min(int x, int y) { return x < y ? x : y; }
int main() {
  int (*comp) (int, int);
  comp = max;
  cout << comp(-42, 42); //42
}</pre>
```

- We can assign a value to our new variable
- The value is the name of a function
- The instructions for executing a function are stored somewhere...a function pointer actually stores the **address** where the function is stored

```
int max(int x, int y) { return x > y ? x : y; }
int min(int x, int y) { return x < y ? x : y; }
int main() {
  int (*comp)(int,int);
  comp = max;
  cout << comp(-42, 42); //42
}</pre>
```

- Then, we can use our comp variable just like a function
- It will do the same thing as max in this example

```
int max(int x, int y) { return x > y ? x : y; }
int min(int x, int y) { return x < y ? x : y; }
int list_extreme(list_t list, int so_far, int (*comp)(int,int)) {
   //...
}</pre>
```

Now we'll add an input to list_extreme using our new type

```
int max(int x, int y) { return x > y ? x : y; }
int min(int x, int y) { return x < y ? x : y; }
int list_extreme(list_t list, int so_far, int (*comp)(int,int)) {
  if (list_isEmpty(list)) return so_far;
  return list_extreme(
    list_rest(list),
    comp(list_first(list), so_far)
    );}</pre>
```

• And use comp instead of min or max

```
int max(int x, int y) { return x > y ? x : y; }
int min(int x, int y) { return x < y ? x : y; }
int list extreme (list t list, int so far, int (*comp)(int,int)) {
  if (list isEmpty(list)) return so far;
  return list extreme (
    list rest(list),
    comp(list first(list), so far)
    );
• Finally, we can write list min() and list max()
int list max(list t list) { return list extreme(list, 1, max); }
int list min(list t list) { return list extreme(list, 1, min); }
```

Why function pointers?

- Avoid code duplication
- We can use a simple function, like min or max to change the behavior of a more complicated function, like

```
list_extreme
```

Exercise

```
bool all of(list t list, bool(*predicate)(int)) {
  //EFFECTS: returns true if predicate returns true
 // for all elements in list
  if (list isEmpty(list)) //base case
    return true;
  if (!predicate(list first(list))) //check first
    return false;
  return all_of(list rest(list), fn); //rec. case
```

• Write these two functions. Use all_of() and helper functions

```
bool all_even(list_t list);
bool all_odd(list_t list);
```

Solution

```
bool is_even(int i) {
   //EFFECTS: returns true if i is even
   return (i % 2) == 0;
}

bool all_even(list_t list) {
   //EFFECTS: returns true if all elements are even
   return all_of(list, is_even);
}
```

Solution

```
bool is_odd(int i) {
   //EFFECTS: returns true if i is odd
  return (i % 2) == 1;
}

bool all_odd(list_t list) {
   //EFFECTS: returns true if all elements are odd
  return all_of(list, is_odd);
}
```

Testing

- The goals of testing:
- Ensure that code works as expected
- Meet the requirements of the spec

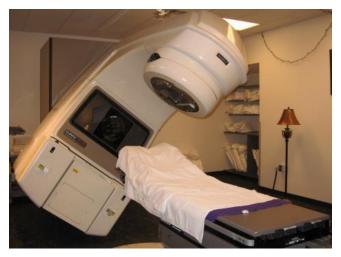
When Testing Goes Wrong

- Patriot missile bug
 - Persian Gulf war
 - Missile defense system to detect an incoming missile, then deploy a counter-missile to destroy it.
 - Needs to happen fast
 - Bug in floating point time conversion code
- System failed to intercept a missile
 - 28 killed, 98 injured



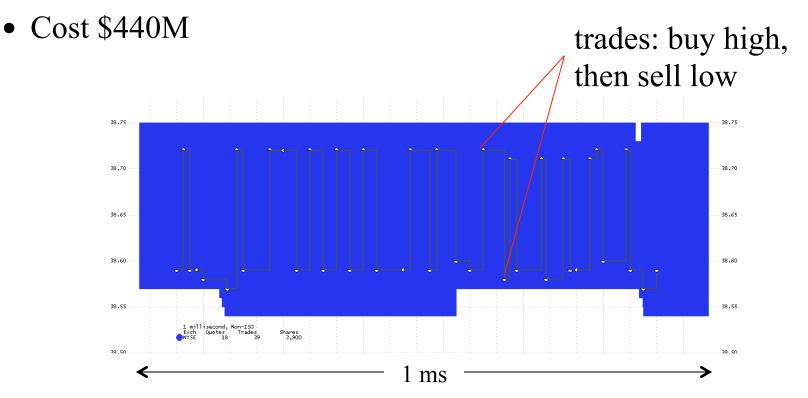
When Testing Goes Wrong

- Therac-25 Radiation Therapy Machine (1980's)
- Radiation therapy
 - Low-power direct electron beam or megavolt X-ray
- Bug in software caused megavolt X-ray instead of low-power
 - Race condition
- Patients died of radiation poisoning



When Testing Goes Wrong

- Knight Capital Group: high-speed stock trading
 - Goal: buy low, sell high
- August 1, 2012: upgraded software algorithm
 - Bug: buy high, sell low!



Motivation for testing

- Good testing yields correctly working software products
- Good testing yields good scores on projects
 - Typically, the difference between a good and bad score on a project doesn't have very much to do with your talent as a programmer.

It has much more to do with your talents as a tester!

- Testing is not the same as debugging
 - Testing: Discovering that something is broken
 - Debugging: Fixing something once you know it's broken

Psychology of a Good Tester

- Adversarial frame of mind.
- You must be convinced that the code you are testing is broken and your task is to find out where.
- You can NEVER REST and must ALWAYS BE DILIGENT, because the code is NEVER FINISHED!
- Everyone makes mistakes, and one essential nature of a mistake is that the person who made it (i.e. YOU) didn't realize it was wrong in the first place you thought it was perfect!

Types of testing

- Unit testing
 - One piece at a time (e.g., a function)
 - Find and fix bugs early! Less work.
 - Test smaller, less complex, easier to understand units.
 - You just wrote the code: easier to debug
- System testing
 - Entire project (code base)
 - Do this *after* unit testing
- Regression testing
 - Automatically run all unit and system tests after a code change

How to test

- 1. Understand the specification
 - Identify the required behaviors
- 2. Write tests
 - Code to run your other code
 - Try to test only one behavior per test
- 3. Check the results
 - Know the answers in advance

Debugging Hint for Recursion

- Add a cout statement to the top of your recursive function
 - Watch the recursion in action
 - Very helpful for trees!

```
int countdown4(int i, int n) {
  cout << "countdown4("</pre>
       << i << ","
       << n << ")" << endl;
  if (i > n) return 0;
  countdown4(i+1, n);
  cout << i << endl;
  return i;
```

```
$ ./a.out
countdown4(1,3)
countdown4(2,3)
countdown4(3,3)
countdown4(4,3)
3
2
1
```

Kinds of test cases

Imagine we are writing test cases for the chop function from project 2.¹

```
// REQUIRES n >= 0 and list has at least n elements
// EFFECTS: returns the list equal to list without its last n elements
list_t chop(list_t list, int n);
```

Don't write these	Type Prohibited	("cat", 2)	(1, 5)	(3, [1,2])	(true, 0)
	REQUIRES Prohibited	([], 1)	([42], -1)	([1,2,3,4,5], 4)	
	Simple	([1,2,3],2)		([1,2,3,4,5],1)	
	(Edge) Special	([], 0)	([42], 0)	([42], 1)	([1,2], 2)
	Nonsense	Sorry, no good examples for chop.			
	Stress	([1,2,3,4,5,6,7,8,9,,100000], 50000)			

"Guard against Murphy, not Machiavelli!

- Do: Try to write test cases to catch bugs that people would realistically make
 - ([1,2,3,4,5],5)
 - "Tricky because it chops everything. Could expose a bug."
- Don't: Try to write test cases to catch bugs introduced by a devious coder
 - ([4,4,4,4,4],3)
 - "Tricky because maybe it crashes when the list has all fours."
- Think about what makes test cases meaningfully different.
 - chop doesn't use element values, so these aren't meaningfully different.
- ([V,W,X,Y,Z],5) vs. ([A,B,C,D,E],5)

Thorough testing with "small" test cases is sufficient to find all bugs within a system.

The Small Scope Hypothesis

- Think about what makes two test cases meaningfully different for the function's behavior
- Beyond a small size, just making test cases bigger doesn't make them meaningfully different
 - ([1,2,3,4,5],3) vs. ([1,2,3,4,5,6],3)