# MIPS Functions

A MIPS function is called by the jal instruction, which does two things: 1) going to the address of the first instruction in the function, 2) passing the arguments in \$a0 to \$a3.

A MIPS function must be ended by the jr \$ra instruction.

MIPS functions are stored in a different part in the memory and not in the same part as the main function.

A MIPS function has explicitly declare a name along with the list of arguments to be passed to it, including the names and the types.

A function in MIPS cannot have loops.

In MIPS, the name of a function is a special data type.

The jal L1 instruction jumps to L1, and saves the address of L1 into \$ra.

Suppose \$s0, \$v0, and \$a0 are holding 60, 0, and 35, respectively. After the program executes till p9L2, what will be the value in \$v0?

- (a) 100
- (b) 70
- (c) The program will never run to p9L2.
- (d) None of the above.

In MIPS, "la \$t0, L1" is a pseudo instruction which loads the address of the instruction associated with L1 into \$t0. [f "jal" is not supported by hardware and is a pseudo instruction, which of the following correctly implements instruction "jal f1?"

```
(a)
     la $ra, L1
      j f1
                  # or any instruction after calling f1
L1:
     nop
(b)
     la $t0, f1
      jr f1
                  # or any instruction after calling f1
L1:
     nop
(c)
     la $ra, f1
      i fl
L1:
                  # or any instruction after calling f1
     nop
(d)
     None of the above.
```

#### MIPS Calling Conventions

- MIPS assembly follows the following convention in using registers
  - \$a0 \$a3: four argument registers in which to pass parameters
  - \$v0 \$v1: two value registers in which to return values
  - \$ra: one return address register to return to the point of origin

#### MIPS Conventions

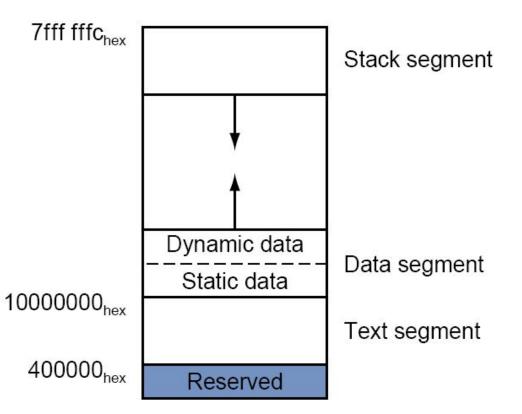
- Quite often, our function needs to use some registers to do dome calculation. So we will modify the values of them.
- We can use \$t0-\$t9 freely inside a function, because the caller does not expect the values inside \$t0-\$t9 to stay the same after the function call.
- But, the caller do expect the values in \$s0 to \$s7 to be the same after a function call.

#### MIPS Conventions

- So, just try to avoid using \$s0 and \$s7 inside a function whenever possible.
- But what if do need it? Such occasions will arise...

#### Stack

- So, if we do have to use \$s0 -\$s7, we MUST save it somewhere before entering the main part of the function, and restore it before we return (before we execute "jr \$ra").
- In MIPS, we save them in the stack.
- Stack is a part in the memory allocated for functions. It starts at 0x7ffffffc and grows down as we add more stuffs to it.
- Stack is "first in last out."



## \$sp

- The top address of the stack, the address of the first word that is storing value, is (should be) always stored in \$sp.
- So, adding a word into the stack (pushing a word onto the stack) is a two-step thing, because **you** have to maintain the correctness of \$sp:
  - •addi \$sp, \$sp, -4
  - •sw \$s0, 0(\$sp)

### Suppose we want to

```
| int weirdfun(int a, int b)
    int res;
    res = a + a + b - a / 2;
    return res;
int t = 0;
for (int i=0;i<10;i+=2)
    t += weirdfun(A[i],A[i+1]);
```

# Stack and \$sp

- Suppose we want to store a/2 in \$s0.
  - How do we get a/2?
- At the beginning, we do
  - addi \$sp, \$sp, -4
  - •sw \$s0, 0(\$sp)
- At the end, we do
  - •lw \$s0, 0(\$sp)
  - •addi \$sp, \$sp, 4

.data A: .word 12, 34, 67, 1, 45, 90, 11, 33, 67, 19 .text .globl main main: la \$s7, A li \$s0, 0 #i li \$s1, 0 #res li \$s6, 9 loop: sll \$t0, \$s0, 2 add \$t0, \$t0, \$s7 lw \$a0, 0(\$t0) lw \$a1, 4(\$t0) jal weirdfun add \$s1, \$s1, \$v0 addi \$s0, \$s0, 2 blt \$s0, \$s6, loop done: li \$v0,10 syscall weirdfun: addi \$sp, \$sp, -4 sw \$s0, 0(\$sp) srl \$s0, \$a0, 1 add \$t0, \$a0, \$a0 add \$t0, \$t0, \$a1 sub \$t0, \$t0, \$s0 ori \$v0, \$t0, 0 lw \$s0, 0(\$sp) addi \$sp, \$sp, 4 jr \$ra