

## 1 Cognition, computation, computers, programs

1. Cognitive Science: mind/brain is analogous to software/hardware.
  - 1.1. Various names: “the computational view of mind,” “information processing psychology,” “the computational theory of mind,” or simply “computationalism.”
  - 1.2. The analogy is usually far from strict and there are many varieties.
  - 1.3. For instance: “mind/cognition is computation” versus “mind/cognition is computable.”
2. Why we need computationalism?
  - 2.1. Compare the tasks:
    - Predicting the trajectory of celestial bodies, say the motion of the earth in the next six hours.
    - Predicting the next move of a chess player at a given state of the game.
  - 2.2. In the case of chess, physical description and physical laws are helpless. The source of helplessness is twofold: (i) The state of a chess game has infinitely many physical realizations, and whether a physical state is a game state is dependent on the whether or not some cognitive agents interpret the “scene” as a chess game or not. Therefore a function from physical description to game state is at best extremely complex. (ii) Even if we find a way from physical description to game state, the physical description of rule-based behavior would be a function of the particular realization function (the mapping from physics-to-chess states) at that particular occasion.
3. Computation and cognition share an essential property: both operate on rules and representations, yet are based on (instantiated by) physical causal systems. What is happening in a computer and mind/brain are quite similar.

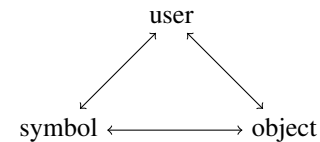
**Levels** (more on this below):

- i. physical (device)
- i. symbolic (syntactic)
- iii. intentional (semantic)

### 1.1 What is (symbolic) computation?

1. Symbols (or signs) and signification are central concepts in language, logic, and computation.

- 1.1. Take signification as a three-part relation:



- A user uses a sign to **refer** to an object.
  - A sign **denotes** an object.
  - A user has certain **intentions** about an object.
- 1.2. Example: 001011010100111000000010.
  2. A computational process is a sequence of manipulations performed over symbolic representations.
    - 2.1. Example: the process by which you flip the digits of the representation above, one at a time is a computational process.

### 1.2 Hardware/software

1. Most computers around are based on von Neumann<sup>1</sup> architecture.
2. The main components of VNA is a **central processing unit** and a **memory**. You can think of these as the homes of **processes** and **data**.<sup>2</sup>
3. Memory is simpler, so let's start with it. It consists of a sequence of slots with the informational capacity of a **byte**<sup>3</sup>. Bytes are organized into **words**. The size of a word depends on the design; 4-bytes is a typical value. Therefore memory can be thought of as a sequence of words. Sometimes words are also called **cells**. Each cell of memory has a unique **address**.

<sup>1</sup>Named after the mathematician and physicist John von Neumann.

<sup>2</sup>But beware that the clarity of this distinction is not present in every model or architecture of computation.

<sup>3</sup>A byte consists of 8 bits. A bit is a binary digit which can either be 0 or 1. A single byte can hold 2<sup>8</sup> different values, can you see why?

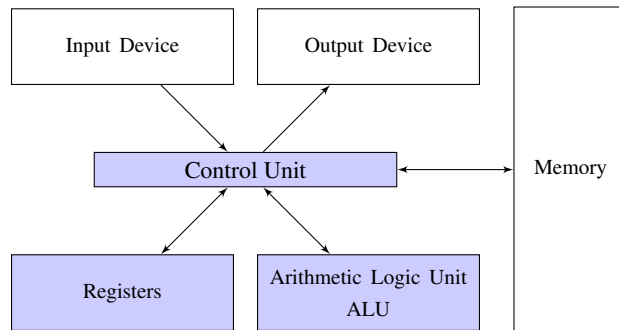


Figure 1: Von Neumann architecture

4. The two basic components, CPU and memory, communicate through three channels:
  - i. A collection of wires called **address bus**;
  - ii. Another collection of wires called **data bus**;
  - iii. A single wire called **R/W line**, the status of which signals whether CPU wants to write to or read from the memory.
5. The address bus and the R/W line are one way channels. The value is always dictated by the CPU. The data bus is a two-way channel.
6. The two basic interactions between CPU and Memory go as follows:
  - CPU sets the address bus and R/W line to W. In that case it also sets the data bus. Obviously, this amounts to dictating to write the data to the given address in memory.
  - CPU sets the address bus and R/W line to R. In this case it is the memory that sets the data bus in accord with the data located at the address provided by CPU; this is reading from memory.
7. CPU itself also has some local memory slots. These are called **registers**. Access to registers is faster than access to memory. But there is a reason to keep the size of the CPU small, therefore there are a limited number of registers, which are used to store intermediate results of computations and some frequently used information.
8. The computation unit of CPU is **arithmetic logic unit** (ALU, for short.) ALU is responsible for arithmetic and logical operations.
9. CPU feeds on **instructions**. Some typical types of instructions are:
  - i. Store the number at the register  $X$  at the memory address  $Y$ ;
  - ii. Fetch the number at the address  $X$  and store it at the register  $Y$ ;
  - iii. Add the number of address  $X$  to the number at address  $Y$ , and store the result at address  $Z$ ;
  - iv. Compute the bitwise *and* of the number at the address  $X$  and  $Y$ , and store the result at the address  $Z$ ;
  - v. If the contents of the register  $X$  and  $Y$  are not identical, go to address  $Z$ ;
  - vi. Jump to the address  $X$ ;
  - vii. and so on.
10. **Stored-program** concept
  - a. Not only data but also instructions are represented as numbers;
  - b. Programs are stored in memory; they can be read and written just like data.
11. A central aspect of a computational system is **flow of control**. Some instructions have *go to* or *jump*, which sends the control to another instruction. But other instructions lack such a mechanism for flow of control.
12. There is a special register called **program counter**, where the address of the next instruction is kept.
13. CPU operates through a sequence of cycles. Each cycle begins by fetching a binary description of what to do, an **instruction** from an address in Memory. Following this CPU understands, or technically speaking, **decodes** the instruction. The next step is to **execute** the instruction. This completes a single cycle, which is followed by reading the next instruction from Memory.<sup>4</sup>
14. Beside instructions Memory is the store for any sort of information that needs to be kept for reuse.
15. Most of our commercial computers are based on this architecture.

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<sup>4</sup>Cycles are called 'fetch-decode-execute' or 'instruction cycle'.

### 1.3 Levels of programming

1. Computers operate on numbers, in the sense that the functional architecture of the machines get their instructions as binary numbers organized into expressions of **machine code**.
2. Given a functional architecture, say VNA, a straightforward specification would be to code the memory byte-by-byte (writing **machine code**); so that CPU “knows” what to do in each possible state of the process.
3. Example of an addition instruction:  
  
000000 10001 10010 01000 00000 1000000
4. A machine code instruction is organized into **fields** with different meanings.
5. One level up the machine code is the **assembly** language. The above machine code takes the form:

```
add $s1,$t1,$t2
```

6. Assembly level lies between high-level languages and machine code.

A simple while loop in C:

```
while (save[i] == k)
    i = i + 1;
```

Assuing that i and k are stored in registers \$s3 and \$s5, and the array save starts at the memory address stored in \$s6, the above C code is **compiled** to the following assembly code (for MIPS):

```
Loop:  sll $t1,$s3,2
       add $t1,$t1,$s6
       lw  $t0,0($t1)
       bne $t0,$s5,Exit
       add $s3,$s3,1
       j   Loop
Exit:
```