

1 An import property of XOR

Y is a rand var over $\{0, 1\}^n$, X is an independent uniform var over $\{0, 1\}^n$ Then $Z := X \oplus Y$ is uniform var over $\{0, 1\}^n$

2 Symmetric ciphers: definition

A cipher defined over (K, M, C) is a pair of "efficient" algs (E, D) Where $E : K \times M \rightarrow C, D : C \rightarrow M$ s.t. $\forall m \in M D(k, E(k, m)) = m$

3 Information theoretic security

A cipher (E, D) over (K, M, C) has perfect secrecy if $\forall m_0, m_1 \in M, (|m_0| = |m_1|)$, and $\forall c \in C Pr[E(k, m_0) = c] = Pr[E(k, m_1) = c]$ where $k \leftarrow \text{rand } K$

4 PRG must be unpredictable

We say $G : K \rightarrow \{0, 1\}^n$ is predictable if: \exists efficient alg α and $\exists i, 1 \leq i \leq n-1$ s.t. $Pr[\alpha(G(k)|_{1,\dots,i}) = G(k)|_{i+1}] \geq \frac{1}{2} + \epsilon$

5 Secure PRGs: crypto definition

Def we say that $G: K \rightarrow \{0, 1\}^n$ is a secure PRG if \forall "eff" stat test $A : Adv_{PRG}[A, G]$ is "negligible"

6 PRF: pseudo random function

Def: let $F: K \times X \rightarrow Y$ be a PRF $Funs[X, Y]$: the set of all functions from X to Y $S_F = \{F(K, \bullet) \text{ s.t. } k \in K\} \subseteq Funs[X, Y]$

Intuition: a PRF is secure if a random function in $Funs[X, Y]$ is indistinguishable from a random function in S_F