

## **EECS 280**

Programming and Introductory Data Structures

Recursion and Iteration

#### Review: recursion

- A recursive problem is one that is defined in terms of itself.
- A recursive problem has two important features:
  - A base case(s) that stops the recursion
  - A recursive step that solves a smaller version of the same problem
- Unfortunately, general recursion requires many stack frames, which can use a lot of memory

```
int factorial(int n) {
  if (n == 0) return 1;
  return n * factorial(n-1);
}
```

#### Tail call review

• A function call is a **tail call** if it is the very last thing in its containing function

```
int factorial(int n) {
   //...
  return factorial(n-1);
}
```

- The calling function has **no pending work** to do after a tail call (in the passive flow)
- Avoids saving the stack frame!
- Tail call optimization can take a recursive algorithm from linear to constant space complexity as long as it only makes tail calls

#### Review: Tail Recursion

- The many stacks frames problem was solved with tail recursion
- Tail recursion reuses the stack frame

```
int fact helper(int n, int result) {
  // REQUIRES: n >= 0
  // EFFECTS: returns result * n!
  if (n == 0) return result;
  return fact helper(n-1, result*n);
int factorial(int num) {
  // REQUIRES: n >= 0
  // EFFECTS: returns num!
  return fact helper(num, 1);
```

#### Tail Recursion and Iteration

• Both of these implementation require only 1 stack frame

```
int factorial(int n) {
  int result = 1;
  while (n != 0) {
    result *= n;
    n -= 1;
  }
  return result;
}
```

```
int factorial(int n, int result) {
  if (n == 0) return result;
  return factorial(n-1, result*n);
}
```

## Stack frame example

• Stack frames example

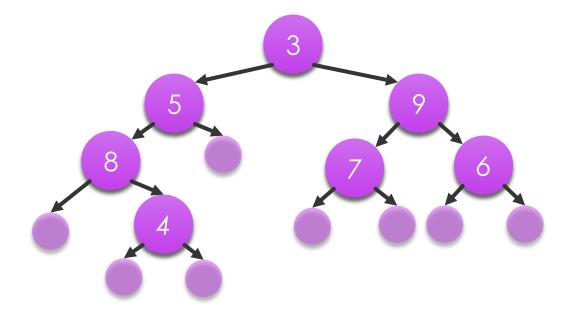
```
int factorial helper(int n, int result){
  if (n == 0) return result;
 return factorial helper(n-1, n * result);
int factorial(int n) {
 return factorial helper(n ,1);
int main() {
 factorial(3);
 return 0;
```

#### Recursive data structures

• List



• Tree



#### List definition

• List



- A list is:
  - The empty list, or
  - An integer followed by a list

## List functions provided in project 2

```
list t list make();
// returns an empty list.
list t list make(int elt, list t list);
// given the list (list) make a new list consisting
// of the new element followed by the elements of
// the original list.
void list print(list t list);
    // MODIFIES: cout
    // EFFECTS: prints list to cout.
```

## list\_make()

• list t empty = list make(); list t list1 = list make(4, empty); list t list2 = list make(1, list make (5,list make(6, list1)));

## list\_first() and list\_rest()

```
empty list1
                     list2
list first(list2); // 1
list first(empty); // BAD!
list rest(list1);
list rest(list2);
list rest(list rest(list1)); // BAD!
list first(list rest(list2)); // 5
```

# list\_isEmpty()

```
list_isEmpty(list2); // false
list_isEmpty(empty); // true
list_isEmpty(list_rest(list1)); // true
```

#### List properties

- Lists are *immutable* 
  - There is no list\_setElement
- Think of lists as values
  - To "modify" a list, you just have to "compute" a new one with desired changes
- What does this line do?

```
list2 = list_make(0, list_rest(list_2));
```

## Exercise: list\_max()

```
empty list1 list2

(4→) (1→6)→6→4→

// REQUIRES: list must not be empty
// EFFECTS: Returns the largest element in list int list_max (list_t list){
```

#### Use "general" recursion

```
max(empty); // BAD!
max(list1); // 4
max(list2); // 6
```

# Solution: list\_max()

```
static int max(int x, int y) {
  return x > y ? x : y;
// REQUIRES: list must not be empty
// EFFECTS: Returns the largest element in list
int list max (list t list) {
  if (list isEmpty(list)) {
    return -9999; //HACK!
  return max (
    list first(list),
    list max(list rest(list))
    );
```

# Solution: list\_max()

```
static int max(int x, int y) {
  return x > y ? x : y;
  REQUIRES: list must not be empty
// EFFECTS: Returns the largest element in list
int list max (list t list) {
  assert(!list isEmpty(list);
  if (list isEmpty(list rest(list)) { //hack fixed :)
    return list first(list); //single element
  return max (
    list first(list),
    list max(list rest(list))
    );
```

## Exercise: list\_max() tail

max(list2); // 6

```
empty list1
                      list2
// REQUIRES: list must not be empty
// EFFECTS: Returns the largest element in list
int list max (list t list) {
  Use tail recursion recursion
  Hint: make a helper function with
  an extra parameter if needed
max(empty); // BAD!
max(list1); // 4
```

#### Solution: list\_max() tail

```
static int max(int x, int y) { return x > y ? x : y; }
int list max h(list t list, int soFar) {
  if (list isEmpty(list)) return soFar;
  return list max h(
    list rest(list),
   max(list first(list), list max(list_rest(list))
   );
int list max (list t list) {
  assert(!list isEmpty(list);
 return list max h(list, list first(list));
```

```
int list_max_h(list_t list, int soFar){
Recursive
    int list_max (list_ist){
       return list_max list_first(list));
     int list /
                 (list_t list){
Iterative
       int soFar = list_first(list);
```

```
int list_max_h(list_t list, int soFar){
      if (list_isEmpty(list)){ // BASE CASE
        return soFar; // final result
          // RECURSIVE CASE
       list_max (list_t list){ ... }
      list_max(list_t list){
          soFar = list_first(list);
Iterative
      while (!list_isEmpty(list)){
      return soFar;
```

int list\_max\_h(list\_t list, int soFar){ ... // BASE CASE Tail Recursive Compiler handles these! // RECURSIVE CASE return list\_max\_h(list\_rest(list), max(list\_first(list), soFar)); int list\_max (list\_t list){ ... } omputation int list\_max(list\_t list){ int soFar = list\_first(list); Iterative while (!list\_isEmpty(list)){ list = list\_rest(list); soFar = max(list\_first(list), soFar); return soFar; Update ordering matters!

```
while (!list_isEmpty(list)){
   list = list_rest(list);
   soFar = max(list_first(list), soFar);
}
Ordering 1
```

```
while (!list_isEmpty(list)){
   soFar = max(list_first(list), soFar);
   list = list_rest(list);
}
Ordering 2
```

- The new value of soFar depends on list
- The new value of list does NOT depend on soFar
- Ordering 2 is safe, but ordering 1 is not

## Dependency Graphs

- Allow us to determine an ordering to safely update variables.
- An arrow from one variable to another indicates a dependency.
  - e.g. new soFar depends on current list

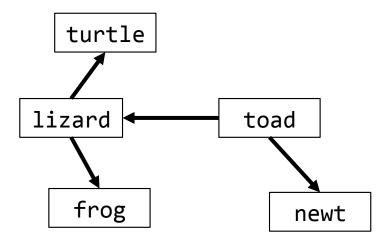
So we must update soFar before list.



• In general, we "topologically sort" the dependency graph to determine update order

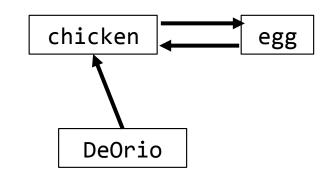
## Exercise: Dependency Graphs

• Find a suitable ordering for updating the variables with this dependency graph



# Dependency Graphs: Cycles

• What if you have a cycle in the dependency graph?



Shadow variables

```
int cTemp = chicken;
int eTemp = egg;
chicken = // use eTemp;
egg = // use cTemp;
```

#### So...where are we now?

- Tail recursion and iteration are the same
  - We just walked through a constructive proof
  - Should you ever use tail recursion?

Pro: Compiler handles update dependencies

Pro: Sometimes base/recursive cases easy

Con: Sometimes iteration is more intuitive

Con: Can be tricky to ensure TCO

- "General" recursion is expensive
  - Stack frames sit and store pending work
  - Why would you ever use this?
  - Sometimes you need it (example: project 2 tree functions)