



**BSI Standards Publication**

# **Information technology — Security techniques — Encryption algorithms**

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## **Part 4: Stream ciphers**

# National foreword

This British Standard is the UK implementation of ISO/IEC 18033-4:2011+A1:2020. It supersedes BS ISO/IEC 18033-4:2011, which is withdrawn.

The start and finish of text introduced or altered by amendment is indicated in the text by tags. Tags indicating changes to ISO/IEC text carry the number of the ISO/IEC amendment. For example, text altered by ISO/IEC amendment 1 is indicated by A1 A1.

The UK participation in its preparation was entrusted to Technical Committee IST/33, Information security, cybersecurity and privacy protection.

A list of organizations represented on this committee can be obtained on request to its committee manager.

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© The British Standards Institution 2020  
Published by BSI Standards Limited 2020

ISBN 978 0 539 04207 8

ICS 35.030

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This British Standard was published under the authority of the Standards Policy and Strategy Committee on 31 January 2012.

## Amendments/corrigenda issued since publication

Date	Text affected
30 September 2020	Implementation of ISO/IEC amendment 1:2020

# INTERNATIONAL STANDARD

# ISO/IEC 18033-4

Second edition  
2011-12-15

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## Information technology — Security techniques — Encryption algorithms —

### Part 4: Stream ciphers

*Technologies de l'information — Techniques de sécurité — Algorithmes  
de chiffrement —*

*Partie 4: Chiffrements en flot*

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Reference number  
ISO/IEC 18033-4:2011(E)



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Published in Switzerland

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## Foreword

ISO (the International Organization for Standardization) and IEC (the International Electrotechnical Commission) form the specialized system for worldwide standardization. National bodies that are members of ISO or IEC participate in the development of International Standards through technical committees established by the respective organization to deal with particular fields of technical activity. ISO and IEC technical committees collaborate in fields of mutual interest. Other international organizations, governmental and non-governmental, in liaison with ISO and IEC, also take part in the work. In the field of information technology, ISO and IEC have established a joint technical committee, ISO/IEC JTC 1.

International Standards are drafted in accordance with the rules given in the ISO/IEC Directives, Part 2.

The main task of the joint technical committee is to prepare International Standards. Draft International Standards adopted by the joint technical committee are circulated to national bodies for voting. Publication as an International Standard requires approval by at least 75 % of the national bodies casting a vote.

Attention is drawn to the possibility that some of the elements of this document may be the subject of patent rights. ISO and IEC shall not be held responsible for identifying any or all such patent rights.

ISO/IEC 18033-4 was prepared by Joint Technical Committee ISO/IEC JTC 1, *Information technology*, Subcommittee SC 27, *IT Security techniques*.

This second edition cancels and replaces the first edition (ISO/IEC 18033-4:2005), which has been technically revised. It also incorporates the Amendment ISO/IEC 18033-4:2005/Amd.1:2009.

ISO/IEC 18033 consists of the following parts, under the general title *Information technology — Security techniques — Encryption algorithms*:

- *Part 1: General*
- *Part 2: Asymmetric ciphers*
- *Part 3: Block ciphers*
- *Part 4: Stream ciphers*

## Introduction

This part of ISO/IEC 18033 includes stream cipher algorithms. A stream cipher is an encryption mechanism that uses a keystream to encrypt a plaintext in a bitwise or a block-wise manner. There are two types of stream ciphers: a synchronous stream cipher, in which the keystream is generated from only the secret key (and an initialization vector) and a self-synchronizing stream cipher, in which the keystream is generated from the secret key and some past ciphertexts (and an initialization vector). This part of ISO/IEC 18033 describes both pseudorandom number generators for producing keystream and output functions to combine a keystream with plaintext.

This part of ISO/IEC 18033 includes two output functions:

- Binary-additive output function; and
- MULTI-S01 output function.

**A1** This document includes six dedicated keystream generators:

- MUGI keystream generator;
- SNOW 2.0 keystream generator;
- Rabbit keystream generator;
- Decim<sup>v2</sup> keystream generator;
- KCipher-2 (K2) keystream generator; and
- ZUC keystream generator. **A1**

# Information technology — Security techniques — Encryption algorithms —

## Part 4: Stream ciphers

### 1 Scope

This part of ISO/IEC 18033 specifies

- a) output functions to combine a keystream with plaintext,
- b) keystream generators for producing keystream, and
- c) object identifiers assigned to dedicated keystream generators in accordance with ISO/IEC 9834.

NOTE 1 The list of assigned object identifiers is given in Annex A.

NOTE 2 Any change to the specification of these algorithms resulting in a change of functional behaviour will result in a change of the object identifier assigned to the algorithms concerned.

### 2 Normative references

The following referenced documents are indispensable for the application of this document. For dated references, only the edition cited applies. For undated references, the latest edition of the referenced document (including any amendments) applies.

ISO/IEC 18033-1, *Information technology — Security techniques — Encryption algorithms — Part 1: General*

### 3 Terms and definitions

For the purposes of this document, the terms and definitions given in ISO/IEC 18033-1 and the following apply.

#### 3.1

##### **big-endian**

method of storage of multi-byte numbers with the most significant bytes at the lowest memory addresses

[ISO/IEC 10118-1:2000]

#### 3.2

##### **ciphertext**

data which has been transformed to hide its information content

[ISO/IEC 10116:2006]



**3.3 confidentiality**  
property that information is not made available or disclosed to unauthorized individuals, entities, or processes

**3.4 data integrity**  
property that data has not been altered or destroyed in an unauthorized manner

[ISO/IEC 9797-1:2011]

**3.5 decryption**  
reversal of a corresponding encryption

[ISO/IEC 10116:2006]

**3.6 encryption**  
reversible transformation of data by a cryptographic algorithm to produce ciphertext, i.e., to hide the information content of the data

[ISO/IEC 9797-1:2011]

**3.7 initialization value**  
value used in defining the starting point of an encryption process

**3.8 key**  
sequence of symbols that controls the operation of a cryptographic transformation (e.g., encryption, decryption, cryptographic check function computation, signature generation, or signature verification)

[ISO/IEC 11770-1:2010]

**3.9 keystream function**  
function that takes as input, the current state of the keystream generator and (optionally) part of the previously generated ciphertext, and gives as output the next part of the keystream

**3.10 keystream generator**  
state-based process (i.e., a finite state machine) that takes as input, a key, an initialization vector, and if necessary the ciphertext, and gives as output a keystream (i.e., a sequence of bits or blocks of bits) of arbitrary length

**3.11  $n$ -bit block cipher**  
block cipher with the property that plaintext blocks and ciphertext blocks are  $n$  bits in length

[ISO/IEC 10116:2006]

**3.12 next-state function**  
function that takes as input, the current state of the keystream generator and (optionally) part of the previously generated ciphertext, and gives as output a new state for the keystream generator

**3.13 output function**  
function that combines the keystream and the plaintext to produce the ciphertext

NOTE This function is often bitwise XOR.

### 3.14

#### **padding**

appending extra bits to a data string

[ISO/IEC 10118-1:2000]

### 3.15

#### **plaintext**

unencrypted information

[ISO/IEC 9797-1:2011]

### 3.16

#### **secret key**

key used with symmetric cryptographic techniques by a specified set of entities

[ISO/IEC 11770-3:2008]

### 3.17

#### **state**

current internal state of a keystream generator

## 4 Symbols and abbreviated terms

### 4.1 Symbols

$0x$	Prefix for hexadecimal values.
$0^{(n)}$	$n$ -bit variable where 0 is assigned to every bit.
$AND$	Bitwise logical $AND$ operation.
$Am^{(i)}[Y]$	The $Y$ -th bit of the register $Am^{(i)}$ in KCipher-2 (K2).
$a_i$	Variables in an internal state of a keystream generator.
$b_i$	Variables in an internal state of a keystream generator.
CFB	Cipher FeedBack mode of a block cipher.
CTR	Counter mode of a block cipher.
$C_i$	Ciphertext block.
$D_i$	64-bit constants used for MUGL.
$e_K$	Symmetric block cipher encryption function using secret key $K$ .
$F$	Subfunction used for MUGL.
$FSM$	Subfunction used for SNOW 2.0.
$GF(2^n)$	Finite field of exactly $2^n$ elements.
$GF(2^n)[x]$	The polynomial ring over the finite field $GF(2^n)$ .

<i>Init</i>	Function which generates the initial internal state of a keystream generator.
<i>IV</i>	Initialization vector.
<i>IK</i>	Internal key used for KCipher-2 (K2).
<i>K</i>	Key.
$\boxed{A_1}$ <i>L</i> <sub>1</sub>	Linear transform with index 1 used for ZUC. $\langle A_1 \rangle$
$\boxed{A_1}$ <i>L</i> <sub>2</sub>	Linear transform with index 2 used for ZUC. $\langle A_1 \rangle$
<i>M</i>	Subfunction used for MUGI.
<i>Next</i>	Next-state function of a keystream generator.
<i>NLF</i>	Nonlinear function used for KCipher-2 (K2).
<i>n</i>	Block length.
OFB	Output FeedBack mode of a block cipher.
<i>OR</i>	Bitwise logical OR operation.
<i>Out</i>	Output function combining keystream and plaintext in order to generate ciphertext.
<i>P</i>	Plaintext.
<i>P</i> <sub><i>i</i></sub>	Plaintext block.
<i>R</i>	Additional input to Out.
<i>S</i> <sub><i>R</i></sub>	Subfunction used for MUGI.
$\boxed{A_1}$ <i>SS</i>	Subfunction used for ZUC. $\langle A_1 \rangle$
<i>Strm</i>	Keystream function of a keystream generator.
<i>SUB</i>	Lookup table used for MUGI and SNOW 2.0.
$\boxed{A_1}$ <i>SUB</i> 1	Lookup table with index 1 used for ZUC $\langle A_1 \rangle$
$\boxed{A_1}$ <i>SUB</i> 2	Lookup table with index 2 used for ZUC. $\langle A_1 \rangle$
<i>Sub</i> <sub>K2</sub>	Subfunction used for KCipher-2 (K2).
<i>S</i> <sub><i>i</i></sub>	Internal state of a keystream generator.

**NOTE** During normal operation of the cipher, *i* will increase monotonically starting from zero. However, during initialization of the ciphers, it is convenient from a notational point of view to let *i* take negative values and define the starting state *S*<sub>0</sub> in terms of values of *S*<sub>*i*</sub> for *i* < 0.

<i>T</i>	Subfunction used for SNOW 2.0.
<i>Z</i>	Keystream.
<i>Z</i> <sub><i>i</i></sub>	Keystream block.
$\alpha_{\text{MUL}}$	Lookup table used for SNOW 2.0.
$\alpha_{\text{MUL}0}$	Lookup table with index 0 used for KCipher-2 (K2).
$\alpha_{\text{MUL}1}$	Lookup table with index 1 used for KCipher-2 (K2).
$\alpha_{\text{MUL}2}$	Lookup table with index 2 used for KCipher-2 (K2).

$\alpha_{\text{MUL3}}$	Lookup table with index 3 used for KCipher-2 (K2).
$\alpha_{\text{inv\_MUL}}$	Inverse lookup table used for SNOW 2.0.
$\rho_1$	Subfunction used for MUGL.
$\lambda_1$	Subfunction used for MUGL.
$\lceil x \rceil$	The smallest integer greater than or equal to the real number $x$ .
$\neg x$	Bitwise complement operation.
$\bullet$	Polynomial multiplication.
$\parallel$	Bit concatenation.
$+_m$	Integer addition modulo $2^m$ .
$\oplus$	Bitwise XOR (eXclusive OR) operation.
$\otimes$	Operation of multiplication of elements in the finite field $GF(2^n)$ .

**EXAMPLE**  $C = A \otimes B$ : In this operation, the finite field is represented as a selected irreducible polynomial  $F(x)$  of degree  $n$  with binary coefficients, the  $n$ -bit blocks  $A = (a_0, a_1, \dots, a_{n-1})$  and  $B = (b_0, b_1, \dots, b_{n-1})$  (where the  $a_i$  and  $b_i$  are bits) are represented as the polynomials,  $A(x) = a_{n-1}x^{n-1} + a_{n-2}x^{n-2} + \dots + a_1$  and  $B(x) = b_{n-1}x^{n-1} + b_{n-2}x^{n-2} + \dots + b_0$  respectively, then let  $C(x) = A(x) \bullet B(x) \bmod F(x)$ , i.e.,  $C(x)$  is the polynomial of degree at most  $n-1$  obtained by multiplying  $A(x)$  and  $B(x)$ , dividing the result by  $F(x)$ , and then taking the remainder. If  $C(x) = c_{n-1}x^{n-1} + c_{n-2}x^{n-2} + \dots + c_0$  (where the  $c_i$  are bits) then let  $C$  be the  $n$ -bit block  $(c_0, c_1, \dots, c_{n-1})$ .

$\boxplus$	Modular addition operation
$\llcorner_n t$	$t$ -bit left shift in an $n$ -bit register.
$\ggcorner_n t$	$t$ -bit right shift in an $n$ -bit register.
$\lllcorner_n t$	$t$ -bit left circular rotation in an $n$ -bit register.
$\gggcorner_n t$	$t$ -bit right circular rotation in an $n$ -bit register.

## 4.2 Functions

### 4.2.1 Left-truncation of bits

The operation of selecting the  $j$  leftmost bits of an array  $A = (a_0, a_1, \dots, a_{m-1})$  to generate a  $j$ -bit array is written

$$(j \sim A) = (a_0, a_1, \dots, a_{j-1})$$

This operation is defined only when  $1 \leq j \leq m$ .

See ISO/IEC 10116:2006.

#### 4.2.2 Shift operation

The operation *Shift* is defined as follows: Given an  $n$ -bit variable  $X$  and a  $k$ -bit variable  $V$  where  $1 \leq k \leq n$ , the effect of the shift function *Shift* is to produce the  $n$ -bit variable

$$\text{Shift}_k(X \mid V) = (x_k, x_{k+1}, \dots, x_{n-1}, v_0, v_1, \dots, v_{k-1}) \quad (k < n)$$

$$\text{Shift}_k(X \mid V) = (v_0, v_1, \dots, v_{k-1}) \quad (k = n)$$

The effect is to shift the bits of array  $X$  left by  $k$  places, discarding  $x_0, x_1, \dots, x_{k-1}$  and to place the array  $V$  in the rightmost  $k$  places of  $X$ . When  $k = n$  the effect is to totally replace  $X$  by  $V$ .

See ISO/IEC 10116:2006.

#### 4.2.3 Variable $I(k)$

The variable  $I(k)$  is a  $k$ -bit variable where 1 is assigned to every bit.

### 5 Framework for stream ciphers

This clause contains a high-level description of a framework for the stream ciphers specified in this part of ISO/IEC 18033. A detailed description of the general model for a stream cipher is provided in Clause 6. A stream cipher specified in this part of ISO/IEC 18033 is defined by the specification of the following processes:

- a) The keystream generator, which may be either
  - a Synchronous keystream generator, or
  - a Self-synchronizing keystream generator.

NOTE 1 Block cipher modes of operation are methods by which a block cipher can be used to construct a keystream generator. These modes are standardised in ISO/IEC 10116, and the meaning of the functions used in the specification is defined in 6.2.1 and 6.2.2.

NOTE 2 Block ciphers are defined in this part of ISO/IEC 18033.

- b) The output function, which may be either
  - the Binary-additive output function, or
  - the MULTI-S01 output function.

### 6 General models for stream ciphers

#### 6.1 Keystream generators

##### 6.1.1 Synchronous keystream generators

A synchronous keystream generator is a finite-state machine. It is defined by:

- a) An initialization function, *Init*, which takes as input a key  $K$  and an initialization vector  $IV$ , and outputs an initial state  $S_0$  for the keystream generator. The initialization vector should be chosen so that no two messages are ever encrypted using the same key and the same  $IV$ .

- b) A next-state function, *Next*, which takes as input the current state of the keystream generator  $S_i$ , and outputs the next state of the keystream generator  $S_{i+1}$ .
- c) A keystream function, *Strm*, which takes as input a state of the keystream generator  $S_i$ , and outputs a keystream block  $Z_i$ .

When the synchronous keystream generator is first initialized, it will enter an initial state  $S_0$  defined by:

$$S_0 = \text{Init}(IV, K).$$

On demand the synchronous keystream generator will, for  $i=0,1,\dots$ :

- a) Output a keystream block  $Z_i = \text{Strm}(S_i, K)$ .
- b) Update the state of the machine  $S_{i+1} = \text{Next}(S_i, K)$ .

Therefore to define a synchronous keystream generator it is only necessary to specify the functions *Init*, *Next* and *Strm*, including the lengths and alphabets of the key, the initialization vector, the state, and the output block.

### 6.1.2 Self-synchronizing keystream generators

Generation of a keystream for a self-synchronizing stream cipher is dependent only on previous ciphertexts, the key, and the initialization vector. A general model for a keystream generator for a self-synchronizing stream cipher is now defined:

- a) An initialization function, *Init*, which takes as input a key  $K$  and an initialization vector  $IV$  and outputs an internal input for the keystream generator  $S$  and  $r$  dummy ciphertext blocks  $C_{-1}, C_{-2}, \dots, C_{-r}$ .
- b) A keystream function, *Strm*, that takes as input  $S$  and  $r$  ciphertext blocks  $C_{i-1}, C_{i-2}, \dots, C_{i-r}$ , and outputs a keystream block  $Z_i$ .

To define a self-synchronizing keystream generator it is only necessary to specify the number of feedback blocks  $r$  and the functions *Init* and *Strm*.

**NOTE** A self-synchronizing stream cipher differs from a synchronous stream cipher in that the keystream depends only on previous ciphertext, the initialization vector and the key, i.e., the keystream generator operates in a stateless fashion. As a result, a decryptor for such a cipher can recover from loss of synchronization after receiving sufficient ciphertext blocks. This also means that the method of keystream generation is dependent upon the selected output function *Out*, which is typically the bitwise XOR operation.

## 6.2 Output functions

### 6.2.1 General model of output function

6.2 specifies two stream cipher output functions, i.e., techniques to be used in a stream cipher to combine a keystream with plaintext to derive ciphertext.

An output function for a synchronous or a self-synchronizing stream cipher is a function *Out* that combines a plaintext block  $P_i$ , a keystream block  $Z_i$ , and some other input  $R$  if necessary to give a ciphertext block  $C_i$  ( $i \geq 0$ ). A general model for a stream cipher output function is now defined:

Encryption of a plaintext block  $P_i$  by a keystream block  $Z_i$  is given by:

$$C_i = \text{Out}(P_i, Z_i, R),$$

and decryption of a ciphertext block  $C_i$  by a keystream block  $Z_i$  is given by:

$$P_i = \text{Out}^{-1}(C_i, Z_i, R).$$

The output function shall satisfy that for any keystream block  $Z_i$ , plaintext block  $P_i$ , and other input  $R$ ,

$$P_i = \text{Out}^{-1}(\text{Out}(P_i, Z_i, R), Z_i, R).$$

## 6.2.2 Binary-additive output function

A binary-additive stream cipher is a stream cipher in which the keystream, plaintext, and ciphertext blocks are strings of binary digits, and the operation to combine plaintext with keystream is bitwise XOR. The operation  $\text{Out}$  takes two inputs and does not use any additional information  $R$  for calculation. Let  $n$  to be the bit length of  $P_i$ . This function is specified by

$$\text{Out}(P_i, Z_i, R) = P_i \oplus Z_i.$$

The operation  $\text{Out}^{-1}$  is specified by

$$\text{Out}^{-1}(C_i, Z_i, R) = C_i \oplus Z_i.$$

**NOTE** The binary-additive stream cipher does not provide any integrity protection for encrypted data. If data integrity is required, either the MULTI-S01 output function or a separate integrity mechanism should be used, such as a MAC, i.e., a Message Authentication Code (such mechanisms are specified in ISO/IEC 9797).

## 6.2.3 MULTI-S01 output function

### a) General model of MULTI-S01

MULTI-S01 is an output function for a synchronous stream cipher that supports both data confidentiality and data integrity. The MULTI-S01 encryption operation is suitable for use in an online environment. However, the decryption operation of MULTI-S01 can only be performed in an offline situation, as the integrity check is only performed after receiving all the ciphertext blocks. MULTI-S01 has a security parameter  $n$ . The computation of  $\text{Out}$  depends on the choice of a field  $GF(2^n)$ , i.e., on the choice of an irreducible polynomial over  $GF(2)$  of degree  $n$ . The MULTI-S01 function only accepts messages whose length is a multiple of  $n$ . To encrypt messages whose length is not a multiple of  $n$ , a padding mechanism  $\text{Pad}(M)$  is required.

**NOTE 1** The redundancy  $R$  is generated in such a way that the sender and the receiver share it.  $R$  can be a fixed public value like 0x00...0.

### b) The encryption function $\text{Out}(P, Z, R)$

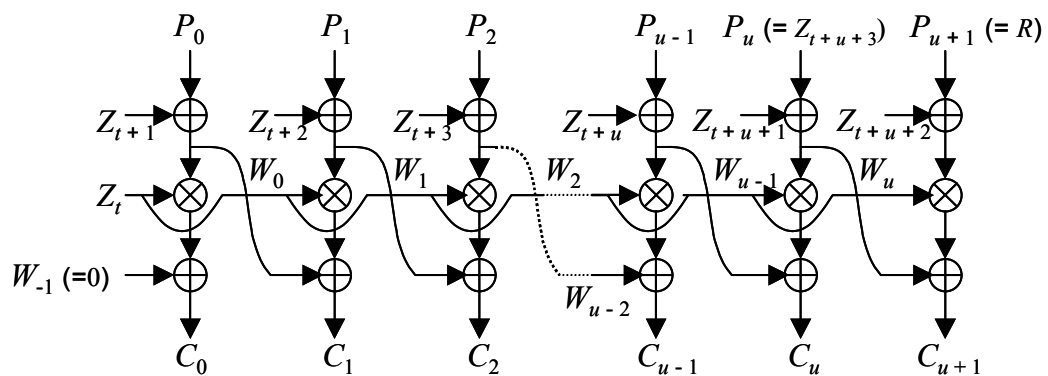
Input:  $n \cdot u$ -bit plaintext  $P$ , keystream  $Z = (Z_0, Z_1, \dots)$ , where  $Z_i$  are  $n$ -bit blocks,  $n$ -bit redundancy  $R$ .

Output: Ciphertext  $C$ .

- 1) Let  $t$  be the lowest value of  $i$  ( $i \geq 0$ ) such that  $Z_i \neq 0^{(n)}$ .
- 2) Let  $(P_0, P_1, \dots, P_{u-1}) = P$ , where  $P_i$  is an  $n$ -bit block.
- 3) Set  $P_u = Z_t + u + 3$ .
- 4) Set  $P_{u+1} = R$ .

- 5) For each  $P_i$ , do the following calculations (for  $i = 0, 1, \dots, u + 1$ ):
  - Let  $W_i = P_i \oplus Z_{t+i+1}$ .
  - Let  $X_i = Z_t \otimes W_i$  (in  $GF(2^n)$ ).
  - Let  $C_i = X_i \oplus W_{i-1}$ , where  $W_{i-1}$  is the  $W$  value of the previous block  $i - 1$ , and  $W_{-1} = 0^{(n)}$ .
  - Set  $C = C_0 \parallel C_1 \parallel \dots \parallel C_{u+1}$ .
  - Output  $C$ .

Figure 1 shows the block diagram of *Out* function.



**Figure 1 — *Out* function of MULTI-S01 mode**

NOTE 2 The irreducible polynomial used to define multiplication in the field depends on  $n$ . For instance, in the case of  $n = 64$  and  $128$ , the irreducible polynomial  $x^{64} + x^4 + x^3 + x + 1$  and  $x^{128} + x^7 + x^2 + x + 1$  can be used.

- c) The decryption function  $Out^{-1}(P, Z, R)$

Input:  $n \cdot v$ -bit ciphertext  $C$ , keystream  $Z$ ,  $n$ -bit redundancy  $R$ .

Output: Plaintext  $P$  or “*reject*”.

- 1) Let  $t$  be the lowest value of  $i$  ( $i \geq 0$ ) such that  $Z_i \neq 0^{(n)}$ .
- 2) Let  $(C_0, C_1, \dots, C_{v-1}) = C$ , where  $C_i$  is an  $n$ -bit block.
- 3) For each  $C_i$ , do the following calculations (for  $i = 0, 1, \dots, v - 1$ ):
  - Let  $X_i = C_i \oplus W_{i-1}$ , where  $W_{-1} = 0^{(n)}$ .
  - Let  $W_i = Z_t^{-1} \otimes X_i$  (in  $GF(2^n)$ ).
  - Let  $P_i = W_i \oplus Z_{t+i+1}$ .
- 4) If  $P_{v-2} = Z_{t+v+1}$  and  $P_{v-1} = R$ , output  $P = P_0 \parallel P_1 \parallel \dots \parallel P_{v-3}$  as plaintext. Otherwise, output the special symbol meaning “*reject*” without any text.



Figure 2 shows the block diagram of  $Out^{-1}$  function.

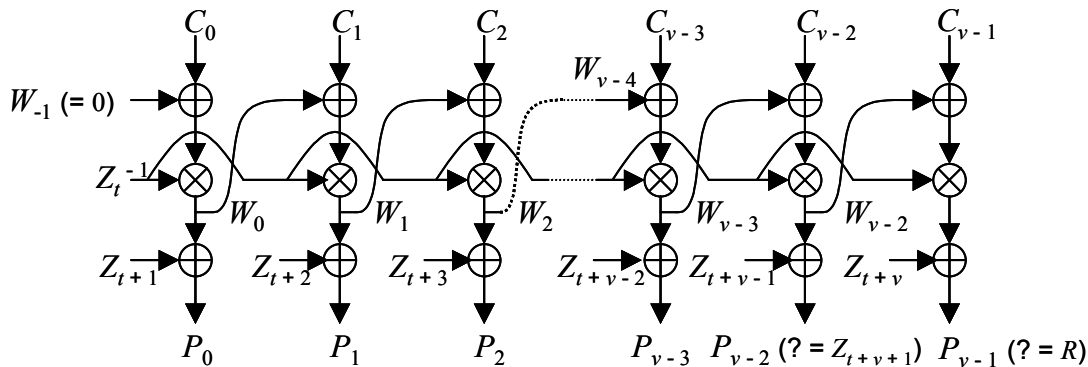


Figure 2 —  $Out^{-1}$  function of MULTI-S01 mode

d) Padding mechanism  $Pad(M)$

Only when lengths of input messages are not multiples of  $n$ , the following padding mechanism  $Pad(M)$  is executed:

Input:  $(n \nu + c)$ -bit string  $M$ , where  $\nu$  is a non-negative integer and  $0 \leq c < n$ .

Output: Padded plaintext  $P$ .

- 1) Pad a bit string "1" at the end of the message.
- 2) Pad  $(n - c - 1)$ -bit string  $0^{(n - c - 1)}$  to the string generated by step a).
- 3) Output the whole data string in length of  $(n \nu + n)$  bits.

NOTE 3 If the length of the message is a multiple of  $n$  in an environment where the length is not certain to be so, this padding mechanism is recommended.

NOTE 4 In order to unpad the message, remove consecutive 0 bits at the end of the data, and remove another bit "1".

## 7 Constructing keystream generators from block ciphers

### 7.1 Block cipher modes for a synchronous keystream generator

#### 7.1.1 The OFB (Output FeedBack) mode and the CTR (Counter) mode

Subclause 7.1 specifies two  $n$ -bit block cipher modes for a synchronous keystream generator. They are the OFB (Output FeedBack) mode and the CTR (Counter) mode of an  $n$ -bit block cipher  $e_K$ .

#### 7.1.2 OFB mode

The OFB mode is defined by one parameter  $r$ ,  $1 \leq r \leq n$ , which is the size of a plaintext and ciphertext block.

The initialization vector  $IV$  is an  $n$ -bit string.  $IV$  shall be generated differently for two encryptions with the same key  $K$ . The functions  $Init$ ,  $Next$  and  $Strm$  are specified as follows:

—  $Init(IV, K) = IV$ .

- $Next(S_i, K) = e_K(S_i)$ .
- $Strm(S_i) = (r \sim S_i)$ .

NOTE  $Init(IV, K) = IV$ , is equivalent to  $S_0 = IV$ .

In case of the OFB mode, the binary-additive output function defined in 6.2.2 is used. Figure 3 shows the block diagram of a keystream generator based on CFB mode.

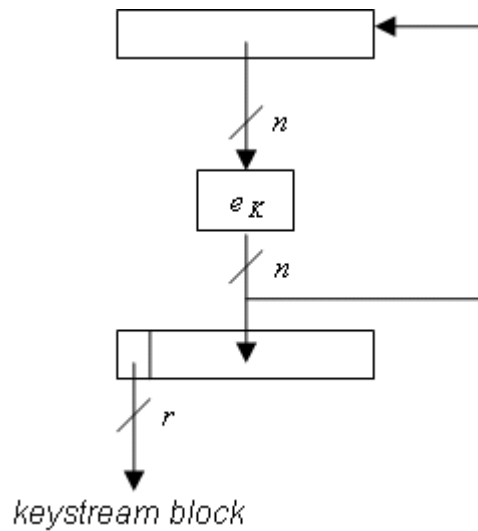


Figure 3 — Keystream generation based on OFB mode

### 7.1.3 CTR mode

The CTR mode is defined by one parameter  $r$ ,  $1 \leq r \leq n$ , which is the size of a plaintext and ciphertext block.

The initialization vector  $IV$  is an  $n$ -bit string. It shall be assured that  $S_i \neq S_j$  for two keystreams  $S_0, S_1, S_2, \dots$  and  $S_0', S_1', S_2', \dots$  generated with the same key  $K$ . The functions  $Init$ ,  $Next$  and  $Strm$  are specified as follows:

- $Init(IV, K) = IV$ .
- $Next(S_i, K) = S_i + 1 \bmod 2^n$ .
- $Strm(S_i, K) = (r \sim e_K(S_i))$ .

NOTE  $Init(IV, K) = IV$ , is equivalent to  $S_0 = IV$ .

In case of the CTR mode, the binary-additive output function defined in 6.2.2 is used. Figure 4 shows the block diagram of a keystream generator based on CFB mode.

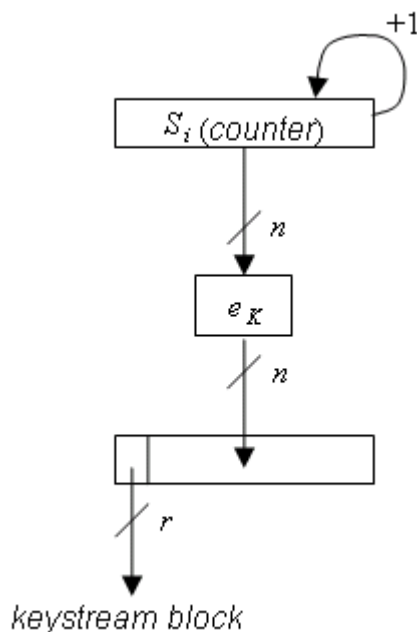


Figure 4 — Keystream generation based on CTR mode

## 7.2 Block cipher mode for a self-synchronizing keystream generator

### 7.2.1 Introduction to the CFB mode

The CFB mode of an  $n$ -bit block cipher is a self-synchronizing stream cipher.

#### 7.2.2 CFB mode

The CFB (Cipher FeedBack) mode is defined by three parameters, i.e., the size  $j$  of feedback buffer  $S_i$ , where  $n \leq j \leq 1024n$ , the size  $b$  of feedback variable, where  $1 \leq b \leq n$  and the size  $r$  of the output block, where  $1 \leq r \leq b$ .

NOTE 1 The value  $b-r$  shall be small compared to  $b$ .

The initialization vector  $IV$  shall be a randomly generated  $j$ -bit string and also shall be generated differently for two encryptions with the same key  $K$ . The functions *Init*, *Next* and *Strm* are specified as follows:

- $Init(IV, K) = IV$ .
- $Next(S) = Shift_b(S \parallel Shift_r(I(b) \parallel C_i))$ .
- $Strm(S, K) = (r \sim e_K((n \sim S)))$ .

NOTE 2  $Init(IV, K) = IV$ , is equivalent to  $S_0 = IV$ .

In case of the CFB mode, the binary-additive output function defined in 6.2.2 is used. Figure 5 shows the block diagram of a keystream generator based on CFB mode.

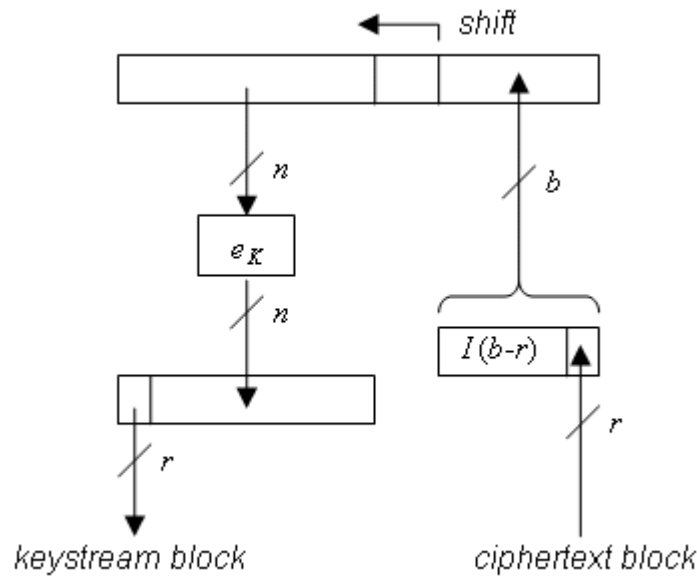


Figure 5 — Keystream generation based on CFB mode

## 8 Dedicated keystream generators

### 8.1 MUGI keystream generator

#### 8.1.1 Introduction to MUGI

MUGI is a keystream generator which uses a 128-bit secret key  $K$ , a 128-bit initialization vector  $IV$ , and a state variable  $S_i$  ( $i \geq 0$ ) consisting of 19 64-bit blocks (note that the term block is used through the specification of MUGI for a 64-bit block), and outputs a keystream block  $Z_i$  at every iteration of the function  $Strm$ .

NOTE This keystream generator is originally proposed in [17].

The state variable  $S_i$  is sub-divided into a combination of a 3-block variable:

$$a^{(i)} = (a_0^{(i)}, a_1^{(i)}, a_2^{(i)}),$$

where  $a_j^{(i)}$  is a block (for  $j = 0, 1, 2$ ), and a 16-block variable:

$$b^{(i)} = (b_0^{(i)}, b_1^{(i)}, \dots, b_{15}^{(i)}),$$

where  $b_j^{(i)}$  is a block (for  $j = 0, 1, \dots, 15$ ).

The  $Init$  function, defined in detail in 8.1.2, takes as input the 128-bit key  $K$  and the 128-bit initializing vector  $IV$ , and produces the initial value of the state variable  $S_0 = (a^{(0)}, b^{(0)})$ .

The  $Next$  function, defined in detail in 8.1.3, takes as input the 19-block state variable  $S_i = (a^{(i)}, b^{(i)})$  and produces as output the next value of the state variable  $S_{i+1} = (a^{(i+1)}, b^{(i+1)})$ .

The  $Strm$  function, defined in detail in 8.1.4, takes as input the 19-block state variable  $S_i = (a^{(i)}, b^{(i)})$  and produces as output the keystream block  $Z_i$ .

Note that the *Next* function is defined in terms of the functions  $\rho_1$  and  $\lambda_1$  which are defined in 8.1.5 and 8.1.6, respectively. The function  $\rho_1$  is defined in terms of a function  $F$  which is defined in 8.1.7

There are three constants used in MUGI,  $D_0$  in the initialization function *Init*, and  $D_1, D_2$  in  $\rho_1$ . These are given by:

$$D_0 = 0 \times 6A09E667F3BCC908,$$

$$D_1 = 0 \times BB67AE8584CAA73B,$$

$$D_2 = 0 \times 3C6EF372FE94F82B.$$

### 8.1.2 Initialization function *Init*

The initialization of MUGI is divided into eight steps. The left- and right-half blocks of  $K$  are denoted by  $K_0$  and  $K_1$  respectively.  $IV_0$  and  $IV_1$  are defined in the same manner. The initialization function *Init* is as follows:

Input: 128-bit key  $K$ , 128-bit initialization vector  $IV$ .

Output: Initial value of the state variable  $S_0 = (a^{(0)}, b^{(0)})$ .

a) Set the key  $K$  into the part of the state variable  $a^{(-49)}$  as follows:

- Set  $(K_0, K_1) = K$ , where  $K_i$  is 64 bits for  $i = 0, 1$ .
- Set  $a_0^{(-49)} = K_0$ .
- Set  $a_1^{(-49)} = K_1$ .
- Set  $a_2^{(-49)} = (K_0 \lll_{64} 7) \oplus (K_1 \ggg_{64} 7) \oplus D_0$ .

$D_0$  in the above equation is a constant (see 8.1).

b) For  $i = -49, -48, \dots, -34$ , set  $a^{(i+1)} = \rho_1(a^{(i)}, 0^{(64)}, 0^{(64)})$ . For the description of  $\rho_1$  see 8.1.5.

c) For  $i = 0, 1, \dots, 15$ , set  $b_{15-i}^{(-16)} = a_0^{(i-48)}$ .

d) Add the initialization vector  $IV$  into the state  $a$  as follows:

- Set  $IV_0 \parallel IV_1 = IV$ , where  $IV_i$  is a block.
- Set  $a_0^{(-32)} = a_0^{(-33)} \oplus IV_0$ .
- Set  $a_1^{(-32)} = a_1^{(-33)} \oplus IV_1$ .
- Set  $a_2^{(-32)} = a_2^{(-33)} \oplus (IV_0 \lll_{64} 7) \oplus (IV_1 \ggg_{64} 7) \oplus D_0$ .

e) For  $i = -32, -31, \dots, -17$ , set  $a^{(i+1)} = \rho_1(a^{(i)}, 0^{(64)}, 0^{(64)})$ .

f) Set  $S_{-16} = (a^{(-16)}, b^{(-16)})$ .

g) Iterate the update function *Next* 16 times:

Set  $S_0 = \text{Next}^{16}(S_{-16})$ ,

where  $\text{Next}^{16}$  stands for 16 iterations of the next-state function *Next*.

h) Output  $S_0$ .

### 8.1.3 Next-state function *Next*

The next-state function of MUGI is described as a combination of  $\rho_1$  and  $\lambda_1$ . The next-state function *Next* of MUGI is as follows:

Input: State variable  $S_i = (a^{(i)}, b^{(i)})$ .

Output: Next value of the state variable  $S_{i+1} = (a^{(i+1)}, b^{(i+1)})$ .

- Set  $a^{(i+1)} = \rho_1(a^{(i)}, b_4^{(i)}, b_{10}^{(i)})$ . The detailed description of the function  $\rho_1$  is given in 8.1.5.
- Set  $b^{(i+1)} = \lambda_1(b^{(i)}, a_0^{(i)})$ . The detailed description of the function  $\lambda_1$  is given in 8.1.6.
- Set  $S_{i+1} = (a^{(i+1)}, b^{(i+1)})$ .
- Output  $S_{i+1}$ .

### 8.1.4 Keystream function *Strm*

The keystream function *Strm* is as follows:

Input: State variable  $S_i$ .

Output: Keystream block  $Z_i$ .

- Set  $Z_i = a_2^{(i)}$ .
- Output  $Z_i$ .

### 8.1.5 Function $\rho_1$

The function  $\rho_1$  is as follows:

Input: State variable  $a^{(i)}$ , two 64-bit parameters  $w_1, w_2$ .

Output: The next value of the state variable  $a^{(i+1)}$ .

- Set  $a_0^{(i+1)} = a_1^{(i)}$ .
- Set  $a_1^{(i+1)} = a_2^{(i)} \oplus F(a_1^{(i)}, w_1) \oplus D_1$ .
- Set  $a_2^{(i+1)} = a_0^{(i)} \oplus F(a_1^{(i)}, (w_2 \lll 17)) \oplus D_2$ .
- Output  $a^{(i+1)}$ .

$D_1, D_2$  are constants (see 8.1 for details).

Figure 6 shows the block diagram of the function  $\rho_1$ . The detailed description of the function  $F$  is given in 8.1.7.

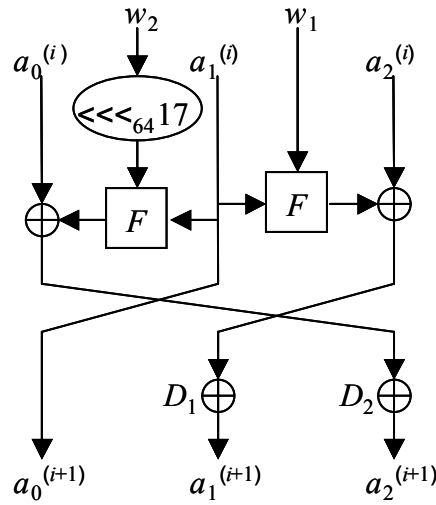


Figure 6 —  $\rho_1$  function of MUGI

### 8.1.6 Function $\lambda_1$

The function  $\lambda_1$  is as follows:

Input: State variable  $b^{(i)}$ , 64-bit parameter  $a'$ .

Output: The next value of the state variable  $b^{(i+1)}$ .

- Set  $b_j^{(i+1)} = b_{j-1}^{(i)}$  for  $j \neq 0, 4, 10$ .
- Set  $b_0^{(i+1)} = b_{15}^{(i)} \oplus a'$ .
- Set  $b_4^{(i+1)} = b_3^{(i)} \oplus b_7^{(i)}$ .
- Set  $b_{10}^{(i+1)} = b_9^{(i)} \oplus (b_{13}^{(i)} \lll_{64} 32)$ .

Output  $b^{(i+1)}$ .

### 8.1.7 Function $F$

Function  $F$  uses operations over the finite field  $GF(2^8)$ . In the polynomial representation,  $GF(2^8)$  is realized as  $GF(2)[x] / f(x)$ , where  $f(x)$  is an irreducible polynomial of degree 8 defined over  $GF(2)$ . The MUGI keystream generator uses the following irreducible polynomial:

$$f(x) = x^8 + x^4 + x^3 + x + 1.$$

The function  $F$  is the composition of a key addition (the data addition from the part of state variable  $b$ ), a non-linear transformation using the function  $S_R$ , a linear transformation using the matrix  $M$  and byte shuffling (see Figure 7).

Let us denote the input and the output to the  $F$  function as  $X$  and  $Y$  respectively. Then the function  $F$  is as follows:

Input: Two 64-bit strings  $X$  and  $T$ .

Output: 64-bit string  $Y$ .

- $X' = X \oplus T$ .
- Set  $(X_0, X_1, X_2, X_3, X_4, X_5, X_6, X_7) = X'$ , where  $X_i$  is an 8-bit string.
- Set  $P_i = S_R(X_i)$  for  $i = 0, 1, \dots, 7$ .
- Set  $P_L = P_0 \parallel P_1 \parallel P_2 \parallel P_3$ .
- Set  $P_R = P_4 \parallel P_5 \parallel P_6 \parallel P_7$ .
- Set  $Q_L = M(P_L)$ .
- Set  $Q_R = M(P_R)$ .
- Set  $(Q_0, Q_1, Q_2, Q_3) = Q_L$ .
- Set  $(Q_4, Q_5, Q_6, Q_7) = Q_R$ .
- Set  $Y = Q_4 \parallel Q_5 \parallel Q_2 \parallel Q_3 \parallel Q_0 \parallel Q_1 \parallel Q_6 \parallel Q_7$ .
- Output  $Y$ .

Figure 7 shows the block diagram of the function  $F$ .

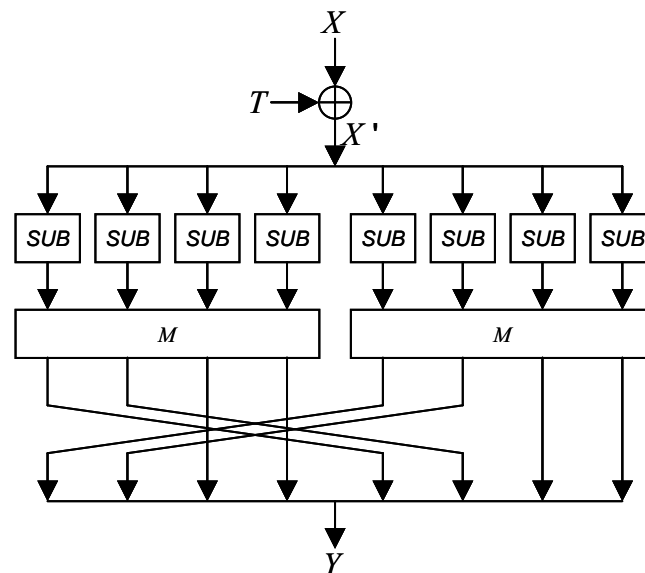


Figure 7 — function  $F$  of MUGI

### 8.1.8 Function $S_R$

The function  $S_R$  is the internal function of  $F$ . The function  $S_R$  can be described by using a substitution table. In this case, the function  $S_R$  is as follows:



Input: 8-bit string  $x$ .

Output: 8-bit string  $y$ .

— Set  $y = SUB[x]$ .

— Output  $y$ .

The  $SUB$  used in function  $S_R$  is a substitution as follows:

$SUB[256] = \{$

0x63, 0x7c, 0x77, 0x7b, 0xf2, 0x6b, 0x6f, 0xc5, 0x30, 0x01, 0x67, 0x2b, 0xfe, 0xd7, 0xab, 0x76,  
0xca, 0x82, 0xc9, 0x7d, 0xfa, 0x59, 0x47, 0xf0, 0xad, 0xd4, 0xa2, 0xaf, 0x9c, 0xa4, 0x72, 0xc0,  
0xb7, 0xfd, 0x93, 0x26, 0x36, 0x3f, 0xf7, 0xcc, 0x34, 0xa5, 0xe5, 0xf1, 0x71, 0xd8, 0x31, 0x15,  
0x04, 0xc7, 0x23, 0xc3, 0x18, 0x96, 0x05, 0x9a, 0x07, 0x12, 0x80, 0xe2, 0xeb, 0x27, 0xb2, 0x75,  
0x09, 0x83, 0x2c, 0x1a, 0x1b, 0x6e, 0x5a, 0xa0, 0x52, 0x3b, 0xd6, 0xb3, 0x29, 0xe3, 0x2f, 0x84,  
0x53, 0xd1, 0x00, 0xed, 0x20, 0xfc, 0xb1, 0x5b, 0x6a, 0xcb, 0xbe, 0x39, 0x4a, 0x4c, 0x58, 0xcf,  
0xd0, 0xef, 0xaa, 0xfb, 0x43, 0x4d, 0x33, 0x85, 0x45, 0xf9, 0x02, 0x7f, 0x50, 0x3c, 0x9f, 0xa8,  
0x51, 0xa3, 0x40, 0x8f, 0x92, 0x9d, 0x38, 0xf5, 0xbc, 0xb6, 0xda, 0x21, 0x10, 0xff, 0xf3, 0xd2,  
0xcd, 0x0c, 0x13, 0xec, 0x5f, 0x97, 0x44, 0x17, 0xc4, 0xa7, 0x7e, 0x3d, 0x64, 0x5d, 0x19, 0x73,  
0x60, 0x81, 0x4f, 0xdc, 0x22, 0x2a, 0x90, 0x88, 0x46, 0xee, 0xb8, 0x14, 0xde, 0x5e, 0x0b, 0xdb,  
0xe0, 0x32, 0x3a, 0x0a, 0x49, 0x06, 0x24, 0x5c, 0xc2, 0xd3, 0xac, 0x62, 0x91, 0x95, 0xe4, 0x79,  
0xe7, 0xc8, 0x37, 0x6d, 0x8d, 0xd5, 0x4e, 0xa9, 0x6c, 0x56, 0xf4, 0xea, 0x65, 0x7a, 0xae, 0x08,  
0xba, 0x78, 0x25, 0x2e, 0x1c, 0xa6, 0xb4, 0xc6, 0xe8, 0xdd, 0x74, 0x1f, 0x4b, 0xbd, 0x8b, 0x8a,  
0x70, 0x3e, 0xb5, 0x66, 0x48, 0x03, 0xf6, 0x0e, 0x61, 0x35, 0x57, 0xb9, 0x86, 0xc1, 0x1d, 0x9e,  
0xe1, 0xf8, 0x98, 0x11, 0x69, 0xd9, 0x8e, 0x94, 0x9b, 0x1e, 0x87, 0xe9, 0xce, 0x55, 0x28, 0xdf,  
0x8c, 0xa1, 0x89, 0x0d, 0xbf, 0xe6, 0x42, 0x68, 0x41, 0x99, 0x2d, 0x0f, 0xb0, 0x54, 0xbb, 0x16}

### 8.1.9 Function $M$

The function  $M$  is the internal function of the  $F$  function. The function  $M$  is as follows:

Input: 32-bit string  $X$ .

Output: 32-bit string  $Y$ .

— Set  $(x_0, x_1, x_2, x_3) = X$ , where  $x_i$  is an 8-bit string and an element of  $GF(2^8)$ .

— Set

$$\begin{pmatrix} y_0 \\ y_1 \\ y_2 \\ y_3 \end{pmatrix} = \begin{pmatrix} 0x02 & 0x03 & 0x01 & 0x01 \\ 0x01 & 0x02 & 0x03 & 0x01 \\ 0x01 & 0x01 & 0x02 & 0x03 \\ 0x03 & 0x01 & 0x01 & 0x02 \end{pmatrix} \begin{pmatrix} x_0 \\ x_1 \\ x_2 \\ x_3 \end{pmatrix},$$

where 0x01, 0x02, and 0x03 are the hexadecimal expressions of the elements of  $GF(2^8)$ .

— Set  $Y = y_0 \parallel y_1 \parallel y_2 \parallel y_3$ .

— Output  $Y$ .

## 8.2 SNOW 2.0 keystream generator

### 8.2.1 Introduction to SNOW 2.0

SNOW 2.0, in the sequel simply denoted SNOW, is a keystream generator which uses as input a 128 or 256-bit secret key  $K$ , and a 128-bit initialization vector  $IV$ . These are used to initiate a state variable  $S_i$  ( $i \geq 0$ )

consisting of eighteen  $n = 32$  bit blocks. Bit/byte order is big-endian, i.e., if the key and initialization vector are given as a sequence of bits/bytes, the first/leftmost bit/byte is the most significant of the corresponding data. For every iteration of the *Strm* function, a 32-bit keystream  $Z_i$  is produced as output.

SNOW's state variable  $S_i$  consists of two components. First, 16 32-bit variables:

$$a^{(i)} = (a_{15}^{(i)}, a_{14}^{(i)}, \dots, a_0^{(i)}),$$

implements a linear feedback shift register (LFSR). Secondly, 2 32-bit variables:

$$b^{(i)} = (b_2^{(i)}, b_1^{(i)}),$$

maintains the state of a finite state machine (FSM). SNOW is best understood with reference to Figure 8, which shows a snapshot, at time  $i$ , omitting the time dependence variable  $(i)$  from the notation.

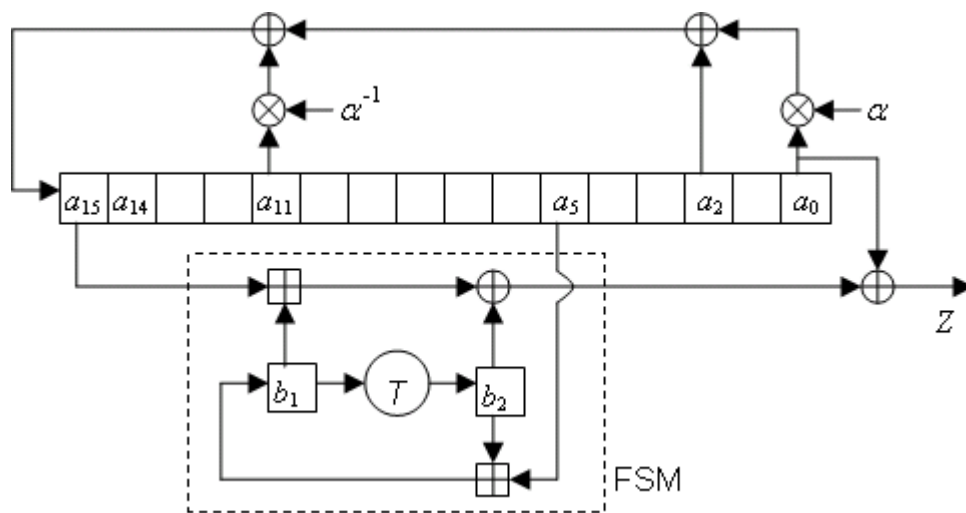


Figure 8 — Schematic drawing of SNOW

SNOW operation is defined by:

The *Init* function, defined in detail in 8.2.2, takes as input the 128 or 256-bit key  $K$  and the 128-bit  $IV$ , and produces the initial value of the state variable  $S_0 = (a^{(0)}, b^{(0)})$ .

The *Next* function, defined in detail in 8.2.3, takes as input the 18 32-bit state variable  $S_i = (a^{(i)}, b^{(i)})$  and produces as output the next value of the state variable  $S_{i+1} = (a^{(i+1)}, b^{(i+1)})$ . The *Next* function runs in two modes, depending on whether the iteration performed is part of the initialization, or, of the normal mode of generating output, see below.

The *Strm* function, defined in detail in 8.2.4, takes as input the 18 32-bit state variable  $S_i = (a^{(i)}, b^{(i)})$  and produces as output the 32-bit keystream  $Z_i$ .

NOTE 1 For SNOW, the maximum recommended amount of keystream produced from a given  $(K, IV)$  is  $2^{50}$  32 bits. This bound has been selected to provide good security margin against cryptanalysis, and implies no practical limitation in applicability of the algorithm.

NOTE 2 The paper [10] is referred for theoretical background on the design rationale for SNOW.

### 8.2.2 Initialization function *Init*

The Initialization function *Init* is as follows.

Input: 128- or 256-bit key  $K$ , 128-bit initialization vector  $IV$ .

Output: Initial value of state variable  $S_0 = (a^{(0)}, b^{(0)})$ .

a) Initialize the registers by the key information

- For a 128-bit key, set  $(K_3, K_2, K_1, K_0) = K$ ,  $a_{15-j}^{(-34)} = a_{15-j-8}^{(-34)} = K_{3-j}$ , and  $a_{15-j-4}^{(-34)} = a_{15-j-12}^{(-34)} = \neg K_{3-j}$  for  $j = 0, 1, 2, 3$ .
- For a 256-bit key, set  $(K_7, K_6, \dots, K_0) = K$ ,  $a_{15-j}^{(-34)} = K_{7-j}$ , and  $a_{15-j-8}^{(-34)} = \neg(K_{7-j})$  for  $j = 0, 1, \dots, 7$ .

b) Set  $S_{-33} = (a^{(-33)}, b^{(-33)})$  by:

- Set  $(IV_3, IV_2, IV_1, IV_0) = IV$ .
- Set  $a_i^{(-33)} = a_i^{(-34)}$  for  $i = 0, 1, 2, 3, 4, 5, 6, 7, 8, 11, 13, 14$ .
- Set  $a_{15}^{(-33)} = a_{15}^{(-34)} \oplus IV_0$ ;  $a_{12}^{(-33)} = a_{12}^{(-34)} \oplus IV_1$ ;  $a_{10}^{(-33)} = a_{10}^{(-34)} \oplus IV_2$ ;  $a_9^{(-33)} = a_9^{(-34)} \oplus IV_3$ .
- Set  $b_1^{(-33)} = b_2^{(-33)} = 0^{(32)}$ .

c) Set  $S_{-1} = \text{Next}^{32}(S_{-33}, \text{INIT})$ , where  $\text{Next}^{32}$  denotes 32 iterations of the *Next* function.

d)  $S_0 = \text{Next}(S_{-1})$ .

e) Output  $S_0$ .

### 8.2.3 Next-state function *Next*

SNOW has two modes for *Next* function.

Input: State variable  $S_i = (a^{(i)}, b^{(i)})$ , mode = {INIT, null}.

Output: Next value of the state variable  $S_{i+1} = (a^{(i+1)}, b^{(i+1)})$ .

- a) Set  $b_2^{(i+1)} = T(b_1^{(i)})$ .
- b) Set  $b_1^{(i+1)} = b_2^{(i)} +_{32} a_5^{(i)}$ .
- c) For  $j = 0, 1, \dots, 14$  set  $a_j^{(i+1)} = a_{j+1}^{(i)}$
- d) If INIT mode, set  $a_{15}^{(i+1)} = (a_0^{(i)} \otimes \alpha) \oplus a_2^{(i)} \oplus (a_{11}^{(i)} \oplus \alpha^{-1}) \oplus \text{FSM}(a_{15}^{(i)}, b_1^{(i)}, b_2^{(i)})$ . Otherwise, set  $a_{15}^{(i+1)} = (a_0^{(i)} \otimes \alpha) \oplus a_2^{(i)} \oplus (a_{11}^{(i)} \otimes \alpha^{-1})$ .
- e)  $S_{i+1} = (a^{(i+1)}, b^{(i+1)})$ .
- f) Output  $S_{i+1}$ .

The description of the  $T$  function and the finite field arithmetic involving the fixed element  $\alpha$  refers to 8.2.5 and 8.2.6, respectively.

NOTE Figure 9 shows the block diagram of the INIT mode of *Next* function.

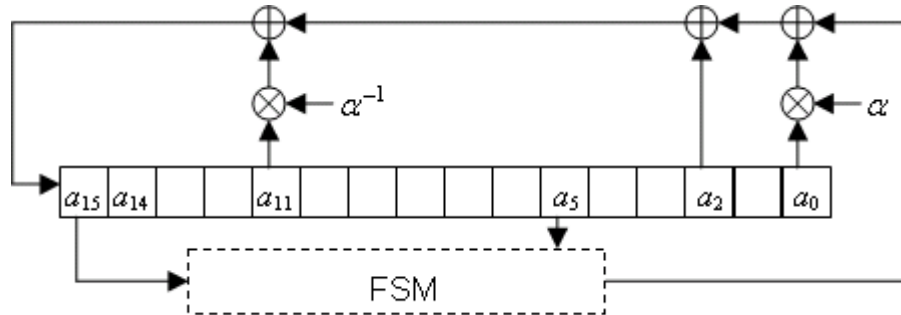


Figure 9 — INIT mode of *Next* function

The definition of the *FSM* function refers to 8.2.8.

#### 8.2.4 Keystream function *Strm*

The keystream function *Strm* is as follows:

Input: State variable  $S_i$ .

Output: 32-bit keystream  $Z_i$ .

- Set  $Z_i = FSM(a_{15}^{(i)}, b_1^{(i)}, b_2^{(i)}) \oplus a_0^{(i)}$ .
- Output  $Z_i$ .

#### 8.2.5 Function $T$

The  $T$  function is a substitution, specifically a permutation of  $GF(2^{32})$ , based on components from the Advanced Encryption Standard (AES), ISO/IEC 18033-3. To this end, the finite field  $GF(2^8)$  is used, which is viewed as  $GF(2)[x]$  modulo the irreducible polynomial

$$f(x) = x^8 + x^4 + x^3 + x + 1.$$

and the polynomial ring  $GF(2^8)[y]$  modulo  $(y^4 + 1)$ .

Input: 32-bit string  $w$ .

Output: 32-bit string  $q = T(w)$ .

- Set  $(w_3, w_2, w_1, w_0) = w$ , where each  $w_j$  is 8 bit.
- For  $j = 0, 1, 2, 3$  set  $t_j = SUB[w_j]$ .
- Let  $t(y)$  be the polynomial  $t(y) = t_3 y^3 + t_2 y^2 + t_1 y + t_0$  in  $GF(2^8)[y]$ , where  $t_j$  is interpreted as an element of  $GF(2^8)$  in the natural way:  $t_j = t_{j,7} x^7 + \dots + t_{j,1} x + t_{j,0}$ ,  $t_{j,k}$  in  $GF(2)$ .
- set  $q(y) = c(y) \cdot t(y)$  modulo  $(y^4 + 1)$ , where  $c(y) = (x+1)y^3 + y^2 + y + x$  in  $GF(2^8)[y]$ .

e) associate the 32-bit string  $q = (q_3, q_2, q_1, q_0)$  with the result of the above,  $q(y) = q_3 y^3 + q_2 y^2 + q_1 y + q_0$ .

f) Output  $q$ .

Note that in step c), the two polynomials are multiplied where coefficient-by-coefficient operations are carried out in  $GF(2^8)$  as defined by  $f(x)$  above. The result is then reduced modulo  $y^4 + 1$ .

NOTE 1 The AES S-box can be found in 8.1.8.

NOTE 2 The paper [10] is referred for details of this (and other) optimisation(s).

## 8.2.6 Multiplication of $\alpha$ in finite field arithmetic

Input: 32-bit string  $w$ , representing an element of  $GF(2^{32})$ .

Output: 32-bit string  $w'$ , representing  $\alpha \otimes w$  in  $GF(2^{32})$ .

a) Set  $w' = (w \ll_{32} 8) \oplus \alpha_{\text{MUL}}[w \gg_{32} 24]$ .

b) Output  $w'$ .

The function  $\alpha_{\text{MUL}}$  is defined in the following:

```
 $\alpha_{\text{MUL}} [256] = \{$ 
0x00000000, 0xE19FCF13, 0x6B973726, 0x8A08F835, 0xD6876E4C, 0x3718A15F, 0xBD10596A, 0x5C8F9679,
0x05A7DC98, 0xE438138B, 0x6E30EBBE, 0x8FAF24AD, 0xD320B2D4, 0x32BF7DC7, 0xB8B785F2, 0x59284AE1,
0x0AE71199, 0xEB78DE8A, 0x617026BF, 0x80EFE9AC, 0xDC607FD5, 0x3DFFB0C6, 0xB7F748F3, 0x566887E0,
0x0F40CD01, 0xEEDF0212, 0x64D7FA27, 0x85483534, 0xD9C7A34D, 0x38586C5E, 0xB250946B, 0x53CF5B78,
0x1467229B, 0xF5F8ED88, 0x7FF015BD, 0x9E6FDAEE, 0xC2E04CD7, 0x237F83C4, 0xA9777BF1, 0x48E8B4E2,
0x11C0FE03, 0xF05F3110, 0x7A57C925, 0x9BC80636, 0xC747904F, 0x26D85F5C, 0xADC0A769, 0x4D4F687A,
0x1E803302, 0xFF1FFC11, 0x75170424, 0x9488CB37, 0xC8075D4E, 0x2998925D, 0xA3906A68, 0x420FA57B,
0x1B27EF9A, 0xFAB82089, 0x70B0D8BC, 0x912F17AF, 0xCDA081D6, 0x2C3F4EC5, 0xA637B6F0, 0x47A879E3,
0x28CE449F, 0xC9518B8C, 0x435973B9, 0xA2C6BCAA, 0xFE492AD3, 0x1FD6E5C0, 0x95DE1DF5, 0x7441D2E6,
0x2D699807, 0xCCF65714, 0x46FEAF21, 0xA7616032, 0xFBEEF64B, 0x1A713958, 0x9079C16D, 0x71E60E7E,
0x22295506, 0xC3B69A15, 0x49BE6220, 0xA821AD33, 0xF4AE3B4A, 0x1531F459, 0x9F390C6C, 0x7EA6C37F,
0x278E899E, 0xC611468D, 0x4C19BEB8, 0xAD8671AB, 0xF109E7D2, 0x109628C1, 0x9A9ED0F4, 0x7B011FE7,
0x3CA96604, 0xDD36A917, 0x573E5122, 0xB6A19E31, 0xEA2E0848, 0x0BB1C75B, 0x81B93F6E, 0x6026F07D,
0x390EBA9C, 0xD891758F, 0x52998DBA, 0xB30642A9, 0xEF89D4D0, 0x0E161BC3, 0x841EE3F6, 0x65812CE5,
0x364E779D, 0xD7D1B88E, 0x5DD940BB, 0xBC468FA8, 0xE0C919D1, 0x0156D6C2, 0x8B5E2EF7, 0x6AC1E1E4,
0x33E9AB05, 0xD2766416, 0x587E9C23, 0xB9E15330, 0xE56EC549, 0x04F10A5A, 0x8EF9F26F, 0x6F663D7C,
0x50358897, 0xB1AA4784, 0x3BA2BFB1, 0xDA3D70A2, 0x86B2E6DB, 0x672D29C8, 0xED25D1FD, 0x0CBA1EEE,
0x5592540F, 0xB40D9B1C, 0x3E056329, 0xDF9AAC3A, 0x83153A43, 0x628AF550, 0xE8820D65, 0x091DC276,
0x5AD2990E, 0xBB4D561D, 0x3145AE28, 0xD0DA613B, 0x8C55F742, 0x6DCA3851, 0xE7C2C064, 0x065D0F77,
0x5F754596, 0xBEEA8A85, 0x34E272B0, 0xD57DBDA3, 0x89F22BDA, 0x686DE4C9, 0xE2651CFC, 0x03FAD3EF,
0x4452AA0C, 0xA5CD651F, 0x2FC59D2A, 0xCE5A5239, 0x92D5C440, 0x734A0B53, 0xF942F366, 0x18DD3C75,
0x41F57694, 0xA06AB987, 0x2A6241B2, 0xCBFD8EA1, 0x977218D8, 0x76EDD7CB, 0xFCE52FFE, 0x1D7AE0ED,
0x4EB5BB95, 0xAF2A7486, 0x25228CB3, 0xC4BD43A0, 0x9832D5D9, 0x79AD1ACA, 0xF3A5E2FF, 0x123A2DEC,
0x4B12670D, 0xAA8DA81E, 0x2085502B, 0xC11A9F38, 0x9D950941, 0x7C0AC652, 0xF6023E67, 0x179DF174,
0x78FBCC08, 0x9964031B, 0x136CFB2E, 0xF2F3343D, 0xAE7CA244, 0x4FE36D57, 0xC5EB9562, 0x24745A71,
0x7D5C1090, 0x9CC3DF83, 0x16CB27B6, 0xF754E8A5, 0xABDB7EDC, 0x4A44B1CF, 0xC04C49FA, 0x21D386E9,
0x721CDD91, 0x93831282, 0x198BEAB7, 0xF81425A4, 0xA49BB3DD, 0x45047CCE, 0xCF0C84FB, 0x2E934BE8,
0x77BB0109, 0x9624CE1A, 0x1C2C362F, 0xFDB3F93C, 0xA13C6F45, 0x40A3A056, 0xCAAB5863, 0x2B349770,
0x6C9CEE93, 0x8D032180, 0x070BD9B5, 0xE69416A6, 0xBA1B80DF, 0x5B844FCC, 0xD18CB7F9, 0x301378EA,
0x693B320B, 0x88A4FD18, 0x02AC052D, 0xE333CA3E, 0xBFBC5C47, 0x5E239354, 0xD42B6B61, 0x35B4A472,
0x667BFF0A, 0x87E43019, 0x0DECC82C, 0xEC73073F, 0xB0FC9146, 0x51635E55, 0xDB6BA660, 0x3AF46973,
0x63DC2392, 0x8243EC81, 0x084B14B4, 0xE9D4DBA7, 0xB55B4DDE, 0x54C482CD, 0xDECC7AF8, 0x3F53B5EB};
```

## 8.2.7 Multiplication of $\alpha^{-1}$ in finite field arithmetic

Input: 32-bit string  $y$ , representing an element of  $GF(2^{32})$ .

Output: 32-bit string  $y$ , representing  $\alpha^{-1} \otimes y$  in  $GF(2^{32})$ .

a) Set  $y' = (y \gg_{32} 8) \otimes \alpha_{\text{inv\_MUL}}[y \bmod 256]$ .

b) Output  $y'$ .

The function  $\alpha_{\text{inv\_MUL}}$  is defined in the following:

```
 $\alpha_{\text{inv\_MUL}}[256] = \{$ 
0x00000000, 0x180F40CD, 0x301E8033, 0x2811C0FE, 0x603CA966, 0x7833E9AB, 0x50222955, 0x482D6998,
0xC078FBCC, 0xD877BB01, 0xF0667BFF, 0xE8693B32, 0xA04452AA, 0xB84B1267, 0x905AD299, 0x88559254,
0x29F05F31, 0x31FF1FFC, 0x19EEDF02, 0x01E19FCF, 0x49CCF657, 0x51C3B69A, 0x79D27664, 0x61DD36A9,
0xE988A4FD, 0xF187E430, 0xD99624CE, 0xC1996403, 0x89B40D9B, 0x91BB4D56, 0xB9AA8DA8, 0xA1A5CD65,
0x5249BE62, 0x4A46FEAF, 0x62573E51, 0x7A587E9C, 0x32751704, 0x2A7A57C9, 0x026B9737, 0x1A64D7FA,
0x923145AE, 0x8A3E0563, 0xA22FC59D, 0xBA208550, 0xF20DECC8, 0xEA02AC05, 0xC2136CFB, 0xDA1C2C36,
0x7BB9E153, 0x63B6A19E, 0x4BA76160, 0x53A821AD, 0x1B854835, 0x038A08F8, 0x2B9BC806, 0x339488CB,
0xBBC11A9F, 0xA3CE5A52, 0x8BDF9AAC, 0x93D0DA61, 0xDBFDB3F9, 0xC3F2F334, 0xEBE333CA, 0xF3EC7307,
0xA492D5C4, 0xBC9D9509, 0x948C55F7, 0x8C83153A, 0xC4AE7CA2, 0xDCA13C6F, 0xF4B0FC91, 0xECBFC5C,
0x64EA2E08, 0x7CE56EC5, 0x54F4AE3B, 0x4CFBEEF6, 0x04D6876E, 0x1CD9C7A3, 0x34C8075D, 0x2CC74790,
0x8D628AF5, 0x956DCA38, 0xBD7C0AC6, 0xA5734A0B, 0xED5E2393, 0xF551635E, 0xDD40A3A0, 0xC54FE36D,
0x4D1A7139, 0x551531F4, 0x7D04F10A, 0x650BB1C7, 0x2D26D85F, 0x35299892, 0x1D38586C, 0x053718A1,
0xF6DB6BA6, 0xEED42B6B, 0xC6C5EB95, 0xDECAB58, 0x96E7C2C0, 0x8EE8820D, 0xA6F942F3, 0xBEF6023E,
0x36A3906A, 0x2EACD0A7, 0x06BD1059, 0x1EB25094, 0x569F390C, 0x4E9079C1, 0x6681B93F, 0x7E8EF9F2,
0xDF2B3497, 0xC724745A, 0xEF35B4A4, 0xF73AF469, 0xBF179DF1, 0xA718DD3C, 0x8F091DC2, 0x97065D0F,
0x1F53CF5B, 0x075C8F96, 0x2F4D4F68, 0x37420FA5, 0x7F6F663D, 0x676026F0, 0x4F71E60E, 0x577EA6C3,
0xE18D0321, 0xF98243EC, 0xD1938312, 0xC99CC3DF, 0x81B1AA47, 0x99BEEA8A, 0xB1AF2A74, 0xA9A06AB9,
0x21F5F8ED, 0x39FAB820, 0x11EB78DE, 0x09E43813, 0x41C9518B, 0x59C61146, 0x71D7D1B8, 0x69D89175,
0xC87D5C10, 0xD0721CDD, 0xF863DC23, 0xE06C9CEE, 0xA841F576, 0xB04EB5BB, 0x985F7545, 0x80503588,
0x0805A7DC, 0x100AE711, 0x381B27EF, 0x20146722, 0x68390EBA, 0x70364E77, 0x58278E89, 0x4028CE44,
0xB3C4BD43, 0xABCBD8E, 0x83DA3D70, 0x9BD57DBD, 0xD3F81425, 0xCBF754E8, 0xE3E69416, 0xFBE9D4DB,
0x73BC468F, 0x6BB30642, 0x43A2C6BC, 0x5BAD8671, 0x1380EFE9, 0x0B8FAF24, 0x239E6FDA, 0x3B912F17,
0x9A34E272, 0x823BA2BF, 0xAA2A6241, 0xB225228C, 0xFA084B14, 0xE2070BD9, 0xCA16CB27, 0xD2198BEA,
0x5A4C19BE, 0x42435973, 0x6A52998D, 0x725DD940, 0x3A70B0D8, 0x227FF015, 0xA6E30EB, 0x12617026,
0x451FD6E5, 0x5D109628, 0x750156D6, 0x6D0E161B, 0x25237F83, 0x3D2C3F4E, 0x153DFFB0, 0xD32BF7D,
0x85672D29, 0x9D68DE4, 0xB579AD1A, 0xAD76EDD7, 0xE55B844F, 0xFD54C482, 0xD545047C, 0xCD4A44B1,
0x6CEF89D4, 0x74E0C919, 0x5CF109E7, 0x44FE492A, 0x0CD320B2, 0x14DC607F, 0x3CCDA081, 0x24C2E04C,
0xAC977218, 0xB49832D5, 0x9C89F22B, 0x8486B2E6, 0xCCABDB7E, 0xD4A49BB3, 0xFCB55B4D, 0xE4BA1B80,
0x17566887, 0x0F59284A, 0x2748E8B4, 0x3F47A879, 0x776AC1E1, 0x6F65812C, 0x477441D2, 0x5F7B011F,
0xD72E934B, 0xCF21D386, 0xE7301378, 0xFF3F53B5, 0xB7123A2D, 0xAF1D7AE0, 0x870CBA1E, 0x9F03FAD3,
0x3EA637B6, 0x26A9777B, 0x0EB8B785, 0x16B7F748, 0x5E9A9ED0, 0x4695DE1D, 0xE841EE3, 0x768B5E2E,
0xFEDECC7A, 0xE6D18CB7, 0xCEC04C49, 0xD6CF0C84, 0x9EE2651C, 0x86ED25D1, 0xAEFCE52F, 0xB6F3A5E2};
```

## 8.2.8 Function $FSM(x, y, z)$

Input: Three 32-bit strings,  $x$ ,  $y$ , and  $z$ .

Output: 32-bit string  $q$ .

a) Set  $q = (x +_{32} y) \oplus z$ .

b) Output  $q$ .

## 8.3 Rabbit keystream generator

### 8.3.1 Introduction to Rabbit

Rabbit is a keystream generator which uses a 128-bit secret key  $K$ , a 64-bit initialization vector  $IV$ , and a 513-bit internal state variable  $S_i$  ( $i \geq 0$ ). It outputs a 128-bit keystream block  $Z_i$  at every iteration of the function  $Strm$ .

The 513 bits of the internal state  $S_i$  are divided between eight 32-bit state variables  $X_0^{(i)}, \dots, X_7^{(i)}$ , eight 32-bit counter variables  $C_0^{(i)}, \dots, C_7^{(i)}$ , and one counter carry bit  $b^{(i)}$ .

The description uses the notation laid out in Clause 4 of this part of ISO/IEC 18033. In addition, a special notation for bit arrays is used to enhance readability: When labeling the bits of a variable  $A$ , the least significant bit is denoted by  $A^{(0)}$ . The notation  $A^{[h:g]}$  represents bits  $h$  through  $g$  of variable  $A$ , where bit position  $h$  is more significant than bit position  $g$ .

NOTE 1 For Rabbit, the maximum recommended amount of keystream produced from a given key  $K$  is  $2^{64}$  keystream blocks. This provides a large security margin against cryptanalysis, while at the same time implying no practical limitations on the applicability of the algorithm.

NOTE 2 The paper [8] is referred for the original proposal of the cipher and the paper [9] is referred for an overview of its cryptographic security.

### 8.3.2 Additional variables and notation

For the Rabbit keystream generator, the following notation is added:

$A$	Constant for Rabbit
$b$	Carry bit for Rabbit
$C$	Counter variable for Rabbit
$g$	Subfunction used for Rabbit
$X$	inner state variable for Rabbit

In addition, a number of other symbols are used for auxiliary local variables in algorithm descriptions. These symbols occur only within a given function specification and do not have a global meaning. They are thus described in the function declaration.

### 8.3.3 Initialization function *Init*

In the following, the initialization function *Init* of Rabbit is specified.

Input: 128-bit key  $K$ , 64-bit initialization vector  $IV$ .

Output: Initial value of the state variable  $S_0 = (b^{(0)}, X_0^{(0)}, \dots, X_7^{(0)}, C_0^{(0)}, \dots, C_7^{(0)})$ .

Local variables: counters  $i, j$

- a) Let  $K_0 = K^{[15..0]}$ ,  $K_1 = K^{[31..16]}$ , ..., and  $K_7 = K^{[127..112]}$ .
- b) Set  $S_{-9}$  as follows:
  - 1) Set  $b^{(-9)} = 0$ .
  - 2) For  $j = 0, 1, \dots, 7$ :
    - If  $j$  is even, set  $X_j^{(-9)} = K_{(j+1 \bmod 8)} \parallel K_j$  and  $C_j^{(-9)} = K_{(j+4 \bmod 8)} \parallel K_{(j+5 \bmod 8)}$ .
    - Else, set  $X_j^{(-9)} = K_{(j+5 \bmod 8)} \parallel K_{(j+4 \bmod 8)}$  and  $C_j^{(-9)} = K_j \parallel K_{(j+1 \bmod 8)}$ .
- c) Iterate the next-state function *Next* four times: set  $S_i = \text{Next}(S_{i-1})$  for  $i = -8, -7, -6, -5$ .
- d) Set  $S_{-4}$  as follows:
  - 1) Modify the counters as follows:

$C_0^{(-4)} = C_0^{(-5)} \oplus X_4^{(-5)} \oplus IV^{[31..0]}$	$C_1^{(-4)} = C_1^{(-5)} \oplus X_5^{(-5)} \oplus (IV^{[63..48]} \parallel IV^{[31..16]})$
$C_2^{(-4)} = C_2^{(-5)} \oplus X_6^{(-5)} \oplus IV^{[63..32]}$	$C_3^{(-4)} = C_3^{(-5)} \oplus X_7^{(-5)} \oplus (IV^{[47..32]} \parallel IV^{[15..0]})$
$C_4^{(-4)} = C_4^{(-5)} \oplus X_0^{(-5)} \oplus IV^{[31..0]}$	$C_5^{(-4)} = C_5^{(-5)} \oplus X_1^{(-5)} \oplus (IV^{[63..48]} \parallel IV^{[31..16]})$
$C_6^{(-4)} = C_6^{(-5)} \oplus X_2^{(-5)} \oplus IV^{[63..32]}$	$C_7^{(-4)} = C_7^{(-5)} \oplus X_3^{(-5)} \oplus (IV^{[47..32]} \parallel IV^{[15..0]})$

2) Set  $X_0^{(-4)} = X_0^{(-5)}, \dots, X_7^{(-4)} = X_7^{(-5)}, b^{(-4)} = b^{(-5)}$ .

e) Iterate the next-state function *Next* four times: set  $S_i = \text{Next}(S_{i-1})$  for  $i = -3, -2, -1, 0$ .

f) Output  $S_0 = (b^{(0)}, X_0^{(0)}, \dots, X_7^{(0)}, C_0^{(0)}, \dots, C_7^{(0)})$ .

NOTE The *IV* is mixed into the internal state in steps d) and e) of the algorithm. If the application requires frequent re-initialization under the same key, it makes sense to store the internal state after step c) as master state and to perform only steps d) through f) for re-initialization.

#### 8.3.4 Next-state function *Next*

The next-state function *Next* of Rabbit is specified as follows:

Input: State variable  $S_i = (b^{(i)}, X_0^{(i)}, \dots, X_7^{(i)}, C_0^{(i)}, \dots, C_7^{(i)})$ .

Output: State variable  $S_{i+1} = (b^{(i+1)}, X_0^{(i+1)}, \dots, X_7^{(i+1)}, C_0^{(i+1)}, \dots, C_7^{(i+1)})$ .

Local variables: counter  $j$ , 33-bit positive integer *temp*

a) Set constants  $A_0, \dots, A_7$  as follows:

$$\begin{array}{llll} A_0 = 0 \times 4D34D34D & A_1 = 0 \times D34D34D3 & A_2 = 0 \times 34D34D34 & A_3 = 0 \times 4D34D34D \\ A_4 = 0 \times D34D34D3 & A_5 = 0 \times 34D34D34 & A_6 = 0 \times 4D34D34D & A_7 = 0 \times D34D34D3 \end{array}$$

b) Let  $b_0^{(i+1)} = b^{(i)}$

c) For  $j = 0, 1, \dots, 7$ :

— Let  $\text{temp} = C_j^{(i)} + A_j + b_j^{(i+1)}$ ; this results in a 33-bit value.

— Let  $b_{j+1}^{(i+1)} = \text{temp}^{[32]}$ .

— Let  $C_j^{(i+1)} = \text{temp}^{[31..0]}$ .

d) Let  $b^{(i+1)} = b_8^{(i+1)}$

e) For  $j = 0, 1, \dots, 7$ , let  $G_j = g(X_j^{(i)}, C_j^{(i+1)})$ , where the function  $g$  is given in 8.3.6.

f) Modify internal state as follows:

$$X_0^{(i+1)} = G_0 +_{32} (G_7 \lll_{32} 16) +_{32} (G_6 \lll_{32} 16)$$

$$X_1^{(i+1)} = G_1 +_{32} (G_0 \lll_{32} 8) +_{32} G_7$$

$$X_2^{(i+1)} = G_2 +_{32} (G_1 \lll_{32} 16) +_{32} (G_0 \lll_{32} 16)$$

$$X_3^{(i+1)} = G_3 +_{32} (G_2 \lll_{32} 8) +_{32} G_1$$

$$X_4^{(i+1)} = G_4 +_{32} (G_3 \lll_{32} 16) +_{32} (G_2 \lll_{32} 16)$$

$$X_5^{(i+1)} = G_5 +_{32} (G_4 \lll_{32} 8) +_{32} G_3$$



$$X_6^{(i+1)} = G_6 +_{32} (G_5 \ll_{32} 16) +_{32} (G_4 \ll_{32} 16)$$

$$X_7^{(i+1)} = G_7 +_{32} (G_6 \ll_{32} 8) +_{32} G_5$$

g) Output  $S_{i+1} = (b^{(i+1)}, X_0^{(i+1)}, \dots, X_7^{(i+1)}, C_0^{(i+1)}, \dots, C_7^{(i+1)})$ .

### 8.3.5 Keystream function *Strm*

The keystream function *Strm* of Rabbit is specified as follows:

Input: State variable  $S_i = (b^{(i)}, X_0^{(i)}, \dots, X_7^{(i)}, C_0^{(i)}, \dots, C_7^{(i)})$ .

Output: Keystream block  $Z_i$ .

a) Set  $Z_i$  as follows:

$$Z_i^{[15..0]} = X_0^{(i)[15..0]} \oplus X_5^{(i)[31..16]}$$

$$Z_i^{[31..16]} = X_0^{(i)[31..16]} \oplus X_3^{(i)[15..0]}$$

$$Z_i^{[47..32]} = X_2^{(i)[15..0]} \oplus X_7^{(i)[31..16]}$$

$$Z_i^{[63..48]} = X_2^{(i)[31..16]} \oplus X_5^{(i)[15..0]}$$

$$Z_i^{[79..64]} = X_4^{(i)[15..0]} \oplus X_1^{(i)[31..16]}$$

$$Z_i^{[95..80]} = X_4^{(i)[31..16]} \oplus X_7^{(i)[15..0]}$$

$$Z_i^{[111..96]} = X_6^{(i)[15..0]} \oplus X_3^{(i)[31..16]}$$

$$Z_i^{[127..112]} = X_6^{(i)[31..16]} \oplus X_1^{(i)[15..0]}$$

b) Output  $Z_i$ .

### 8.3.6 Function *g*

The function *g* is specified as follows:

Input: Two 32-bit parameters  $u$  and  $v$ .

Output: 32-bit result  $g(u, v)$ .

Local variables: 64-bit positive integer *temp*

a) Let  $temp = (u +_{32} v)^2$ ; this results in a 64-bit value.

b) Let  $g(u, v) = temp^{[31..0]} \oplus temp^{[63..32]}$ .

c) Output  $g(u, v)$ .

## 8.4 Decim<sup>v2</sup> keystream generator

### 8.4.1 Introduction to Decim<sup>v2</sup>

DECIM<sup>v2</sup> is a keystream generator which uses an 80-bit secret key  $K$  and a 64-bit initialization vector  $IV$ . DECIM<sup>v2</sup> is composed of a 192-bit maximum length linear feedback shift register  $A$ , filtered by a 14-variable Boolean function  $LF$ . In keystream generation mode, the output of  $LF$  is used to feed a compression block which is a function called  $ABSG$ , whose output finally passes through a 32-bit long buffer  $B$  to regulate the keystream output rate.

DECIM<sup>v2</sup> is described in Figure 10, which shows a snapshot, at time  $i$ , omitting the time dependence variable ( $i$ ) from the notation.

NOTE 1 The paper [6] is referred for theoretical background on the design rationale of DECIM<sup>v2</sup>.

The state variable  $S_i$  of DECIM<sup>v2</sup> consists of the 192-bit value  $a^{(i)} = (a_0^{(i)}, a_1^{(i)}, \dots, a_{191}^{(i)})$  of register  $A$ , a 3-bit variable  $T^{(i)}$  which corresponds to the state of the compression function  $ABSG$ , the 32 bits  $b^{(i)} = (b_0^{(i)}, b_1^{(i)}, \dots, b_{31}^{(i)})$  in buffer  $B$ , and the number  $I^{(i)}$  of bits in buffer  $B$  that are ready to be output.

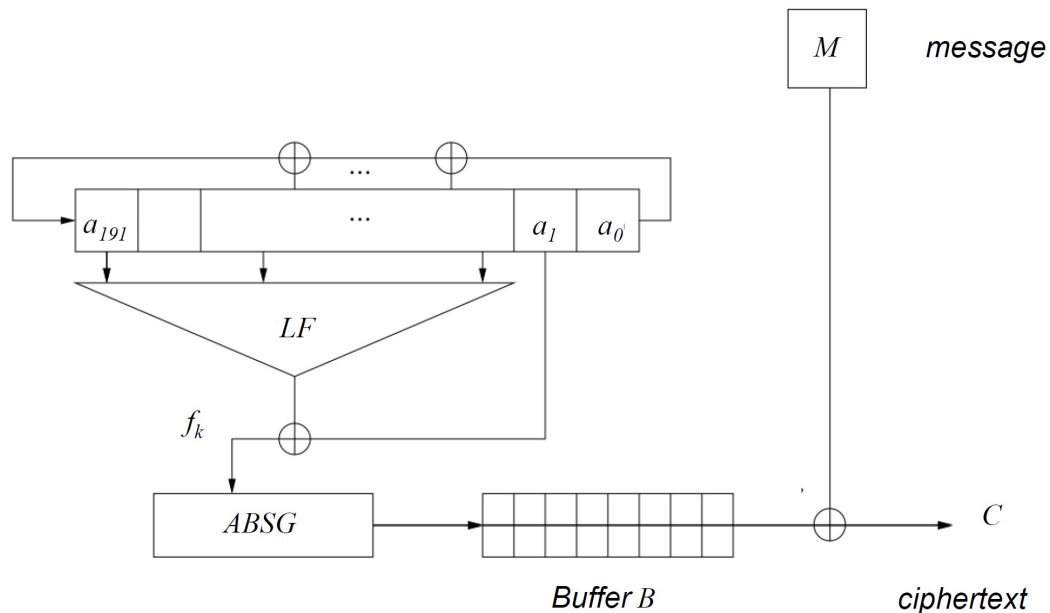


Figure 10 — Schematic drawing of DECIM<sup>v2</sup>

The *Init* function, defined in detail in 8.4.3, takes as input the 80-bit key  $K$  and the 64-bit initialization vector  $IV$ , and produces the initial value of the state variable  $S_0 = (a^{(0)}, T^{(0)}, b^{(0)}, I^{(0)})$ .

The *Next* function, defined in detail in 8.4.5, takes as an input the value of the state variable  $S_i = (a^{(i)}, T^{(i)}, b^{(i)}, I^{(i)})$  and produces as output the next value of the state variable  $S_{i+1} = (a^{(i+1)}, T^{(i+1)}, b^{(i+1)}, I^{(i+1)})$ . The *Next* function runs in three modes, depending on whether the iteration performed is part of the initialization of the register, the initialization of the buffer, or the subsequent keystream generation.

The *Strm* function, defined in detail in 8.4.6, takes as an input the value of the state variable  $S_i = (a^{(i)}, T^{(i)}, b^{(i)}, I^{(i)})$ , and produces as output a keystream bit  $Z_i$ .

NOTE 2 The standard output rate of  $\text{DECIM}^{v2}$  is 1/4. Therefore, in order to synchronize the state variable and the keystream output, the *Next* function performs four standard iterations of  $\text{DECIM}^{v2}$  as specified in [6].

NOTE 3 The compression function of  $\text{DECIM}^{v2}$  has a variable output rate, equal to 1/3 on average. Therefore, a buffer mechanism is used to ensure a constant output rate. The differences between the buffer output rate and the compression function output rate, as well as the buffer length, have been chosen to ensure that the buffer always functions as expected with overwhelming probability, as described in 8.4.3.

NOTE 4  $\text{DECIM}^{v2}$  is immune to the attacks as described in [18].

## 8.4.2 Additional variables and notation

For the  $\text{Decim}^{v2}$  keystream generator, the following notation is added:

$a$	Inner state variable for $\text{Decim}^{v2}$
$ABSG$	Compression function used for $\text{Decim}^{v2}$
$b, b'$	Inner state variables for $\text{Decim}^{v2}$
$B$	Buffering function used for $\text{Decim}^{v2}$
$F$	Linear feedback function used for $\text{Decim}^{v2}$
$I, I'$	Inner state variables for $\text{Decim}^{v2}$
$LF$	Filtering function used for $\text{Decim}^{v2}$
$T, T'$	Inner state variables for $\text{Decim}^{v2}$
$Y$	Boolean function used for $\text{Decim}^{v2}$

In addition, a number of other symbols are used for auxiliary local variables in algorithm descriptions. These symbols occur only within a given function specification and do not have a global meaning. They are thus described in the function declaration.

## 8.4.3 Initialization function *Init*

The Initialization function *Init* is defined as follows.

The Initialization function *Init* is defined as follows.

Input: 80-bit key  $K$ , 64-bit initialization vector  $IV$ .

Output: Initial value of the state variable  $S_0 = (a^{(0)}, T^{(0)}, b^{(0)}, I^{(0)})$ .

Local variables: counters  $i, j$

a) Initialize the register with the key  $K$  and the initialization vector  $IV$ .

- Set  $a_j^{(-256)} = K_j$  for  $j = 0, 1, \dots, 79$ .
- Set  $a_j^{(-256)} = K_{j-80} \oplus IV_{j-80}$  for  $j = 80, 81, \dots, 143$ .
- Set  $a_j^{(-256)} = K_{j-80} \oplus IV_{j-144} \oplus IV_{j-128} \oplus IV_{j-112} \oplus IV_{j-96}$  for  $j = 144, 145, \dots, 159$ .

- Set  $a_j^{(-256)} = IV_{j-160} \oplus IV_{j-128} \oplus 1$  for  $j = 160, 161, \dots, 191$ .
- b) Initialize the buffer and the compression function:
  - Set  $T^{(-256)} = 000$ .
  - Set  $b_j^{(-256)} = 0$  for  $j = 0, 1, \dots, 31$ .
  - Set  $I^{(-256)} = 0$ .
- c) Set  $S_{-64} = \text{InitNext}^{192}(S_{-256}, \text{LFSR})$ .
- d) Set  $i = -64$ .
- e) While  $I^{(i)} < 32$  and  $i < 0$ : set  $S_{i+1} = \text{InitNext}(S_i, \text{BUFF})$  and  $i = i + 1$ . The test  $I^{(i)} < 32$  can be removed, if a fixed, constant number of steps in the *Init* function are needed for implementation.
- f) Set  $S_0 = S_i$ .
- g) Output  $S_0$ .

NOTE Steps d), e) and f) of the DECIM<sup>v2</sup> initialization involve filling the buffer before starting the keystream output. As the output rate of the compression function varies, the number of steps required to fill the buffer may vary. In step e), the *InitNext*(BUFF) function is iterated 64 times at most, which guarantees that the buffer is full with probability more than  $1-2^{-97}$ . On average, the buffer is full after 24 iterations.

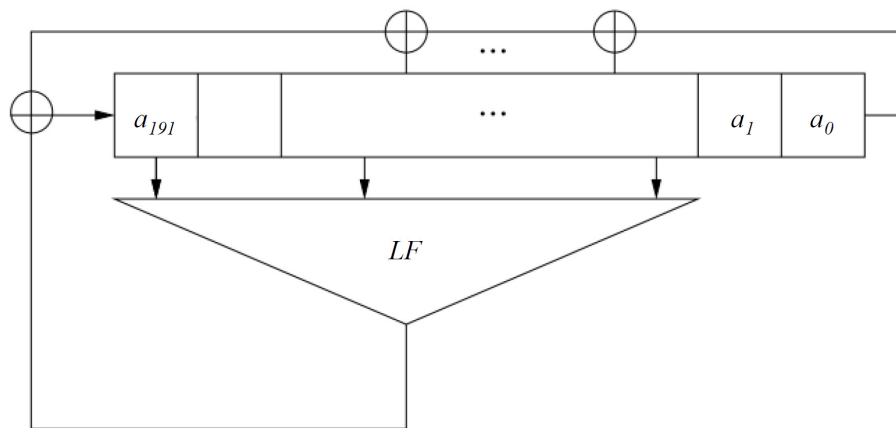


Figure 11 — LFSR mode of Initialization Next-state function *InitNext*

#### 8.4.4 Initialization Next-state function *InitNext*

Decim<sup>v2</sup> has two modes for the *InitNext* function: one mode is used during the initialization of the register *A* and the second during the initial filling of the buffer.

Input: State variable  $S_i = (a^{(i)}, T^{(i)}, b^{(i)}, I^{(i)})$ , mode  $\in \{\text{LFSR}, \text{BUFF}\}$ .

Output: Next value of the state variable  $S_{i+1} = (a^{(i+1)}, T^{(i+1)}, b^{(i+1)}, I^{(i+1)})$ .

Local variables: counters  $j, k$ , buffers  $f_k, r, c$ , state buffers  $\alpha^{(0)}, \dots, \alpha^{(4)}, \tau^{(0)}, \dots, \tau^{(4)}, \beta^{(0)}, \dots, \beta^{(4)}, \iota^{(0)}, \dots, \iota^{(4)}$ .

**LFSR mode (execute if mode = LFSR):**

a) Update the state of the register  $A$  with the following steps:

- 1) Set  $\alpha^{(0)} = a^{(i)}$ .
- 2) For  $k = 0, 1, 2, 3$ :
  - Set  $f_k = LF(\alpha^{(k)})$  and  $r = L(\alpha^{(k)}) \oplus f_k$ .
  - For  $j = 0, 1, \dots, 190$  set  $\alpha_j^{(k+1)} = \alpha_{j+1}^{(k)}$ .
  - Set  $\alpha_{191}^{(k+1)} = r$ .
- 3) Set  $a^{(i+1)} = \alpha^{(4)}$ .

Figure 11 shows the block diagram of the LFSR mode of *InitNext* function.

**BUFF mode (execute if mode = BUFF):**

a) Update the state of the register  $A$  with the following steps:

- 4) Set  $\alpha^{(0)} = a^{(i)}$ .
- 5) For  $k = 0, 1, 2, 3$ :
  - Set  $f_k = \alpha_1^{(k)} \oplus LF(\alpha^{(k)})$  and  $r = L(\alpha^{(k)})$ .
  - For  $j = 0, 1, \dots, 190$  set  $\alpha_j^{(k+1)} = \alpha_{j+1}^{(k)}$ .
  - Set  $\alpha_{191}^{(k+1)} = r$ .
- 6) Set  $a^{(i+1)} = \alpha^{(4)}$ .

b) Set  $\tau^{(0)} = T^{(i)}$ ,  $\beta^{(0)} = b^{(i)}$ ,  $t^{(0)} = I^{(i)}$ .

c) For  $k = 0, 1, 2, 3$ :

- 1) Update the state of the compression block with the following steps:
  - Set  $c = f_k \oplus \tau_2^{(k)}$ .
  - Set  $\tau^{(k+1)} = ABSG(\tau^{(k)}, f_k)$ .
  - If  $\tau_0^{(k+1)} = 0$ , set *output* = TRUE. Otherwise set *output* = FALSE.

2) Update the state of the buffer by  $(\beta^{(k+1)}, t^{(k+1)}) = B(\beta^{(k)}, t^{(k)}, output, c)$ .

d) Set  $T^{(i+1)} = \tau^{(4)}$ .

e) Set  $b^{(i+1)} = \beta^{(4)}$  and  $I^{(i+1)} = t^{(4)}$ .

#### 8.4.5 Next-state function *Next*

Input: State variable  $S_i = (a^{(i)}, T^{(i)}, b^{(i)}, I^{(i)})$ .

Output: Next value of the state variable  $S_{i+1} = (a^{(i+1)}, T^{(i+1)}, b^{(i+1)}, I^{(i+1)})$ .

Local variables: counters  $j, k$ , buffers  $f_k, r, c$ , state buffers  $\alpha^{(0)}, \dots, \alpha^{(4)}, \tau^{(0)}, \dots, \tau^{(4)}, \beta^{(0)}, \dots, \beta^{(4)}, \iota^{(0)}, \dots, \iota^{(4)}$ .

a) Update the state of the register  $A$  with the following steps:

- 7) Set  $\alpha^{(0)} = a^{(i)}$ .
- 8) For  $k = 0, 1, 2, 3$ :
  - Set  $f_k = \alpha_1^{(k)} \oplus LF(\alpha^{(k)})$  and  $r = L(\alpha^{(k)})$ .
  - For  $j = 0, 1, \dots, 190$  set  $\alpha_j^{(k+1)} = \alpha_{j+1}^{(k)}$ .
  - Set  $\alpha_{191}^{(k+1)} = r$ .

9) Set  $a^{(i+1)} = \alpha^{(4)}$ .

b) Set  $\tau^{(0)} = T^{(i)}, \beta^{(0)} = b^{(i)}, \iota^{(0)} = I^{(i)} - 1$ .

c) For  $j = 0, 1, \dots, \tau^{(0)} - 1$ , set  $\beta_j^{(0)} = b_{j+1}^{(i)}$

d) For  $k = 0, 1, 2, 3$ :

- 1) If  $\iota^{(0)} = 0$ , set  $\tau^{(k+1)} = \tau^{(k)}$ , *output* = TRUE and  $c = f_k$ . Otherwise update the state of the compression block with the following steps:
  - Set  $c = f_k \oplus \tau_2^{(k)}$ .
  - Set  $\tau^{(k+1)} = \text{ABSG}(\tau^{(k)}, f_k)$ .
  - If  $\tau_0^{(k+1)} = 0$ , set *output* = TRUE. Otherwise set *output* = FALSE.

2) Update the state of the buffer by  $(\beta^{(k+1)}, \tau^{(k+1)}) = B(\beta^{(k)}, \tau^{(k)}, \text{output}, c)$ .

e) Set  $T^{(i+1)} = \tau^{(4)}, b^{(i+1)} = \beta^{(4)}$  and  $I^{(i+1)} = \iota^{(4)}$ .

NOTE 1 The condition  $\iota^{(0)} = 0$  in step 1) of step d) should never be satisfied; if it is, this means that the buffer has become empty during the keystream generation. This happens with probability less than  $2^{-80}$  at every state update, see [8] for details. Also, this probability is higher if the buffer is not full after the *Init* function, but, as mentioned in 8.4.3 (NOTE), this also happens with negligible probability.

NOTE 2 The *InitNext* function and the *Next* function share many computational steps. Indeed, the LFSR mode of the *InitNext* function mainly consists of the LFSR update of the BUFF mode and of the *Next* function, the only difference being that the Boolean function output is added to the feedback bit. The BUFF mode of the *InitNext* function and the *Next* function differ only in that the buffer B is shifted only in the latter.

#### 8.4.6 Keystream function *Strm*

Input: State variable  $S_i = (a^{(i)}, T^{(i)}, b^{(i)}, I^{(i)})$ .

Output: Keystream bit  $Z_i$ .

- a) Set  $Z_i = b_0^{(i)}$ .
- b) Output  $Z_i$ .

#### 8.4.7 Linear feedback function $L$

Input: 192-bit string  $w = (w_0, w_1, \dots, w_{191})$ .

Output: Bit  $q = L(w)$ .

Set  $q = w_0 \oplus w_3 \oplus w_4 \oplus w_{23} \oplus w_{36} \oplus w_{37} \oplus w_{60} \oplus w_{61} \oplus w_{98} \oplus w_{115} \oplus w_{146} \oplus w_{175} \oplus w_{176} \oplus w_{187}$ .

#### 8.4.8 Filtering function $LF$

Input: 192-bit string  $w = (w_0, w_1, \dots, w_{191})$ .

Output: Bit  $q = LF(w)$ . Set  $q = Y((w_{13}, w_{28}, w_{45}, w_{54}, w_{65}, w_{104}, w_{111}, w_{144}, w_{162}, w_{172}, w_{178}, w_{186}, w_{191}))$ .

#### 8.4.9 Boolean function $Y$

Input: 13-bit string  $w = (w_0, w_1, \dots, w_{12})$ .

Output: Bit  $q = Y(w)$ . Set  $q = (\oplus_{0 \leq j \leq 12} w_j) \oplus (\oplus_{0 \leq j \leq 12} w_j w_k)$ .

NOTE Equivalently,  $q$  is given by  $q = 0$  if  $X = 0$  or  $X = 3$ , and  $q = 1$  otherwise, with  $X = w_0 + w_1 + \dots + w_{12} \bmod 4$ .

#### 8.4.10 Compression function $ABSG$

Input: 3-bit state  $T$ , input bit  $c$ .

Output: 3-bit state  $T' = ABSG(T, c)$ .

- a) If  $T_0 = 1$ , set  $T'_1 = T_1$ , otherwise set  $T'_1 = c$ .
- b) Set  $T'_2 = T_0 \text{ AND } (T_1 \oplus c)$ .
- c) Set  $T'_0 = (T_0 \oplus 1) \text{ OR } T'_2$ .

#### 8.4.11 Buffering function $B$

Input: 32-bit string  $b = (b_0, b_1, \dots, b_{31})$ , index  $I$ , Boolean *output*, input bit  $c$ .

Output: 32-bit string  $b' = (b'_0, b'_1, \dots, b'_{31})$ , index  $I'$ .

- a) Set  $I' = I$ ,  $b' = b$ .
- b) If *output* = TRUE and  $I' < 32$ , do the following:
  - Set  $b'_{I'} = c$ .
  - Set  $I' = I' + 1$ .
- c) Output  $B(b, I, \text{output}, c) = (b', I')$ .

## 8.5 KCipher-2 (K2) keystream generator

### 8.5.1 Introduction to KCipher-2 (K2)

KCIPHER-2 (K2) is a keystream generator which uses as input a 128-bit secret key  $K$  and a 128-bit initial vector  $IV$ . These are used to initialize state variables  $S_i$  ( $i \geq 0$ ) consisting of twenty 32-bit blocks, where  $S_i$  represents the internal state of K2 at clock  $i$ . Bit/byte order is big-endian, i.e., if the key and initialization vector are given as a sequence of bits/bytes, the first/leftmost bit/byte is the most significant of the corresponding data. For every iteration of the *Strm* function, a 64-bit keystream  $Z_i$  is produced as output.

K2's state variable  $S_i$  consists of three components. The first component  $A^{(i)}$  consists of a sequence of five 32-bit variables:

$$A^{(i)} = (A_4^{(i)}, A_3^{(i)}, A_2^{(i)}, A_1^{(i)}, A_0^{(i)}) \quad (A_m^{(i)} \text{ in } GF(2^{32}), m \geq 0)$$

which form the state for a feedback shift register (FSR)  $A$ . The second component  $B^{(i)}$  consists of a sequence of eleven 32-bit variables:

$$B^{(i)} = (B_{10}^{(i)}, B_9^{(i)}, \dots, B_0^{(i)}) \quad (B_m^{(i)} \text{ in } GF(2^{32}), m \geq 0)$$

which form the state for an FSR  $B$ . The third component consists of a set of four 32-bit variables:

$$R1^{(i)}, L1^{(i)}, R2^{(i)}, L2^{(i)} \text{ in } GF(2^{32})$$

which maintain the state of a non-linear function. The operation of K2 is summarized in Figures 12 and 13, which show a snapshot of operation, at time  $i$ , omitting the time dependent variable ( $i$ ) from the notation.



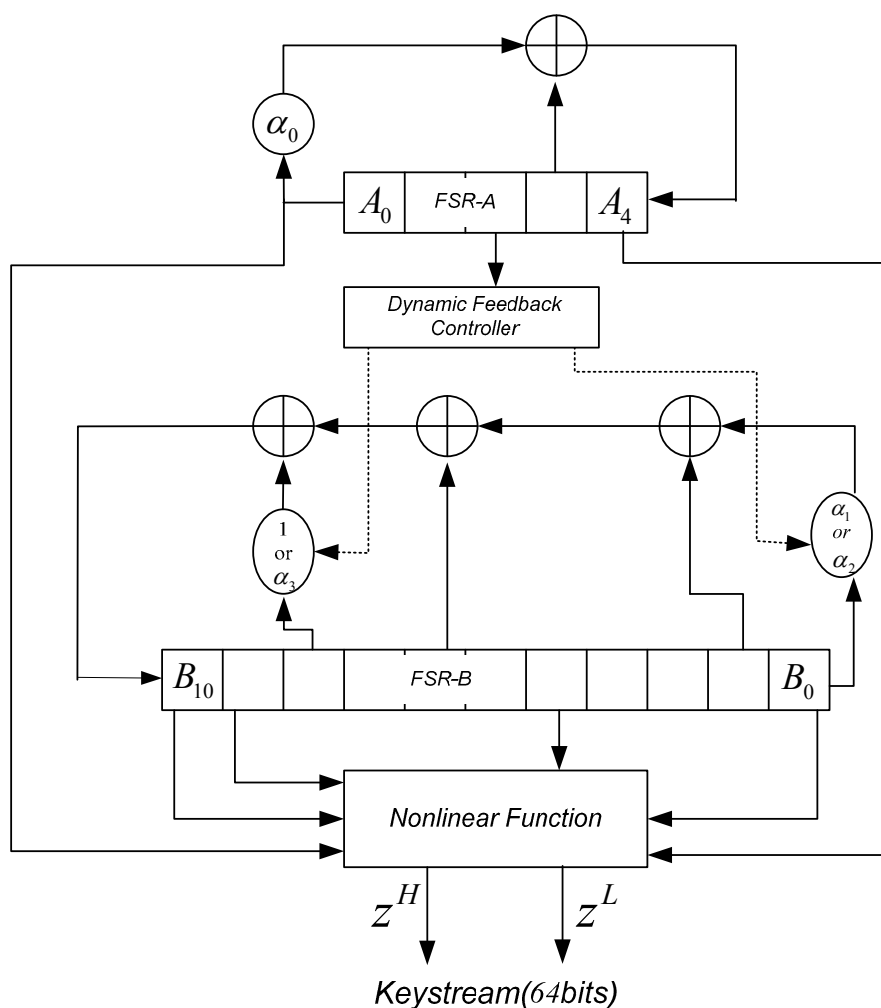


Figure 12 — Schematic drawing of K2

The operation of K2 is defined by the following three functions:

The *Init* function, defined in 8.5.2, takes as input the 128-bit key  $K$  and the 128-bit  $IV$  to produce the initial state  $S_0=(A^{(0)}, B^{(0)}, R1^{(0)}, L1^{(0)}, R2^{(0)}, L2^{(0)})$ .

The *Next* function, defined in 8.5.3, takes as input the internal state,  $S_i=(A^{(i)}, B^{(i)}, R1^{(i)}, L1^{(i)}, R2^{(i)}, L2^{(i)})$  and produces as output the next value of the state variable  $S_{i+1}=(A^{(i+1)}, B^{(i+1)}, R1^{(i+1)}, L1^{(i+1)}, R2^{(i+1)}, L2^{(i+1)})$ . The *Next* function runs in two modes, depending on whether the iteration performed is part of the initialization, or, of the normal mode of generating output.

The *Strm* function, defined in 8.5.4, takes as input the internal state,  $S_i=(A^{(i)}, B^{(i)}, R1^{(i)}, L1^{(i)}, R2^{(i)}, L2^{(i)})$  and produces as output the 64-bit keystream  $Z_i=(Z_i^H, Z_i^L)$ .

NOTE 1 The recommended maximum number of keystream bits without either re-keying or re-initializing with new  $IV$  is  $2^{64}$  bits.

NOTE 2 For the design rational for K2, refer to [12], [13], [14].



The initialization function *Init* works as follows.

Output; Initial value of state variable  $S_0=(A^{(0)},B^{(0)},R1^{(0)},L1^{(0)},R2^{(0)},L2^{(0)})$ .

Local variables: counter  $m$ .

- a) Expand the 128-bit key  $K=(K_0, K_1, K_2, K_3)$  into the 384-bit internal key  $IK=(IK_0, IK_1, \dots, IK_{11})$  as follows:
- 1) For  $m=0, 1, 2, 3$ , set  $IK_m = K_m$ .

2) For  $m=4,5,\dots,11$ ,

- If  $m \neq 4$  or 8, set  $IK_m = IK_{m-4} \oplus IK_{m-1}$ .
- If  $m=4$  set  $Rcon[0]=(0 \times 01, 0 \times 00, 0 \times 00, 0 \times 00)$  and  $IK_m = IK_{m-4} \oplus Sub_{K2}((IK_{m-1} \ll_{32} 8) \oplus (IK_{m-1} \gg_{32} 24)) \oplus Rcon[m/4-1]$ . The  $Sub_{K2}$  function is referred to in 8.5.5.
- If  $m=8$  set  $Rcon[1]=(0 \times 02, 0 \times 00, 0 \times 00, 0 \times 00)$  and  $IK_m = IK_{m-4} \oplus Sub_{K2}((IK_{m-1} \ll_{32} 8) \oplus (IK_{m-1} \gg_{32} 24)) \oplus Rcon[m/4-1]$ . The  $Sub_{K2}$  function is referred to in 8.5.5.

b) Initialize the registers with the internal key  $IK$  and  $IV = (IV_0, IV_1, IV_2, IV_3)$ .

— For  $m=0, 1, 2, 3, 4$ , set  $A_m^{(-24)} = IK_{4-m}$ .

— Set the registers in  $FSR-B$  as follows.

$$\begin{aligned} B_0^{(-24)} &= IK_{10}, B_1^{(-24)} = IK_{11}, B_2^{(-24)} = IV_0, B_3^{(-24)} = IV_1, B_4^{(-24)} = IK_8, \\ B_5^{(-24)} &= IK_9, B_6^{(-24)} = IV_2, B_7^{(-24)} = IV_3, B_8^{(-24)} = IK_7, B_9^{(-24)} = IK_5, B_{10}^{(-24)} = IK_6. \end{aligned}$$

— Set the registers in the nonlinear function as follows.

$$R1^{(-24)} = 0 \times 00000000, L1^{(-24)} = 0 \times 00000000, R2^{(-24)} = 0 \times 00000000, L2^{(-24)} = 0 \times 00000000.$$

c) Set  $S_0 = Next^{24}(S_{-24}, INIT)$ , where  $Next^{24}$  denotes 24 iterations of the  $Next$  function.

d) Output  $S_0$ .

We refer to 8.5.5 for description of the  $Sub_{K2}$  function.

### 8.5.3 Next-state function $Next$

$K2$  has two modes for the  $Next$  function.

Input: State variable  $S_i = (A^{(i)}, B^{(i)}, R1^{(i)}, L1^{(i)}, R2^{(i)}, L2^{(i)})$ , mode = {INIT, null}.

Output: Next value of the state variable  $S_{i+1} = (A^{(i+1)}, B^{(i+1)}, R1^{(i+1)}, L1^{(i+1)}, R2^{(i+1)}, L2^{(i+1)})$ .

Local variables: counter  $m$ .

a) Set the variables in the nonlinear function as follows.

$$\begin{aligned} R1^{(i+1)} &= Sub_{K2}(L2^{(i)} +_{32} B_9^{(i)}), \\ L1^{(i+1)} &= Sub_{K2}(R2^{(i)} +_{32} B_4^{(i)}), \\ R2^{(i+1)} &= Sub_{K2}(R1^{(i)}), \\ L2^{(i+1)} &= Sub_{K2}(L1^{(i)}), \end{aligned}$$

b) For  $m=0, 1, 2, 3$ , set  $A_m^{(i+1)} = A_{m+1}^{(i)}$ .

c) For  $m=0, 1, \dots, 9$ , set  $B_m^{(i+1)} = B_{m+1}^{(i)}$ .

d) For INIT mode, set  $A_4^{(i+1)} = (\alpha_0 \otimes A_0^{(i)}) \oplus A_3^{(i)} \oplus NLF(B_0^{(i)}, R2^{(i)}, R1^{(i)}, A_4^{(i)})$ ,  
For null mode, set  $A_4^{(i+1)} = (\alpha_0 \otimes A_0^{(i)}) \oplus A_3^{(i)}$ .

e) For INIT mode,

$$\begin{aligned} \text{set } B_{10}^{(i+1)} &= ((\alpha_1^{A_2^{(i)}[30]} + \alpha_2^{L_2^{(i)}[30]} - 1) \otimes B_0^{(i)}) \oplus B_1^{(i)} \oplus B_6^{(i)} \oplus (\alpha_3^{A_2^{(i)}[31]} \otimes B_8^{(i)}) \oplus NLF(B_{10}^{(i)}, L2^{(i)}, L1^{(i)}, \\ &A_0^{(i)}), \end{aligned}$$

For null mode,

$$\text{set } B_{10}^{(i+1)} = ((\alpha_1^{A_2^{(i)}[30]} + \alpha_2^{1-A_2^{(i)}[30]} - 1) \otimes B_0^{(i)}) \oplus B_1^{(i)} \oplus B_6^{(i)} \oplus (\alpha_3^{A_2^{(i)}[31]} \otimes B_8^{(i)}).$$

f) set  $S_{i+1} = (A^{(i+1)}, B^{(i+1)}, R1^{(i+1)}, L1^{(i+1)}, R2^{(i+1)}, L2^{(i+1)})$ .

g) Output  $S_{i+1}$ .

$A_m^{(i)}[Y]$  in  $\{0,1\}$  denotes the  $Y$ -th bit of the register  $A_m^{(i)}$ , where  $A_m^{(i)}[31]$  is the most significant bit of  $A_m^{(i)}$ . The description of the  $Sub_{K2}$  function and the finite field arithmetic involving the fixed elements,  $\alpha_0$ ,  $\alpha_1$ ,  $\alpha_2$  and  $\alpha_3$  refers to 8.5.5, 8.5.6, 8.5.7, 8.5.8, and 8.5.9, respectively. Also, the definition of the  $NLF$  function refers to 8.5.10.

Figure 14 is a block diagram of the INIT mode of the *Next* function.

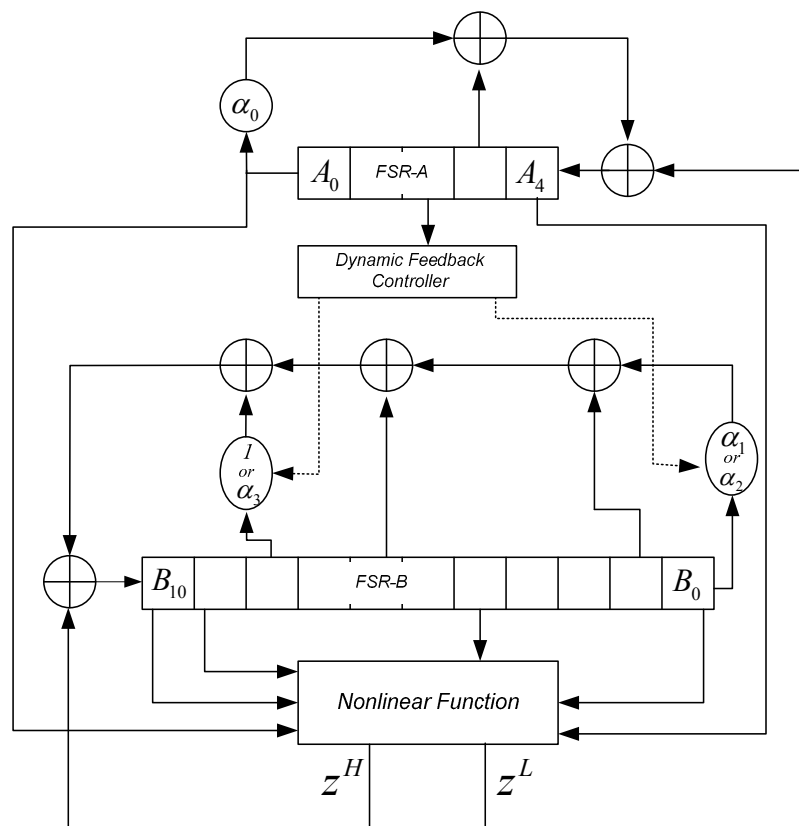


Figure 14 — INIT mode of the *Next* function

### 8.5.4 Keystream function *Strm*

The keystream function *Strm* works as follows.

Input: State variable  $S_i = (A^{(i)}, B^{(i)}, R1^{(i)}, L1^{(i)}, R2^{(i)}, L2^{(i)})$ .

Output: 64-bit keystream  $Z_i = (Z_i^H, Z_i^L)$ .

a) Set  $Z_i^H = NLF(B_{10}^{(i)}, L2^{(i)}, L1^{(i)}, A_0^{(i)})$ .

b) Set  $Z_i^L = NLF(B_0^{(i)}, R2^{(i)}, R1^{(i)}, A_4^{(i)})$ .

c) Set  $Z_i = (Z_i^H, Z_i^L)$ .

d) Output  $Z_i$ .

The function  $NLF$  is defined in 8.5.10.

### 8.5.5 Function $Sub_{K2}$

The  $Sub_{K2}$  function is a permutation of  $GF(2^{32})$ , based on components from the Advanced Encryption Standard (AES) [ISO/IEC 18033-3]. In  $Sub_{K2}$  function, the 32-bit input value is divided into four 1-byte strings and a nonlinear permutation is applied to each byte using an 8x8 bit substitution function ( $SBox$ ) followed by a 32x32 bit linear permutation. The  $SBox$  function is the same as  $SBox$  of AES, and the permutation is the same as AES *Mix Column* operation.

NOTE 1 The AES  $SBox$  function,  $SBox$ , can be found in 8.1.8 as the function  $SUB$ .

NOTE 2 Function  $Sub_{K2}$  produces the same output as function  $T$  of 8.2.5.

Input: A 32-bit value  $w$ .

Output: A 32-bit string  $q = Sub_{K2}(w)$ .

Local variables: counter  $m$ .

- Set  $w = (w_3, w_2, w_1, w_0)$ , where each  $w_m$  is 8-bit.
- For  $m=0,1,2,3$ , set  $t_m = SBox(w_m)$ .
- Set  $q = (q_3, q_2, q_1, q_0)$  as follows.

$$\begin{pmatrix} q_0 \\ q_1 \\ q_2 \\ q_3 \end{pmatrix} = \begin{pmatrix} 02 & 03 & 01 & 01 \\ 01 & 02 & 03 & 01 \\ 01 & 01 & 02 & 03 \\ 03 & 01 & 01 & 02 \end{pmatrix} \begin{pmatrix} t_0 \\ t_1 \\ t_2 \\ t_3 \end{pmatrix}$$

Multiplication of the elements  $t_i$  is performed in  $GF(2^8)$  making use of the irreducible polynomial  $f(x) = x^8 + x^4 + x^3 + x + 1$ .

d) Output  $q$ .

### 8.5.6 Multiplication of $\alpha_0$ in $GF(2^{32})$

Input: A 32-bit value  $w$ , which represents an element of  $GF(2^{32})$ .

Output: A 32-bit string  $w'$ , which represents  $\alpha_0 \otimes w$  in  $GF(2^{32})$ .

- Set  $w' = (w \ll_{32} 8) \oplus \alpha_{MUL0}[w \gg_{32} 24]$ .
- Output  $w'$ .

The function  $\alpha_{MUL0}$  is defined in the following:

$\alpha_{MUL0}[256] = \{$   
 $0 \times 00000000, 0 \times B6086D1A, 0 \times AF10DA34, 0 \times 1918B72E, 0 \times 9D207768, 0 \times 2B281A72, 0 \times 3230AD5C,$   
 $0 \times 8438C046, 0 \times F940EED0, 0 \times 4F4883CA, 0 \times 565034E4, 0 \times E05859FE, 0 \times 646099B8, 0 \times D268F4A2,$   
 $0 \times CB70438C, 0 \times 7D782E96, 0 \times 31801F63, 0 \times 87887279, 0 \times 9E90C557, 0 \times 2898A84D, 0 \times ACA0680B,$   
 $0 \times 1AA80511, 0 \times 03B0B23F, 0 \times B5B8DF25, 0 \times C8C0F1B3, 0 \times 7EC89CA9, 0 \times 67D02B87, 0 \times D1D8469D,$   
 $0 \times 55E086DB, 0 \times E3E8EBC1, 0 \times FAF05CEF, 0 \times 4CF831F5, 0 \times 62C33EC6, 0 \times D4CB53DC, 0 \times CDD3E4F2,$   
 $0 \times 7BDB89E8, 0 \times FFE349AE, 0 \times 49EB24B4, 0 \times 50F3939A, 0 \times E6FBFE80, 0 \times 9B83D016, 0 \times 2D8BBD0C,$

```

0x34930A22, 0x829B6738, 0x06A3A77E, 0xB0ABCA64, 0xA9B37D4A, 0x1FBB1050, 0x534321A5,
0xE54B4CBF, 0xFC53FB91, 0x4A5B968B, 0xCE6356CD, 0x786B3BD7, 0x61738CF9, 0xD77BE1E3,
0xAA03CF75, 0x1C0BA26F, 0x05131541, 0xB31B785B, 0x3723B81D, 0x812BD507, 0x98336229,
0x2E3B0F33, 0xC4457C4F, 0x724D1155, 0x6B55A67B, 0xDD5DCB61, 0x59650B27, 0xEF6D663D,
0xF675D113, 0x407DBC09, 0x3D05929F, 0x8B0DFF85, 0x921548AB, 0x241D25B1, 0xA025E5F7,
0x162D88ED, 0x0F353FC3, 0xB93D52D9, 0xF5C5632C, 0x43CD0E36, 0x5AD5B918, 0xECDDD402,
0x68E51444, 0xDEED795E, 0xC7F5CE70, 0x71FDA36A, 0x0C858DFC, 0xBA8DE0E6, 0xA39557C8,
0x159D3AD2, 0x91A5FA94, 0x27AD978E, 0x3EB520A0, 0x88BD4DBA, 0xA6864289, 0x108E2F93,
0x099698BD, 0xBF9EF5A7, 0x3BA635E1, 0x8DAE58FB, 0x94B6EFD5, 0x22BE82CF, 0x5FC6AC59,
0xE9CEC143, 0xF0D6766D, 0x46DE1B77, 0xC2E6DB31, 0x74EEB62B, 0x6DF60105, 0xDBFE6C1F,
0x97065DEA, 0x210E30F0, 0x381687DE, 0x8E1EEAC4, 0x0A262A82, 0xBC2E4798, 0xA536F0B6,
0x133E9DAC, 0x6E46B33A, 0xD84EDE20, 0xC156690E, 0x775E0414, 0xF366C452, 0x456EA948,
0x5C761E66, 0xEA7E737C, 0x4B8AF89E, 0xFD829584, 0xE49A22AA, 0x52924FB0, 0xD6AA8FF6,
0x60A2E2EC, 0x79BA55C2, 0xCFB238D8, 0xB2CA164E, 0x04C27B54, 0x1DDACC7A, 0xABD2A160,
0x2FEA6126, 0x99E20C3C, 0x80FABB12, 0x36F2D608, 0x7A0AE7FD, 0xCC028AE7, 0xD51A3DC9,
0x631250D3, 0xE72A9095, 0x5122FD8F, 0x483A4AA1, 0xFE3227BB, 0x834A092D, 0x35426437,
0x2C5AD319, 0x9A52BE03, 0x1E6A7E45, 0xA862135F, 0xB17AA471, 0x0772C96B, 0x2949C658,
0x9F41AB42, 0x86591C6C, 0x30517176, 0xB469B130, 0x0261DC2A, 0x1B796B04, 0xAD71061E,
0xD0092888, 0x66014592, 0x7F19F2BC, 0xC9119FA6, 0x4D295FE0, 0xFB2132FA, 0xE23985D4,
0x5431E8CE, 0x18C9D93B, 0xAEC1B421, 0xB7D9030F, 0x01D16E15, 0x85E9AE53, 0x33E1C349,
0x2AF97467, 0x9CF1197D, 0xE18937EB, 0x57815AF1, 0x4E99EDDF, 0xF89180C5, 0x7CA94083,
0xCA12D99, 0xD3B99AB7, 0x65B1F7AD, 0x8FCF84D1, 0x39C7E9CB, 0x20DF5EE5, 0x96D733FF,
0x12EFF3B9, 0xA4E79EA3, 0xBDFD298D, 0x0BF74497, 0x768F6A01, 0xC087071B, 0xD99FB035,
0x6F97DD2F, 0xEBAF1D69, 0x5DA77073, 0x44BFC75D, 0xF2B7AA47, 0xBE4F9BB2, 0x0847F6A8,
0x115F4186, 0xA7572C9C, 0x236FECDA, 0x956781C0, 0x8C7F36EE, 0x3A775BF4, 0x470F7562,
0xF1071878, 0xE81FAF56, 0x5E17C24C, 0xDA2F020A, 0x6C276F10, 0x753FD83E, 0xC337B524,
0xEDCBA17, 0x5B04D70D, 0x421C6023, 0xF4140D39, 0x702CCD7F, 0xC624A065, 0xDF3C174B,
0x69347A51, 0x144C54C7, 0xA24439DD, 0xBB5C8EF3, 0x0D54E3E9, 0x896C23AF, 0x3F644EB5,
0x267CF99B, 0x90749481, 0xDC8CA574, 0x6A84C86E, 0x739C7F40, 0xC594125A, 0x41ACD21C,
0xF7A4BF06, 0xEEBC0828, 0x58B46532, 0x25CC4BA4, 0x93C426BE, 0x8ADC9190, 0x3CD4FC8A,
0xB8EC3CCC, 0x0EE451D6, 0x17FCE6F8, 0xA1F48BE2};

```

### 8.5.7 Multiplication of $\alpha_i$ in $GF(2^{32})$

Input: A 32-bit value  $w$ , which represents an element of  $GF(2^{32})$ .

Output: A 32-bit string  $w'$ , which represents  $\alpha_i \otimes w$  in  $GF(2^{32})$ .

a) Set  $w' = (w \ll_{32} 8) \oplus_{\alpha_{MUL1}} [w \gg_{32} 24]$ .

b) Output  $w'$ .

The function  $\alpha_{MUL1}$  is defined in the following:

```

 $\alpha_{MUL1}[256] = \{$ 
0x00000000, 0xA0F5FC2E, 0x6DC7D55C, 0xCD322972, 0xDAA387B8, 0x7A567B96, 0xB76452E4,
0x1791AECA, 0x996B235D, 0x399EDF73, 0xF4ACF601, 0x54590A2F, 0x43C8A4E5, 0xE33D58CB,
0x2E0F71B9, 0x8EFA8D97, 0x1FD646BA, 0xBF23BA94, 0x721193E6, 0xD2E46FC8, 0xC575C102,
0x65803D2C, 0xA8B2145E, 0x0847E870, 0x86BD65E7, 0x264899C9, 0xEB7AB0BB, 0x4B8F4C95,
0x5C1EE25F, 0xFCEB1E71, 0x31D93703, 0x912CCB2D, 0x3E818C59, 0x9E747077, 0x53465905,
0xF3B3A52B, 0xE4220BE1, 0x44D7F7CF, 0x89E5DEBD, 0x29102293, 0xA7EAAF04, 0x071F532A,
0xCA2D7A58, 0x6AD88676, 0x7D4928BC, 0xDDBCD492, 0x108EFDE0, 0xB07B01CE, 0x2157CAE3,
0x81A236CD, 0x4C901FBF, 0xEC65E391, 0xFBF44D5B, 0x5B01B175, 0x96339807, 0x36C66429,
0xB83CE9BE, 0x18C91590, 0xD5FB3CE2, 0x750EC0CC, 0x629F6E06, 0xC26A9228, 0x0F58BB5A,
0xAFAD4774, 0x7C2F35B2, 0xDCDAC99C, 0x11E8E0EE, 0xB11D1CC0, 0xA68CB20A, 0x06794E24,
0xCB4B6756, 0x6BBE9B78, 0xE54416EF, 0x45B1EAC1, 0x8883C3B3, 0x28763F9D, 0x3FE79157,
0x9F126D79, 0x5220440B, 0xF2D5B825, 0x63F97308, 0xC30C8F26, 0x0E3EA654, 0xAECB5A7A,
0xB95AF4B0, 0x19AF089E, 0xD49D21EC, 0x7468DDC2, 0xFA925055, 0x5A67AC7B, 0x97558509,
0x37A07927, 0x2031D7ED, 0x80C42BC3, 0x4DF602B1, 0xED03FE9F, 0x42AEB9EB, 0xE25B45C5,

```

```
0x2F696CB7, 0x8F9C9099, 0x980D3E53, 0x38F8C27D, 0xF5CAEB0F, 0x553F1721, 0xDBC59AB6,
0x7B306698, 0xB6024FEA, 0x16F7B3C4, 0x01661D0E, 0xA193E120, 0x6CA1C852, 0xCC54347C,
0x5D78FF51, 0xFD8D037F, 0x30BF2A0D, 0x904AD623, 0x87DB78E9, 0x272E84C7, 0xEA1CADB5,
0x4AE9519B, 0xC413DC0C, 0x64E62022, 0xA9D40950, 0x0921F57E, 0x1EB05BB4, 0xBE45A79A,
0x73778EE8, 0xD38272C6, 0xF85E6A49, 0x58AB9667, 0x9599BF15, 0x356C433B, 0x22FDEDF1,
0x820811DF, 0x4F3A38AD, 0xEFCFC483, 0x61354914, 0xC1C0B53A, 0x0CF29C48, 0xAC076066,
0xBB96CEAC, 0x1B633282, 0xD6511BF0, 0x76A4E7DE, 0xE7882CF3, 0x477DD0DD, 0x8A4FF9AF,
0x2ABA0581, 0x3D2BAB4B, 0x9DDE5765, 0x50EC7E17, 0xF0198239, 0x7EE30FAE, 0xDE16F380,
0x1324DAF2, 0xB3D126DC, 0xA4408816, 0x04B57438, 0xC9875D4A, 0x6972A164, 0xC6DFE610,
0x662A1A3E, 0xAB18334C, 0x0BEDCF62, 0x1C7C61A8, 0xBC899D86, 0x71BBB4F4, 0xD14E48DA,
0x5FB4C54D, 0xFF413963, 0x32731011, 0x9286EC3F, 0x851742F5, 0x25E2BEDB, 0xE8D097A9,
0x48256B87, 0xD909A0AA, 0x79FC5C84, 0xB4CE75F6, 0x143B89D8, 0x03AA2712, 0xA35FDB3C,
0x6E6DF24E, 0xCE980E60, 0x406283F7, 0xE0977FD9, 0x2DA556AB, 0x8D50AA85, 0x9AC1044F,
0x3A34F861, 0xF706D113, 0x57F32D3D, 0x84715FFB, 0x2484A3D5, 0xE9B68AA7, 0x49437689,
0x5ED2D843, 0xFE27246D, 0x33150D1F, 0x93E0F131, 0x1D1A7CA6, 0xBDEF8088, 0x70DDA9FA,
0xD02855D4, 0xC7B9FB1E, 0x674C0730, 0xAA7E2E42, 0x0A8BD26C, 0x9BA71941, 0x3B52E56F,
0xF660CC1D, 0x56953033, 0x41049EF9, 0xE1F162D7, 0x2CC34BA5, 0x8C36B78B, 0x02CC3A1C,
0xA239C632, 0x6F0BEF40, 0xCFFE136E, 0xD86FBDA4, 0x789A418A, 0xB5A868F8, 0x155D94D6,
0xBAF0D3A2, 0x1A052F8C, 0xD73706FE, 0x77C2FAD0, 0x6053541A, 0xC0A6A834, 0x0D948146,
0xAD617D68, 0x239BF0FF, 0x836E0CD1, 0x4E5C25A3, 0xEEA9D98D, 0xF9387747, 0x59CD8B69,
0x94FFA21B, 0x340A5E35, 0xA5269518, 0x05D36936, 0xC8E14044, 0x6814BC6A, 0x7F8512A0,
0xDF70EE8E, 0x1242C7FC, 0xB2B73BD2, 0x3C4DB645, 0x9CB84A6B, 0x518A6319, 0xF17F9F37,
0xEE31FD, 0x461BCDD3, 0x8B29E4A1, 0x2BDC188F};
```

### 8.5.8 Multiplication of $\alpha_2$ in $GF(2^{32})$

Input: A 32-bit value  $w$ , which represents an element of  $GF(2^{32})$ .

Output: A 32-bit string  $w'$ , which represents  $\alpha_2 \otimes w$  in  $GF(2^{32})$ .

- Set  $w' = (w \ll_{32} 8) \oplus \alpha_{MUL2}[w \gg_{32} 24]$ .
- Output  $w'$ .

The function  $\alpha_{MUL2}$  is defined in the following:

```
 $\alpha_{MUL2}[256] = \{$ 
0x00000000, 0x5BF87F93, 0xB6BDFE6B, 0xED4581F8, 0x2137B1D6, 0x7ACFCE45, 0x978A4FBD,
0xCC72302E, 0x426E2FE1, 0x19965072, 0xF4D3D18A, 0xAF2BAE19, 0x63599E37, 0x38A1E1A4,
0xD5E4605C, 0x8E1C1FCF, 0x84DC5E8F, 0xDF24211C, 0x3261A0E4, 0x6999DF77, 0xA5EBEF59,
0xFE1390CA, 0x13561132, 0x48AE6EA1, 0xC6B2716E, 0x9D4A0EFD, 0x700F8F05, 0x2BF7F096,
0xE785C0B8, 0xBC7DBF2B, 0x51383ED3, 0x0AC04140, 0x45F5BC53, 0x1E0DC3C0, 0xF3484238,
0xA8B03DAB, 0x64C20D85, 0x3F3A7216, 0xD27FF3EE, 0x89878C7D, 0x079B93B2, 0x5C63EC21,
0xB1266DD9, 0xEADE124A, 0x26AC2264, 0x7D545DF7, 0x9011DC0F, 0xCBE9A39C, 0xC129E2DC,
0x9AD19D4F, 0x77941CB7, 0x2C6C6324, 0xE01E530A, 0xBBE62C99, 0x56A3AD61, 0x0D5BD2F2,
0x8347CD3D, 0xD8BFB2AE, 0x35FA3356, 0x6E024CC5, 0xA2707CEB, 0xF9880378, 0x14CD8280,
0x4F35FD13, 0x8AA735A6, 0xD15F4A35, 0x3C1ACBCD, 0x67E2B45E, 0xAB908470, 0xF068FBE3,
0x1D2D7A1B, 0x46D50588, 0xC8C91A47, 0x933165D4, 0x7E74E42C, 0x258C9BBF, 0xE9FEAB91,
0xB206D402, 0x5F4355FA, 0x04BB2A69, 0x0E7B6B29, 0x558314BA, 0xB8C69542, 0xE33EEAD1,
0x2F4CDAFF, 0x74B4A56C, 0x99F12494, 0xC2095B07, 0x4C1544C8, 0x17ED3B5B, 0xFAA8BAA3,
0xA150C530, 0x6D22F51E, 0x36DA8A8D, 0xDB9F0B75, 0x806774E6, 0xCF5289F5, 0x94AAF666,
0x79EF779E, 0x2217080D, 0xEE653823, 0xB59D47B0, 0x58D8C648, 0x0320B9DB, 0x8D3CA614,
0xD6C4D987, 0x3B81587F, 0x607927EC, 0xAC0B17C2, 0xF7F36851, 0x1AB6E9A9, 0x414E963A,
0x4B8ED77A, 0x1076A8E9, 0xFD332911, 0xA6CB5682, 0x6AB966AC, 0x3141193F, 0xDC0498C7,
0x87FCE754, 0x09E0F89B, 0x52188708, 0xBF5D06F0, 0xE4A57963, 0x28D7494D, 0x732F36DE,
0x9E6AB726, 0xC592C8B5, 0x59036A01, 0x02FB1592, 0xEFBE946A, 0xB446EBF9, 0x7834DBD7,
0x23CCA444, 0xCE8925BC, 0x95715A2F, 0x1B6D45E0, 0x40953A73, 0xADD0BB8B, 0xF628C418,
0x3A5AF436, 0x61A28BA5, 0x8CE70A5D, 0xD71F75CE, 0xDDDF348E, 0x86274B1D, 0x6B62CAE5,
0x309AB576, 0xFCE88558, 0xA710FACB, 0x4A557B33, 0x11AD04A0, 0x9FB11B6F, 0xC44964FC,
```

```
0x290CE504, 0x72F49A97, 0xBE86AAB9, 0xE57ED52A, 0x083B54D2, 0x53C32B41, 0x1CF6D652,
0x470EA9C1, 0xAA4B2839, 0xF1B357AA, 0x3DC16784, 0x66391817, 0x8B7C99EF, 0xD084E67C,
0x5E98F9B3, 0x05608620, 0xE82507D8, 0xB3DD784B, 0x7FAF4865, 0x245737F6, 0xC912B60E,
0x92EAC99D, 0x982A88DD, 0xC3D2F74E, 0x2E9776B6, 0x756F0925, 0xB91D390B, 0xE2E54698,
0x0FA0C760, 0x5458B8F3, 0xDA44A73C, 0x81BCD8AF, 0x6CF95957, 0x370126C4, 0xFB7316EA,
0xA08B6979, 0x4DCEE881, 0x16369712, 0xD3A45FA7, 0x885C2034, 0x6519A1CC, 0x3EE1DE5F,
0xF293EE71, 0xA96B91E2, 0x442E101A, 0x1FD66F89, 0x91CA7046, 0xCA320FD5, 0x27778E2D,
0x7C8FF1BE, 0xB0FDC190, 0xEB05BE03, 0x06403FFB, 0x5DB84068, 0x57780128, 0x0C807EBB,
0xE1C5FF43, 0xBA3D80D0, 0x764FB0FE, 0x2DB7CF6D, 0xC0F24E95, 0x9B0A3106, 0x15162EC9,
0x4EEE515A, 0xA3ABD0A2, 0xF853AF31, 0x34219F1F, 0x6FD9E08C, 0x829C6174, 0xD9641EE7,
0x9651E3F4, 0xCDA99C67, 0x20EC1D9F, 0x7B14620C, 0xB7665222, 0xEC9E2DB1, 0x01DBAC49,
0x5A23D3DA, 0xD43FCC15, 0x8FC7B386, 0x6282327E, 0x397A4DED, 0xF5087DC3, 0xAEF00250,
0x43B583A8, 0x184DFC3B, 0x128DBD7B, 0x4975C2E8, 0xA4304310, 0xFFC83C83, 0x33BA0CAD,
0x6842733E, 0x8507F2C6, 0xDEFF8D55, 0x50E3929A, 0x0B1BED09, 0xE65E6CF1, 0xBDA61362,
0x71D4234C, 0x2A2C5CDF, 0xC769DD27, 0x9C91A2B4};
```

### 8.5.9 Multiplication of $\alpha_3$ in $GF(2^{32})$

Input: A 32-bit value  $w$ , which represents an element of  $GF(2^{32})$ .

Output: A 32-bit string  $w'$ , which represents  $\alpha_3 \otimes w$  in  $GF(2^{32})$ .

a) Set  $w' = (w \ll_{32} 8) \oplus \alpha_{\text{MUL3}}[w \gg_{32} 24]$ .

b) Output  $w'$ .

The function  $\alpha_{\text{MUL3}}$  is defined in the following:

```
 $\alpha_{\text{MUL3}}[256] = \{$ 
0x00000000, 0x4559568B, 0x8AB2AC73, 0xCFEBFAF8, 0x71013DE6, 0x34586B6D, 0xFBB39195,
0xBEEAC71E, 0xE2027AA9, 0xA75B2C22, 0x68B0D6DA, 0x2DE98051, 0x9303474F, 0xD65A11C4,
0x19B1EB3C, 0x5CE8BDB7, 0xA104F437, 0xE45DA2BC, 0x2BB65844, 0x6EEF0ECF, 0xD005C9D1,
0x955C9F5A, 0x5AB765A2, 0x1FEE3329, 0x43068E9E, 0x065FD815, 0xC9B422ED, 0x8CED7466,
0x3207B378, 0x775EE5F3, 0xB8B51F0B, 0xFDEC4980, 0x27088D6E, 0x6251DBE5, 0xADBA211D,
0xE8E37796, 0x5609B088, 0x1350E603, 0xDCBB1CFB, 0x99E24A70, 0xC50AF7C7, 0x8053A14C,
0x4FB85BB4, 0x0AE10D3F, 0xB40BCA21, 0xF1529CAA, 0x3EB96652, 0x7BE030D9, 0x860C7959,
0xC3552FD2, 0x0CBED52A, 0x49E783A1, 0xF70D44BF, 0xB2541234, 0x7DBFE8CC, 0x38E6BE47,
0x640E03F0, 0x2157557B, 0xEEBCAF83, 0xABE5F908, 0x150F3E16, 0x5056689D, 0x9FBD9265,
0xDAE4C4EE, 0x4E107FDC, 0x0B492957, 0xC4A2D3AF, 0x81FB8524, 0x3F11423A, 0x7A4814B1,
0xB5A3EE49, 0xF0FAB8C2, 0xAC120575, 0xE94B53FE, 0x26A0A906, 0x63F9FF8D, 0xDD133893,
0x984A6E18, 0x57A194E0, 0x12F8C26B, 0xEF148BEB, 0xAA4DD60, 0x65A62798, 0x20FF7113,
0x9E15B60D, 0xDB4CE086, 0x14A71A7E, 0x51FE4CF5, 0x0D16F142, 0x484FA7C9, 0x87A45D31,
0xC2FD0BBA, 0x7C17CCA4, 0x394E9A2F, 0xF6A560D7, 0xB3FC365C, 0x6918F2B2, 0x2C41A439,
0xE3AA5EC1, 0xA6F3084A, 0x1819CF54, 0x5D4099DF, 0x92AB6327, 0xD7F235AC, 0x8B1A881B,
0xCE43DE90, 0x01A82468, 0x44F172E3, 0xFA1BB5FD, 0xBF42E376, 0x70A9198E, 0x35F04F05,
0xC81C0685, 0x8D45500E, 0x42AEAAF6, 0x07F7FC7D, 0xB91D3B63, 0xFC446DE8, 0x33AF9710,
0x76F6C19B, 0x2A1E7C2C, 0x6F472AA7, 0xA0ACD05F, 0xE5F586D4, 0x5B1F41CA, 0x1E461741,
0xD1ADEDB9, 0x94F4BB32, 0x9C20FEDD, 0xD979A856, 0x169252AE, 0x53CB0425, 0xED21C33B,
0xA87895B0, 0x67936F48, 0x22CA39C3, 0x7E228474, 0x3B7BD2FF, 0xF4902807, 0xB1C97E8C,
0x0F23B992, 0x4A7AEF19, 0x859115E1, 0xC0C8436A, 0x3D240AEA, 0x787D5C61, 0xB796A699,
0xF2CFF012, 0x4C25370C, 0x097C6187, 0xC6979B7F, 0x83CECDF4, 0xDF267043, 0x9A7F26C8,
0x5594DC30, 0x10CD8ABB, 0xAE274DA5, 0xEB7E1B2E, 0x2495E1D6, 0x61CCB75D, 0xBB2873B3,
0xFE712538, 0x319ADFC0, 0x74C3894B, 0xCA294E55, 0x8F7018DE, 0x409BE226, 0x05C2B4AD,
0x592A091A, 0x1C735F91, 0xD398A569, 0x96C1F3E2, 0x282B34FC, 0x6D726277, 0xA299988F,
0xE7C0CE04, 0x1A2C8784, 0x5F75D10F, 0x909E2BF7, 0xD5C77D7C, 0x6B2DBA62, 0x2E74ECE9,
0xE19F1611, 0xA4C6409A, 0xF82EFD2D, 0xBD77ABA6, 0x729C515E, 0x37C507D5, 0x892FC0CB,
0xCC769640, 0x039D6CB8, 0x46C43A33, 0xD2308101, 0x9769D78A, 0x58822D72, 0x1DDB7BF9,
0xA331BCE7, 0xE668EA6C, 0x29831094, 0x6CDA461F, 0x3032FBA8, 0x756BAD23, 0xBA8057DB,
0xFFD90150, 0x4133C64E, 0x046A90C5, 0xCB816A3D, 0xED83CB6, 0x73347536, 0x366D23BD,
```



0xF986D945, 0xBCDF8FCE, 0x023548D0, 0x476C1E5B, 0x8887E4A3, 0xCDDEB228, 0x91360F9F,  
0xD46F5914, 0x1B84A3EC, 0x5EDDF567, 0xE0373279, 0xA56E64F2, 0x6A859E0A, 0x2FDCC881,  
0xF5380C6F, 0xB0615AE4, 0x7F8AA01C, 0x3AD3F697, 0x84393189, 0xC1606702, 0x0E8B9DFA,  
0x4BD2CB71, 0x173A76C6, 0x5263204D, 0x9D88DAB5, 0xD8D18C3E, 0x663B4B20, 0x23621DAB,  
0xEC89E753, 0xA9D0B1D8, 0x543CF858, 0x1165AED3, 0xDE8E542B, 0x9BD702A0, 0x253DC5BE,  
0x60649335, 0xAF8F69CD, 0xEAD63F46, 0xB63E82F1, 0xF367D47A, 0x3C8C2E82, 0x79D57809,  
0xC73FBF17, 0x8266E99C, 0x4D8D1364, 0x08D445EF};

#### 8.5.10 Function $NLF(a,b,c,d)$

Input: Four 32-bit values,  $a$ ,  $b$ ,  $c$  and  $d$ .

Output: A 32-bit string  $q$ .

a) Set  $q = (a \oplus_{32} b) \oplus c \oplus d$ .

b) Output  $q$ .

### A1 8.6 ZUC KEYSTREAM GENERATOR

#### 8.6.1 INTRODUCTION TO ZUC

ZUC is a keystream generator which uses as input a 128-bit secret key  $K$  and a 128-bit initialization vector  $IV$ . These are used to initialize state variables  $S_i$  ( $i \geq 0$ ). The bit/byte order is big-endian, i.e., if the key and initialization vector are given as a sequence of bits/bytes, the first/leftmost bit/byte is the most significant bit/byte of the corresponding data. It outputs a 32-bit keystream  $Z_i$  at every iteration of the function  $\text{Strm}$ .

The state variable  $S_i$  consists of two components. The first consists of sixteen 31-bit variables:

$$A^{(i)} = (A_{15}^{(i)}, A_{14}^{(i)}, \dots, A_0^{(i)}),$$

and maintains the state of a linear feedback shift register. The second consists of two 32-bit variables:

$$R^{(i)} = (R_2^{(i)}, R_1^{(i)}),$$

that maintains the state of a finite state machine. ZUC is summarised in Figure 15, which shows a snapshot if its operation, at time  $i$ , omitting the time-dependent variable ( $i$ ) from the notation.

The *Init* function, defined in detail in 8.6.2, takes as input the 128-bit key  $K$  and the 128-bit initialization vector  $IV$ , and produces the initial value of the state variable  $S_0 = (A^{(0)}, R^{(0)})$ .

The *Next* function, defined in detail in 8.6.3, takes as input the state variable  $S_i = (A^{(i)}, R^{(i)})$  and produces as output the next value of the state variable  $S_{i+1} = (A^{(i+1)}, R^{(i+1)})$ . The *Next* function runs in two modes, depending on whether the iteration performed is part of the initialization mode or of the normal mode of generating output.

The *Strm* function, defined in detail in 8.6.4, takes as input the state variable  $S_i = (A^{(i)}, R^{(i)})$  and produces as output the 32-bit keystream  $Z_i$ .

NOTE See document [20] for theoretical background on the design rationale for ZUC.

A 240-bit constant  $D = d_0 \parallel d_1 \parallel \dots \parallel d_{15}$  used in ZUC:

$$\begin{aligned} d_0 &= 100010011010111, d_1 = 010011010111100, d_2 = 110001001101011, d_3 = 001001101011110, \\ d_4 &= 101011110001001, d_5 = 011010111100010, d_6 = 111000100110101, d_7 = 000100110101111, \\ d_8 &= 100110101111000, d_9 = 010111100010011, d_{10} = 110101111000100, d_{11} = 001101011110001, \\ d_{12} &= 101111000100110, d_{13} = 011110001001101, d_{14} = 111100010011010, d_{15} = 100011110101100, \end{aligned}$$

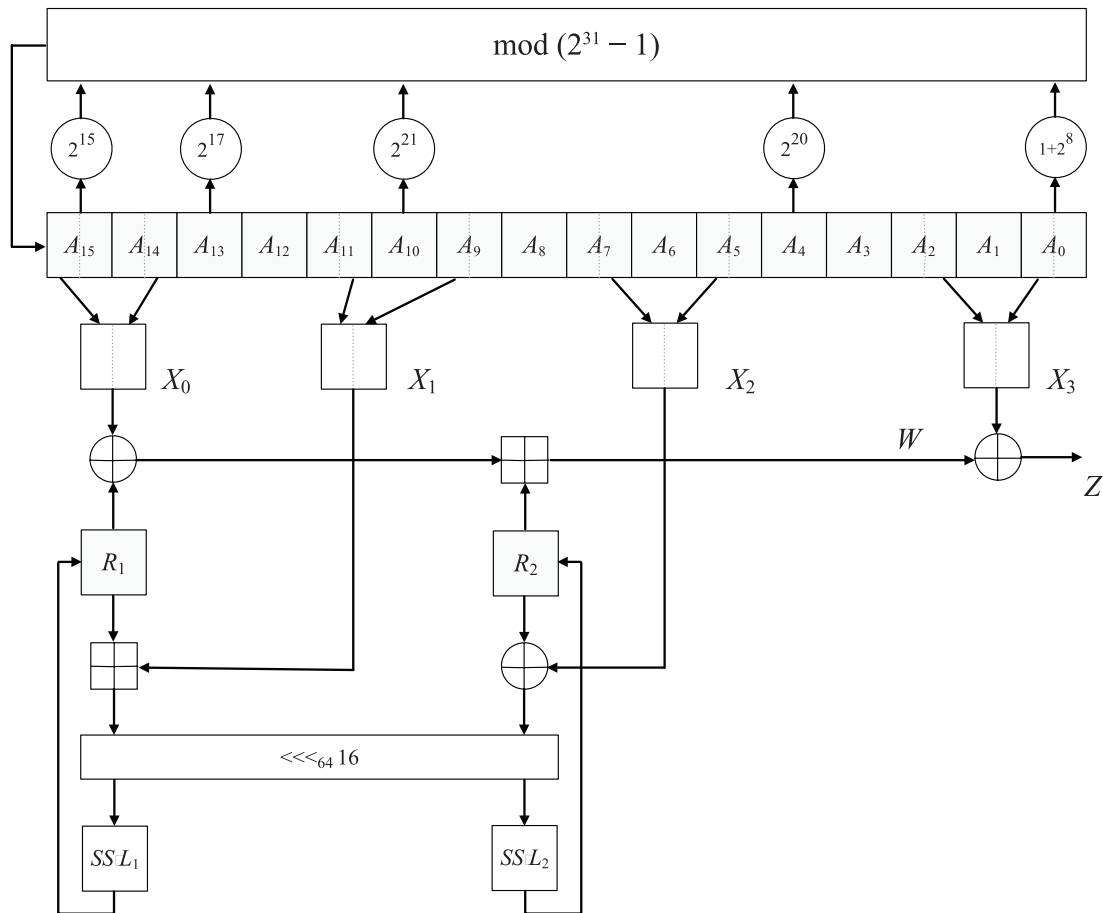


Figure 15 — Schematic drawing of ZUC

where for  $i = 0, 1, \dots, 15$ ,  $d_i$  is a 15-bit variable in binary notation.

The description uses notations defined in Clause 4 of this part of ISO/IEC 18033. For a string  $A$  which has at least 16 bits, the notation  $A_H$  represents the leftmost 16 bits of  $A$  and the notation  $A_L$  represents the rightmost 16 bits of  $A$ . For example, if  $A = 100010011011111011110101111001$  is a 31-bit string, then  $A_H = 1000100110111110$  and  $A_L = 0111110101111001$ .

### 8.6.2 INITIALIZATION FUNCTION *INIT*

The Initialization function *Init* is as follows.

Input: 128-bit key  $K$ , 128-bit initialization vector  $IV$ .

Output: Initial value of state variable  $S_0 = (A^{(0)}, R^{(0)})$ .

- Initialize the state variable  $S_{-33}$  with the key  $K$ , the 128-bit initialization vector  $IV$  and the constant  $D$ .
  - Set  $(k_0, k_1, \dots, k_{15}) = K$ ;  $(iv_0, iv_1, \dots, iv_{15}) = IV$ , where  $k_i$  and  $iv_i$  are bytes for  $i = 0, 1, \dots, 15$ .
  - Set  $A_i^{(-33)} = k_i \parallel d_i \parallel iv_i$  for  $i = 0, 1, \dots, 15$ .
  - Set  $R_1^{(-33)} = R_2^{(-33)} = 0^{(32)}$ .
- Set  $S_{-1} = \text{Next}^{32}(S_{-33}, \text{INIT})$ , where  $\text{Next}^{32}$  denotes 32 iterations of the *Next* function.
- Set  $S_0 = \text{Next}(S_{-1}, \text{null})$ .
- Output  $S_0$ .

### 8.6.3 NEXT-STATE FUNCTION NEXT

The *Next* function has two modes, and is defined as follows.

Input: State variable  $S_i = (A^{(i)}, R^{(i)})$ , mode  $\in \{\text{INIT}, \text{null}\}$ .

Output: Next value of the state variable  $S_{i+1} = (A^{(i+1)}, R^{(i+1)})$ .

Local variables: 32-bit strings  $W$ ,  $W_1$ ,  $W_2$ ,  $X_0$ ,  $X_1$ ,  $X_2$  and 31-bit string  $V$ .

- Set  $X_0 = A_{15}^{(i)} \parallel A_{14}^{(i)}; X_1 = A_{11}^{(i)} \parallel A_9^{(i)}; X_2 = A_7^{(i)} \parallel A_5^{(i)}$ .
- Set  $W = (X_0 \oplus R_1^{(i)}) + 32 R_2^{(i)}; W_1 = R_1^{(i)} + 32 X_1; W_2 = R_2^{(i)} \oplus X_2; R_1^{(i+1)} = SS(L_1(W_{1L} \parallel W_{2H})); R_2^{(i+1)} = SS(L_2(W_{2L} \parallel W_{1H}))$ .
- Set  $V = 2^{15}A_{15}^{(i)} + 2^{17}A_{13}^{(i)} + 2^{21}A_{10}^{(i)} + 2^{20}A_4^{(i)} + (1+2^8)A_0^{(i)} \bmod (2^{31}-1)$ .
- If mode = INIT, set  $A_{15}^{(i+1)} = V + (31 \sim W) \bmod (2^{31}-1)$ . Otherwise, set  $A_{15}^{(i+1)} = V$ . If  $A_{15}^{(i+1)} = 0$ , set  $A_{15}^{(i+1)} = 2^{31}-1$ .
- Set  $A_j^{(i+1)} = A_{j+1}^{(i)}$  for  $j = 0, 1, \dots, 14$ .
- Set  $S_{i+1} = (A^{(i+1)}, R^{(i+1)})$ .
- Output  $S_{i+1}$ .

NOTE For two 31-bit strings  $a$  and  $b$ , if  $b = 2^i$ , then  $ab \bmod (2^{31}-1) = a \lll_{31} i \bmod (2^{31}-1)$ ; if  $b = 2^i + 2^j$ , then  $ab \bmod (2^{31}-1) = (a \lll_{31} i) + (a \lll_{31} j) \bmod (2^{31}-1)$ . Reference C code for ZUC is given in document [21].

### 8.6.4 KEYSTREAM FUNCTION STRM

The keystream function *Strm* is as follows:

Input: State variable  $S_i$ .

Output: 32-bit keystream  $Z_i$ .

Local variables: 32-bit strings  $X_0$ ,  $X_3$ .

- Set  $X_0 = A_{15}^{(i)} \parallel A_{14}^{(i)}; X_3 = A_2^{(i)} \parallel A_0^{(i)}$ .
- Set  $Z_i = ((X_0 \oplus R_1^{(i)}) +_{32} R_2^{(i)}) \oplus X_3$ .
- Output  $Z_i$ .

### 8.6.5 FUNCTION SS

The function *SS* is as follows:

Input: 32-bit string  $X$ .

Output: 32-bit string  $Y$ .

- Define  $X = x_3 \parallel x_2 \parallel x_1 \parallel x_0$ , where  $x_i$  is a byte for  $i = 0, 1, 2, 3$ .
- Set  $Y = SUB1[x_3] \parallel SUB2[x_2] \parallel SUB1[x_1] \parallel SUB2[x_0]$ .
- Output  $Y$ .

The functions *SUB1* and *SUB2* are defined by the following substitution tables:

*SUB1* [256] = {  
0x3e, 0x72, 0x5b, 0x47, 0xca, 0xe0, 0x00, 0x33, 0x04, 0xd1, 0x54, 0x98, 0x09, 0xb9, 0x6d, 0xcb,  
0x7b, 0x1b, 0xf9, 0x32, 0xaf, 0x9d, 0x6a, 0xa5, 0xb8, 0x2d, 0xfc, 0x1d, 0x08, 0x53, 0x03, 0x90,  
0x4d, 0x4e, 0x84, 0x99, 0xe4, 0xce, 0xd9, 0x91, 0xdd, 0xb6, 0x85, 0x48, 0x8b, 0x29, 0x6e, 0xac,  
0xcd, 0xc1, 0xf8, 0x1e, 0x73, 0x43, 0x69, 0xc6, 0xb5, 0xbd, 0xfd, 0x39, 0x63, 0x20, 0xd4, 0x38,  
0x76, 0x7d, 0xb2, 0xa7, 0xcf, 0xed, 0x57, 0xc5, 0xf3, 0x2c, 0xbb, 0x14, 0x21, 0x06, 0x55, 0x9b,  
0xe3, 0xef, 0x5e, 0x31, 0x4f, 0x7f, 0x5a, 0xa4, 0x0d, 0x82, 0x51, 0x49, 0x5f, 0xba, 0x58, 0x1c,  
0x4a, 0x16, 0xd5, 0x17, 0xa8, 0x92, 0x24, 0x1f, 0x8c, 0xff, 0xd8, 0xae, 0x2e, 0x01, 0xd3, 0xad,  
0x3b, 0x4b, 0xda, 0x46, 0xeb, 0xc9, 0xde, 0x9a, 0x8f, 0x87, 0xd7, 0x3a, 0x80, 0x6f, 0x2f, 0xc8,  
0xb1, 0xb4, 0x37, 0xf7, 0x0a, 0x22, 0x13, 0x28, 0x7c, 0xcc, 0x3c, 0x89, 0xc7, 0xc3, 0x96, 0x56,  
0x07, 0xbf, 0x7e, 0xf0, 0x0b, 0x2b, 0x97, 0x52, 0x35, 0x41, 0x79, 0x61, 0xa6, 0x4c, 0x10, 0xfe,  
0xbc, 0x26, 0x95, 0x88, 0x8a, 0xb0, 0xa3, 0xfb, 0xc0, 0x18, 0x94, 0xf2, 0xe1, 0xe5, 0xe9, 0x5d,  
0xd0, 0xdc, 0x11, 0x66, 0x64, 0x5c, 0xec, 0x59, 0x42, 0x75, 0x12, 0xf5, 0x74, 0x9c, 0xaa, 0x23,  
0x0e, 0x86, 0xab, 0xbe, 0x2a, 0x02, 0xe7, 0x67, 0xe6, 0x44, 0xa2, 0x6c, 0xc2, 0x93, 0x9f, 0xf1,  
0xf6, 0xfa, 0x36, 0xd2, 0x50, 0x68, 0x9e, 0x62, 0x71, 0x15, 0x3d, 0xd6, 0x40, 0xc4, 0xe2, 0x0f,  
0x8e, 0x83, 0x77, 0x6b, 0x25, 0x05, 0x3f, 0x0c, 0x30, 0xea, 0x70, 0xb7, 0xa1, 0xe8, 0xa9, 0x65,  
0x8d, 0x27, 0x1a, 0xdb, 0x81, 0xb3, 0xa0, 0xf4, 0x45, 0x7a, 0x19, 0xdf, 0xee, 0x78, 0x34, 0x60};

*SUB2* [256] = {  
0x55, 0xc2, 0x63, 0x71, 0x3b, 0xc8, 0x47, 0x86, 0x9f, 0x3c, 0xda, 0x5b, 0x29, 0xaa, 0xfd, 0x77,  
0x8c, 0xc5, 0x94, 0x0c, 0xa6, 0x1a, 0x13, 0x00, 0xe3, 0xa8, 0x16, 0x72, 0x40, 0xf9, 0xf8, 0x42,  
0x44, 0x26, 0x68, 0x96, 0x81, 0xd9, 0x45, 0x3e, 0x10, 0x76, 0xc6, 0xa7, 0x8b, 0x39, 0x43, 0xe1,  
0x3a, 0xb5, 0x56, 0x2a, 0xc0, 0x6d, 0xb3, 0x05, 0x22, 0x66, 0xbf, 0xdc, 0x0b, 0xfa, 0x62, 0x48,  
0xdd, 0x20, 0x11, 0x06, 0x36, 0xc9, 0xc1, 0xcf, 0xf6, 0x27, 0x52, 0xbb, 0x69, 0xf5, 0xd4, 0x87,  
0x7f, 0x84, 0x4c, 0xd2, 0x9c, 0x57, 0xa4, 0xbc, 0x4f, 0x9a, 0xdf, 0xfe, 0xd6, 0x8d, 0x7a, 0xeb,  
0x2b, 0x53, 0xd8, 0x5c, 0xa1, 0x14, 0x17, 0xfb, 0x23, 0xd5, 0x7d, 0x30, 0x67, 0x73, 0x08, 0x09,  
0xee, 0xb7, 0x70, 0x3f, 0x61, 0xb2, 0x19, 0x8e, 0x4e, 0xe5, 0x4b, 0x93, 0x8f, 0x5d, 0xdb, 0xa9,  
0xad, 0xf1, 0xae, 0x2e, 0xcb, 0x0d, 0xfc, 0xf4, 0x2d, 0x46, 0x6e, 0x1d, 0x97, 0xe8, 0xd1, 0xe9,  
0x4d, 0x37, 0xa5, 0x75, 0x5e, 0x83, 0x9e, 0xab, 0x82, 0x9d, 0xb9, 0x1c, 0xe0, 0xcd, 0x49, 0x89,  
0x01, 0xb6, 0xbd, 0x58, 0x24, 0xa2, 0x5f, 0x38, 0x78, 0x99, 0x15, 0x90, 0x50, 0xb8, 0x95, 0xe4,  
0xd0, 0x91, 0xc7, 0xce, 0xed, 0x0f, 0xb4, 0x6f, 0xa0, 0xcc, 0xf0, 0x02, 0x4a, 0x79, 0xc3, 0xde,  
0xa3, 0xef, 0xea, 0x51, 0xe6, 0x6b, 0x18, 0xec, 0x1b, 0x2c, 0x80, 0xf7, 0x74, 0xe7, 0xff, 0x21,  
0x5a, 0x6a, 0x54, 0x1e, 0x41, 0x31, 0x92, 0x35, 0xc4, 0x33, 0x07, 0x0a, 0xba, 0x7e, 0x0e, 0x34,  
0x88, 0xb1, 0x98, 0x7c, 0xf3, 0x3d, 0x60, 0x6c, 0x7b, 0xca, 0xd3, 0x1f, 0x32, 0x65, 0x04, 0x28,  
0x64, 0xbe, 0x85, 0x9b, 0x2f, 0x59, 0x8a, 0xd7, 0xb0, 0x25, 0xac, 0xaf, 0x12, 0x03, 0xe2, 0xf2}.

### 8.6.6 LINEAR TRANSFORMS $L_1$ AND $L_2$

Both  $L_1$  and  $L_2$  are linear transforms of 32-bit strings, defined as follows:

$$L_1(X) = X \oplus (X \lll_{32} 2) \oplus (X \lll_{32} 10) \oplus (X \lll_{32} 18) \oplus (X \lll_{32} 24),$$

$$L_2(X) = \oplus (X \lll_{32} 8) \oplus (X \lll_{32} 14) \oplus (X \lll_{32} 22) \oplus (X \lll_{32} 30).$$

A1

## Annex A (normative)

### Object Identifiers

This annex lists the object identifiers assigned to algorithms specified in this part of ISO/IEC 18033 and defines algorithm parameter structures. Please refer to ISO/IEC 18033-3 for Object IDs for modes of a block cipher.

```
A1 EncryptionAlgorithms-4 {
    iso(1) standard(0) encryption-algorithms(18033) part(4)
        asn1-module(0) algorithm-object-identifiers(0) }
    DEFINITIONS EXPLICIT TAGS ::= BEGIN

-- EXPORTS All; --

-- IMPORTS None; --

OID ::= OBJECT IDENTIFIER -- Alias

-- Synonyms --

is18033-4 OID ::= { iso(1) standard(0) is18033(18033) part4(4) }

id-kg OID ::= { is18033-4 keystream-generator(1) }
id-scmode OID ::= { is18033-4 stream-cipher-mode(2) }

-- Assignments --

id-kg-mugi OID ::= { id-kg mugi(1) }
id-kg-snow OID ::= { id-kg snow(2) }
id-kg-rabbit OID ::= { id-kg rabbit(3) }
id-kg-decim2 OID ::= { id-kg decim2(4) }
id-kg-k2 OID ::= { id-kg k2(5) }
id-kg-zuc OID ::= { id-kg zuc(6) }

id-scmode-additive OID ::= { id-scmode additive(1) }
id-scmode-multis01 OID ::= { id-scmode multis01(2) }

-- Algorithms and parameters --

StreamCipher ::= AlgorithmIdentifier {{ StreamCipherAlgorithms }}

StreamCipherAlgorithms ALGORITHM ::= {
    additiveStreamCipher |
    multiS01StreamCipher,

... -- Expect additional algorithms --
}

additiveStreamCipher ALGORITHM ::= {
    OID id-scmode-additive PARMS AdditiveStreamCipherParameters
}

AdditiveStreamCipherParameters ::= KeyGenerator

multiS01StreamCipher ALGORITHM ::= {
    OID id-scmode-multis01 PARMS MultiS01StreamCipherParameters
}

MultiS01StreamCipherParameters ::= SEQUENCE {
    keyGenerator KeyGenerator,
```

```

securityParameter INTEGER DEFAULT 64,
irreduciblePolynomial BIT STRING,
redundancy BIT STRING,
publicParameterR BIT STRING
    -- length determined by securityParameter
    -- for full interoperability multis01 parameters should
    -- include the padding method but they do not have object
    -- identifiers. for the time being they will have to be
    -- negotiated in an application-dependent way
}

KeyGenerator ALGORITHM ::= {
    mugiKeyGenerator |
    snowKeyGenerator |
    rabbitKeyGenerator |
    decim2KeyGenerator |
    k2KeyGenerator |
    zucKeyGenerator,

    ... -- Expect additional algorithms --
}

mugiKeyGenerator ALGORITHM ::= {
    OID id-kg-mugi PARMS NullParameters
}

snowKeyGenerator ALGORITHM ::= {
    OID id-kg-snow PARMS NullParameters
}

rabbitKeyGenerator ALGORITHM ::= {
    OID id-kg-rabbit PARMS NullParameters
}

decim2KeyGenerator ALGORITHM ::= {
    OID id-kg-decim2 PARMS NullParameters
}

k2KeyGenerator ALGORITHM ::= {
    OID id-kg-k2 PARMS NullParameters
}

zucKeyGenerator ALGORITHM ::= {
    OID id-kg-zuc PARMS NullParameters
}

NullParameters ::= NULL

-- Cryptographic algorithm identification --

ALGORITHM ::= CLASS {
    &id OBJECT IDENTIFIER UNIQUE,
    &Type OPTIONAL
}
    WITH SYNTAX { OID &id [PARMS &Type] }

AlgorithmIdentifier { ALGORITHM:IOSet } ::= SEQUENCE {
    algorithm ALGORITHM.&id( {IOSet} ),
    parameters ALGORITHM.&Type( {IOSet} {@algorithm} ) OPTIONAL
}

END -- EncryptionAlgorithms-4 -- A1

```

## Annex B (informative)

### Operations over the finite field $GF(2^n)$

For any positive integer  $n$  there exists a finite field containing exactly  $2^n$  elements. This field is unique up to isomorphism, and in this part of ISO/IEC 18033 it is referred to as the finite field  $GF(2^n)$ .

In the polynomial representation, each element of  $GF(2^n)$  is represented by a binary polynomial of degree less than  $n$ . More explicitly the bit string  $a = a_{n-1} \dots a_2 a_1 a_0$  is taken to represent the binary polynomial  $a(x) = a_{n-1}x^{n-1} + \dots + a_2x^2 + a_1x + a_0$ . The polynomial basis is the set  $B = (x^{n-1}, \dots, x^2, x, 1)$ . For two bit strings  $a = a_{n-1} \dots a_2 a_1 a_0$  and  $b = b_{n-1} \dots b_2 b_1 b_0$ , the sum is  $c = a \oplus b = c_{n-1} \dots c_2 c_1 c_0$ , where  $c_i = a_i \oplus b_i$ .

Multiplication in the finite field, written  $a \otimes b$ , corresponds to the multiplication of two polynomials  $a(x)b(x)$  modulo a binary irreducible polynomial  $p(x)$  of degree  $n$ . A polynomial is irreducible if it has no non-trivial divisors.

$GF(2^n) \setminus \{0\}$  denoted as  $GF(2^n)^*$  is an abelian group with respect to multiplication and the identity is 1. For any non-zero binary polynomial  $b(x)$  of degree less than  $n$ , the multiplicative inverse of  $b(x)$ , denoted  $b^{-1}(x)$ , can be computed as follows: the extended Euclidean algorithm is used to compute polynomials  $a(x)$  and  $c(x)$  such that  $b(x) \cdot a(x) + p(x) \cdot c(x) = 1$ . Hence,  $a(x) \cdot b(x) \bmod p(x) = 1$ , which means  $b^{-1}(x) = a(x) \bmod p(x)$ . The extended Euclidean algorithm is described in [15].

## Annex C (informative)

### Examples

#### C.1 Example for MUGI

##### C.1.1 Key, initialization vector, and keystream triplets

K = 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00  
IV= 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00  
Z = c7 6e 14 e7 08 36 e6 b6 cb 0e 9c 5a 0b f0 3e 1e 0a cf 9a f4 9e be 6d 67 d5 72 6e 37 4b 13 97 ac.

K = 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00  
IV= 34 61 69 88 51 81 21 39 01 55 00 a5 3b 7e 59 87  
Z = 2a a1 c5 c7 20 73 b1 b3 a9 d1 0d c6 85 50 66 10 28 30 56 0d 9a 24 65 c9 9c 29 1c 13 81 4e 08 8d.

K = 51 34 00 b1 04 a0 59 91 30 ad 00 fc 48 d7 59 e0  
IV= 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00  
Z = bd df ad 5f 04 b8 86 25 c3 ad ac e1 56 d1 c1 99 36 ff a4 e9 a7 fd f7 5a aa b8 29 13 42 85 aa 4b.

K = 69 e7 06 ee 52 95 37 2c 75 13 01 47 30 23 79 93  
IV= 2a 00 45 c8 49 27 49 d5 3a 9b 16 4a 25 e4 49 15  
Z = e3 cc 67 a0 25 5b 0f 28 2d 9a 5b 1b bd f7 f2 df 84 eb 46 f6 07 d6 e6 dd 32 86 13 43 94 dd 95 fb.

##### C.1.2 Sample internal states

K = 00 01 02 03 04 05 06 07 08 09 0a 0b 0c 0d 0e 0f  
IV= f0 e0 d0 c0 b0 a0 90 80 70 60 50 40 30 20 10 00  
Z = bc 62 43 06 14 b7 9b 71 71 a6 66 81 c3 55 42 de 7a ba 5b 4f b8 0e 82 d7 0b 96 98 28 90 b6 e1 43

Intermediate values of the internal state

rho function 0

a: 0001020304050607 08090a0b0c0d0e0f 7498f5f1e727d094  
b: 0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 0000000000000000 0000000000000000 0000000000000000

rho function 1

a: 08090a0b0c0d0e0f 9724d9144c5d8926 64b47311d52100a5  
b: 0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 0000000000000000 0000000000000000 08090a0b0c0d0e0f

rho function 2

a: 9724d9144c5d8926 09671cfbcfaa95fb e2d338166cd8c441  
b: 0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 0000000000000000 9724d9144c5d8926 08090a0b0c0d0e0f



rho function 3

a: 09671cfbcfaa95fb 9c0c2097edb20067 6ef29c62b7691210  
b: 0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 09671cfbcfaa95fb 9724d9144c5d8926 08090a0b0c0d0e0f

rho function 4

a: 9c0c2097edb20067 c08ee4dcb2d08591 201239b2b04d5d6a  
b: 0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 0000000000000000 0000000000000000 0000000000000000  
9c0c2097edb20067 09671cfbcfaa95fb 9724d9144c5d8926 08090a0b0c0d0e0f

rho function 5

a: c08ee4dcb2d08591 738177859f3210f6 48963357b89312eb  
b: 0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 0000000000000000 0000000000000000 c08ee4dcb2d08591  
9c0c2097edb20067 09671cfbcfaa95fb 9724d9144c5d8926 08090a0b0c0d0e0f

rho function 6

a: 738177859f3210f6 b36b4d944f5d04cb bc7ac7e83f40cca1  
b: 0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 0000000000000000 738177859f3210f6 c08ee4dcb2d08591  
9c0c2097edb20067 09671cfbcfaa95fb 9724d9144c5d8926 08090a0b0c0d0e0f

rho function 7

a: b36b4d944f5d04cb 2d13c00221057d8d 65e12d98fb29feca  
b: 0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 b36b4d944f5d04cb 738177859f3210f6 c08ee4dcb2d08591  
9c0c2097edb20067 09671cfbcfaa95fb 9724d9144c5d8926 08090a0b0c0d0e0f

rho function 8

a: 2d13c00221057d8d 20ead0479e63cdc3 7169edbc504968d2  
b: 0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 0000000000000000 0000000000000000 0000000000000000  
2d13c00221057d8d b36b4d944f5d04cb 738177859f3210f6 c08ee4dcb2d08591  
9c0c2097edb20067 09671cfbcfaa95fb 9724d9144c5d8926 08090a0b0c0d0e0f

rho function 9

a: 20ead0479e63cdc3 591a6857e3112cee 8269181ee80366a1  
b: 0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 0000000000000000 0000000000000000 20ead0479e63cdc3  
2d13c00221057d8d b36b4d944f5d04cb 738177859f3210f6 c08ee4dcb2d08591  
9c0c2097edb20067 09671cfbcfaa95fb 9724d9144c5d8926 08090a0b0c0d0e0f

rho function 10

a: 591a6857e3112cee dfbbb88c02c9c80a fa312d220ef73c78  
b: 0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 0000000000000000 591a6857e3112cee 20ead0479e63cdc3  
2d13c00221057d8d b36b4d944f5d04cb 738177859f3210f6 c08ee4dcb2d08591  
9c0c2097edb20067 09671cfbcfaa95fb 9724d9144c5d8926 08090a0b0c0d0e0f

rho function 11

a: dfbbb88c02c9c80a 5cc4835080bc5321 78e69bd217041ca7  
b: 0000000000000000 0000000000000000 0000000000000000 0000000000000000  
0000000000000000 dfbbb88c02c9c80a 591a6857e3112cee 20ead0479e63cdc3

2d13c00221057d8d b36b4d944f5d04cb 738177859f3210f6 c08ee4dcb2d08591  
9c0c2097edb20067 09671cfbcfaa95fb 9724d9144c5d8926 08090a0b0c0d0e0f

rho function 12

a: 5cc4835080bc5321 fd5755df9cc0ceb9 dd032b76f3534504  
b: 0000000000000000 0000000000000000 0000000000000000 0000000000000000  
5cc4835080bc5321 dfbbb88c02c9c80a 591a6857e3112cee 20ead0479e63cdc3  
2d13c00221057d8d b36b4d944f5d04cb 738177859f3210f6 c08ee4dcb2d08591  
9c0c2097edb20067 09671cfbcfaa95fb 9724d9144c5d8926 08090a0b0c0d0e0f

rho function 13

a: fd5755df9cc0ceb9 c905d08f50fa71db cfc255e594b38ee  
b: 0000000000000000 0000000000000000 0000000000000000 fd5755df9cc0ceb9  
5cc4835080bc5321 dfbbb88c02c9c80a 591a6857e3112cee 20ead0479e63cdc3  
2d13c00221057d8d b36b4d944f5d04cb 738177859f3210f6 c08ee4dcb2d08591  
9c0c2097edb20067 09671cfbcfaa95fb 9724d9144c5d8926 08090a0b0c0d0e0f

rho function 14

a: c905d08f50fa71db bfe2485ac2696cc7 0a77652c7dbcc580  
b: 0000000000000000 0000000000000000 c905d08f50fa71db fd5755df9cc0ceb9  
5cc4835080bc5321 dfbbb88c02c9c80a 591a6857e3112cee 20ead0479e63cdc3  
2d13c00221057d8d b36b4d944f5d04cb 738177859f3210f6 c08ee4dcb2d08591  
9c0c2097edb20067 09671cfbcfaa95fb 9724d9144c5d8926 08090a0b0c0d0e0f

rho function 15

a: bfe2485ac2696cc7 7dea261cb61d4fea 3991ce48e105a4a1  
b: 0000000000000000 bfe2485ac2696cc7 c905d08f50fa71db fd5755df9cc0ceb9  
5cc4835080bc5321 dfbbb88c02c9c80a 591a6857e3112cee 20ead0479e63cdc3  
2d13c00221057d8d b36b4d944f5d04cb 738177859f3210f6 c08ee4dcb2d08591  
9c0c2097edb20067 09671cfbcfaa95fb 9724d9144c5d8926 08090a0b0c0d0e0f

buffer init

a: 7dea261cb61d4fea eafb528479bb687d eb8189612089ff0b  
b: 7dea261cb61d4fea bfe2485ac2696cc7 c905d08f50fa71db fd5755df9cc0ceb9  
5cc4835080bc5321 dfbbb88c02c9c80a 591a6857e3112cee 20ead0479e63cdc3  
2d13c00221057d8d b36b4d944f5d04cb 738177859f3210f6 c08ee4dcb2d08591  
9c0c2097edb20067 09671cfbcfaa95fb 9724d9144c5d8926 08090a0b0c0d0e0f

rho function 0

a: 8d0af6dc06bddf6a 9a9b02c4499b787d f100cffe031d365b

rho function 1

a: 9a9b02c4499b787d 435407f3bbc2c760 b8576326c43c7141

rho function 2

a: 435407f3bbc2c760 b5117172dcf5e507 10d44d672b0cb32b

rho function 3

a: b5117172dcf5e507 9157292760b2892f 45de3e448a22a274

rho function 4

a: 9157292760b2892f aee0542493e7889e d92646e5bf6e90fd

rho function 5

a: aee0542493e7889e a9f2f7fac6cff1ff 668ac5cf634db73d

rho function 6

a: a9f2f7fac6cff1ff 9cb8969f9fc84dc6 d3db5a83153c2d75

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rho function 7

a: 9cb8969f9fc84dc6 b1260b2ec980a340 4c06fba0602d20da

rho function 8

a: b1260b2ec980a340 192a6fd877969848 4e9d5f10f22daa44

rho function 9

a: 192a6fd877969848 bbac287d38601209 c31e21b47993441d

rho function 10

a: bbac287d38601209 d58486545129be34 88b995cf25723d71

rho function 11

a: d58486545129be34 c8af8f1422e98119 7cb36f5145a5f171

rho function 12

a: c8af8f1422e98119 00bba312081aa445 2e8517e066c8b410

rho function 13

a: 00bba312081aa445 2f3864a9c279a14c 4e1balaafc06cb55

rho function 14

a: 2f3864a9c279a14c 6551f5e9cbc1e0d7 acf8aaa64583d0d7

rho function 15

a: 6551f5e9cbc1e0d7 4e466dffcb92db48 4a8ffe073636f5c3

state init

a: 4e466dffcb92db48 f5eb67b928359d8b 5d3c31a0af9cd78f  
b: 7dea261cb61d4fea bfe2485ac2696cc7 c905d08f50fa71db fd5755df9cc0ceb9  
5cc4835080bc5321 dfbbb88c02c9c80a 591a6857e3112cee 20ead0479e63cdc3  
2d13c00221057d8d b36b4d944f5d04cb 738177859f3210f6 c08ee4dcb2d08591  
9c0c2097edb20067 09671cfbcfaa95fb 9724d9144c5d8926 08090a0b0c0d0e0f

update 1

a: f5eb67b928359d8b ace6a90bde0af786 529108c358fa4ada  
b: 464f67f4c79fd547 7dea261cb61d4fea bfe2485ac2696cc7 c905d08f50fa71db  
ddbd859802a3037a 5cc4835080bc5321 dfbbb88c02c9c80a 591a6857e3112cee  
20ead0479e63cdc3 2d13c00221057d8d 7cc1d86f463a1830 738177859f3210f6  
c08ee4dcb2d08591 9c0c2097edb20067 09671cfbcfaa95fb 9724d9144c5d8926

update 2

a: ace6a90bde0af786 9fa7a15367ae1667 5f241cf311a0bfa7  
b: 62cfbead646814ad 464f67f4c79fd547 7dea261cb61d4fea bfe2485ac2696cc7  
901fb8d8b3eb5d35 ddbd859802a3037a 5cc4835080bc5321 dfbbb88c02c9c80a  
591a6857e3112cee 20ead0479e63cdc3 c0a1c065bd095d1a 7cc1d86f463a1830  
738177859f3210f6 c08ee4dcb2d08591 9c0c2097edb20067 09671cfbcfaa95fb

update 3

a: 9fa7a15367ae1667 75195c2e249e4399 8bd43dd671ad8b05  
b: a581b5f011a0627d 62cfbead646814ad 464f67f4c79fd547 7dea261cb61d4fea  
6059f0d6c0a0a4cd 901fb8d8b3eb5d35 ddbd859802a3037a 5cc4835080bc5321  
dfbbb88c02c9c80a 591a6857e3112cee 923a55d65eed291f c0a1c065bd095d1a  
7cc1d86f463a1830 738177859f3210f6 c08ee4dcb2d08591 9c0c2097edb20067

update 4

a: 75195c2e249e4399 9b2239a28cc5a4e3 5554ab4e803a0a19  
b: 03ab81c48a1c1600 a581b5f011a0627d 62cfbead646814ad 464f67f4c79fd547  
212ea54c36a11ccb 6059f0d6c0a0a4cd 901fb8d8b3eb5d35 ddbd859802a3037a

5cc4835080bc5321 dfbbb88c02c9c80a c62878a190905b6b 923a55d65eed291f  
c0a1c065bd095d1a 7cc1d86f463a1830 738177859f3210f6 c08ee4dcb2d08591

#### update 5

a: 9b2239a28cc5a4e3 80c1d0bc18b29e62 4b4f363e4a322d0e  
b: b597b8f2964ec608 03ab81c48a1c1600 a581b5f011a0627d 62cfbead646814ad  
9bf2e26cc53cd63d 212ea54c36a11ccb 6059f0d6c0a0a4cd 901fb8d8b3eb5d35  
ddb859802a3037a 5cc4835080bc5321 9981a0bc7e081065 c62878a190905b6b  
923a55d65eed291f c0a1c065bd095d1a 7cc1d86f463a1830 738177859f3210f6

#### update 6

a: 80c1d0bc18b29e62 dfe2225186dfbca3 1277ae469887f627  
b: e8a34e2713f7b415 b597b8f2964ec608 03ab81c48a1c1600 a581b5f011a0627d  
f2d00675d7834998 9bf2e26cc53cd63d 212ea54c36a11ccb 6059f0d6c0a0a4cd  
901fb8d8b3eb5d35 ddb859802a3037a e1cdde4a401d9344 9981a0bc7e081065  
c62878a190905b6b 923a55d65eed291f c0a1c065bd095d1a 7cc1d86f463a1830

#### update 7

a: dfe2225186dfbca3 ca86cdc9d2bf2007 7a5c92de7be1811f  
b: fc0008d35e888652 e8a34e2713f7b415 b597b8f2964ec608 03ab81c48a1c1600  
c5d84526d100c6b0 f2d00675d7834998 9bf2e26cc53cd63d 212ea54c36a11ccb  
6059f0d6c0a0a4cd 901fb8d8b3eb5d35 8350ac87909956ac e1cdde4a401d9344  
9981a0bc7e081065 c62878a190905b6b 923a55d65eed291f c0a1c065bd095d1a

#### update 8

a: ca86cdc9d2bf2007 0870fbf8e065b266 067f3c2be88481d4  
b: 1f43e2343bd6e1b9 fc0008d35e888652 e8a34e2713f7b415 b597b8f2964ec608  
22852488bcbd0acb c5d84526d100c6b0 f2d00675d7834998 9bf2e26cc53cd63d  
212ea54c36a11ccb 6059f0d6c0a0a4cd 008fe3b375c32594 8350ac87909956ac  
e1cdde4a401d9344 9981a0bc7e081065 c62878a190905b6b 923a55d65eed291f

#### update 9

a: 0870fbf8e065b266 73b145404394710d d756724ed3994273  
b: 58bc981f8c520918 1f43e2343bd6e1b9 fc0008d35e888652 e8a34e2713f7b415  
2e655a9e53721035 22852488bcbd0acb c5d84526d100c6b0 f2d00675d7834998  
9bf2e26cc53cd63d 212ea54c36a11ccb 1e51e0b359210471 008fe3b375c32594  
8350ac87909956ac e1cdde4a401d9344 9981a0bc7e081065 c62878a190905b6b

#### update 10

a: 73b145404394710d 82d164adcac96d62 0607785b7d152b8b  
b: ce58835970f5e90d 58bc981f8c520918 1f43e2343bd6e1b9 fc0008d35e888652  
1a734852c474fd8d 2e655a9e53721035 22852488bcbd0acb c5d84526d100c6b0  
f2d00675d7834998 9bf2e26cc53cd63d 61333608d76cc281 1e51e0b359210471  
008fe3b375c32594 8350ac87909956ac e1cdde4a401d9344 9981a0bc7e081065

#### update 11

a: 82d164adcac96d62 c14072735c68e7e9 f61c61bbde49ed28  
b: ea30e5fc3d9c6168 ce58835970f5e90d 58bc981f8c520918 1f43e2343bd6e1b9  
39d84df58f8840e2 1a734852c474fd8d 2e655a9e53721035 22852488bcbd0acb  
c5d84526d100c6b0 f2d00675d7834998 0b6bb4c0466c7aba 61333608d76cc281  
1e51e0b359210471 008fe3b375c32594 8350ac87909956ac e1cdde4a401d9344

#### update 12

a: c14072735c68e7e9 ce0bee4623950852 af052447a7444e65  
b: 631cbae78ad4fe26 ea30e5fc3d9c6168 ce58835970f5e90d 58bc981f8c520918  
3dc6c6bc876beb72 39d84df58f8840e2 1a734852c474fd8d 2e655a9e53721035  
22852488bcbd0acb c5d84526d100c6b0 871323e1d70caa2b 0b6bb4c0466c7aba  
61333608d76cc281 1e51e0b359210471 008fe3b375c32594 8350ac87909956ac

#### update 13

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a: ce0bee4623950852 f6c22506fc93fb5a 9eb296971244bcb3  
b: 4210def4ccf1b145 631cbae78ad4fe26 ea30e5fc3d9c6168 ce58835970f5e90d  
76d9c281df20192d 3dc6c6bc876beb72 39d84df58f8840e2 1a734852c474fd8d  
2e655a9e53721035 22852488bcbd0acb 9cf94157cf512603 871323e1d70caa2b  
0b6bb4c0466c7aba 61333608d76cc281 1e51e0b359210471 008fe3b375c32594

update 14  
a: f6c22506fc93fb5a b36f504b7eb67fe6 a66ba7dd058722d3  
b: ce840df556562dc6 4210def4ccf1b145 631cbae78ad4fe26 ea30e5fc3d9c6168  
d42bcb0bb4811480 76d9c281df20192d 3dc6c6bc876beb72 39d84df58f8840e2  
1a734852c474fd8d 2e655a9e53721035 f5e9e609dd8e3cc3 9cf94157cf512603  
871323e1d70caa2b 0b6bb4c0466c7aba 61333608d76cc281 1e51e0b359210471

update 15  
a: b36f504b7eb67fe6 0ce5a4d1a0cbc0f7 bd0c30563f8ee4f7  
b: e893c5b5a5b2ff2b ce840df556562dc6 4210def4ccf1b145 631cbae78ad4fe26  
d3e8a809b214218a d42bcb0bb4811480 76d9c281df20192d 3dc6c6bc876beb72  
39d84df58f8840e2 1a734852c474fd8d 680920245819a4f5 f5e9e609dd8e3cc3  
9cf94157cf512603 871323e1d70caa2b 0b6bb4c0466c7aba 61333608d76cc281

update 16  
a: 0ce5a4d1a0cbc0f7 316993816117e50f bc62430614b79b71  
b: d25c6643a9dabd67 e893c5b5a5b2ff2b ce840df556562dc6 4210def4ccf1b145  
5eda7c5b0dbf1554 d3e8a809b214218a d42bcb0bb4811480 76d9c281df20192d  
3dc6c6bc876beb72 39d84df58f8840e2 cd7fe2794367de6c 680920245819a4f5  
f5e9e609dd8e3cc3 9cf94157cf512603 871323e1d70caa2b 0b6bb4c0466c7aba

update 1  
a: 316993816117e50f 4f7c747ce422e686 71a66681c35542de  
b: 078e1011e6a7ba4d d25c6643a9dabd67 e893c5b5a5b2ff2b ce840df556562dc6  
34c91c7513d1a868 5eda7c5b0dbf1554 d3e8a809b214218a d42bcb0bb4811480  
76d9c281df20192d 3dc6c6bc876beb72 f6896bf6137101b5 cd7fe2794367de6c  
680920245819a4f5 f5e9e609dd8e3cc3 9cf94157cf512603 871323e1d70caa2b

update 2  
a: 4f7c747ce422e686 0aeab5f525c1a62f 7aba5b4fb80e82d7  
b: b67ab060b61b4f24 078e1011e6a7ba4d d25c6643a9dabd67 e893c5b5a5b2ff2b  
1aafc6fee2d73946 34c91c7513d1a868 5eda7c5b0dbf1554 d3e8a809b214218a  
d42bcb0bb4811480 76d9c281df20192d e048fa7f72820d7b f6896bf6137101b5  
cd7fe2794367de6c 680920245819a4f5 f5e9e609dd8e3cc3 9cf94157cf512603

update 3  
a: 0aeab5f525c1a62f bd1a2938a57319c8 0b96982890b6e143  
b: d385352b2b73c085 b67ab060b61b4f24 078e1011e6a7ba4d d25c6643a9dabd67  
3b7b6dbc17a6deal 1aafc6fee2d73946 34c91c7513d1a868 5eda7c5b0dbf1554  
d3e8a809b214218a d42bcb0bb4811480 2ec06674b7293909 e048fa7f72820d7b  
f6896bf6137101b5 cd7fe2794367de6c 680920245819a4f5 f5e9e609dd8e3cc3

update 4  
a: bd1a2938a57319c8 e4684a2bf28ff50d 4930b5d033157f46  
b: ff0353fcf84f9aec d385352b2b73c085 b67ab060b61b4f24 078e1011e6a7ba4d  
8c861a18a465a833 3b7b6dbc17a6deal 1aafc6fee2d73946 34c91c7513d1a868  
5eda7c5b0dbf1554 d3e8a809b214218a 974c156779fef6f9 2ec06674b7293909  
e048fa7f72820d7b f6896bf6137101b5 cd7fe2794367de6c 680920245819a4f5

## C.2 128-bit key example for SNOW 2.0

### C.2.1 Key, initialization vector, and keystream triplets

(IV<sub>3</sub>, IV<sub>2</sub>, IV<sub>1</sub>, IV<sub>0</sub>) = (0, 0, 0, 0), key = 80000000000000000000000000000000  
Keystream output: 8D590AE9A74A7D056DC9CA74B72D1A4599B0A083FB45D13FCF9411BD9A503783.

(IV<sub>3</sub>, IV<sub>2</sub>, IV<sub>1</sub>, IV<sub>0</sub>) = (0, 0, 0, 0), key = AAAAAAAAAAAAAAAAAAAAAAAAAAAAAA  
Keystream output: E00982F525F02054214992D8706F2B20DA585E5B85E2746D09F22681B2749407.

(IV<sub>3</sub>, IV<sub>2</sub>, IV<sub>1</sub>, IV<sub>0</sub>) = (4, 3, 2, 1), key = 80000000000000000000000000000000  
Keystream output: D6403358E0354A6957F43FCE44B4B13FF78E24C246618A0767AC83C10BFC45F0.

(IV<sub>3</sub>, IV<sub>2</sub>, IV<sub>1</sub>, IV<sub>0</sub>) = (4, 3, 2, 1), key = AAAAAAAAAAAAAAAAAAAAAAAAAAAAAA  
Keystream output: C355385DB31D6CBDF774AF5366C2E8774DEADAC7DC7229DFED171D7B.

### C.2.2 Sample internal states

K = 80 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00  
IV= 00 00 00 04 00 00 00 03 00 00 00 02 00 00 00 01  
Z = d6 40 33 58 e0 35 4a 69 57 f4 3f ce 44 b4 b1 3f f7 8e 24 c2 46 61 8a 07 67 ac 83 c1 0b fc 45 f0

Snow 2.0 Internal state at time -34  
15:80000000 14:00000000 13:00000000 12:00000000 11:7fffffff 10:ffffffff 09:ffffffff 08:ffffffff  
07:80000000 06:00000000 05:00000000 04:00000000 03:7fffffff 02:ffffffff 01:ffffffff 00:ffffffff

Snow 2.0 Internal state at time -33  
15:80000001 14:00000000 13:00000000 12:00000002 11:7fffffff 10:fffffffc 09:fffffffb 08:ffffffff  
07:80000000 06:00000000 05:00000000 04:00000000 03:7fffffff 02:ffffffff 01:ffffffff 00:ffffffff  
R1:00000000 R2:00000000

Snow 2.0 Internal state at time -32  
15:09dfef08 14:80000001 13:00000000 12:00000000 11:00000002 10:7fffffff 09:fffffffc 08:fffffffb  
07:ffffffff 06:80000000 05:00000000 04:00000000 03:00000000 02:7fffffff 01:ffffffff 00:ffffffff  
R1:00000000 R2:63636363

Snow 2.0 Internal state at time -31  
15:e5f1b94c 14:09dfef08 13:80000001 12:00000000 11:00000000 10:00000002 09:7fffffff 08:fffffffc  
07:fffffffb 06:ffffffff 05:80000000 04:00000000 03:00000000 02:00000000 01:7fffffff 00:ffffffff  
R1:63636363 R2:63636363

Snow 2.0 Internal state at time -30  
15:ea9a3527 14:e5f1b94c 13:09dfef08 12:80000001 11:00000000 10:00000000 09:00000002 08:7fffffff  
07:fffffffc 06:fffffffb 05:ffffffff 04:80000000 03:00000000 02:00000000 01:00000000 00:7fffffff  
R1:e3636363 R2:fbfbfbfb

Snow 2.0 Internal state at time -29  
15:a69fa10d 14:ea9a3527 13:e5f1b94c 12:09dfef08 11:80000001 10:00000000 09:00000000 08:00000002  
07:7fffffff 06:fffffffc 05:fffffffb 04:ffffffff 03:80000000 02:00000000 01:00000000 00:00000000  
R1:fbfbfbfa R2:34de1111

Snow 2.0 Internal state at time -28  
15:8ecacddb 14:a69fa10d 13:ea9a3527 12:e5f1b94c 11:09dfef08 10:80000001 09:00000000 08:00000000  
07:00000002 06:7fffffff 05:fffffffc 04:fffffffb 03:ffffffff 02:80000000 01:00000000 00:00000000  
R1:34de110c R2:692d2d4b

Snow 2.0 Internal state at time -27  
15:eaf4d48f 14:8ecacddb 13:a69fa10d 12:ea9a3527 11:e5f1b94c 10:09dfef08 09:80000001 08:00000000  
07:00000000 06:00000002 05:7fffffff 04:fffffffc 03:fffffffb 02:ffffffff 01:80000000 00:00000000

# BS ISO/IEC 18033-4:2011+A1:2020

## ISO/IEC 18033-4:2011+A1:2020(E)

R1:692d2d47 R2:b66ede7f

Snow 2.0 Internal state at time -26

15:19805681 14:eaf4d48f 13:8ecacddb 12:a69fa10d 11:ea9a3527 10:e5f1b94c 09:09dfef08 08:80000001  
07:00000000 06:00000000 05:00000002 04:7fffffff 03:fffffffc 02:fffffffb 01:fffffffc 00:80000000  
R1:366ede7e R2:12c38109

Snow 2.0 Internal state at time -25

15:e8688f55 14:19805681 13:eaf4d48f 12:8ecacddb 11:a69fa10d 10:ea9a3527 09:e5f1b94c 08:09dfef08  
07:80000001 06:00000000 05:00000000 04:00000002 03:7fffffff 02:fffffffc 01:fffffffb 00:fffffffc  
R1:12c3810b R2:86c47640

Snow 2.0 Internal state at time -24

15:fa565ef1 14:e8688f55 13:19805681 12:eaf4d48f 11:8ecacddb 10:a69fa10d 09:ea9a3527 08:e5f1b94c  
07:09dfef08 06:80000001 05:00000000 04:00000000 03:00000002 02:7fffffff 01:fffffffc 00:fffffffb  
R1:86c47640 R2:d63b88a5

Snow 2.0 Internal state at time -23

15:6c7a94aa 14:fa565ef1 13:e8688f55 12:19805681 11:eaf4d48f 10:8ecacddb 09:a69fa10d 08:ea9a3527  
07:e5f1b94c 06:09dfef08 05:80000001 04:00000000 03:00000000 02:00000002 01:7fffffff 00:fffffffc  
R1:d63b88a5 R2:b7c51902

Snow 2.0 Internal state at time -22

15:5ced2805 14:6c7a94aa 13:fa565ef1 12:e8688f55 11:19805681 10:eaf4d48f 09:8ecacddb 08:a69fa10d  
07:ea9a3527 06:e5f1b94c 05:09dfef08 04:80000001 03:00000000 02:00000000 01:00000002 00:7fffffff  
R1:37c51903 R2:db1c5e4f

Snow 2.0 Internal state at time -21

15:26ac1e81 14:5ced2805 13:6c7a94aa 12:fa565ef1 11:e8688f55 10:19805681 09:eaf4d48f 08:8ecacddb  
07:a69fa10d 06:ea9a3527 05:e5f1b94c 04:09dfef08 03:80000001 02:00000000 01:00000000 00:00000002  
R1:e4fc4d57 R2:d04da3ad

Snow 2.0 Internal state at time -20

15:2e5cc1a4 14:26ac1e81 13:5ced2805 12:6c7a94aa 11:fa565ef1 10:e8688f55 09:19805681 08:eaf4d48f  
07:8ecacddb 06:a69fa10d 05:ea9a3527 04:e5f1b94c 03:09dfef08 02:80000001 01:00000000 00:00000000  
R1:b63f5cf9 R2:6c782451

Snow 2.0 Internal state at time -19

15:2eb71be8 14:2e5cc1a4 13:26ac1e81 12:5ced2805 11:6c7a94aa 10:fa565ef1 09:e8688f55 08:19805681  
07:eaf4d48f 06:8ecacddb 05:a69fa10d 04:ea9a3527 03:e5f1b94c 02:09dfef08 01:80000001 00:00000000  
R1:57125978 R2:13ebdccc

Snow 2.0 Internal state at time -18

15:dc33fa8c 14:2eb71be8 13:2e5cc1a4 12:26ac1e81 11:5ced2805 10:6c7a94aa 09:fa565ef1 08:e8688f55  
07:19805681 06:eaf4d48f 05:8ecacddb 04:a69fa10d 03:ea9a3527 02:e5f1b94c 01:09dfef08 00:80000001  
R1:ba8b7dd9 R2:6b132ab7

Snow 2.0 Internal state at time -17

15:3007668a 14:dc33fa8c 13:2eb71be8 12:2e5cc1a4 11:26ac1e81 10:5ced2805 09:6c7a94aa 08:fa565ef1  
07:e8688f55 06:19805681 05:eaf4d48f 04:8ecacddb 03:a69fa10d 02:ea9a3527 01:e5f1b94c 00:09dfef08  
R1:f9ddf792 R2:6eb763b9

Snow 2.0 Internal state at time -16

15:6fbbfcbf 14:3007668a 13:dc33fa8c 12:2eb71be8 11:2e5cc1a4 10:26ac1e81 09:5ced2805 08:6c7a94aa  
07:fa565ef1 06:e8688f55 05:19805681 04:eaf4d48f 03:8ecacddb 02:a69fa10d 01:ea9a3527 00:e5f1b94c  
R1:59ac3848 R2:510e5e7e

Snow 2.0 Internal state at time -15

15:47128118 14:6fbbfcbf 13:3007668a 12:dc33fa8c 11:2eb71be8 10:2e5cc1a4 09:26ac1e81 08:5ced2805

07:6c7a94aa 06:fa565ef1 05:e8688f55 04:19805681 03:eaf4d48f 02:8ecacddb 01:a69fa10d 00:ea9a3527  
R1:6a8eb4ff R2:ed2a3ff7

Snow 2.0 Internal state at time -14

15:9dd8c346 14:47128118 13:6fbbfcfb 12:3007668a 11:dc33fa8c 10:2eb71be8 09:2e5cc1a4 08:26ac1e81  
07:5ced2805 06:6c7a94aa 05:fa565ef1 04:e8688f55 03:19805681 02:eaf4d48f 01:8ecacddb 00:a69fa10d  
R1:d592cf4c R2:aaaf3ebb

Snow 2.0 Internal state at time -13

15:14c6e4b1 14:9dd8c346 13:47128118 12:6fbbfcfb 11:3007668a 10:dc33fa8c 09:2eb71be8 08:2e5cc1a4  
07:26ac1e81 06:5ced2805 05:6c7a94aa 04:fa565ef1 03:e8688f55 02:19805681 01:eaf4d48f 00:8ecacddb  
R1:a5059dac R2:b838f49b

Snow 2.0 Internal state at time -12

15:2be1899a 14:14c6e4b1 13:9dd8c346 12:47128118 11:6fbbfcfb 10:3007668a 09:dc33fa8c 08:2eb71be8  
07:2e5cc1a4 06:26ac1e81 05:5ced2805 04:6c7a94aa 03:fa565ef1 02:e8688f55 01:19805681 00:eaf4d48f  
R1:24b38945 R2:911396b6

Snow 2.0 Internal state at time -11

15:09363669 14:2be1899a 13:14c6e4b1 12:9dd8c346 11:47128118 10:6fbbfcfb 09:3007668a 08:dc33fa8c  
07:2eb71be8 06:2e5cc1a4 05:26ac1e81 04:5ced2805 03:6c7a94aa 02:fa565ef1 01:e8688f55 00:19805681  
R1:ee00bebb R2:1449ba75

Snow 2.0 Internal state at time -10

15:9e6f24ce 14:09363669 13:2be1899a 12:14c6e4b1 11:9dd8c346 10:47128118 09:6fbbfcfb 08:3007668a  
07:dc33fa8c 06:2eb71be8 05:2e5cc1a4 04:26ac1e81 03:5ced2805 02:6c7a94aa 01:fa565ef1 00:e8688f55  
R1:3af5d8f6 R2:b8fa206d

Snow 2.0 Internal state at time -9

15:f87d0a5a 14:9e6f24ce 13:09363669 12:2be1899a 11:14c6e4b1 10:9dd8c346 09:47128118 08:6fbbfcfb  
07:3007668a 06:dc33fa8c 05:2eb71be8 04:2e5cc1a4 03:26ac1e81 02:5ced2805 01:6c7a94aa 00:fa565ef1  
R1:e756e211 R2:5a6f3141

Snow 2.0 Internal state at time -8

15:056b74c0 14:f87d0a5a 13:9e6f24ce 12:09363669 11:2be1899a 10:14c6e4b1 09:9dd8c346 08:47128118  
07:6fbbfcfb 06:3007668a 05:dc33fa8c 04:2eb71be8 03:2e5cc1a4 02:26ac1e81 01:5ced2805 00:6c7a94aa  
R1:89264d29 R2:87c4f589

Snow 2.0 Internal state at time -7

15:82d49257 14:056b74c0 13:f87d0a5a 12:9e6f24ce 11:09363669 10:2be1899a 09:14c6e4b1 08:9dd8c346  
07:47128118 06:6fbbfcfb 05:3007668a 04:dc33fa8c 03:2eb71be8 02:2e5cc1a4 01:26ac1e81 00:5ced2805  
R1:63f8f015 R2:b541dd3f

Snow 2.0 Internal state at time -6

15:4f549ab4 14:82d49257 13:056b74c0 12:f87d0a5a 11:9e6f24ce 10:09363669 09:2be1899a 08:14c6e4b1  
07:9dd8c346 06:47128118 05:6fbbfcfb 04:3007668a 03:dc33fa8c 02:2eb71be8 01:2e5cc1a4 00:26ac1e81  
R1:e54943c9 R2:cb416287

Snow 2.0 Internal state at time -5

15:01d936bb 14:4f549ab4 13:82d49257 12:056b74c0 11:f87d0a5a 10:9e6f24ce 09:09363669 08:2be1899a  
07:14c6e4b1 06:9dd8c346 05:47128118 04:6fbbfcfb 03:3007668a 02:dc33fa8c 01:2eb71be8 00:2e5cc1a4  
R1:3afd5f82 R2:f4c17d6d

Snow 2.0 Internal state at time -4

15:99c99eb5 14:01d936bb 13:4f549ab4 12:82d49257 11:056b74c0 10:f87d0a5a 09:9e6f24ce 08:09363669  
07:2be1899a 06:14c6e4b1 05:9dd8c346 04:47128118 03:6fbbfcfb 02:3007668a 01:dc33fa8c 00:2eb71be8  
R1:3bd3fe85 R2:b5efeab8



# BS ISO/IEC 18033-4:2011+A1:2020

## ISO/IEC 18033-4:2011+A1:2020(E)

Snow 2.0 Internal state at time -3

15:0ea4e3f0 14:99c99eb5 13:01d936bb 12:4f549ab4 11:82d49257 10:056b74c0 09:f87d0a5a 08:9e6f24ce  
07:09363669 06:2be1899a 05:14c6e4b1 04:9dd8c346 03:47128118 02:6fbbfcfb 01:3007668a 00:dc33fa8c  
R1:53c8adfe R2:a0ddb267

Snow 2.0 Internal state at time -2

15:fa000bc8 14:0ea4e3f0 13:99c99eb5 12:01d936bb 11:4f549ab4 10:82d49257 09:056b74c0 08:f87d0a5a  
07:9e6f24ce 06:09363669 05:2be1899a 04:14c6e4b1 03:9dd8c346 02:47128118 01:6fbbfcfb 00:3007668a  
R1:b5a49718 R2:6ac944cc

Snow 2.0 Internal state at time -1

15:61dec1b8 14:fa000bc8 13:0ea4e3f0 12:99c99eb5 11:01d936bb 10:4f549ab4 09:82d49257 08:056b74c0  
07:f87d0a5a 06:9e6f24ce 05:09363669 04:2be1899a 03:14c6e4b1 02:9dd8c346 01:47128118 00:6fbbfcfb  
R1:96aace66 R2:9cd3a85e

Snow 2.0 Internal state at time 0

15:31f914d5 14:61dec1b8 13:fa000bc8 12:0ea4e3f0 11:99c99eb5 10:01d936bb 09:4f549ab4 08:82d49257  
07:056b74c0 06:f87d0a5a 05:9e6f24ce 04:09363669 03:2be1899a 02:14c6e4b1 01:9dd8c346 00:47128118  
R1:a609dec7 R2:495041dc

Snow 2.0 Internal state at time 1

15:9098ec10 14:31f914d5 13:61dec1b8 12:fa000bc8 11:0ea4e3f0 10:99c99eb5 09:01d936bb 08:4f549ab4  
07:82d49257 06:056b74c0 05:f87d0a5a 04:9e6f24ce 03:09363669 02:2be1899a 01:14c6e4b1 00:9dd8c346  
R1:e7bf66aa R2:05b5db95

Snow 2.0 Internal state at time 2

15:a5e7b806 14:9098ec10 13:31f914d5 12:61dec1b8 11:fa000bc8 10:0ea4e3f0 09:99c99eb5 08:01d936bb  
07:4f549ab4 06:82d49257 05:056b74c0 04:f87d0a5a 03:9e6f24ce 02:09363669 01:2be1899a 00:14c6e4b1  
R1:fe32e5ef R2:e728468a

Snow 2.0 Internal state at time 3

15:962fd59e 14:a5e7b806 13:9098ec10 12:31f914d5 11:61dec1b8 10:fa000bc8 09:0ea4e3f0 08:99c99eb5  
07:01d936bb 06:4f549ab4 05:82d49257 04:056b74c0 03:f87d0a5a 02:9e6f24ce 01:09363669 00:2be1899a  
R1:ec93bb4a R2:ed96a84d

Snow 2.0 Internal state at time 4

15:be037f87 14:962fd59e 13:a5e7b806 12:9098ec10 11:31f914d5 10:61dec1b8 09:fa000bc8 08:0ea4e3f0  
07:99c99eb5 06:01d936bb 05:4f549ab4 04:82d49257 03:056b74c0 02:f87d0a5a 01:9e6f24ce 00:09363669  
R1:706b3aa4 R2:d0d6a880

Snow 2.0 Internal state at time 5

15:3e9ee9ba 14:be037f87 13:962fd59e 12:a5e7b806 11:9098ec10 10:31f914d5 09:61dec1b8 08:fa000bc8  
07:0ea4e3f0 06:99c99eb5 05:01d936bb 04:4f549ab4 03:82d49257 02:056b74c0 01:f87d0a5a 00:9e6f24ce  
R1:202b4334 R2:86c48227

Snow 2.0 Internal state at time 6

15:a14a61e1 14:3e9ee9ba 13:be037f87 12:962fd59e 11:a5e7b806 10:9098ec10 09:31f914d5 08:61dec1b8  
07:fa000bc8 06:0ea4e3f0 05:99c99eb5 04:01d936bb 03:4f549ab4 02:82d49257 01:056b74c0 00:f87d0a5a  
R1:889db8e2 R2:b6399358

Snow 2.0 Internal state at time 7

15:cc852528 14:a14a61e1 13:3e9ee9ba 12:be037f87 11:962fd59e 10:a5e7b806 09:9098ec10 08:31f914d5  
07:61dec1b8 06:fa000bc8 05:0ea4e3f0 04:99c99eb5 03:01d936bb 02:4f549ab4 01:82d49257 00:056b74c0  
R1:5003320d R2:121f6605

Snow 2.0 Internal state at time 8

15:4b895ab7 14:cc852528 13:a14a61e1 12:3e9ee9ba 11:be037f87 10:962fd59e 09:a5e7b806 08:9098ec10  
07:31f914d5 06:61dec1b8 05:fa000bc8 04:0ea4e3f0 03:99c99eb5 02:01d936bb 01:4f549ab4 00:82d49257  
R1:20c449f5 R2:9cf74ff8

## C.3 256-bit key example for SNOW 2.0

### C.3.1 Key, initialization vector, and keystream triplets

$(IV_3, IV_2, IV_1, IV_0) = (0, 0, 0, 0),$   
key = 8000  
Keystream output: 0B5BCCE20323E28E0FC203809C66AB73CA35A680F2A5DD197E0C5C02287BE822.

$(IV_3, IV_2, IV_1, IV_0) = (0, 0, 0, 0),$   
key = AA  
Keystream output: D9CC22FD861492D0AE6F43FB0F072012078C5AEEE479DE8CF0E555F458EED858.

$(IV_3, IV_2, IV_1, IV_0) = (4, 3, 2, 1),$   
key = 8000  
Keystream output: 7861080D5755E90B736F10916ED519B12C1A3A4255297FC2246AB7FA6C089526.

$(IV_3, IV_2, IV_1, IV_0) = (4, 3, 2, 1),$   
key = AA  
Keystream output: 29261FCE5ED038201D6AFAF8B87E74FED49ECB10197EAC025D024EB45E0C7655.

### C.3.2 Sample internal state

K = 80 00  
IV= 00 00 00 04 00 00 00 03 00 00 00 02 00 00 00 01  
Z0= 78 61 08 0d 57 55 e9 0b 73 6f 10 91 6e d5 19 b1 2c 1a 3a 42 55 29 7f c2 24 6a b7 fa 6c 08 95 26

Snow 2.0 Internal state at time -34  
15:80000000 14:00000000 13:00000000 12:00000000 11:00000000 10:00000000 09:00000000 08:00000000  
07:7fffffff 06:ffffffff 05:ffffffff 04:ffffffff 03:ffffffff 02:ffffffff 01:ffffffff 00:ffffffff  
R1:0804a9cc R2:00000001

Snow 2.0 Internal state at time -33  
15:80000001 14:00000000 13:00000000 12:00000002 11:00000000 10:00000003 09:00000004 08:00000000  
07:7fffffff 06:ffffffff 05:ffffffff 04:ffffffff 03:ffffffff 02:ffffffff 01:ffffffff 00:ffffffff  
R1:00000000 R2:00000000

Snow 2.0 Internal state at time -32  
15:bf53b515 14:80000001 13:00000000 12:00000000 11:00000002 10:00000000 09:00000003 08:00000004  
07:00000000 06:7fffffff 05:ffffffff 04:ffffffff 03:ffffffff 02:ffffffff 01:ffffffff 00:ffffffff  
R1:ffffffff R2:63636363

Snow 2.0 Internal state at time -31  
15:d37de350 14:bf53b515 13:80000001 12:00000000 11:00000000 10:00000002 09:00000000 08:00000003  
07:00000004 06:00000000 05:7fffffff 04:ffffffff 03:ffffffff 02:ffffffff 01:ffffffff 00:ffffffff  
R1:63636362 R2:16161616

Snow 2.0 Internal state at time -30  
15:1fa4e5b0 14:d37de350 13:bf53b515 12:80000001 11:00000000 10:00000000 09:00000002 08:00000000  
07:00000003 06:00000004 05:00000000 04:7fffffff 03:ffffffff 02:ffffffff 01:ffffffff 00:ffffffff  
R1:96161615 R2:08aaaa59

Snow 2.0 Internal state at time -29  
15:8243e488 14:1fa4e5b0 13:d37de350 12:bf53b515 11:80000001 10:00000000 09:00000000 08:00000002  
07:00000000 06:00000003 05:00000004 04:00000000 03:7fffffff 02:ffffffff 01:ffffffff 00:ffffffff  
R1:08aaaa59 R2:d03b8eac

Snow 2.0 Internal state at time -28  
15:7d09f594 14:8243e488 13:1fa4e5b0 12:d37de350 11:bf53b515 10:80000001 09:00000000 08:00000000

BS ISO/IEC 18033-4:2011+A1:2020  
ISO/IEC 18033-4:2011+A1:2020(E)

07:00000002 06:00000000 05:00000003 04:00000004 03:00000000 02:7fffffff 01:ffffffff 00:ffffffff  
R1:d03b8eb0 R2:267457fe

Snow 2.0 Internal state at time -27  
15:851e8381 14:7d09f594 13:8243e488 12:1fa4e5b0 11:d37de350 10:bf53b515 09:80000001 08:00000000  
07:00000000 06:00000002 05:00000000 04:00000003 03:00000004 02:00000000 01:7fffffff 00:ffffffff  
R1:26745801 R2:29b1986c

Snow 2.0 Internal state at time -26  
15:cf3efe13 14:851e8381 13:7d09f594 12:8243e488 11:1fa4e5b0 10:d37de350 09:bf53b515 08:80000001  
07:00000000 06:00000000 05:00000002 04:00000000 03:00000003 02:00000004 01:00000000 00:7fffffff  
R1:29b1986c R2:892bf223

Snow 2.0 Internal state at time -25  
15:7b69e0b3 14:cf3efe13 13:851e8381 12:7d09f594 11:8243e488 10:1fa4e5b0 09:d37de350 08:bf53b515  
07:80000001 06:00000000 05:00000000 04:00000002 03:00000000 02:00000003 01:00000004 00:00000000  
R1:892bf225 R2:2f693a07

Snow 2.0 Internal state at time -24  
15:0a8b53d5 14:7b69e0b3 13:cf3efe13 12:851e8381 11:7d09f594 10:8243e488 09:1fa4e5b0 08:d37de350  
07:bf53b515 06:80000001 05:00000000 04:00000000 03:00000002 02:00000000 01:00000003 00:00000004  
R1:2f693a07 R2:6cbd99a8

Snow 2.0 Internal state at time -23  
15:fd75ecf7 14:0a8b53d5 13:7b69e0b3 12:cf3efe13 11:851e8381 10:7d09f594 09:8243e488 08:1fa4e5b0  
07:d37de350 06:bf53b515 05:80000001 04:00000000 03:00000000 02:00000002 01:00000000 00:00000003  
R1:6cbd99a8 R2:0793dbe6

Snow 2.0 Internal state at time -22  
15:94a70314 14:fd75ecf7 13:0a8b53d5 12:7b69e0b3 11:cf3efe13 10:851e8381 09:7d09f594 08:8243e488  
07:1fa4e5b0 06:d37de350 05:bf53b515 04:80000001 03:00000000 02:00000000 01:00000002 00:00000000  
R1:8793dbe7 R2:6928db9c

Snow 2.0 Internal state at time -21  
15:743ca456 14:94a70314 13:fd75ecf7 12:0a8b53d5 11:7b69e0b3 10:cf3efe13 09:851e8381 08:7d09f594  
07:8243e488 06:1fa4e5b0 05:d37de350 04:bf53b515 03:80000001 02:00000000 01:00000000 00:00000002  
R1:287c90b1 R2:ecb79528

Snow 2.0 Internal state at time -20  
15:c250e943 14:743ca456 13:94a70314 12:fd75ecf7 11:0a8b53d5 10:7b69e0b3 09:cf3efe13 08:851e8381  
07:7d09f594 06:8243e488 05:1fa4e5b0 04:d37de350 03:bf53b515 02:80000001 01:00000000 00:00000000  
R1:c0357878 R2:5bd40c0f

Snow 2.0 Internal state at time -19  
15:4d848699 14:c250e943 13:743ca456 12:94a70314 11:fd75ecf7 10:0a8b53d5 09:7b69e0b3 08:cf3efe13  
07:851e8381 06:7d09f594 05:8243e488 04:1fa4e5b0 03:d37de350 02:bf53b515 01:80000001 00:00000000  
R1:7b78f1bf R2:9ae2c490

Snow 2.0 Internal state at time -18  
15:9b3a221f 14:4d848699 13:c250e943 12:743ca456 11:94a70314 10:fd75ecf7 09:0a8b53d5 08:7b69e0b3  
07:cf3efe13 06:851e8381 05:7d09f594 04:8243e488 03:1fa4e5b0 02:d37de350 01:bf53b515 00:80000001  
R1:1d26a918 R2:47a9af75

Snow 2.0 Internal state at time -17  
15:35d95fd1 14:9b3a221f 13:4d848699 12:c250e943 11:743ca456 10:94a70314 09:fd75ecf7 08:0a8b53d5  
07:7b69e0b3 06:cf3efe13 05:851e8381 04:7d09f594 03:8243e488 02:1fa4e5b0 01:d37de350 00:bf53b515  
R1:c4b3a509 R2:9b7cb67c

Snow 2.0 Internal state at time -16

15:e7492c66 14:35d95fd1 13:9b3a221f 12:4d848699 11:c250e943 10:743ca456 09:94a70314 08:fd75ecf7  
07:0a8b53d5 06:7b69e0b3 05:cf3efe13 04:851e8381 03:7d09f594 02:8243e488 01:1fa4e5b0 00:d37de350  
R1:209b39fd R2:50f9a679

Snow 2.0 Internal state at time -15

15:dce814e5 14:e7492c66 13:35d95fd1 12:9b3a221f 11:4d848699 10:c250e943 09:743ca456 08:94a70314  
07:fd75ecf7 06:0a8b53d5 05:7b69e0b3 04:cf3efe13 03:851e8381 02:7d09f594 01:8243e488 00:1fa4e5b0  
R1:2038a48c R2:8facfb3d

Snow 2.0 Internal state at time -14

15:e8e83f37 14:dce814e5 13:e7492c66 12:35d95fd1 11:9b3a221f 10:4d848699 09:c250e943 08:743ca456  
07:94a70314 06:fd75ecf7 05:0a8b53d5 04:7b69e0b3 03:cf3efe13 02:851e8381 01:7d09f594 00:8243e488  
R1:0b16dbf0 R2:97e148a3

Snow 2.0 Internal state at time -13

15:387810f3 14:e8e83f37 13:dce814e5 12:e7492c66 11:35d95fd1 10:9b3a221f 09:4d848699 08:c250e943  
07:743ca456 06:94a70314 05:fd75ecf7 04:0a8b53d5 03:7b69e0b3 02:cf3efe13 01:851e8381 00:7d09f594  
R1:a26c9c78 R2:27c607bf

Snow 2.0 Internal state at time -12

15:4bcdadab 14:387810f3 13:e8e83f37 12:dce814e5 11:e7492c66 10:35d95fd1 09:9b3a221f 08:4d848699  
07:c250e943 06:743ca456 05:94a70314 04:fd75ecf7 03:0a8b53d5 02:7b69e0b3 01:cf3efe13 00:851e8381  
R1:253bf4b6 R2:258cd170

Snow 2.0 Internal state at time -11

15:f05c5d45 14:4bcdadab 13:387810f3 12:e8e83f37 11:dce814e5 10:e7492c66 09:35d95fd1 08:9b3a221f  
07:4d848699 06:c250e943 05:743ca456 04:94a70314 03:fd75ecf7 02:0a8b53d5 01:7b69e0b3 00:cf3efe13  
R1:ba33d484 R2:f16f299b

Snow 2.0 Internal state at time -10

15:21e0b756 14:f05c5d45 13:4bcdadab 12:387810f3 11:e8e83f37 10:dce814e5 09:e7492c66 08:35d95fd1  
07:9b3a221f 06:4d848699 05:c250e943 04:743ca456 03:94a70314 02:fd75ecf7 01:0a8b53d5 00:7b69e0b3  
R1:65abdcf1 R2:998d6551

Snow 2.0 Internal state at time -9

15:00098c25 14:21e0b756 13:f05c5d45 12:4bcdadab 11:387810f3 10:e8e83f37 09:dce814e5 08:e7492c66  
07:35d95fd1 06:9b3a221f 05:4d848699 04:c250e943 03:743ca456 02:94a70314 01:fd75ecf7 00:0a8b53d5  
R1:5bde4e94 R2:bd0f2baa

Snow 2.0 Internal state at time -8

15:81a343e1 14:00098c25 13:21e0b756 12:f05c5d45 11:4bcdadab 10:387810f3 09:e8e83f37 08:dce814e5  
07:e7492c66 06:35d95fd1 05:9b3a221f 04:4d848699 03:c250e943 02:743ca456 01:94a70314 00:fd75ecf7  
R1:0a93b243 R2:267c6211

Snow 2.0 Internal state at time -7

15:7b933a92 14:81a343e1 13:00098c25 12:21e0b756 11:f05c5d45 10:4bcdadab 09:387810f3 08:e8e83f37  
07:dce814e5 06:e7492c66 05:35d95fd1 04:9b3a221f 03:4d848699 02:c250e943 01:743ca456 00:94a70314  
R1:c1b68430 R2:0b276cd6

Snow 2.0 Internal state at time -6

15:0339b827 14:7b933a92 13:81a343e1 12:00098c25 11:21e0b756 10:f05c5d45 09:4bcdadab 08:387810f3  
07:e8e83f37 06:dce814e5 05:e7492c66 04:35d95fd1 03:9b3a221f 02:4d848699 01:c250e943 00:743ca456  
R1:4100cca7 R2:ed4f10df

Snow 2.0 Internal state at time -5

15:e5fd1e4e 14:0339b827 13:7b933a92 12:81a343e1 11:00098c25 10:21e0b756 09:f05c5d45 08:4bcdadab  
07:387810f3 06:e8e83f37 05:dce814e5 04:e7492c66 03:35d95fd1 02:9b3a221f 01:4d848699 00:c250e943  
R1:d4983d45 R2:d14fec85

# BS ISO/IEC 18033-4:2011+A1:2020

## ISO/IEC 18033-4:2011+A1:2020(E)

Snow 2.0 Internal state at time -4

15:991f7362 14:e5fd1e4e 13:0339b827 12:7b933a92 11:81a343e1 10:00098c25 09:21e0b756 08:f05c5d45  
07:4bcddadb 06:387810f3 05:e8e83f37 04:dce814e5 03:e7492c66 02:35d95fd1 01:9b3a221f 00:4d848699  
R1:ae38016a R2:431da2bb

Snow 2.0 Internal state at time -3

15:a0bca2f7 14:991f7362 13:e5fd1e4e 12:0339b827 11:7b933a92 10:81a343e1 09:00098c25 08:21e0b756  
07:f05c5d45 06:4bcddadb 05:387810f3 04:e8e83f37 03:dce814e5 02:e7492c66 01:35d95fd1 00:9b3a221f  
R1:2c05elf2 R2:ae471763

Snow 2.0 Internal state at time -2

15:928b5256 14:a0bca2f7 13:991f7362 12:e5fd1e4e 11:0339b827 10:7b933a92 09:81a343e1 08:00098c25  
07:21e0b756 06:f05c5d45 05:4bcddadb 04:387810f3 03:e8e83f37 02:dce814e5 01:e7492c66 00:35d95fd1  
R1:e6bf2856 R2:f134ae00

Snow 2.0 Internal state at time -1

15:be366d56 14:928b5256 13:a0bca2f7 12:991f7362 11:e5fd1e4e 10:0339b827 09:7b933a92 08:81a343e1  
07:00098c25 06:21e0b756 05:f05c5d45 04:4bcddadb 03:387810f3 02:e8e83f37 01:dce814e5 00:e7492c66  
R1:3d0288db R2:f31c4fa3

Snow 2.0 Internal state at time 0

15:a5fadb9e 14:be366d56 13:928b5256 12:a0bca2f7 11:991f7362 10:e5fd1e4e 09:0339b827 08:7b933a92  
07:81a343e1 06:00098c25 05:21e0b756 04:f05c5d45 03:4bcddadb 02:387810f3 01:e8e83f37 00:dce814e5  
R1:e378ace8 R2:2dfa946e

Snow 2.0 Internal state at time 1

15:b70c6e50 14:a5fadb9e 13:be366d56 12:928b5256 11:a0bca2f7 10:991f7362 09:e5fd1e4e 08:0339b827  
07:7b933a92 06:81a343e1 05:00098c25 04:21e0b756 03:f05c5d45 02:4bcddadb 01:387810f3 00:e8e83f37  
R1:4fdb4bc4 R2:b95a6c28

Snow 2.0 Internal state at time 2

15:bce23d5c 14:b70c6e50 13:a5fadb9e 12:be366d56 11:928b5256 10:a0bca2f7 09:991f7362 08:e5fd1e4e  
07:0339b827 06:7b933a92 05:81a343e1 04:00098c25 03:21e0b756 02:f05c5d45 01:4bcddadb 00:387810f3  
R1:b963f84d R2:3d5135cb

Snow 2.0 Internal state at time 3

15:4eb9692d 14:bce23d5c 13:b70c6e50 12:a5fadb9e 11:be366d56 10:928b5256 09:a0bca2f7 08:991f7362  
07:e5fd1e4e 06:0339b827 05:7b933a92 04:81a343e1 03:00098c25 02:21e0b756 01:f05c5d45 00:4bcddadb  
R1:bef479ac R2:28b521b3

Snow 2.0 Internal state at time 4

15:96a599a9 14:4eb9692d 13:bce23d5c 12:b70c6e50 11:a5fadb9e 10:be366d56 09:928b5256 08:a0bca2f7  
07:991f7362 06:e5fd1e4e 05:0339b827 04:7b933a92 03:81a343e1 02:00098c25 01:21e0b756 00:f05c5d45  
R1:a4485c45 R2:e6ab92e9

Snow 2.0 Internal state at time 5

15:62ad427d 14:96a599a9 13:4eb9692d 12:bce23d5c 11:b70c6e50 10:a5fadb9e 09:be366d56 08:928b5256  
07:a0bca2f7 06:991f7362 05:e5fd1e4e 04:0339b827 03:7b933a92 02:81a343e1 01:00098c25 00:21e0b756  
R1:e9e54b10 R2:385b4519

Snow 2.0 Internal state at time 6

15:19397ef2 14:62ad427d 13:96a599a9 12:4eb9692d 11:bce23d5c 10:b70c6e50 09:a5fadb9e 08:be366d56  
07:928b5256 06:a0bca2f7 05:991f7362 04:e5fd1e4e 03:0339b827 02:7b933a92 01:81a343e1 00:00098c25  
R1:1e586367 R2:13f2d986

Snow 2.0 Internal state at time 7

15:5f8525f0 14:19397ef2 13:62ad427d 12:96a599a9 11:4eb9692d 10:bce23d5c 09:b70c6e50 08:a5fadb9e  
07:be366d56 06:928b5256 05:a0bca2f7 04:991f7362 03:e5fd1e4e 02:0339b827 01:7b933a92 00:81a343e1  
R1:ad124ce8 R2:e13ca41f

```
Snow 2.0 Internal state at time 8
15:fb9c0bcf 14:5f8525f0 13:19397ef2 12:62ad427d 11:96a599a9 10:4eb9692d 09:bce23d5c 08:b70c6e50
07:a5fadb9e 06:be366d56 05:928b5256 04:a0bca2f7 03:991f7362 02:e5fd1e4e 01:0339b827 00:7b933a92
R1:81f94716 R2:679f1c0a
```

## C.4 Example for Rabbit

### C.4.1 Introduction

All test vectors for Rabbit are given in the little-endian notation, i.e., for multi-byte numbers, the most significant bytes are stored at the highest memory addresses.

### C.4.2 Key, initialization vector, and keystream triplets

```
K = 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
IV= 00 00 00 00 00 00 00 00 00
Z = ED B7 05 67 37 5D CD 7C D8 95 54 F8 5E 27 A7 C6 8D 4A DC 70 32 29 8F 7B D4 EF F5 04 AC A6 29 5F
    66 8F BF 47 8A DB 2B E5 1E 6C DE 29 2B 82 DE 2A B4 8D 2A C6 56 59 79 22 0E C9 09 A7 E7 57 60 98
```

```
K = 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
IV= 00 01 02 03 04 05 06 07
Z = 98 71 C7 BA 4E A3 08 07 CD AA 49 64 66 39 2D 2F 4A FF 43 55 EF 90 69 56 10 9B 96 65 97 8D AC ED
    9B 7C 6F 7F C8 2C 67 D2 73 22 CB DE 9D B0 16 45 8C 38 2C 9C 7D 30 44 E6 52 0B B9 2A 13 53 C0 FF
```

```
K = 00 01 02 03 04 05 06 07 08 09 0A 0B 0C 0D 0E 0F
IV= 00 00 00 00 00 00 00 00
Z = A8 F7 E6 9B 69 40 A7 8D 13 6A 5C 15 4A 15 79 52 A6 E4 23 58 59 E3 02 20 EA 68 64 36 BB 38 EF 53
    9C 29 40 55 6B 09 EC D7 FE A2 B0 AC 83 07 F1 69 62 65 A3 D6 44 28 1C 39 C9 CD 5E 1E 2F 9B E4 D0
```

```
K = 00 01 02 03 04 05 06 07 08 09 0A 0B 0C 0D 0E 0F
IV= 00 01 02 03 04 05 06 07
Z = F2 89 19 DD A1 28 F8 F9 0A 30 34 6E 97 94 D2 B7 4C 69 A2 D9 91 37 27 BC 5A 30 18 E6 33 2A F7 F3
    BE 3A C3 EF B3 68 F4 3A 4C B8 58 67 B8 1C 91 F9 24 29 0C 81 6B 8B 57 88 98 C5 7F B4 C0 BA 05 BD
```

### C.4.3 Sample internal states

```
K = 00 01 02 03 04 05 06 07 08 09 0A 0B 0C 0D 0E 0F
IV= 00 01 02 03 04 05 06 07
Z = F2 89 19 DD A1 28 F8 F9 0A 30 34 6E 97 94 D2 B7 4C 69 A2 D9 91 37 27 BC 5A 30 18 E6 33 2A F7 F3
    BE 3A C3 EF B3 68 F4 3A 4C B8 58 67 B8 1C 91 F9 24 29 0C 81 6B 8B 57 88 98 C5 7F B4 C0 BA 05 BD
```

After key expansion (Internal state S(-9))

```
x0:03020100 x1:0D0C0B0A x2:07060504 x3:01000F0E x4:0B0A0908 x5:05040302 x6:0F0E0D0C x7:09080706
c0:09080B0A c1:03020504 c2:0D0C0F0E c3:07060908 c4:01000302 c5:0B0A0D0C c6:05040706 c7:0F0E0100
carry:0
```

After key setup iteration 1 (Internal state S(-8))

```
x0:05783933 x1:162113C0 x2:B38F168E x3:F08A919E x4:7F2CDA94 x5:ACBEB878 x6:0D5257A9 x7:4FF46B46
c0:563CDE57 c1:D64F39D7 c2:41DF5C42 c3:543ADC55 c4:D44D37D5 c5:3FDD5A40 c6:5238DA53 c7:E25B35D3
carry:0
```

After key setup iteration 2 (Internal state S(-7))

```
x0:798C2CEC x1:CC05FFD4 x2:50D68324 x3:2C306745 x4:AD519559 x5:81595E7A x6:29A589E2 x7:15212B97
c0:A371B1A4 c1:A99C6EAA c2:76B2A977 c3:A16FAFA2 c4:A79A6CA8 c5:74B0A775 c6:9F6DADA0 c7:B5A86AA6
carry:1
```

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## ISO/IEC 18033-4:2011+A1:2020(E)

After key setup iteration 3 (Internal state S(-6))

x0:CD328957 x1:66D5AB1F x2:0D115824 x3:FCCEB784 x4:12E900D7 x5:36A46997 x6:9F40C5BC x7:AB1C8A08  
c0:F0A684F2 c1:7CE9A37D c2:AB85F6AC c3:EEA482EF c4:7AE7A17B c5:A983F4AA c6:ECA280ED c7:88F59F79  
carry:1

After key setup iteration 4 (Internal state S(-5))

x0:A31515F8 x1:5DFD3AC6 x2:33CD6AD2 x3:4BD778E5 x4:89708269 x5:D93095C1 x6:5E495F60 x7:C197863A  
c0:3DDB5840 c1:5036D851 c2:E05943E1 c3:3BD9563C c4:4E34D64F c5:DE5741DF c6:39D7543A c7:5C42D44D  
carry:1

After counter modification / IV setup (Internal state S(-4))

x0:A31515F8 x1:5DFD3AC6 x2:33CD6AD2 x3:4BD778E5 x4:89708269 x5:D93095C1 x6:5E495F60 x7:C197863A  
c0:B7A9DB29 c1:8E004E92 c2:B9161985 c3:FF4AD106 c4:EE23C2B7 c5:84AC781B c6:0D1C3BEC c7:1291ADA8  
carry:1

After IV setup iteration 1 (Internal state S(-3))

x0:054A3F2F x1:BE444CDE x2:573425A4 x3:9347FAD1 x4:29036A2F x5:DD3C6B50 x6:12CC3803 x7:6F7847C0  
c0:04DEAE77 c1:614D8366 c2:EDE966BA c3:4C7FA453 c4:C170F78B c5:B97FC550 c6:5A510F39 c7:E5DEE27B  
carry:0

After IV setup iteration 2 (Internal state S(-2))

x0:0FDB9A3A x1:334807E8 x2:E66BCC98 x3:0FDA371C x4:9C3E3036 x5:7774E657 x6:C6FCBB4C x7:A8D1AC4F  
c0:521381C4 c1:349AB839 c2:22BCB3EF c3:99B477A1 c4:94BE2C5E c5:EE531285 c6:A785E286 c7:B92C174E  
carry:1

After IV setup iteration 3 (Internal state S(-1))

x0:1A2EF77E x1:FDEEE287 x2:A918F5A1 x3:D6414F76 x4:4848D473 x5:BCE9BD30 x6:3E524094 x7:16242C51  
c0:9F485512 c1:07E7ED0C c2:57900124 c3:E6E94AEE c4:680B6131 c5:23265FBA c6:F4BAB5D4 c7:8C794C21  
carry:1

After IV setup iteration 4 (Internal state S(0))

x0:987651C2 x1:FF5F0007 x2:5C48C79E x3:661B3E75 x4:49247B9A x5:3C7AA744 x6:4AEF3F40 x7:D117584E  
c0:EC7D2860 c1:DB3521DF c2:8C634E58 c3:341E1E3B c4:3B589605 c5:57F9ACEF c6:41EF8921 c7:5FC680F5  
carry:1

After keystream iteration 1 (Internal state S(1))

x0:2A158BE4 x1:D93EC5A4 x2:298B7C1B x3:01F4F70C x4:E241E934 x5:0216D073 x6:72769563 x7:54BA8C75  
c0:39B1FBAE c1:AE8256B3 c2:C1369B8D c3:8152F188 c4:0EA5CAD8 c5:8CCCFA24 c6:8F245C6E c7:3313B5C8  
carry:1  
output F2 89 19 DD A1 28 F8 F9 0A 30 34 6E 97 94 D2 B7

After keystream iteration 2 (Internal state S(2))

x0:46EC0492 x1:A4B5D46E x2:7B374C9E x3:93249F4E x4:E93894EF x5:6DDEC710 x6:2799B917 x7:7B0F0F20  
c0:86E6CEFC c1:81CF8B86 c2:F609E8C2 c3:CE87C4D5 c4:E1F2FFAB c5:C1A04758 c6:DC592FBB c7:0660EA9B  
carry:1  
output 4C 69 A2 D9 91 37 27 BC 5A 30 18 E6 33 2A F7 F3

After keystream iteration 3 (Internal state S(3))

x0:98C27422 x1:0D5B5EC2 x2:FEEC9F8D x3:423F7701 x4:E22AB517 x5:4E9CC418 x6:A7535E87 x7:F73E8572  
c0:D41BA24A c1:551CC059 c2:2ADD35F7 c3:1BBC9823 c4:B540347F c5:F673948D c6:298E0308 c7:D9AE1F6F  
carry:0  
output BE 3A C3 EF B3 68 F4 3A 4C B8 58 67 B8 1C 91 F9

After keystream iteration 4 (Internal state S(4))

x0:3B844C36 x1:AF5CD78B x2:2619A0AC x3:774FBA88 x4:D16C6AC4 x5:6512AE4E x6:6A8ECD8F x7:2BC76513  
c0:21507597 c1:2869F52D c2:5FB0832C c3:68F16B70 c4:888D6952 c5:2B46E1C2 c6:76C2D656 c7:ACFB5442  
carry:1  
output 24 29 0C 81 6B 8B 57 88 98 C5 7F B4 C0 BA 05 BD

## C.5 Example for Decim<sup>v2</sup>

### C.5.1 Introduction to the Decim<sup>v2</sup> example

The byte-values and binary decomposition of bytes follow the big-endian notation, i.e., for multi-byte numbers, the most significant bytes are stored at the lowest memory addresses. In particular, this holds for the key,  $IV$ , keystream, register and buffer byte- and binary values given below.

$$\begin{aligned} K &= K_{79} \dots K_0 \\ IV &= IV_{63} \dots IV_0 \\ Z &= Z_n \dots Z_0 \\ a &= a_{191} \dots a_0 \\ b &= b_{31} \dots b_0 \\ T &= T_2 T_1 T_0 \end{aligned}$$

and, for instance, given the key

$$K = \text{de aa 00 40 00 30 00 0f 08 80},$$

each registers are defined to:

$$[K_{79} \dots K_{72}] = \text{de}, [K_{71} \dots K_{64}] = \text{aa} \dots [K_7 \dots K_0] = \text{80},$$

with bit-decomposition as follows:

$$\begin{aligned} K_{79} \dots K_0 &= 11011110 \ 10101010 \ 00000000 \ 01000000 \ 00000000 \\ &\quad 00110000 \ 00000000 \ 00001111 \ 00001000 \ 10000000 \end{aligned}$$

### C.5.2 Key, initialization vector, and keystream triplets

$$\begin{aligned} K &= 00 \ 00 \ 00 \ 00 \ 00 \ 00 \ 00 \ 00 \ 00 \ 80 \\ IV &= 00 \ 00 \ 00 \ 00 \ 00 \ 00 \ 00 \ 00 \\ Z &= 76 \ \text{e3} \ 89 \ \text{be} \ 1b \ \text{fb} \ \text{ad} \ d5 \ 3c \ \text{ce} \ a0 \ \text{fe} \ 43 \ b8 \ c8 \ \text{fb} \ d3 \ 92 \ b8 \ 0b \ 52 \ 94 \ 60 \ \text{f8} \end{aligned}$$

$$\begin{aligned} K &= 00 \ 00 \ 00 \ 00 \ 00 \ 00 \ 00 \ 00 \ 00 \ 00 \\ IV &= 00 \ 00 \ 00 \ 00 \ 00 \ 00 \ 00 \ 80 \\ Z &= 4c \ \text{ec} \ \text{bd} \ b3 \ 0e \ \text{cd} \ c9 \ c0 \ 8b \ 41 \ 8f \ 7f \ 28 \ \text{ff} \ 83 \ 48 \ 75 \ 40 \ \text{ff} \ c5 \ \text{cb} \ 0a \ 33 \ da \end{aligned}$$

$$\begin{aligned} K &= 09 \ 09 \ 09 \ 09 \ 09 \ 09 \ 09 \ 09 \ 09 \ 09 \\ IV &= 00 \ 00 \ 00 \ 00 \ 00 \ 00 \ 00 \ 00 \\ Z &= 43 \ 9b \ ba \ f8 \ a7 \ 84 \ dc \ f9 \ e6 \ d2 \ 90 \ 1d \ 12 \ 4d \ 43 \ 09 \ 22 \ 33 \ f2 \ 47 \ 60 \ 19 \ 70 \ 53 \end{aligned}$$

$$\begin{aligned} K &= 09 \ 08 \ 07 \ 06 \ 05 \ 04 \ 03 \ 02 \ 01 \ 00 \\ IV &= 00 \ 00 \ 00 \ 00 \ 00 \ 00 \ 00 \ 00 \\ Z &= 52 \ b1 \ 73 \ 10 \ 01 \ 2a \ cd \ 3a \ d2 \ 20 \ 4f \ e2 \ b2 \ 2a \ 5d \ 21 \ 64 \ 41 \ f6 \ 3d \ d3 \ b4 \ 43 \ 6a \end{aligned}$$

$$\begin{aligned} K &= \text{eb} \ 98 \ 45 \ f2 \ 9f \ 4c \ f9 \ a6 \ 53 \ 00 \\ IV &= \text{de} \ 77 \ 10 \ a9 \ 42 \ db \ 74 \ 0d \\ Z &= 62 \ \text{ff} \ c9 \ cc \ 21 \ 0e \ 07 \ ea \ 6e \ 50 \ f0 \ \text{fb} \ 1b \ 60 \ 36 \ 7f \ 88 \ a6 \ a5 \ 27 \ 9b \ 18 \ \text{cb} \ b8 \end{aligned}$$

$$\begin{aligned} K &= fa \ a7 \ 54 \ 01 \ ae \ 5b \ 08 \ b5 \ 62 \ 0f \\ IV &= f9 \ 92 \ 2b \ c4 \ 5d \ f6 \ 8f \ 28 \\ Z &= f0 \ af \ 66 \ 52 \ 2a \ 23 \ 8b \ 29 \ 63 \ 37 \ 8b \ 18 \ ec \ 1f \ 4c \ a8 \ 27 \ 91 \ 3d \ 2c \ f0 \ ad \ 94 \ d9 \end{aligned}$$



### C.5.3 Sample internal states

The binary equivalents of the internal states for key stages, namely at time -256, time -64, time 0 and time 193 are provided.

```
K = 00 00 00 00 00 00 00 00 00 80
IV= 00 00 00 00 00 00 00 00
Z = 76 e3 89 be 1b fb ad d5 3c ce a0 fe 43 b8 c8 fb d3 92 b8 0b 52 94 60 f8
```

For time -256 until -64 (executions of *InitNext* (S,LFSR)), internal state variables *T*, *b* and *I* have the following values:

```
T: 000
b: 00 00 00 00
I: 0
```

```
Decim v2 Internal State at time -256
a: ff ff ff ff 00 00 00 00 00 00 00 00 00 80 00 00 00 00 00 00 00 80
```

```
Decim v2 Binary Internal State at time -256 (Binary notation)
a: 11111111 11111111 11111111 11111111 00000000 00000000 00000000 00000000
   00000000 00000000 00000000 00000000 00000000 10000000 00000000 00000000
   00000000 00000000 00000000 00000000 00000000 00000000 00000000 10000000
```

Executions of *InitNext* (S,LFSR)

```
Decim v2 Internal State at time -255
a: 2f ff ff ff f0 00 00 00 00 00 00 00 00 08 00 00 00 00 00 00 00 08
```

```
Decim v2 Internal State at time -254
a: 42 ff ff ff ff 00 00 00 00 00 00 00 00 00 80 00 00 00 00 00 00 00 00
```

```
Decim v2 Internal State at time -253
a: 84 2f ff ff ff f0 00 00 00 00 00 00 00 08 00 00 00 00 00 00 00 00
```

```
Decim v2 Internal State at time -252
a: 98 42 ff ff ff ff 00 00 00 00 00 00 00 00 80 00 00 00 00 00 00 00
```

```
Decim v2 Internal State at time -251
a: 99 84 2f ff ff ff f0 00 00 00 00 00 00 00 08 00 00 00 00 00 00 00
```

```
Decim v2 Internal State at time -250
a: 89 98 42 ff ff ff ff 00 00 00 00 00 00 00 80 00 00 00 00 00 00 00
```

```
Decim v2 Internal State at time -249
a: e8 99 84 2f ff ff ff f0 00 00 00 00 00 00 00 08 00 00 00 00 00 00
```

```
Decim v2 Internal State at time -248
a: 4e 89 98 42 ff ff ff ff 00 00 00 00 00 00 00 80 00 00 00 00 00 00
```

```
Decim v2 Internal State at time -247
a: 04 e8 99 84 2f ff ff ff f0 00 00 00 00 00 00 00 08 00 00 00 00 00 00
```

```
Decim v2 Internal State at time -246
a: 00 4e 89 98 42 ff ff ff ff 00 00 00 00 00 00 00 80 00 00 00 00 00 00
```

```
Decim v2 Internal State at time -245
a: d0 04 e8 99 84 2f ff ff ff f0 00 00 00 00 00 00 00 08 00 00 00 00 00
```

```
Decim v2 Internal State at time -244
a: 6d 00 4e 89 98 42 ff ff ff ff 00 00 00 00 00 00 00 80 00 00 00 00 00
```

```
Decim v2 Internal State at time -243
a: 16 d0 04 e8 99 84 2f ff ff ff f0 00 00 00 00 00 00 00 08 00 00 00 00
```

```
Decim v2 Internal State at time -242
a: d1 6d 00 4e 89 98 42 ff ff ff ff 00 00 00 00 00 00 00 80 00 00 00 00
```

```
Decim v2 Internal State at time -241
a: fd 16 d0 04 e8 99 84 2f ff ff ff f0 00 00 00 00 00 00 00 08 00 00 00
```

```

Decim v2 Internal State at time -240
a: df d1 6d 00 4e 89 98 42 ff ff ff ff 00 00 00 00 00 00 00 00 80 00 00

Decim v2 Internal State at time -239
a: ed fd 16 d0 04 e8 99 84 2f ff ff ff f0 00 00 00 00 00 00 00 08 00 00

Decim v2 Internal State at time -238
a: fe df d1 6d 00 4e 89 98 42 ff ff ff ff 00 00 00 00 00 00 00 00 80 00

Decim v2 Internal State at time -237
a: ef ed fd 16 d0 04 e8 99 84 2f ff ff ff f0 00 00 00 00 00 00 00 08 00

Decim v2 Internal State at time -236
a: ee fe df d1 6d 00 4e 89 98 42 ff ff ff ff 00 00 00 00 00 00 00 00 80

Decim v2 Internal State at time -235
a: ce ef ed fd 16 d0 04 e8 99 84 2f ff ff ff f0 00 00 00 00 00 00 00 08

Decim v2 Internal State at time -234
a: 0c ee fe df d1 6d 00 4e 89 98 42 ff ff ff ff 00 00 00 00 00 00 00 00 00

Decim v2 Internal State at time -233
a: 60 ce ef ed fd 16 d0 04 e8 99 84 2f ff ff ff f0 00 00 00 00 00 00 00 00

Decim v2 Internal State at time -232
a: 06 0c ee fe df d1 6d 00 4e 89 98 42 ff ff ff ff 00 00 00 00 00 00 00 00

Decim v2 Internal State at time -231
a: b0 60 ce ef ed fd 16 d0 04 e8 99 84 2f ff ff ff f0 00 00 00 00 00 00 00

Decim v2 Internal State at time -230
a: 5b 06 0c ee fe df d1 6d 00 4e 89 98 42 ff ff ff ff 00 00 00 00 00 00 00

Decim v2 Internal State at time -229
a: c5 b0 60 ce ef ed fd 16 d0 04 e8 99 84 2f ff ff ff f0 00 00 00 00 00 00

Decim v2 Internal State at time -228
a: bc 5b 06 0c ee fe df d1 6d 00 4e 89 98 42 ff ff ff ff 00 00 00 00 00 00

Decim v2 Internal State at time -227
a: 6b c5 b0 60 ce ef ed fd 16 d0 04 e8 99 84 2f ff ff ff f0 00 00 00 00 00

Decim v2 Internal State at time -226
a: d6 bc 5b 06 0c ee fe df d1 6d 00 4e 89 98 42 ff ff ff ff 00 00 00 00 00

Decim v2 Internal State at time -225
a: bd 6b c5 b0 60 ce ef ed fd 16 d0 04 e8 99 84 2f ff ff ff f0 00 00 00 00

Decim v2 Internal State at time -224
a: 8b d6 bc 5b 06 0c ee fe df d1 6d 00 4e 89 98 42 ff ff ff ff 00 00 00 00

Decim v2 Internal State at time -223
a: 38 bd 6b c5 b0 60 ce ef ed fd 16 d0 04 e8 99 84 2f ff ff ff f0 00 00 00

Decim v2 Internal State at time -222
a: b3 8b d6 bc 5b 06 0c ee fe df d1 6d 00 4e 89 98 42 ff ff ff ff 00 00 00

Decim v2 Internal State at time -221
a: fb 38 bd 6b c5 b0 60 ce ef ed fd 16 d0 04 e8 99 84 2f ff ff ff f0 00 00

Decim v2 Internal State at time -220
a: 4f b3 8b d6 bc 5b 06 0c ee fe df d1 6d 00 4e 89 98 42 ff ff ff ff 00 00

Decim v2 Internal State at time -219
a: 84 fb 38 bd 6b c5 b0 60 ce ef ed fd 16 d0 04 e8 99 84 2f ff ff ff f0 00

Decim v2 Internal State at time -218
a: 18 4f b3 8b d6 bc 5b 06 0c ee fe df d1 6d 00 4e 89 98 42 ff ff ff ff 00

Decim v2 Internal State at time -217
a: 21 84 fb 38 bd 6b c5 b0 60 ce ef ed fd 16 d0 04 e8 99 84 2f ff ff ff f0

Decim v2 Internal State at time -216
a: f2 18 4f b3 8b d6 bc 5b 06 0c ee fe df d1 6d 00 4e 89 98 42 ff ff ff ff

```

# BS ISO/IEC 18033-4:2011+A1:2020

## ISO/IEC 18033-4:2011+A1:2020(E)

Decim v2 Internal State at time -215  
a: 9f 21 84 fb 38 bd 6b c5 b0 60 ce ef ed fd 16 d0 04 e8 99 84 2f ff ff ff

Decim v2 Internal State at time -214  
a: e9 f2 18 4f b3 8b d6 bc 5b 06 0c ee fe df d1 6d 00 4e 89 98 42 ff ff ff

Decim v2 Internal State at time -213  
a: 6e 9f 21 84 fb 38 bd 6b c5 b0 60 ce ef ed fd 16 d0 04 e8 99 84 2f ff ff

Decim v2 Internal State at time -212  
a: f6 e9 f2 18 4f b3 8b d6 bc 5b 06 0c ee fe df d1 6d 00 4e 89 98 42 ff ff

Decim v2 Internal State at time -211  
a: 5f 6e 9f 21 84 fb 38 bd 6b c5 b0 60 ce ef ed fd 16 d0 04 e8 99 84 2f ff

Decim v2 Internal State at time -210  
a: 45 f6 e9 f2 18 4f b3 8b d6 bc 5b 06 0c ee fe df d1 6d 00 4e 89 98 42 ff

Decim v2 Internal State at time -209  
a: 04 5f 6e 9f 21 84 fb 38 bd 6b c5 b0 60 ce ef ed fd 16 d0 04 e8 99 84 2f

Decim v2 Internal State at time -208  
a: 10 45 f6 e9 f2 18 4f b3 8b d6 bc 5b 06 0c ee fe df d1 6d 00 4e 89 98 42

Decim v2 Internal State at time -207  
a: d1 04 5f 6e 9f 21 84 fb 38 bd 6b c5 b0 60 ce ef ed fd 16 d0 04 e8 99 84

Decim v2 Internal State at time -206  
a: 1d 10 45 f6 e9 f2 18 4f b3 8b d6 bc 5b 06 0c ee fe df d1 6d 00 4e 89 98

Decim v2 Internal State at time -205  
a: 31 d1 04 5f 6e 9f 21 84 fb 38 bd 6b c5 b0 60 ce ef ed fd 16 d0 04 e8 99

Decim v2 Internal State at time -204  
a: 13 1d 10 45 f6 e9 f2 18 4f b3 8b d6 bc 5b 06 0c ee fe df d1 6d 00 4e 89

Decim v2 Internal State at time -203  
a: 81 31 d1 04 5f 6e 9f 21 84 fb 38 bd 6b c5 b0 60 ce ef ed fd 16 d0 04 e8

Decim v2 Internal State at time -202  
a: 88 13 1d 10 45 f6 e9 f2 18 4f b3 8b d6 bc 5b 06 0c ee fe df d1 6d 00 4e

Decim v2 Internal State at time -201  
a: c8 81 31 d1 04 5f 6e 9f 21 84 fb 38 bd 6b c5 b0 60 ce ef ed fd 16 d0 04

Decim v2 Internal State at time -200  
a: 7c 88 13 1d 10 45 f6 e9 f2 18 4f b3 8b d6 bc 5b 06 0c ee fe df d1 6d 00

Decim v2 Internal State at time -199  
a: 77 c8 81 31 d1 04 5f 6e 9f 21 84 fb 38 bd 6b c5 b0 60 ce ef ed fd 16 d0

Decim v2 Internal State at time -198  
a: 17 7c 88 13 1d 10 45 f6 e9 f2 18 4f b3 8b d6 bc 5b 06 0c ee fe df d1 6d

Decim v2 Internal State at time -197  
a: 91 77 c8 81 31 d1 04 5f 6e 9f 21 84 fb 38 bd 6b c5 b0 60 ce ef ed fd 16

Decim v2 Internal State at time -196  
a: 49 17 7c 88 13 1d 10 45 f6 e9 f2 18 4f b3 8b d6 bc 5b 06 0c ee fe df d1

Decim v2 Internal State at time -195  
a: 44 91 77 c8 81 31 d1 04 5f 6e 9f 21 84 fb 38 bd 6b c5 b0 60 ce ef ed fd

Decim v2 Internal State at time -194  
a: c4 49 17 7c 88 13 1d 10 45 f6 e9 f2 18 4f b3 8b d6 bc 5b 06 0c ee fe df

Decim v2 Internal State at time -193  
a: bc 44 91 77 c8 81 31 d1 04 5f 6e 9f 21 84 fb 38 bd 6b c5 b0 60 ce ef ed

Decim v2 Internal State at time -192  
a: 5b c4 49 17 7c 88 13 1d 10 45 f6 e9 f2 18 4f b3 8b d6 bc 5b 06 0c ee fe

Decim v2 Internal State at time -191  
a: c5 bc 44 91 77 c8 81 31 d1 04 5f 6e 9f 21 84 fb 38 bd 6b c5 b0 60 ce ef

Decim v2 Internal State at time -190  
a: 5c 5b c4 49 17 7c 88 13 1d 10 45 f6 e9 f2 18 4f b3 8b d6 bc 5b 06 0c ee

Decim v2 Internal State at time -189  
a: 25 c5 bc 44 91 77 c8 81 31 d1 04 5f 6e 9f 21 84 fb 38 bd 6b c5 b0 60 ce

Decim v2 Internal State at time -188  
a: 92 5c 5b c4 49 17 7c 88 13 1d 10 45 f6 e9 f2 18 4f b3 8b d6 bc 5b 06 0c

Decim v2 Internal State at time -187  
a: a9 25 c5 bc 44 91 77 c8 81 31 d1 04 5f 6e 9f 21 84 fb 38 bd 6b c5 b0 60

Decim v2 Internal State at time -186  
a: ca 92 5c 5b c4 49 17 7c 88 13 1d 10 45 f6 e9 f2 18 4f b3 8b d6 bc 5b 06

Decim v2 Internal State at time -185  
a: fc a9 25 c5 bc 44 91 77 c8 81 31 d1 04 5f 6e 9f 21 84 fb 38 bd 6b c5 b0

Decim v2 Internal State at time -184  
a: 9f ca 92 5c 5b c4 49 17 7c 88 13 1d 10 45 f6 e9 f2 18 4f b3 8b d6 bc 5b

Decim v2 Internal State at time -183  
a: 89 fc a9 25 c5 bc 44 91 77 c8 81 31 d1 04 5f 6e 9f 21 84 fb 38 bd 6b c5

Decim v2 Internal State at time -182  
a: f8 9f ca 92 5c 5b c4 49 17 7c 88 13 1d 10 45 f6 e9 f2 18 4f b3 8b d6 bc

Decim v2 Internal State at time -181  
a: af 89 fc a9 25 c5 bc 44 91 77 c8 81 31 d1 04 5f 6e 9f 21 84 fb 38 bd 6b

Decim v2 Internal State at time -180  
a: 6a f8 9f ca 92 5c 5b c4 49 17 7c 88 13 1d 10 45 f6 e9 f2 18 4f b3 8b d6

Decim v2 Internal State at time -179  
a: 46 af 89 fc a9 25 c5 bc 44 91 77 c8 81 31 d1 04 5f 6e 9f 21 84 fb 38 bd

Decim v2 Internal State at time -178  
a: b4 6a f8 9f ca 92 5c 5b c4 49 17 7c 88 13 1d 10 45 f6 e9 f2 18 4f b3 8b

Decim v2 Internal State at time -177  
a: 6b 46 af 89 fc a9 25 c5 bc 44 91 77 c8 81 31 d1 04 5f 6e 9f 21 84 fb 38

Decim v2 Internal State at time -176  
a: 76 b4 6a f8 9f ca 92 5c 5b c4 49 17 7c 88 13 1d 10 45 f6 e9 f2 18 4f b3

Decim v2 Internal State at time -175  
a: e7 6b 46 af 89 fc a9 25 c5 bc 44 91 77 c8 81 31 d1 04 5f 6e 9f 21 84 fb

Decim v2 Internal State at time -174  
a: 2e 76 b4 6a f8 9f ca 92 5c 5b c4 49 17 7c 88 13 1d 10 45 f6 e9 f2 18 4f

Decim v2 Internal State at time -173  
a: d2 e7 6b 46 af 89 fc a9 25 c5 bc 44 91 77 c8 81 31 d1 04 5f 6e 9f 21 84

Decim v2 Internal State at time -172  
a: 8d 2e 76 b4 6a f8 9f ca 92 5c 5b c4 49 17 7c 88 13 1d 10 45 f6 e9 f2 18

Decim v2 Internal State at time -171  
a: f8 d2 e7 6b 46 af 89 fc a9 25 c5 bc 44 91 77 c8 81 31 d1 04 5f 6e 9f 21

Decim v2 Internal State at time -170  
a: 6f 8d 2e 76 b4 6a f8 9f ca 92 5c 5b c4 49 17 7c 88 13 1d 10 45 f6 e9 f2

Decim v2 Internal State at time -169  
a: d6 f8 d2 e7 6b 46 af 89 fc a9 25 c5 bc 44 91 77 c8 81 31 d1 04 5f 6e 9f

Decim v2 Internal State at time -168  
a: cd 6f 8d 2e 76 b4 6a f8 9f ca 92 5c 5b c4 49 17 7c 88 13 1d 10 45 f6 e9

Decim v2 Internal State at time -167  
a: dc d6 f8 d2 e7 6b 46 af 89 fc a9 25 c5 bc 44 91 77 c8 81 31 d1 04 5f 6e

Decim v2 Internal State at time -166  
a: ed cd 6f 8d 2e 76 b4 6a f8 9f ca 92 5c 5b c4 49 17 7c 88 13 1d 10 45 f6

Decim v2 Internal State at time -165  
a: fe dc d6 f8 d2 e7 6b 46 af 89 fc a9 25 c5 bc 44 91 77 c8 81 31 d1 04 5f

Decim v2 Internal State at time -164  
a: 1f ed cd 6f 8d 2e 76 b4 6a f8 9f ca 92 5c 5b c4 49 17 7c 88 13 1d 10 45

Decim v2 Internal State at time -163  
a: 21 fe dc d6 f8 d2 e7 6b 46 af 89 fc a9 25 c5 bc 44 91 77 c8 81 31 d1 04

Decim v2 Internal State at time -162  
a: 22 1f ed cd 6f 8d 2e 76 b4 6a f8 9f ca 92 5c 5b c4 49 17 7c 88 13 1d 10

Decim v2 Internal State at time -161  
a: d2 21 fe dc d6 f8 d2 e7 6b 46 af 89 fc a9 25 c5 bc 44 91 77 c8 81 31 d1

Decim v2 Internal State at time -160  
a: 3d 22 1f ed cd 6f 8d 2e 76 b4 6a f8 9f ca 92 5c 5b c4 49 17 7c 88 13 1d

Decim v2 Internal State at time -159  
a: 53 d2 21 fe dc d6 f8 d2 e7 6b 46 af 89 fc a9 25 c5 bc 44 91 77 c8 81 31

Decim v2 Internal State at time -158  
a: 45 3d 22 1f ed cd 6f 8d 2e 76 b4 6a f8 9f ca 92 5c 5b c4 49 17 7c 88 13

Decim v2 Internal State at time -157  
a: d4 53 d2 21 fe dc d6 f8 d2 e7 6b 46 af 89 fc a9 25 c5 bc 44 91 77 c8 81

Decim v2 Internal State at time -156  
a: ed 45 3d 22 1f ed cd 6f 8d 2e 76 b4 6a f8 9f ca 92 5c 5b c4 49 17 7c 88

Decim v2 Internal State at time -155  
a: 3e d4 53 d2 21 fe dc d6 f8 d2 e7 6b 46 af 89 fc a9 25 c5 bc 44 91 77 c8

Decim v2 Internal State at time -154  
a: 23 ed 45 3d 22 1f ed cd 6f 8d 2e 76 b4 6a f8 9f ca 92 5c 5b c4 49 17 7c

Decim v2 Internal State at time -153  
a: b2 3e d4 53 d2 21 fe dc d6 f8 d2 e7 6b 46 af 89 fc a9 25 c5 bc 44 91 77

Decim v2 Internal State at time -152  
a: 1b 23 ed 45 3d 22 1f ed cd 6f 8d 2e 76 b4 6a f8 9f ca 92 5c 5b c4 49 17

Decim v2 Internal State at time -151  
a: 91 b2 3e d4 53 d2 21 fe dc d6 f8 d2 e7 6b 46 af 89 fc a9 25 c5 bc 44 91

Decim v2 Internal State at time -150  
a: e9 1b 23 ed 45 3d 22 1f ed cd 6f 8d 2e 76 b4 6a f8 9f ca 92 5c 5b c4 49

Decim v2 Internal State at time -149  
a: fe 91 b2 3e d4 53 d2 21 fe dc d6 f8 d2 e7 6b 46 af 89 fc a9 25 c5 bc 44

Decim v2 Internal State at time -148  
a: cf e9 1b 23 ed 45 3d 22 1f ed cd 6f 8d 2e 76 b4 6a f8 9f ca 92 5c 5b c4

Decim v2 Internal State at time -147  
a: 4c fe 91 b2 3e d4 53 d2 21 fe dc d6 f8 d2 e7 6b 46 af 89 fc a9 25 c5 bc

Decim v2 Internal State at time -146  
a: 64 cf e9 1b 23 ed 45 3d 22 1f ed cd 6f 8d 2e 76 b4 6a f8 9f ca 92 5c 5b

Decim v2 Internal State at time -145  
a: 46 4c fe 91 b2 3e d4 53 d2 21 fe dc d6 f8 d2 e7 6b 46 af 89 fc a9 25 c5

Decim v2 Internal State at time -144  
a: 94 64 cf e9 1b 23 ed 45 3d 22 1f ed cd 6f 8d 2e 76 b4 6a f8 9f ca 92 5c

Decim v2 Internal State at time -143  
a: d9 46 4c fe 91 b2 3e d4 53 d2 21 fe dc d6 f8 d2 e7 6b 46 af 89 fc a9 25

Decim v2 Internal State at time -142  
a: 6d 94 64 cf e9 1b 23 ed 45 3d 22 1f ed cd 6f 8d 2e 76 b4 6a f8 9f ca 92

Decim v2 Internal State at time -141  
a: b6 d9 46 4c fe 91 b2 3e d4 53 d2 21 fe dc d6 f8 d2 e7 6b 46 af 89 fc a9

Decim v2 Internal State at time -140  
a: 9b 6d 94 64 cf e9 1b 23 ed 45 3d 22 1f ed cd 6f 8d 2e 76 b4 6a f8 9f ca

Decim v2 Internal State at time -139  
a: f9 b6 d9 46 4c fe 91 b2 3e d4 53 d2 21 fe dc d6 f8 d2 e7 6b 46 af 89 fc

Decim v2 Internal State at time -138  
a: 1f 9b 6d 94 64 cf e9 1b 23 ed 45 3d 22 1f ed cd 6f 8d 2e 76 b4 6a f8 9f

Decim v2 Internal State at time -137  
a: 51 f9 b6 d9 46 4c fe 91 b2 3e d4 53 d2 21 fe dc d6 f8 d2 e7 6b 46 af 89

Decim v2 Internal State at time -136  
a: 85 1f 9b 6d 94 64 cf e9 1b 23 ed 45 3d 22 1f ed cd 6f 8d 2e 76 b4 6a f8

Decim v2 Internal State at time -135  
a: 18 51 f9 b6 d9 46 4c fe 91 b2 3e d4 53 d2 21 fe dc d6 f8 d2 e7 6b 46 af

Decim v2 Internal State at time -134  
a: e1 85 1f 9b 6d 94 64 cf e9 1b 23 ed 45 3d 22 1f ed cd 6f 8d 2e 76 b4 6a

Decim v2 Internal State at time -133  
a: 6e 18 51 f9 b6 d9 46 4c fe 91 b2 3e d4 53 d2 21 fe dc d6 f8 d2 e7 6b 46

Decim v2 Internal State at time -132  
a: f6 e1 85 1f 9b 6d 94 64 cf e9 1b 23 ed 45 3d 22 1f ed cd 6f 8d 2e 76 b4

Decim v2 Internal State at time -131  
a: 9f 6e 18 51 f9 b6 d9 46 4c fe 91 b2 3e d4 53 d2 21 fe dc d6 f8 d2 e7 6b

Decim v2 Internal State at time -130  
a: 39 f6 e1 85 1f 9b 6d 94 64 cf e9 1b 23 ed 45 3d 22 1f ed cd 6f 8d 2e 76

Decim v2 Internal State at time -129  
a: c3 9f 6e 18 51 f9 b6 d9 46 4c fe 91 b2 3e d4 53 d2 21 fe dc d6 f8 d2 e7

Decim v2 Internal State at time -128  
a: 8c 39 f6 e1 85 1f 9b 6d 94 64 cf e9 1b 23 ed 45 3d 22 1f ed cd 6f 8d 2e

Decim v2 Internal State at time -127  
a: 38 c3 9f 6e 18 51 f9 b6 d9 46 4c fe 91 b2 3e d4 53 d2 21 fe dc d6 f8 d2

Decim v2 Internal State at time -126  
a: 53 8c 39 f6 e1 85 1f 9b 6d 94 64 cf e9 1b 23 ed 45 3d 22 1f ed cd 6f 8d

Decim v2 Internal State at time -125  
a: 55 38 c3 9f 6e 18 51 f9 b6 d9 46 4c fe 91 b2 3e d4 53 d2 21 fe dc d6 f8

Decim v2 Internal State at time -124  
a: 75 53 8c 39 f6 e1 85 1f 9b 6d 94 64 cf e9 1b 23 ed 45 3d 22 1f ed cd 6f

Decim v2 Internal State at time -123  
a: 77 55 38 c3 9f 6e 18 51 f9 b6 d9 46 4c fe 91 b2 3e d4 53 d2 21 fe dc d6

Decim v2 Internal State at time -122  
a: 07 75 53 8c 39 f6 e1 85 1f 9b 6d 94 64 cf e9 1b 23 ed 45 3d 22 1f ed cd

Decim v2 Internal State at time -121  
a: 10 77 55 38 c3 9f 6e 18 51 f9 b6 d9 46 4c fe 91 b2 3e d4 53 d2 21 fe dc

Decim v2 Internal State at time -120  
a: f1 07 75 53 8c 39 f6 e1 85 1f 9b 6d 94 64 cf e9 1b 23 ed 45 3d 22 1f ed

Decim v2 Internal State at time -119  
a: 8f 10 77 55 38 c3 9f 6e 18 51 f9 b6 d9 46 4c fe 91 b2 3e d4 53 d2 21 fe

Decim v2 Internal State at time -118  
a: d8 f1 07 75 53 8c 39 f6 e1 85 1f 9b 6d 94 64 cf e9 1b 23 ed 45 3d 22 1f

Decim v2 Internal State at time -117  
a: 9d 8f 10 77 55 38 c3 9f 6e 18 51 f9 b6 d9 46 4c fe 91 b2 3e d4 53 d2 21

Decim v2 Internal State at time -116  
a: 69 d8 f1 07 75 53 8c 39 f6 e1 85 1f 9b 6d 94 64 cf e9 1b 23 ed 45 3d 22

# BS ISO/IEC 18033-4:2011+A1:2020

## ISO/IEC 18033-4:2011+A1:2020(E)

Decim v2 Internal State at time -115  
a: 86 9d 8f 10 77 55 38 c3 9f 6e 18 51 f9 b6 d9 46 4c fe 91 b2 3e d4 53 d2

Decim v2 Internal State at time -114  
a: 48 69 d8 f1 07 75 53 8c 39 f6 e1 85 1f 9b 6d 94 64 cf e9 1b 23 ed 45 3d

Decim v2 Internal State at time -113  
a: 64 86 9d 8f 10 77 55 38 c3 9f 6e 18 51 f9 b6 d9 46 4c fe 91 b2 3e d4 53

Decim v2 Internal State at time -112  
a: 06 48 69 d8 f1 07 75 53 8c 39 f6 e1 85 1f 9b 6d 94 64 cf e9 1b 23 ed 45

Decim v2 Internal State at time -111  
a: 60 64 86 9d 8f 10 77 55 38 c3 9f 6e 18 51 f9 b6 d9 46 4c fe 91 b2 3e d4

Decim v2 Internal State at time -110  
a: b6 06 48 69 d8 f1 07 75 53 8c 39 f6 e1 85 1f 9b 6d 94 64 cf e9 1b 23 ed

Decim v2 Internal State at time -109  
a: 1b 60 64 86 9d 8f 10 77 55 38 c3 9f 6e 18 51 f9 b6 d9 46 4c fe 91 b2 3e

Decim v2 Internal State at time -108  
a: 71 b6 06 48 69 d8 f1 07 75 53 8c 39 f6 e1 85 1f 9b 6d 94 64 cf e9 1b 23

Decim v2 Internal State at time -107  
a: 47 1b 60 64 86 9d 8f 10 77 55 38 c3 9f 6e 18 51 f9 b6 d9 46 4c fe 91 b2

Decim v2 Internal State at time -106  
a: a4 71 b6 06 48 69 d8 f1 07 75 53 8c 39 f6 e1 85 1f 9b 6d 94 64 cf e9 1b

Decim v2 Internal State at time -105  
a: ba 47 1b 60 64 86 9d 8f 10 77 55 38 c3 9f 6e 18 51 f9 b6 d9 46 4c fe 91

Decim v2 Internal State at time -104  
a: bb a4 71 b6 06 48 69 d8 f1 07 75 53 8c 39 f6 e1 85 1f 9b 6d 94 64 cf e9

Decim v2 Internal State at time -103  
a: 5b ba 47 1b 60 64 86 9d 8f 10 77 55 38 c3 9f 6e 18 51 f9 b6 d9 46 4c fe

Decim v2 Internal State at time -102  
a: 55 bb a4 71 b6 06 48 69 d8 f1 07 75 53 8c 39 f6 e1 85 1f 9b 6d 94 64 cf

Decim v2 Internal State at time -101  
a: 35 5b ba 47 1b 60 64 86 9d 8f 10 77 55 38 c3 9f 6e 18 51 f9 b6 d9 46 4c

Decim v2 Internal State at time -100  
a: a3 55 bb a4 71 b6 06 48 69 d8 f1 07 75 53 8c 39 f6 e1 85 1f 9b 6d 94 64

Decim v2 Internal State at time -99  
a: 8a 35 5b ba 47 1b 60 64 86 9d 8f 10 77 55 38 c3 9f 6e 18 51 f9 b6 d9 46

Decim v2 Internal State at time -98  
a: 58 a3 55 bb a4 71 b6 06 48 69 d8 f1 07 75 53 8c 39 f6 e1 85 1f 9b 6d 94

Decim v2 Internal State at time -97  
a: 15 8a 35 5b ba 47 1b 60 64 86 9d 8f 10 77 55 38 c3 9f 6e 18 51 f9 b6 d9

Decim v2 Internal State at time -96  
a: d1 58 a3 55 bb a4 71 b6 06 48 69 d8 f1 07 75 53 8c 39 f6 e1 85 1f 9b 6d

Decim v2 Internal State at time -95  
a: 9d 15 8a 35 5b ba 47 1b 60 64 86 9d 8f 10 77 55 38 c3 9f 6e 18 51 f9 b6

Decim v2 Internal State at time -94  
a: 69 d1 58 a3 55 bb a4 71 b6 06 48 69 d8 f1 07 75 53 8c 39 f6 e1 85 1f 9b

Decim v2 Internal State at time -93  
a: 96 9d 15 8a 35 5b ba 47 1b 60 64 86 9d 8f 10 77 55 38 c3 9f 6e 18 51 f9

Decim v2 Internal State at time -92  
a: 49 69 d1 58 a3 55 bb a4 71 b6 06 48 69 d8 f1 07 75 53 8c 39 f6 e1 85 1f

Decim v2 Internal State at time -91  
a: 14 96 9d 15 8a 35 5b ba 47 1b 60 64 86 9d 8f 10 77 55 38 c3 9f 6e 18 51

Decim v2 Internal State at time -90  
a: a1 49 69 d1 58 a3 55 bb a4 71 b6 06 48 69 d8 f1 07 75 53 8c 39 f6 e1 85

Decim v2 Internal State at time -89  
a: 5a 14 96 9d 15 8a 35 5b ba 47 1b 60 64 86 9d 8f 10 77 55 38 c3 9f 6e 18

Decim v2 Internal State at time -88  
a: a5 a1 49 69 d1 58 a3 55 bb a4 71 b6 06 48 69 d8 f1 07 75 53 8c 39 f6 e1

Decim v2 Internal State at time -87  
a: da 5a 14 96 9d 15 8a 35 5b ba 47 1b 60 64 86 9d 8f 10 77 55 38 c3 9f 6e

Decim v2 Internal State at time -86  
a: 8d a5 a1 49 69 d1 58 a3 55 bb a4 71 b6 06 48 69 d8 f1 07 75 53 8c 39 f6

Decim v2 Internal State at time -85  
a: d8 da 5a 14 96 9d 15 8a 35 5b ba 47 1b 60 64 86 9d 8f 10 77 55 38 c3 9f

Decim v2 Internal State at time -84  
a: ad 8d a5 a1 49 69 d1 58 a3 55 bb a4 71 b6 06 48 69 d8 f1 07 75 53 8c 39

Decim v2 Internal State at time -83  
a: 7a d8 da 5a 14 96 9d 15 8a 35 5b ba 47 1b 60 64 86 9d 8f 10 77 55 38 c3

Decim v2 Internal State at time -82  
a: 57 ad 8d a5 a1 49 69 d1 58 a3 55 bb a4 71 b6 06 48 69 d8 f1 07 75 53 8c

Decim v2 Internal State at time -81  
a: 75 7a d8 da 5a 14 96 9d 15 8a 35 5b ba 47 1b 60 64 86 9d 8f 10 77 55 38

Decim v2 Internal State at time -80  
a: a7 57 ad 8d a5 a1 49 69 d1 58 a3 55 bb a4 71 b6 06 48 69 d8 f1 07 75 53

Decim v2 Internal State at time -79  
a: 7a 75 7a d8 da 5a 14 96 9d 15 8a 35 5b ba 47 1b 60 64 86 9d 8f 10 77 55

Decim v2 Internal State at time -78  
a: 87 a7 57 ad 8d a5 a1 49 69 d1 58 a3 55 bb a4 71 b6 06 48 69 d8 f1 07 75

Decim v2 Internal State at time -77  
a: c8 7a 75 7a d8 da 5a 14 96 9d 15 8a 35 5b ba 47 1b 60 64 86 9d 8f 10 77

Decim v2 Internal State at time -76  
a: 4c 87 a7 57 ad 8d a5 a1 49 69 d1 58 a3 55 bb a4 71 b6 06 48 69 d8 f1 07

Decim v2 Internal State at time -75  
a: 34 c8 7a 75 7a d8 da 5a 14 96 9d 15 8a 35 5b ba 47 1b 60 64 86 9d 8f 10

Decim v2 Internal State at time -74  
a: 63 4c 87 a7 57 ad 8d a5 a1 49 69 d1 58 a3 55 bb a4 71 b6 06 48 69 d8 f1

Decim v2 Internal State at time -73  
a: 46 34 c8 7a 75 7a d8 da 5a 14 96 9d 15 8a 35 5b ba 47 1b 60 64 86 9d 8f

Decim v2 Internal State at time -72  
a: 94 63 4c 87 a7 57 ad 8d a5 a1 49 69 d1 58 a3 55 bb a4 71 b6 06 48 69 d8

Decim v2 Internal State at time -71  
a: 49 46 34 c8 7a 75 7a d8 da 5a 14 96 9d 15 8a 35 5b ba 47 1b 60 64 86 9d

Decim v2 Internal State at time -70  
a: 54 94 63 4c 87 a7 57 ad 8d a5 a1 49 69 d1 58 a3 55 bb a4 71 b6 06 48 69

Decim v2 Internal State at time -69  
a: 95 49 46 34 c8 7a 75 7a d8 da 5a 14 96 9d 15 8a 35 5b ba 47 1b 60 64 86

Decim v2 Internal State at time -68  
a: f9 54 94 63 4c 87 a7 57 ad 8d a5 a1 49 69 d1 58 a3 55 bb a4 71 b6 06 48

Decim v2 Internal State at time -67  
a: 3f 95 49 46 34 c8 7a 75 7a d8 da 5a 14 96 9d 15 8a 35 5b ba 47 1b 60 64

Decim v2 Internal State at time -66  
a: 23 f9 54 94 63 4c 87 a7 57 ad 8d a5 a1 49 69 d1 58 a3 55 bb a4 71 b6 06



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## ISO/IEC 18033-4:2011+A1:2020(E)

Decim v2 Internal State at time -65  
a: 82 3f 95 49 46 34 c8 7a 75 7a d8 da 5a 14 96 9d 15 8a 35 5b ba 47 1b 60

Decim v2 Internal State at time -64  
a: 38 23 f9 54 94 63 4c 87 a7 57 ad 8d a5 a1 49 69 d1 58 a3 55 bb a4 71 b6

Decim v2 Internal State at time -64 (Binary notation)  
a: 00111000 00100011 11111001 01010100 10010100 01100011 01011000 11010001  
01101001 01010111 10101101 10001101 10100101 10100001 01001001 10100111  
10000111 01001100 10100011 01010101 10111011 10100100 01110001 10110110  
T: 000 b: 00000000 00000000 00000000 00000000 I: 0

Executions of *InitNext*(S,BUFF)

Decim v2 Internal State at time -63  
a: 23 82 3f 95 49 46 34 c8 7a 75 7a d8 da 5a 14 96 9d 15 8a 35 5b ba 47 1b  
T: 111 b: 00 00 00 00 I: 1

Decim v2 Internal State at time -62  
a: a2 38 23 f9 54 94 63 4c 87 a7 57 ad 8d a5 a1 49 69 d1 58 a3 55 bb a4 71  
T: 111 b: 00 00 00 00 I: 2

Decim v2 Internal State at time -61  
a: 7a 23 82 3f 95 49 46 34 c8 7a 75 7a d8 da 5a 14 96 9d 15 8a 35 5b ba 47  
T: 101 b: 00 00 00 00 I: 3

Decim v2 Internal State at time -60  
a: 07 a2 38 23 f9 54 94 63 4c 87 a7 57 ad 8d a5 a1 49 69 d1 58 a3 55 bb a4  
T: 101 b: 00 00 00 08 I: 4

Decim v2 Internal State at time -59  
a: b0 7a 23 82 3f 95 49 46 34 c8 7a 75 7a d8 da 5a 14 96 9d 15 8a 35 5b ba  
T: 101 b: 00 00 00 18 I: 5

Decim v2 Internal State at time -58  
a: 3b 07 a2 38 23 f9 54 94 63 4c 87 a7 57 ad 8d a5 a1 49 69 d1 58 a3 55 bb  
T: 001 b: 00 00 00 38 I: 6

Decim v2 Internal State at time -57  
a: 33 b0 7a 23 82 3f 95 49 46 34 c8 7a 75 7a d8 da 5a 14 96 9d 15 8a 35 5b  
T: 101 b: 00 00 00 78 I: 7

Decim v2 Internal State at time -56  
a: 53 3b 07 a2 38 23 f9 54 94 63 4c 87 a7 57 ad 8d a5 a1 49 69 d1 58 a3 55  
T: 001 b: 00 00 00 f8 I: 9

Decim v2 Internal State at time -55  
a: 05 33 b0 7a 23 82 3f 95 49 46 34 c8 7a 75 7a d8 da 5a 14 96 9d 15 8a 35  
T: 001 b: 00 00 00 f8 I: 11

Decim v2 Internal State at time -54  
a: 20 53 3b 07 a2 38 23 f9 54 94 63 4c 87 a7 57 ad 8d a5 a1 49 69 d1 58 a3  
T: 010 b: 00 00 00 f8 I: 13

Decim v2 Internal State at time -53  
a: a2 05 33 b0 7a 23 82 3f 95 49 46 34 c8 7a 75 7a d8 da 5a 14 96 9d 15 8a  
T: 010 b: 00 00 60 f8 I: 15

Decim v2 Internal State at time -52  
a: 3a 20 53 3b 07 a2 38 23 f9 54 94 63 4c 87 a7 57 ad 8d a5 a1 49 69 d1 58  
T: 011 b: 00 00 60 f8 I: 16

Decim v2 Internal State at time -51  
a: 93 a2 05 33 b0 7a 23 82 3f 95 49 46 34 c8 7a 75 7a d8 da 5a 14 96 9d 15  
T: 011 b: 00 00 60 f8 I: 17

Decim v2 Internal State at time -50  
a: d9 3a 20 53 3b 07 a2 38 23 f9 54 94 63 4c 87 a7 57 ad 8d a5 a1 49 69 d1  
T: 010 b: 00 00 60 f8 I: 18

Decim v2 Internal State at time -49  
a: 2d 93 a2 05 33 b0 7a 23 82 3f 95 49 46 34 c8 7a 75 7a d8 da 5a 14 96 9d  
T: 001 b: 00 04 60 f8 I: 19

Decim v2 Internal State at time -48  
a: 42 d9 3a 20 53 3b 07 a2 38 23 f9 54 94 63 4c 87 a7 57 ad 8d a5 a1 49 69

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T: 001          b: 00 14 60 f8          I: 21

Decim v2 Internal State at time -47
a: 64 2d 93 a2 05 33 b0 7a 23 82 3f 95 49 46 34 c8 7a 75 7a d8 da 5a 14 96
T: 001          b: 00 14 60 f8          I: 23

Decim v2 Internal State at time -46
a: 56 42 d9 3a 20 53 3b 07 a2 38 23 f9 54 94 63 4c 87 a7 57 ad 8d a5 a1 49
T: 000          b: 00 94 60 f8          I: 25

Decim v2 Internal State at time -45
a: 45 64 2d 93 a2 05 33 b0 7a 23 82 3f 95 49 46 34 c8 7a 75 7a d8 da 5a 14
T: 000          b: 02 94 60 f8          I: 27

Decim v2 Internal State at time -44
a: e4 56 42 d9 3a 20 53 3b 07 a2 38 23 f9 54 94 63 4c 87 a7 57 ad 8d a5 a1
T: 010          b: 12 94 60 f8          I: 29

Decim v2 Internal State at time -43
a: 7e 45 64 2d 93 a2 05 33 b0 7a 23 82 3f 95 49 46 34 c8 7a 75 7a d8 da 5a
T: 010          b: 12 94 60 f8          I: 30

Decim v2 Internal State at time -42
a: 07 e4 56 42 d9 3a 20 53 3b 07 a2 38 23 f9 54 94 63 4c 87 a7 57 ad 8d a5
T: 000          b: 52 94 60 f8          I: 32

Decim v2 Internal State at time 0
a: 07 e4 56 42 d9 3a 20 53 3b 07 a2 38 23 f9 54 94 63 4c 87 a7 57 ad 8d a5
T: 000          b: 52 94 60 f8          I: 32

Decim v2 Internal State at time 0 (Binary notation)
a: 00000111 11100100 01010110 01000010 11011001 00111010 01001100 01100011
   10010100 00000111 10100010 00111000 00100011 11111001 01010100 00111011
   01010011 00100000 10000111 10100111 01010111 10101101 10001101 10100101
T: 000          b: 01010010 10010100 01100000 11111000          I: 32

Executions of Next(S)

Decim v2 Internal State at time 1
a: a0 7e 45 64 2d 93 a2 05 33 b0 7a 23 82 3f 95 49 46 34 c8 7a 75 7a d8 da
T: 101  b: a9 4a 30 7c  I: 32

Decim v2 Internal State at time 2
a: 0a 07 e4 56 42 d9 3a 20 53 3b 07 a2 38 23 f9 54 94 63 4c 87 a7 57 ad 8d
T: 111  b: d4 a5 18 3e  I: 32

Decim v2 Internal State at time 3
a: b0 a0 7e 45 64 2d 93 a2 05 33 b0 7a 23 82 3f 95 49 46 34 c8 7a 75 7a d8
T: 010  b: 6a 52 8c 1f  I: 32

Decim v2 Internal State at time 4
a: bb 0a 07 e4 56 42 d9 3a 20 53 3b 07 a2 38 23 f9 54 94 63 4c 87 a7 57 ad
T: 111  b: b5 29 46 0f  I: 32

Decim v2 Internal State at time 5
a: eb b0 a0 7e 45 64 2d 93 a2 05 33 b0 7a 23 82 3f 95 49 46 34 c8 7a 75 7a
T: 111  b: 5a 94 a3 07  I: 32

Decim v2 Internal State at time 6
a: ce bb 0a 07 e4 56 42 d9 3a 20 53 3b 07 a2 38 23 f9 54 94 63 4c 87 a7 57
T: 111  b: 2d 4a 51 83  I: 31

Decim v2 Internal State at time 7
a: fc eb b0 a0 7e 45 64 2d 93 a2 05 33 b0 7a 23 82 3f 95 49 46 34 c8 7a 75
T: 110  b: 16 a5 28 c1  I: 32

Decim v2 Internal State at time 8
a: 3f ce bb 0a 07 e4 56 42 d9 3a 20 53 3b 07 a2 38 23 f9 54 94 63 4c 87 a7
T: 111  b: 0b 52 94 60  I: 32

Decim v2 Internal State at time 9
a: f3 fc eb b0 a0 7e 45 64 2d 93 a2 05 33 b0 7a 23 82 3f 95 49 46 34 c8 7a
T: 010  b: 05 a9 4a 30  I: 32

Decim v2 Internal State at time 10
a: 3f 3f ce bb 0a 07 e4 56 42 d9 3a 20 53 3b 07 a2 38 23 f9 54 94 63 4c 87

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# BS ISO/IEC 18033-4:2011+A1:2020

## ISO/IEC 18033-4:2011+A1:2020(E)

T: 010 b: 02 d4 a5 18 I: 32

Decim v2 Internal State at time 11

a: f3 f3 fc eb b0 a0 7e 45 64 2d 93 a2 05 33 b0 7a 23 82 3f 95 49 46 34 c8

T: 101 b: 01 6a 52 8c I: 32

Decim v2 Internal State at time 12

a: 6f 3f 3f ce bb 0a 07 e4 56 42 d9 3a 20 53 3b 07 a2 38 23 f9 54 94 63 4c

T: 100 b: 80 b5 29 46 I: 32

Decim v2 Internal State at time 13

a: 76 f3 f3 fc eb b0 a0 7e 45 64 2d 93 a2 05 33 b0 7a 23 82 3f 95 49 46 34

T: 101 b: c0 5a 94 a3 I: 31

Decim v2 Internal State at time 14

a: 37 6f 3f 3f ce bb 0a 07 e4 56 42 d9 3a 20 53 3b 07 a2 38 23 f9 54 94 63

T: 110 b: e0 2d 4a 51 I: 31

Decim v2 Internal State at time 15

a: 53 76 f3 f3 fc eb b0 a0 7e 45 64 2d 93 a2 05 33 b0 7a 23 82 3f 95 49 46

T: 110 b: 70 16 a5 28 I: 32

Decim v2 Internal State at time 16

a: f5 37 6f 3f 3f ce bb 0a 07 e4 56 42 d9 3a 20 53 3b 07 a2 38 23 f9 54 94

T: 000 b: b8 0b 52 94 I: 32

Decim v2 Internal State at time 17

a: 8f 53 76 f3 f3 fc eb b0 a0 7e 45 64 2d 93 a2 05 33 b0 7a 23 82 3f 95 49

T: 111 b: dc 05 a9 4a I: 31

Decim v2 Internal State at time 18

a: 48 f5 37 6f 3f 3f ce bb 0a 07 e4 56 42 d9 3a 20 53 3b 07 a2 38 23 f9 54

T: 100 b: ae 02 d4 a5 I: 32

Decim v2 Internal State at time 19

a: 74 8f 53 76 f3 f3 fc eb b0 a0 7e 45 64 2d 93 a2 05 33 b0 7a 23 82 3f 95

T: 010 b: 57 01 6a 52 I: 32

Decim v2 Internal State at time 20

a: b7 48 f5 37 6f 3f 3f ce bb 0a 07 e4 56 42 d9 3a 20 53 3b 07 a2 38 23 f9

T: 101 b: 2b 80 b5 29 I: 32

Decim v2 Internal State at time 21

a: eb 74 8f 53 76 f3 f3 fc eb b0 a0 7e 45 64 2d 93 a2 05 33 b0 7a 23 82 3f

T: 101 b: 15 c0 5a 94 I: 31

Decim v2 Internal State at time 22

a: be b7 48 f5 37 6f 3f 3f ce bb 0a 07 e4 56 42 d9 3a 20 53 3b 07 a2 38 23

T: 110 b: 4a e0 2d 4a I: 31

Decim v2 Internal State at time 23

a: 9b eb 74 8f 53 76 f3 f3 fc eb b0 a0 7e 45 64 2d 93 a2 05 33 b0 7a 23 82

T: 000 b: 25 70 16 a5 I: 32

Decim v2 Internal State at time 24

a: 39 be b7 48 f5 37 6f 3f 3f ce bb 0a 07 e4 56 42 d9 3a 20 53 3b 07 a2 38

T: 100 b: 92 b8 0b 52 I: 32

Decim v2 Internal State at time 25

a: 43 9b eb 74 8f 53 76 f3 f3 fc eb b0 a0 7e 45 64 2d 93 a2 05 33 b0 7a 23

T: 101 b: c9 5c 05 a9 I: 32

Decim v2 Internal State at time 26

a: 24 39 be b7 48 f5 37 6f 3f 3f ce bb 0a 07 e4 56 42 d9 3a 20 53 3b 07 a2

T: 010 b: e4 ae 02 d4 I: 32

Decim v2 Internal State at time 27

a: 52 43 9b eb 74 8f 53 76 f3 f3 fc eb b0 a0 7e 45 64 2d 93 a2 05 33 b0 7a

T: 111 b: 72 57 01 6a I: 32

Decim v2 Internal State at time 28

a: b5 24 39 be b7 48 f5 37 6f 3f 3f ce bb 0a 07 e4 56 42 d9 3a 20 53 3b 07

T: 100 b: 39 2b 80 b5 I: 32

Decim v2 Internal State at time 29

a: 6b 52 43 9b eb 74 8f 53 76 f3 f3 fc eb b0 a0 7e 45 64 2d 93 a2 05 33 b0

T: 101 b: 1c 95 c0 5a I: 31

Decim v2 Internal State at time 30

a: d6 b5 24 39 be b7 48 f5 37 6f 3f 3f ce bb 0a 07 e4 56 42 d9 3a 20 53 3b  
T: 000 b: 4e 4a e0 2d I: 32

Decim v2 Internal State at time 31

a: bd 6b 52 43 9b eb 74 8f 53 76 f3 f3 fc eb b0 a0 7e 45 64 2d 93 a2 05 33  
T: 010 b: a7 25 70 16 I: 32

Decim v2 Internal State at time 32

a: 4b d6 b5 24 39 be b7 48 f5 37 6f 3f 3f ce bb 0a 07 e4 56 42 d9 3a 20 53  
T: 110 b: d3 92 b8 0b I: 32

Decim v2 Internal State at time 33

a: c4 bd 6b 52 43 9b eb 74 8f 53 76 f3 f3 fc eb b0 a0 7e 45 64 2d 93 a2 05  
T: 110 b: e9 c9 5c 05 I: 32

Decim v2 Internal State at time 34

a: 3c 4b d6 b5 24 39 be b7 48 f5 37 6f 3f 3f ce bb 0a 07 e4 56 42 d9 3a 20  
T: 111 b: f4 e4 ae 02 I: 32

Decim v2 Internal State at time 35

a: c3 c4 bd 6b 52 43 9b eb 74 8f 53 76 f3 f3 fc eb b0 a0 7e 45 64 2d 93 a2  
T: 101 b: 7a 72 57 01 I: 32

Decim v2 Internal State at time 36

a: dc 3c 4b d6 b5 24 39 be b7 48 f5 37 6f 3f 3f ce bb 0a 07 e4 56 42 d9 3a  
T: 110 b: bd 39 2b 80 I: 32

Decim v2 Internal State at time 37

a: fd c3 c4 bd 6b 52 43 9b eb 74 8f 53 76 f3 f3 fc eb b0 a0 7e 45 64 2d 93  
T: 110 b: de 9c 95 c0 I: 32

Decim v2 Internal State at time 38

a: af dc 3c 4b d6 b5 24 39 be b7 48 f5 37 6f 3f 3f ce bb 0a 07 e4 56 42 d9  
T: 110 b: ef 4e 4a e0 I: 32

Decim v2 Internal State at time 39

a: 2a fd c3 c4 bd 6b 52 43 9b eb 74 8f 53 76 f3 f3 fc eb b0 a0 7e 45 64 2d  
T: 010 b: f7 a7 25 70 I: 32

Decim v2 Internal State at time 40

a: e2 af dc 3c 4b d6 b5 24 39 be b7 48 f5 37 6f 3f 3f ce bb 0a 07 e4 56 42  
T: 000 b: fb d3 92 b8 I: 32

Decim v2 Internal State at time 41

a: fe 2a fd c3 c4 bd 6b 52 43 9b eb 74 8f 53 76 f3 f3 fc eb b0 a0 7e 45 64  
T: 111 b: 7d e9 c9 5c I: 32

Decim v2 Internal State at time 42

a: 9f e2 af dc 3c 4b d6 b5 24 39 be b7 48 f5 37 6f 3f 3f ce bb 0a 07 e4 56  
T: 000 b: 3e f4 e4 ae I: 32

Decim v2 Internal State at time 43

a: 59 fe 2a fd c3 c4 bd 6b 52 43 9b eb 74 8f 53 76 f3 f3 fc eb b0 a0 7e 45  
T: 010 b: 1f 7a 72 57 I: 32

Decim v2 Internal State at time 44

a: 35 9f e2 af dc 3c 4b d6 b5 24 39 be b7 48 f5 37 6f 3f 3f ce bb 0a 07 e4  
T: 000 b: 8f bd 39 2b I: 32

Decim v2 Internal State at time 45

a: d3 59 fe 2a fd c3 c4 bd 6b 52 43 9b eb 74 8f 53 76 f3 f3 fc eb b0 a0 7e  
T: 111 b: 47 de 9c 95 I: 32

Decim v2 Internal State at time 46

a: 1d 35 9f e2 af dc 3c 4b d6 b5 24 39 be b7 48 f5 37 6f 3f 3f ce bb 0a 07  
T: 000 b: 23 ef 4e 4a I: 32

Decim v2 Internal State at time 47

a: 81 d3 59 fe 2a fd c3 c4 bd 6b 52 43 9b eb 74 8f 53 76 f3 f3 fc eb b0 a0  
T: 101 b: 11 f7 a7 25 I: 31

Decim v2 Internal State at time 48

a: f8 1d 35 9f e2 af dc 3c 4b d6 b5 24 39 be b7 48 f5 37 6f 3f 3f ce bb 0a

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T: 101    b: 48 fb d3 92    I: 31

Decim v2 Internal State at time 49  
a: 6f 81 d3 59 fe 2a fd c3 c4 bd 6b 52 43 9b eb 74 8f 53 76 f3 f3 fc eb b0  
T: 110    b: 64 7d e9 c9    I: 32

Decim v2 Internal State at time 50  
a: 86 f8 1d 35 9f e2 af dc 3c 4b d6 b5 24 39 be b7 48 f5 37 6f 3f 3f ce bb  
T: 111    b: 32 3e f4 e4    I: 32

Decim v2 Internal State at time 51  
a: e8 6f 81 d3 59 fe 2a fd c3 c4 bd 6b 52 43 9b eb 74 8f 53 76 f3 f3 fc eb  
T: 111    b: 19 1f 7a 72    I: 31

Decim v2 Internal State at time 52  
a: fe 86 f8 1d 35 9f e2 af dc 3c 4b d6 b5 24 39 be b7 48 f5 37 6f 3f 3f ce  
T: 101    b: 0c 8f bd 39    I: 31

Decim v2 Internal State at time 53  
a: 0f e8 6f 81 d3 59 fe 2a fd c3 c4 bd 6b 52 43 9b eb 74 8f 53 76 f3 f3 fc  
T: 000    b: 46 47 de 9c    I: 31

Decim v2 Internal State at time 54  
a: 70 fe 86 f8 1d 35 9f e2 af dc 3c 4b d6 b5 24 39 be b7 48 f5 37 6f 3f 3f  
T: 100    b: 63 23 ef 4e    I: 31

Decim v2 Internal State at time 55  
a: 27 0f e8 6f 81 d3 59 fe 2a fd c3 c4 bd 6b 52 43 9b eb 74 8f 53 76 f3 f3  
T: 111    b: 71 91 f7 a7    I: 31

Decim v2 Internal State at time 56  
a: 32 70 fe 86 f8 1d 35 9f e2 af dc 3c 4b d6 b5 24 39 be b7 48 f5 37 6f 3f  
T: 100    b: b8 c8 fb d3    I: 32

Decim v2 Internal State at time 57  
a: 73 27 0f e8 6f 81 d3 59 fe 2a fd c3 c4 bd 6b 52 43 9b eb 74 8f 53 76 f3  
T: 100    b: dc 64 7d e9    I: 32

Decim v2 Internal State at time 58  
a: 47 32 70 fe 86 f8 1d 35 9f e2 af dc 3c 4b d6 b5 24 39 be b7 48 f5 37 6f  
T: 111    b: ee 32 3e f4    I: 32

Decim v2 Internal State at time 59  
a: b4 73 27 0f e8 6f 81 d3 59 fe 2a fd c3 c4 bd 6b 52 43 9b eb 74 8f 53 76  
T: 000    b: 77 19 1f 7a    I: 32

Decim v2 Internal State at time 60  
a: 0b 47 32 70 fe 86 f8 1d 35 9f e2 af dc 3c 4b d6 b5 24 39 be b7 48 f5 37  
T: 010    b: 3b 8c 8f bd    I: 32

Decim v2 Internal State at time 61  
a: 50 b4 73 27 0f e8 6f 81 d3 59 fe 2a fd c3 c4 bd 6b 52 43 9b eb 74 8f 53  
T: 100    b: 1d c6 47 de    I: 32

Decim v2 Internal State at time 62  
a: f5 0b 47 32 70 fe 86 f8 1d 35 9f e2 af dc 3c 4b d6 b5 24 39 be b7 48 f5  
T: 101    b: 0e e3 23 ef    I: 32

Decim v2 Internal State at time 63  
a: af 50 b4 73 27 0f e8 6f 81 d3 59 fe 2a fd c3 c4 bd 6b 52 43 9b eb 74 8f  
T: 111    b: 87 71 91 f7    I: 32

Decim v2 Internal State at time 64  
a: ba f5 0b 47 32 70 fe 86 f8 1d 35 9f e2 af dc 3c 4b d6 b5 24 39 be b7 48  
T: 111    b: 43 b8 c8 fb    I: 32

Decim v2 Internal State at time 65  
a: bb af 50 b4 73 27 0f e8 6f 81 d3 59 fe 2a fd c3 c4 bd 6b 52 43 9b eb 74  
T: 010    b: 21 dc 64 7d    I: 32

Decim v2 Internal State at time 66  
a: eb ba f5 0b 47 32 70 fe 86 f8 1d 35 9f e2 af dc 3c 4b d6 b5 24 39 be b7  
T: 000    b: 90 ee 32 3e    I: 32

Decim v2 Internal State at time 67  
a: 2e bb af 50 b4 73 27 0f e8 6f 81 d3 59 fe 2a fd c3 c4 bd 6b 52 43 9b eb

T: 110 b: c8 77 19 1f I: 32

Decim v2 Internal State at time 68

a: f2 eb ba f5 0b 47 32 70 fe 86 f8 1d 35 9f e2 af dc 3c 4b d6 b5 24 39 be

T: 100 b: e4 3b 8c 8f I: 32

Decim v2 Internal State at time 69

a: cf 2e bb af 50 b4 73 27 0f e8 6f 81 d3 59 fe 2a fd c3 c4 bd 6b 52 43 9b

T: 101 b: f2 1d c6 47 I: 32

Decim v2 Internal State at time 70

a: bc f2 eb ba f5 0b 47 32 70 fe 86 f8 1d 35 9f e2 af dc 3c 4b d6 b5 24 39

T: 000 b: f9 0e e3 23 I: 32

Decim v2 Internal State at time 71

a: ab cf 2e bb af 50 b4 73 27 0f e8 6f 81 d3 59 fe 2a fd c3 c4 bd 6b 52 43

T: 000 b: fc 87 71 91 I: 32

Decim v2 Internal State at time 72

a: ea bc f2 eb ba f5 0b 47 32 70 fe 86 f8 1d 35 9f e2 af dc 3c 4b d6 b5 24

T: 101 b: fe 43 b8 c8 I: 31

Decim v2 Internal State at time 73

a: ce ab cf 2e bb af 50 b4 73 27 0f e8 6f 81 d3 59 fe 2a fd c3 c4 bd 6b 52

T: 010 b: 7f 21 dc 64 I: 32

Decim v2 Internal State at time 74

a: 6c ea bc f2 eb ba f5 0b 47 32 70 fe 86 f8 1d 35 9f e2 af dc 3c 4b d6 b5

T: 010 b: 3f 90 ee 32 I: 32

Decim v2 Internal State at time 75

a: 36 ce ab cf 2e bb af 50 b4 73 27 0f e8 6f 81 d3 59 fe 2a fd c3 c4 bd 6b

T: 110 b: 1f c8 77 19 I: 32

Decim v2 Internal State at time 76

a: 63 6c ea bc f2 eb ba f5 0b 47 32 70 fe 86 f8 1d 35 9f e2 af dc 3c 4b d6

T: 110 b: 0f e4 3b 8c I: 32

Decim v2 Internal State at time 77

a: b6 36 ce ab cf 2e bb af 50 b4 73 27 0f e8 6f 81 d3 59 fe 2a fd c3 c4 bd

T: 000 b: 07 f2 1d c6 I: 32

Decim v2 Internal State at time 78

a: 3b 63 6c ea bc f2 eb ba f5 0b 47 32 70 fe 86 f8 1d 35 9f e2 af dc 3c 4b

T: 100 b: 83 f9 0e e3 I: 32

Decim v2 Internal State at time 79

a: a3 b6 36 ce ab cf 2e bb af 50 b4 73 27 0f e8 6f 81 d3 59 fe 2a fd c3 c4

T: 101 b: 41 fc 87 71 I: 32

Decim v2 Internal State at time 80

a: aa 3b 63 6c ea bc f2 eb ba f5 0b 47 32 70 fe 86 f8 1d 35 9f e2 af dc 3c

T: 010 b: a0 fe 43 b8 I: 32

Decim v2 Internal State at time 81

a: 8a a3 b6 36 ce ab cf 2e bb af 50 b4 73 27 0f e8 6f 81 d3 59 fe 2a fd c3

T: 111 b: d0 7f 21 dc I: 31

Decim v2 Internal State at time 82

a: 68 aa 3b 63 6c ea bc f2 eb ba f5 0b 47 32 70 fe 86 f8 1d 35 9f e2 af dc

T: 010 b: a8 3f 90 ee I: 31

Decim v2 Internal State at time 83

a: f6 8a a3 b6 36 ce ab cf 2e bb af 50 b4 73 27 0f e8 6f 81 d3 59 fe 2a fd

T: 110 b: d4 1f c8 77 I: 31

Decim v2 Internal State at time 84

a: df 68 aa 3b 63 6c ea bc f2 eb ba f5 0b 47 32 70 fe 86 f8 1d 35 9f e2 af

T: 110 b: ea 0f e4 3b I: 32

Decim v2 Internal State at time 85

a: bd f6 8a a3 b6 36 ce ab cf 2e bb af 50 b4 73 27 0f e8 6f 81 d3 59 fe 2a

T: 111 b: 75 07 f2 1d I: 32

Decim v2 Internal State at time 86

a: 8b df 68 aa 3b 63 6c ea bc f2 eb ba f5 0b 47 32 70 fe 86 f8 1d 35 9f e2

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T: 110    b: 3a 83 f9 0e    I: 32

Decim v2 Internal State at time 87  
a: e8 bd f6 8a a3 b6 36 ce ab cf 2e bb af 50 b4 73 27 0f e8 6f 81 d3 59 fe  
T: 101    b: 9d 41 fc 87    I: 32

Decim v2 Internal State at time 88  
a: 2e 8b df 68 aa 3b 63 6c ea bc f2 eb ba f5 0b 47 32 70 fe 86 f8 1d 35 9f  
T: 000    b: ce a0 fe 43    I: 32

Decim v2 Internal State at time 89  
a: 92 e8 bd f6 8a a3 b6 36 ce ab cf 2e bb af 50 b4 73 27 0f e8 6f 81 d3 59  
T: 010    b: 67 50 7f 21    I: 32

Decim v2 Internal State at time 90  
a: 09 2e 8b df 68 aa 3b 63 6c ea bc f2 eb ba f5 0b 47 32 70 fe 86 f8 1d 35  
T: 110    b: 33 a8 3f 90    I: 32

Decim v2 Internal State at time 91  
a: 20 92 e8 bd f6 8a a3 b6 36 ce ab cf 2e bb af 50 b4 73 27 0f e8 6f 81 d3  
T: 000    b: 99 d4 1f c8    I: 32

Decim v2 Internal State at time 92  
a: 92 09 2e 8b df 68 aa 3b 63 6c ea bc f2 eb ba f5 0b 47 32 70 fe 86 f8 1d  
T: 000    b: cc ea 0f e4    I: 32

Decim v2 Internal State at time 93  
a: f9 20 92 e8 bd f6 8a a3 b6 36 ce ab cf 2e bb af 50 b4 73 27 0f e8 6f 81  
T: 101    b: e6 75 07 f2    I: 32

Decim v2 Internal State at time 94  
a: df 92 09 2e 8b df 68 aa 3b 63 6c ea bc f2 eb ba f5 0b 47 32 70 fe 86 f8  
T: 000    b: f3 3a 83 f9    I: 32

Decim v2 Internal State at time 95  
a: 7d f9 20 92 e8 bd f6 8a a3 b6 36 ce ab cf 2e bb af 50 b4 73 27 0f e8 6f  
T: 111    b: f9 9d 41 fc    I: 31

Decim v2 Internal State at time 96  
a: 77 df 92 09 2e 8b df 68 aa 3b 63 6c ea bc f2 eb ba f5 0b 47 32 70 fe 86  
T: 111    b: bc ce a0 fe    I: 31

Decim v2 Internal State at time 97  
a: 27 7d f9 20 92 e8 bd f6 8a a3 b6 36 ce ab cf 2e bb af 50 b4 73 27 0f e8  
T: 010    b: 9e 67 50 7f    I: 32

Decim v2 Internal State at time 98  
a: c2 77 df 92 09 2e 8b df 68 aa 3b 63 6c ea bc f2 eb ba f5 0b 47 32 70 fe  
T: 110    b: 4f 33 a8 3f    I: 32

Decim v2 Internal State at time 99  
a: 7c 27 7d f9 20 92 e8 bd f6 8a a3 b6 36 ce ab cf 2e bb af 50 b4 73 27 0f  
T: 110    b: a7 99 d4 1f    I: 32

Decim v2 Internal State at time 100  
a: c7 c2 77 df 92 09 2e 8b df 68 aa 3b 63 6c ea bc f2 eb ba f5 0b 47 32 70  
T: 000    b: 53 cc ea 0f    I: 32

Decim v2 Internal State at time 101  
a: 0c 7c 27 7d f9 20 92 e8 bd f6 8a a3 b6 36 ce ab cf 2e bb af 50 b4 73 27  
T: 000    b: a9 e6 75 07    I: 32

Decim v2 Internal State at time 102  
a: e0 c7 c2 77 df 92 09 2e 8b df 68 aa 3b 63 6c ea bc f2 eb ba f5 0b 47 32  
T: 010    b: 54 f3 3a 83    I: 32

Decim v2 Internal State at time 103  
a: 4e 0c 7c 27 7d f9 20 92 e8 bd f6 8a a3 b6 36 ce ab cf 2e bb af 50 b4 73  
T: 010    b: aa 79 9d 41    I: 32

Decim v2 Internal State at time 104  
a: 34 e0 c7 c2 77 df 92 09 2e 8b df 68 aa 3b 63 6c ea bc f2 eb ba f5 0b 47  
T: 000    b: d5 3c ce a0    I: 32

Decim v2 Internal State at time 105  
a: 53 4e 0c 7c 27 7d f9 20 92 e8 bd f6 8a a3 b6 36 ce ab cf 2e bb af 50 b4

T: 110 b: ea 9e 67 50 I: 32

Decim v2 Internal State at time 106

a: 85 34 e0 c7 c2 77 df 92 09 2e 8b df 68 aa 3b 63 6c ea bc f2 eb ba f5 0b

T: 010 b: 75 4f 33 a8 I: 32

Decim v2 Internal State at time 107

a: a8 53 4e 0c 7c 27 7d f9 20 92 e8 bd f6 8a a3 b6 36 ce ab cf 2e bb af 50

T: 010 b: ba a7 99 d4 I: 32

Decim v2 Internal State at time 108

a: 6a 85 34 e0 c7 c2 77 df 92 09 2e 8b df 68 aa 3b 63 6c ea bc f2 eb ba f5

T: 100 b: dd 53 cc ea I: 32

Decim v2 Internal State at time 109

a: 36 a8 53 4e 0c 7c 27 7d f9 20 92 e8 bd f6 8a a3 b6 36 ce ab cf 2e bb af

T: 010 b: 6e a9 e6 75 I: 32

Decim v2 Internal State at time 110

a: 83 6a 85 34 e0 c7 c2 77 df 92 09 2e 8b df 68 aa 3b 63 6c ea bc f2 eb ba

T: 110 b: b7 54 f3 3a I: 32

Decim v2 Internal State at time 111

a: 38 36 a8 53 4e 0c 7c 27 7d f9 20 92 e8 bd f6 8a a3 b6 36 ce ab cf 2e bb

T: 010 b: 5b aa 79 9d I: 32

Decim v2 Internal State at time 112

a: 13 83 6a 85 34 e0 c7 c2 77 df 92 09 2e 8b df 68 aa 3b 63 6c ea bc f2 eb

T: 101 b: ad d5 3c ce I: 32

Decim v2 Internal State at time 113

a: 91 38 36 a8 53 4e 0c 7c 27 7d f9 20 92 e8 bd f6 8a a3 b6 36 ce ab cf 2e

T: 000 b: d6 ea 9e 67 I: 32

Decim v2 Internal State at time 114

a: 49 13 83 6a 85 34 e0 c7 c2 77 df 92 09 2e 8b df 68 aa 3b 63 6c ea bc f2

T: 100 b: eb 75 4f 33 I: 32

Decim v2 Internal State at time 115

a: 04 91 38 36 a8 53 4e 0c 7c 27 7d f9 20 92 e8 bd f6 8a a3 b6 36 ce ab cf

T: 100 b: 75 ba a7 99 I: 32

Decim v2 Internal State at time 116

a: c0 49 13 83 6a 85 34 e0 c7 c2 77 df 92 09 2e 8b df 68 aa 3b 63 6c ea bc

T: 010 b: ba dd 53 cc I: 32

Decim v2 Internal State at time 117

a: ac 04 91 38 36 a8 53 4e 0c 7c 27 7d f9 20 92 e8 bd f6 8a a3 b6 36 ce ab

T: 100 b: dd 6e a9 e6 I: 32

Decim v2 Internal State at time 118

a: ba c0 49 13 83 6a 85 34 e0 c7 c2 77 df 92 09 2e 8b df 68 aa 3b 63 6c ea

T: 010 b: ee b7 54 f3 I: 32

Decim v2 Internal State at time 119

a: 4b ac 04 91 38 36 a8 53 4e 0c 7c 27 7d f9 20 92 e8 bd f6 8a a3 b6 36 ce

T: 000 b: f7 5b aa 79 I: 32

Decim v2 Internal State at time 120

a: 14 ba c0 49 13 83 6a 85 34 e0 c7 c2 77 df 92 09 2e 8b df 68 aa 3b 63 6c

T: 010 b: fb ad d5 3c I: 32

Decim v2 Internal State at time 121

a: 41 4b ac 04 91 38 36 a8 53 4e 0c 7c 27 7d f9 20 92 e8 bd f6 8a a3 b6 36

T: 000 b: fd d6 ea 9e I: 32

Decim v2 Internal State at time 122

a: 74 14 ba c0 49 13 83 6a 85 34 e0 c7 c2 77 df 92 09 2e 8b df 68 aa 3b 63

T: 010 b: fe eb 75 4f I: 32

Decim v2 Internal State at time 123

a: b7 41 4b ac 04 91 38 36 a8 53 4e 0c 7c 27 7d f9 20 92 e8 bd f6 8a a3 b6

T: 100 b: 7f 75 ba a7 I: 32

Decim v2 Internal State at time 124

a: bb 74 14 ba c0 49 13 83 6a 85 34 e0 c7 c2 77 df 92 09 2e 8b df 68 aa 3b



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T: 000    b: bf ba dd 53    I: 32

Decim v2 Internal State at time 125  
a: 9b b7 41 4b ac 04 91 38 36 a8 53 4e 0c 7c 27 7d f9 20 92 e8 bd f6 8a a3  
T: 111    b: df dd 6e a9    I: 32

Decim v2 Internal State at time 126  
a: 29 bb 74 14 ba c0 49 13 83 6a 85 34 e0 c7 c2 77 df 92 09 2e 8b df 68 aa  
T: 100    b: 6f ee b7 54    I: 32

Decim v2 Internal State at time 127  
a: 22 9b b7 41 4b ac 04 91 38 36 a8 53 4e 0c 7c 27 7d f9 20 92 e8 bd f6 8a  
T: 110    b: 37 f7 5b aa    I: 32

Decim v2 Internal State at time 128  
a: 22 29 bb 74 14 ba c0 49 13 83 6a 85 34 e0 c7 c2 77 df 92 09 2e 8b df 68  
T: 111    b: 1b fb ad d5    I: 32

Decim v2 Internal State at time 129  
a: b2 22 9b b7 41 4b ac 04 91 38 36 a8 53 4e 0c 7c 27 7d f9 20 92 e8 bd f6  
T: 101    b: 0d fd d6 ea    I: 32

Decim v2 Internal State at time 130  
a: 4b 22 29 bb 74 14 ba c0 49 13 83 6a 85 34 e0 c7 c2 77 df 92 09 2e 8b df  
T: 100    b: 86 fe eb 75    I: 32

Decim v2 Internal State at time 131  
a: e4 b2 22 9b b7 41 4b ac 04 91 38 36 a8 53 4e 0c 7c 27 7d f9 20 92 e8 bd  
T: 100    b: c3 7f 75 ba    I: 32

Decim v2 Internal State at time 132  
a: 1e 4b 22 29 bb 74 14 ba c0 49 13 83 6a 85 34 e0 c7 c2 77 df 92 09 2e 8b  
T: 101    b: e1 bf ba dd    I: 31

Decim v2 Internal State at time 133  
a: 51 e4 b2 22 9b b7 41 4b ac 04 91 38 36 a8 53 4e 0c 7c 27 7d f9 20 92 e8  
T: 110    b: f0 df dd 6e    I: 32

Decim v2 Internal State at time 134  
a: 15 1e 4b 22 29 bb 74 14 ba c0 49 13 83 6a 85 34 e0 c7 c2 77 df 92 09 2e  
T: 010    b: f8 6f ee b7    I: 32

Decim v2 Internal State at time 135  
a: 91 51 e4 b2 22 9b b7 41 4b ac 04 91 38 36 a8 53 4e 0c 7c 27 7d f9 20 92  
T: 101    b: 7c 37 f7 5b    I: 32

Decim v2 Internal State at time 136  
a: 89 15 1e 4b 22 29 bb 74 14 ba c0 49 13 83 6a 85 34 e0 c7 c2 77 df 92 09  
T: 101    b: be 1b fb ad    I: 32

Decim v2 Internal State at time 137  
a: e8 91 51 e4 b2 22 9b b7 41 4b ac 04 91 38 36 a8 53 4e 0c 7c 27 7d f9 20  
T: 111    b: df 0d fd d6    I: 32

Decim v2 Internal State at time 138  
a: 5e 89 15 1e 4b 22 29 bb 74 14 ba c0 49 13 83 6a 85 34 e0 c7 c2 77 df 92  
T: 010    b: 6f 86 fe eb    I: 32

Decim v2 Internal State at time 139  
a: 15 e8 91 51 e4 b2 22 9b b7 41 4b ac 04 91 38 36 a8 53 4e 0c 7c 27 7d f9  
T: 101    b: 37 c3 7f 75    I: 32

Decim v2 Internal State at time 140  
a: a1 5e 89 15 1e 4b 22 29 bb 74 14 ba c0 49 13 83 6a 85 34 e0 c7 c2 77 df  
T: 000    b: 9b e1 bf ba    I: 32

Decim v2 Internal State at time 141  
a: 7a 15 e8 91 51 e4 b2 22 9b b7 41 4b ac 04 91 38 36 a8 53 4e 0c 7c 27 7d  
T: 010    b: 4d f0 df dd    I: 32

Decim v2 Internal State at time 142  
a: d7 a1 5e 89 15 1e 4b 22 29 bb 74 14 ba c0 49 13 83 6a 85 34 e0 c7 c2 77  
T: 111    b: 26 f8 6f ee    I: 32

Decim v2 Internal State at time 143  
a: dd 7a 15 e8 91 51 e4 b2 22 9b b7 41 4b ac 04 91 38 36 a8 53 4e 0c 7c 27

T: 110 b: 13 7c 37 f7 I: 32

Decim v2 Internal State at time 144

a: 2d d7 a1 5e 89 15 1e 4b 22 29 bb 74 14 ba c0 49 13 83 6a 85 34 e0 c7 c2

T: 100 b: 89 be 1b fb I: 32

Decim v2 Internal State at time 145

a: a2 dd 7a 15 e8 91 51 e4 b2 22 9b b7 41 4b ac 04 91 38 36 a8 53 4e 0c 7c

T: 110 b: c4 df 0d fd I: 32

Decim v2 Internal State at time 146

a: ea 2d d7 a1 5e 89 15 1e 4b 22 29 bb 74 14 ba c0 49 13 83 6a 85 34 e0 c7

T: 110 b: e2 6f 86 fe I: 32

Decim v2 Internal State at time 147

a: 1e a2 dd 7a 15 e8 91 51 e4 b2 22 9b b7 41 4b ac 04 91 38 36 a8 53 4e 0c

T: 111 b: 71 37 c3 7f I: 32

Decim v2 Internal State at time 148

a: 11 ea 2d d7 a1 5e 89 15 1e 4b 22 29 bb 74 14 ba c0 49 13 83 6a 85 34 e0

T: 000 b: 38 9b e1 bf I: 32

Decim v2 Internal State at time 149

a: 11 1e a2 dd 7a 15 e8 91 51 e4 b2 22 9b b7 41 4b ac 04 91 38 36 a8 53 4e

T: 010 b: 1c 4d f0 df I: 32

Decim v2 Internal State at time 150

a: 31 11 ea 2d d7 a1 5e 89 15 1e 4b 22 29 bb 74 14 ba c0 49 13 83 6a 85 34

T: 101 b: 0e 26 f8 6f I: 31

Decim v2 Internal State at time 151

a: 73 11 1e a2 dd 7a 15 e8 91 51 e4 b2 22 9b b7 41 4b ac 04 91 38 36 a8 53

T: 100 b: c7 13 7c 37 I: 32

Decim v2 Internal State at time 152

a: a7 31 11 ea 2d d7 a1 5e 89 15 1e 4b 22 29 bb 74 14 ba c0 49 13 83 6a 85

T: 010 b: e3 89 be 1b I: 32

Decim v2 Internal State at time 153

a: fa 73 11 1e a2 dd 7a 15 e8 91 51 e4 b2 22 9b b7 41 4b ac 04 91 38 36 a8

T: 100 b: 71 c4 df 0d I: 32

Decim v2 Internal State at time 154

a: df a7 31 11 ea 2d d7 a1 5e 89 15 1e 4b 22 29 bb 74 14 ba c0 49 13 83 6a

T: 100 b: b8 e2 6f 86 I: 32

Decim v2 Internal State at time 155

a: ad fa 73 11 1e a2 dd 7a 15 e8 91 51 e4 b2 22 9b b7 41 4b ac 04 91 38 36

T: 000 b: dc 71 37 c3 I: 32

Decim v2 Internal State at time 156

a: 1a df a7 31 11 ea 2d d7 a1 5e 89 15 1e 4b 22 29 bb 74 14 ba c0 49 13 83

T: 100 b: 6e 38 9b e1 I: 32

Decim v2 Internal State at time 157

a: 41 ad fa 73 11 1e a2 dd 7a 15 e8 91 51 e4 b2 22 9b b7 41 4b ac 04 91 38

T: 101 b: b7 1c 4d f0 I: 32

Decim v2 Internal State at time 158

a: 84 1a df a7 31 11 ea 2d d7 a1 5e 89 15 1e 4b 22 29 bb 74 14 ba c0 49 13

T: 100 b: db 8e 26 f8 I: 32

Decim v2 Internal State at time 159

a: a8 41 ad fa 73 11 1e a2 dd 7a 15 e8 91 51 e4 b2 22 9b b7 41 4b ac 04 91

T: 111 b: ed c7 13 7c I: 32

Decim v2 Internal State at time 160

a: 6a 84 1a df a7 31 11 ea 2d d7 a1 5e 89 15 1e 4b 22 29 bb 74 14 ba c0 49

T: 101 b: 76 e3 89 be I: 32

Decim v2 Internal State at time 161

a: a6 a8 41 ad fa 73 11 1e a2 dd 7a 15 e8 91 51 e4 b2 22 9b b7 41 4b ac 04

T: 000 b: bb 71 c4 df I: 32

Decim v2 Internal State at time 162

a: 0a 6a 84 1a df a7 31 11 ea 2d d7 a1 5e 89 15 1e 4b 22 29 bb 74 14 ba c0

# BS ISO/IEC 18033-4:2011+A1:2020

## ISO/IEC 18033-4:2011+A1:2020(E)

T: 111 b: dd b8 e2 6f I: 31

Decim v2 Internal State at time 163

a: 60 a6 a8 41 ad fa 73 11 1e a2 dd 7a 15 e8 91 51 e4 b2 22 9b b7 41 4b ac

T: 010 b: ae dc 71 37 I: 31

Decim v2 Internal State at time 164

a: 26 0a 6a 84 1a df a7 31 11 ea 2d d7 a1 5e 89 15 1e 4b 22 29 bb 74 14 ba

T: 100 b: 97 6e 38 9b I: 31

Decim v2 Internal State at time 165

a: f2 60 a6 a8 41 ad fa 73 11 1e a2 dd 7a 15 e8 91 51 e4 b2 22 9b b7 41 4b

T: 110 b: cb b7 1c 4d I: 31

Decim v2 Internal State at time 166

a: df 26 0a 6a 84 1a df a7 31 11 ea 2d d7 a1 5e 89 15 1e 4b 22 29 bb 74 14

T: 101 b: a5 db 8e 26 I: 31

Decim v2 Internal State at time 167

a: 8d f2 60 a6 a8 41 ad fa 73 11 1e a2 dd 7a 15 e8 91 51 e4 b2 22 9b b7 41

T: 100 b: d2 ed c7 13 I: 31

Decim v2 Internal State at time 168

a: 68 df 26 0a 6a 84 1a df a7 31 11 ea 2d d7 a1 5e 89 15 1e 4b 22 29 bb 74

T: 010 b: 29 76 e3 89 I: 32

Decim v2 Internal State at time 169

a: 26 8d f2 60 a6 a8 41 ad fa 73 11 1e a2 dd 7a 15 e8 91 51 e4 b2 22 9b b7

T: 101 b: 14 bb 71 c4 I: 31

Decim v2 Internal State at time 170

a: f2 68 df 26 0a 6a 84 1a df a7 31 11 ea 2d d7 a1 5e 89 15 1e 4b 22 29 bb

T: 010 b: 4a 5d b8 e2 I: 32

Decim v2 Internal State at time 171

a: af 26 8d f2 60 a6 a8 41 ad fa 73 11 1e a2 dd 7a 15 e8 91 51 e4 b2 22 9b

T: 101 b: 25 2e dc 71 I: 32

Decim v2 Internal State at time 172

a: aa f2 68 df 26 0a 6a 84 1a df a7 31 11 ea 2d d7 a1 5e 89 15 1e 4b 22 29

T: 000 b: 92 97 6e 38 I: 32

Decim v2 Internal State at time 173

a: aa af 26 8d f2 60 a6 a8 41 ad fa 73 11 1e a2 dd 7a 15 e8 91 51 e4 b2 22

T: 010 b: 49 4b b7 1c I: 32

Decim v2 Internal State at time 174

a: 3a aa f2 68 df 26 0a 6a 84 1a df a7 31 11 ea 2d d7 a1 5e 89 15 1e 4b 22

T: 111 b: a4 a5 db 8e I: 32

Decim v2 Internal State at time 175

a: a3 aa af 26 8d f2 60 a6 a8 41 ad fa 73 11 1e a2 dd 7a 15 e8 91 51 e4 b2

T: 111 b: 52 52 ed c7 I: 32

Decim v2 Internal State at time 176

a: 6a 3a aa f2 68 df 26 0a 6a 84 1a df a7 31 11 ea 2d d7 a1 5e 89 15 1e 4b

T: 010 b: 29 29 76 e3 I: 32

Decim v2 Internal State at time 177

a: a6 a3 aa af 26 8d f2 60 a6 a8 41 ad fa 73 11 1e a2 dd 7a 15 e8 91 51 e4

T: 010 b: 94 94 bb 71 I: 32

Decim v2 Internal State at time 178

a: 4a 6a 3a aa f2 68 df 26 0a 6a 84 1a df a7 31 11 ea 2d d7 a1 5e 89 15 1e

T: 100 b: ca 4a 5d b8 I: 32

Decim v2 Internal State at time 179

a: 14 a6 a3 aa af 26 8d f2 60 a6 a8 41 ad fa 73 11 1e a2 dd 7a 15 e8 91 51

T: 010 b: e5 25 2e dc I: 32

Decim v2 Internal State at time 180

a: 41 4a 6a 3a aa f2 68 df 26 0a 6a 84 1a df a7 31 11 ea 2d d7 a1 5e 89 15

T: 101 b: f2 92 97 6e I: 31

Decim v2 Internal State at time 181

a: 44 14 a6 a3 aa af 26 8d f2 60 a6 a8 41 ad fa 73 11 1e a2 dd 7a 15 e8 91

T: 000 b: f9 49 4b b7 I: 31

Decim v2 Internal State at time 182

a: 44 41 4a 6a 3a aa f2 68 df 26 0a 6a 84 1a df a7 31 11 ea 2d d7 a1 5e 89  
T: 100 b: fc a4 a5 db I: 31

Decim v2 Internal State at time 183

a: 94 44 14 a6 a3 aa af 26 8d f2 60 a6 a8 41 ad fa 73 11 1e a2 dd 7a 15 e8  
T: 100 b: 3e 52 52 ed I: 32

Decim v2 Internal State at time 184

a: 99 44 41 4a 6a 3a aa f2 68 df 26 0a 6a 84 1a df a7 31 11 ea 2d d7 a1 5e  
T: 010 b: 1f 29 29 76 I: 32

Decim v2 Internal State at time 185

a: 59 94 44 14 a6 a3 aa af 26 8d f2 60 a6 a8 41 ad fa 73 11 1e a2 dd 7a 15  
T: 010 b: 8f 94 94 bb I: 32

Decim v2 Internal State at time 186

a: c5 99 44 41 4a 6a 3a aa f2 68 df 26 0a 6a 84 1a df a7 31 11 ea 2d d7 a1  
T: 111 b: c7 ca 4a 5d I: 32

Decim v2 Internal State at time 187

a: 4c 59 94 44 14 a6 a3 aa af 26 8d f2 60 a6 a8 41 ad fa 73 11 1e a2 dd 7a  
T: 111 b: 63 e5 25 2e I: 32

Decim v2 Internal State at time 188

a: 34 c5 99 44 41 4a 6a 3a aa f2 68 df 26 0a 6a 84 1a df a7 31 11 ea 2d d7  
T: 101 b: 31 f2 92 97 I: 32

Decim v2 Internal State at time 189

a: b3 4c 59 94 44 14 a6 a3 aa af 26 8d f2 60 a6 a8 41 ad fa 73 11 1e a2 dd  
T: 100 b: 98 f9 49 4b I: 32

Decim v2 Internal State at time 190

a: ab 34 c5 99 44 41 4a 6a 3a aa f2 68 df 26 0a 6a 84 1a df a7 31 11 ea 2d  
T: 101 b: 4c 7c a4 a5 I: 32

Decim v2 Internal State at time 191

a: 4a b3 4c 59 94 44 14 a6 a3 aa af 26 8d f2 60 a6 a8 41 ad fa 73 11 1e a2  
T: 101 b: a6 3e 52 52 I: 32

Decim v2 Internal State at time 192

a: 44 ab 34 c5 99 44 41 4a 6a 3a aa f2 68 df 26 0a 6a 84 1a df a7 31 11 ea  
T: 010 b: d3 1f 29 29 I: 32

Decim v2 Internal State at time 192 (Binary notation)

a: 01000100 10101011 00110100 11000101 10011001 01000100 01000001 01001010  
01101010 00111010 10101010 11110010 01101000 11011111 00100110 00001010  
01101010 10000100 00011010 11011111 10100111 00110001 00010001 11101010  
T: 010 b: 11010011 00011111 00101001 00101001 I: 32

## C.6 Example for KCipher-2 (K2)

### C.6.1 Key, initialization vector, and keystream triplets

K = (K0,K1,K2,K3)=(0x00000000, 0x00000000, 0x00000000, 0x00000000)

IV = (IV0, IV1, IV2, IV3)=(0x00000000, 0x00000000, 0x00000000, 0x00000000)

Keystream[0] = 0xF871EBEF945B7272

Keystream[1] = 0xE40C04941DFF0537

Keystream[2] = 0x0B981A59FBC8AC57

Keystream[3] = 0x566D3B02C179DBB4

Keystream[4] = 0x3B46F1F033554C72

Keystream[5] = 0x5DE68BCC9872858F

Keystream[6] = 0x575496024062F0E9

Keystream[7] = 0xF932C998226DB6BA

K = (K0,K1,K2,K3)=(0x0F1E2D3C, 0x4B5A6978, 0x8796A5B4, 0xC3D2E1F0)

IV = (IV0, IV1, IV2, IV3)=(0xF0E0D0C0, 0xB0A09080, 0x70605040, 0x30201000)

```
Keystream[0] = 0x9FB6B580A6A5E7AF
Keystream[1] = 0xD1989DC6A77D5E28
Keystream[2] = 0x4EFCC8CB7BCFB32B
Keystream[3] = 0xF69297F5DD974CE8
Keystream[4] = 0xFBD9139C7A71F41A
Keystream[5] = 0x61382C76D3D2F6CA
Keystream[6] = 0xD5265037659CF838
Keystream[7] = 0x774121C26F6474F3
```

```
K = (K0,K1,K2,K3)=(0xAC2F75C0, 0x43FBC367, 0x09D315F2, 0x245746D8)
IV = (IV0, IV1, IV2, IV3)=(0xF6B29A58, 0x45CCCD8C, 0x6229393A, 0x7A4842C1)
Keystream[0] = 0xDA38138B32864E05
Keystream[1] = 0x24B8B90944E5117A
Keystream[2] = 0xC3E883DCFA22C458
Keystream[3] = 0x1F2C9DDFE98DC5DE
Keystream[4] = 0x33B2FC05064C6FEF
Keystream[5] = 0xA9A3D3ED31660DFF
Keystream[6] = 0xF7DE1857E224E70F
Keystream[7] = 0x4EFE5C36CEB974AC
```

## C.6.2 Sample internal states

```
K = (K0,K1,K2,K3)=(0x80000000, 0x00000000, 0x00000000, 0x00000000)
IV = (IV0, IV1, IV2, IV3)=(0x00000004, 0x00000003, 0x00000002, 0x00000001)
Keystream = (0x9B753FAA, 0x404A0EF5, 0x52919406, 0x18177FDD, 0xA419D11E, 0x47481D1B, 0x2DD49337, 0x640BDEC9)
```

```
K2 internal state at time -24
(A0, A1, A2, A3, A4) = (0xE2636363, 0x00000000, 0x00000000, 0x00000000, 0x80000000)
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =
(0xBEFBFB5E, 0x5C98983D, 0x00000004, 0x00000003, 0xBEFBFB5E, 0x5C98983D, 0x00000002, 0x00000001,
0xE2636363, 0xE2636363, 0xE2636363)
(R1, L1, R2, L2) = (0x00000000, 0x00000000, 0x00000000, 0x00000000)
```

```
K2 internal state at time -23
(A0, A1, A2, A3, A4) = (0x00000000, 0x00000000, 0x00000000, 0x80000000, 0x1F84F87D)
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =
(0x5C98983D, 0x00000004, 0x00000003, 0xBEFBFB5E, 0x5C98983D, 0x00000002, 0x00000001, 0xE2636363,
0xE2636363, 0xE2636363, 0x08CE4DDD)
(R1, L1, R2, L2) = (0x3D5E9898, 0xAFA0F900, 0x63636363, 0x63636363)
```

```
K2 internal state at time -22
(A0, A1, A2, A3, A4) = (0x00000000, 0x00000000, 0x80000000, 0x1F84F87D, 0x1D219B45)
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =
(0x00000004, 0x00000003, 0xBEFBFB5E, 0x5C98983D, 0x00000002, 0x00000001, 0xE2636363, 0xE2636363,
0xE2636363, 0x08CE4DDD, 0xD448E338)
(R1, L1, R2, L2) = (0x1BC16E6E, 0x2BE9E7CD, 0x9AD90539, 0x2EAA08EF)
```

```
K2 internal state at time -21
(A0, A1, A2, A3, A4) = (0x00000000, 0x80000000, 0x1F84F87D, 0x1D219B45, 0x83BD086B)
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =
(0x00000003, 0xBEFBFB5E, 0x5C98983D, 0x00000002, 0x00000001, 0xE2636363, 0xE2636363, 0xE2636363,
0x08CE4DDD, 0xD448E338, 0xD791AF96)
(R1, L1, R2, L2) = (0xFF2C7778, 0x0830D3EF, 0x181A9D48, 0xAF1D5D29)
```

```
K2 internal state at time -20
(A0, A1, A2, A3, A4) = (0x80000000, 0x1F84F87D, 0x1D219B45, 0x83BD086B, 0x79AA791D)
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =
```

(0xBEFBFB5E, 0x5C98983D, 0x00000002, 0x00000001, 0xE2636363, 0xE2636363, 0xE2636363, 0x08CE4DDD, 0xD448E338, 0xD791AF96, 0xDAC909B0)  
(R1, L1, R2, L2) = (0xD3AF4401, 0xF0D6D79B, 0x7791C800, 0x78E12F3B)

K2 internal state at time -19

(A0, A1, A2, A3, A4) = (0x1F84F87D, 0x1D219B45, 0x83BD086B, 0x79AA791D, 0x54BF0EB7)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0x5C98983D, 0x00000002, 0x00000001, 0xE2636363, 0xE2636363, 0xE2636363, 0x08CE4DDD, 0xD448E338, 0xD791AF96, 0xDAC909B0, 0xFFFA4197)  
(R1, L1, R2, L2) = (0xB9569748, 0x8C9BF8C8, 0x2A3FA7CA, 0xC7628540)

K2 internal state at time -18

(A0, A1, A2, A3, A4) = (0x1D219B45, 0x83BD086B, 0x79AA791D, 0x54BF0EB7, 0xDA9BEC10)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0x00000002, 0x00000001, 0xE2636363, 0xE2636363, 0xE2636363, 0x08CE4DDD, 0xD448E338, 0xD791AF96, 0xDAC909B0, 0xFFFA4197, 0xF546678B)  
(R1, L1, R2, L2) = (0x13228CE3, 0xB5FE6E22, 0x6359C7C0, 0xBE2D3278)

K2 internal state at time -17

(A0, A1, A2, A3, A4) = (0x83BD086B, 0x79AA791D, 0x54BF0EB7, 0xDA9BEC10, 0x3C2C0747)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0x00000001, 0xE2636363, 0xE2636363, 0xE2636363, 0x08CE4DDD, 0xD448E338, 0xD791AF96, 0xDAC909B0, 0xFFFA4197, 0xF546678B, 0x152D87ED)  
(R1, L1, R2, L2) = (0x838EA823, 0x3D913F50, 0x3ECF0A60, 0x3B05B5E9)

K2 internal state at time -16

(A0, A1, A2, A3, A4) = (0x79AA791D, 0x54BF0EB7, 0xDA9BEC10, 0x3C2C0747, 0xB46C6DA5)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0xE2636363, 0xE2636363, 0xE2636363, 0x08CE4DDD, 0xD448E338, 0xD791AF96, 0xDAC909B0, 0xFFFA4197, 0xF546678B, 0x152D87ED, 0x2417868F)  
(R1, L1, R2, L2) = (0x2868BEE5, 0x060AB10E, 0x72F97EE4, 0x4F56069F)

K2 internal state at time -15

(A0, A1, A2, A3, A4) = (0x54BF0EB7, 0xDA9BEC10, 0x3C2C0747, 0xB46C6DA5, 0x8743F560)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0xE2636363, 0xE2636363, 0x08CE4DDD, 0xD448E338, 0xD791AF96, 0xDAC909B0, 0xFFFA4197, 0xF546678B, 0x152D87ED, 0x2417868F, 0xA6305225)  
(R1, L1, R2, L2) = (0xDF7B3A4C, 0x6295074A, 0xF3A16531, 0x971CE606)

K2 internal state at time -14

(A0, A1, A2, A3, A4) = (0xDA9BEC10, 0x3C2C0747, 0xB46C6DA5, 0x8743F560, 0xEDBBC959)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0xE2636363, 0x08CE4DDD, 0xD448E338, 0xD791AF96, 0xDAC909B0, 0xFFFA4197, 0xF546678B, 0x152D87ED, 0x2417868F, 0xA6305225, 0x405CEA50)  
(R1, L1, R2, L2) = (0xF96F4856, 0x5680685E, 0xFD52CF76, 0xC1A29363)

K2 internal state at time -13

(A0, A1, A2, A3, A4) = (0x3C2C0747, 0xB46C6DA5, 0x8743F560, 0xEDBBC959, 0x3FD2F9E0)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0x08CE4DDD, 0xD448E338, 0xD791AF96, 0xDAC909B0, 0xFFFA4197, 0xF546678B, 0x152D87ED, 0x2417868F, 0xA6305225, 0x405CEA50, 0x726C92BC)  
(R1, L1, R2, L2) = (0x2AF82CD3, 0x5A241664, 0x1B186FBE, 0x19542F03)

K2 internal state at time -12

(A0, A1, A2, A3, A4) = (0xB46C6DA5, 0x8743F560, 0xEDBBC959, 0x3FD2F9E0, 0xC0535EEC)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0xD448E338, 0xD791AF96, 0xDAC909B0, 0xFFFA4197, 0xF546678B, 0x152D87ED, 0x2417868F, 0xA6305225, 0x405CEA50, 0x726C92BC, 0xD047818E)  
(R1, L1, R2, L2) = (0xBDF4D6A5, 0x5B5798C6, 0x4BA1A2FB, 0xD3B129C7)

# BS ISO/IEC 18033-4:2011+A1:2020

## ISO/IEC 18033-4:2011+A1:2020(E)

K2 internal state at time -11

(A0, A1, A2, A3, A4) = (0x8743F560, 0xEDBBC959, 0x3FD2F9E0, 0xC0535EEC, 0xB41BFCC9)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0xD791AF96, 0xDAC909B0, 0xFFFFA4197, 0xF546678B, 0x152D87ED, 0x2417868F, 0xA6305225, 0x405CEA50, 0x726C92BC, 0xD047818E, 0xA458FE0A)  
(R1, L1, R2, L2) = (0x5A348B92, 0x221535B3, 0xB71B51C8, 0xA80FECDB)

K2 internal state at time -10

(A0, A1, A2, A3, A4) = (0xEDBBC959, 0x3FD2F9E0, 0xC0535EEC, 0xB41BFCC9, 0x2C967031)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0xDAC909B0, 0xFFFFA4197, 0xF546678B, 0x152D87ED, 0x2417868F, 0xA6305225, 0x405CEA50, 0x726C92BC, 0xD047818E, 0xA458FE0A, 0xED0E4419)  
(R1, L1, R2, L2) = (0xB70F8DB4, 0x959902F7, 0x939BA37F, 0x45E722B1)

K2 internal state at time -9

(A0, A1, A2, A3, A4) = (0x3FD2F9E0, 0xC0535EEC, 0xB41BFCC9, 0x2C967031, 0xC54BBBD6)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0xFFFFA4197, 0xF546678B, 0x152D87ED, 0x2417868F, 0xA6305225, 0x405CEA50, 0x726C92BC, 0xD047818E, 0xA458FE0A, 0xED0E4419, 0x67C6A2D5)  
(R1, L1, R2, L2) = (0x8EDD0383, 0x277464BD, 0xEEDC0439, 0x75A6858D)

K2 internal state at time -8

(A0, A1, A2, A3, A4) = (0xC0535EEC, 0xB41BFCC9, 0x2C967031, 0xC54BBBD6, 0x75146287)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0xF546678B, 0x152D87ED, 0x2417868F, 0xA6305225, 0x405CEA50, 0x726C92BC, 0xD047818E, 0xA458FE0A, 0xED0E4419, 0x67C6A2D5, 0xB6E631E9)  
(R1, L1, R2, L2) = (0x731DA313, 0xF37012AC, 0xA7255B96, 0xDC499D6F)

K2 internal state at time -7

(A0, A1, A2, A3, A4) = (0xB41BFCC9, 0x2C967031, 0xC54BBBD6, 0x75146287, 0x83B8D1B2)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0x152D87ED, 0x2417868F, 0xA6305225, 0x405CEA50, 0x726C92BC, 0xD047818E, 0xA458FE0A, 0xED0E4419, 0x67C6A2D5, 0xB6E631E9, 0xA64D395C)  
(R1, L1, R2, L2) = (0xD8B057FC, 0xC761F332, 0x2CAE11CF, 0x2AEDE625)

K2 internal state at time -6

(A0, A1, A2, A3, A4) = (0x2C967031, 0xC54BBBD6, 0x75146287, 0x83B8D1B2, 0xF1D21A26)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0x2417868F, 0xA6305225, 0x405CEA50, 0x726C92BC, 0xD047818E, 0xA458FE0A, 0xED0E4419, 0x67C6A2D5, 0xB6E631E9, 0xA64D395C, 0x7B59587F)  
(R1, L1, R2, L2) = (0xE885CA11, 0x015D89D8, 0xB59D5510, 0x10BAD578)

K2 internal state at time -5

(A0, A1, A2, A3, A4) = (0xC54BBBD6, 0x75146287, 0x83B8D1B2, 0xF1D21A26, 0xD3884C64)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0xA6305225, 0x405CEA50, 0x726C92BC, 0xD047818E, 0xA458FE0A, 0xED0E4419, 0x67C6A2D5, 0xB6E631E9, 0xA64D395C, 0x7B59587F, 0xAFA3772A)  
(R1, L1, R2, L2) = (0x0A63FCEF, 0xB78DD0E9, 0x5375538F, 0xB0DA9C00)

K2 internal state at time -4

(A0, A1, A2, A3, A4) = (0x75146287, 0x83B8D1B2, 0xF1D21A26, 0xD3884C64, 0x3EC047BA)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0x405CEA50, 0x726C92BC, 0xD047818E, 0xA458FE0A, 0xED0E4419, 0x67C6A2D5, 0xB6E631E9, 0xA64D395C, 0x7B59587F, 0xAFA3772A, 0x83AE3259)  
(R1, L1, R2, L2) = (0xF36398D7, 0xA38AB94C, 0xFF2BD5F2, 0x4634B058)

K2 internal state at time -3

(A0, A1, A2, A3, A4) = (0x83B8D1B2, 0xF1D21A26, 0xD3884C64, 0x3EC047BA, 0x89EF93D3)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =

(0x726C92BC, 0xD047818E, 0xA458FE0A, 0xED0E4419, 0x67C6A2D5, 0xB6E631E9, 0xA64D395C, 0x7B59587F, 0xAFA3772A, 0x83AE3259, 0x479E0649)  
(R1, L1, R2, L2) = (0x4F2CD5EE, 0xD8DB21E5, 0xB5B29920, 0x479D0DDC)

K2 internal state at time -2

(A0, A1, A2, A3, A4) = (0xF1D21A26, 0xD3884C64, 0x3EC047BA, 0x89EF93D3, 0x3A5FD7EB)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0xD047818E, 0xA458FE0A, 0xED0E4419, 0x67C6A2D5, 0xB6E631E9, 0xA64D395C, 0x7B59587F, 0xAFA3772A, 0x83AE3259, 0x479E0649, 0xBAC82EA9)  
(R1, L1, R2, L2) = (0x25C35580, 0x36845CF8, 0x195E39A0, 0xF6EE896D)

K2 internal state at time -1

(A0, A1, A2, A3, A4) = (0xD3884C64, 0x3EC047BA, 0x89EF93D3, 0x3A5FD7EB, 0xC74844F8)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0xA458FE0A, 0xED0E4419, 0x67C6A2D5, 0xB6E631E9, 0xA64D395C, 0x7B59587F, 0xAFA3772A, 0x83AE3259, 0x479E0649, 0xBAC82EA9, 0x3AB6E5CC)  
(R1, L1, R2, L2) = (0xBA38B6DF, 0x767E04BF, 0xE02C638F, 0xDCBA3106)

K2 internal state at time 0

(A0, A1, A2, A3, A4) = (0x3EC047BA, 0x89EF93D3, 0x3A5FD7EB, 0xC74844F8, 0xECB10CC9)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0xED0E4419, 0x67C6A2D5, 0xB6E631E9, 0xA64D395C, 0x7B59587F, 0xAFA3772A, 0x83AE3259, 0x479E0649, 0xBAC82EA9, 0x3AB6E5CC, 0xFBE9F17C)  
(R1, L1, R2, L2) = (0x5C134123, 0xC08CCB42, 0x03D9FF06, 0x694FC1D6)

K2 internal state at time 1

(A0, A1, A2, A3, A4) = (0x89EF93D3, 0x3A5FD7EB, 0xC74844F8, 0xECB10CC9, 0x9F3C9CD1)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0x67C6A2D5, 0xB6E631E9, 0xA64D395C, 0x7B59587F, 0xAFA3772A, 0x83AE3259, 0x479E0649, 0xBAC82EA9, 0x3AB6E5CC, 0xFBE9F17C, 0xFE14A575)  
(R1, L1, R2, L2) = (0xEF637AB6, 0x853CADC9, 0x0081F6E5, 0x602E04A7)

K2 internal state at time 2

(A0, A1, A2, A3, A4) = (0x3A5FD7EB, 0xC74844F8, 0xECB10CC9, 0x9F3C9CD1, 0x07E0A49D)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0xB6E631E9, 0xA64D395C, 0x7B59587F, 0xAFA3772A, 0x83AE3259, 0x479E0649, 0xBAC82EA9, 0x3AB6E5CC, 0xFBE9F17C, 0xFE14A575, 0x3ABC9602)  
(R1, L1, R2, L2) = (0x4C41E330, 0xEFA5F58E, 0x560328CD, 0x37275D79)

K2 internal state at time 3

(A0, A1, A2, A3, A4) = (0xC74844F8, 0xECB10CC9, 0x9F3C9CD1, 0x07E0A49D, 0xC5F86290)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0xA64D395C, 0x7B59587F, 0xAFA3772A, 0x83AE3259, 0x479E0649, 0xBAC82EA9, 0x3AB6E5CC, 0xFBE9F17C, 0xFE14A575, 0x3ABC9602, 0x04FD08B2)  
(R1, L1, R2, L2) = (0xD33376B4, 0x991AF343, 0xCC739191, 0x6E891BDA)

K2 internal state at time 4

(A0, A1, A2, A3, A4) = (0xECB10CC9, 0x9F3C9CD1, 0x07E0A49D, 0xC5F86290, 0x4453180A)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0x7B59587F, 0xAFA3772A, 0x83AE3259, 0x479E0649, 0xBAC82EA9, 0x3AB6E5CC, 0xFBE9F17C, 0xFE14A575, 0x3ABC9602, 0x04FD08B2, 0xDE821E38)  
(R1, L1, R2, L2) = (0x8AFC6CE9, 0x1CD53B55, 0xBB82C5EC, 0x4661136F)

K2 internal state at time 5

(A0, A1, A2, A3, A4) = (0x9F3C9CD1, 0x07E0A49D, 0xC5F86290, 0x4453180A, 0xFD98883F)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0xAFA3772A, 0x83AE3259, 0x479E0649, 0xBAC82EA9, 0x3AB6E5CC, 0xFBE9F17C, 0xFE14A575, 0x3ABC9602, 0x04FD08B2, 0xDE821E38, 0xD30B1637)  
(R1, L1, R2, L2) = (0xA51F85B5, 0x676A1660, 0x3EB70B02, 0xDDA7BA41)



K2 internal state at time 6  
(A0, A1, A2, A3, A4) = (0x07E0A49D, 0xC5F86290, 0x4453180A, 0xFD98883F, 0x7FBD0061)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0x83AE3259, 0x479E0649, 0xBAC82EA9, 0x3AB6E5CC, 0xFBE9F17C, 0xFE14A575, 0x3ABC9602, 0x04FD08B2, 0xDE821E38, 0xD30B1637, 0x1AE1594D)  
(R1, L1, R2, L2) = (0xCF29E514, 0x41BE7A08, 0x3FD3BDD5, 0x3F07DDF5)

K2 internal state at time 7  
(A0, A1, A2, A3, A4) = (0xC5F86290, 0x4453180A, 0xFD98883F, 0x7FBD0061, 0x9904D579)  
(B0, B1, B2, B3, B4, B5, B6, B7, B8, B9, B10) =  
(0x479E0649, 0xBAC82EA9, 0x3AB6E5CC, 0xFBE9F17C, 0xFE14A575, 0x3ABC9602, 0x04FD08B2, 0xDE821E38, 0xD30B1637, 0x1AE1594D, 0x38FE9448)  
(R1, L1, R2, L2) = (0x6C079D38, 0xB4614FAA, 0x66F72DB0, 0x3933F538)

**A1 C.7 Example for ZUC**

**C.7.1 Key, initialization vector, and keystream triplets**

k = 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00  
iv= 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00  
z = 27 be de 74 01 80 82 da 87 d4 e5 b6 9f 18 bf 66 32 07 0e 0f 39 b7 b6 92 b4 67 3e dc 31  
84 a4 8e

k = ff ff ff ff ff ff ff ff ff ff ff ff ff ff ff  
iv= ff ff ff ff ff ff ff ff ff ff ff ff ff ff ff  
z = 06 57 cf a0 70 96 39 8b 73 4b 6c b4 88 3e ed f4 25 7a 76 eb 97 59 52 08 d8 84 ad cd b1  
cb ff b8

k = 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00  
iv= ff ff ff ff ff ff ff ff ff ff ff ff ff ff ff  
z = 58 fb 51 5e 39 08 74 6d 7a 91 f2 34 49 4e d8 c8 51 2d 61 eb 69 6c 14 b8 cd 2d 3b fe 69  
4f e8 1d

k = 3d 4c 4b e9 6a 82 fd ae b5 8f 64 1d b1 7b 45 5b  
iv= 84 31 9a a8 de 69 15 ca 1f 6b da 6b fb d8 c7 66  
z = 14 f1 c2 72 32 79 c4 19 4b 8e a4 1d 0c c8 08 63 d2 80 62 e1 e7 1d 3d da e3 c4 d1 58 a7  
f0 67 ac

**C.7.2 Sample internal states**

k = 3d 4c 4b e9 6a 82 fd ae b5 8f 64 1d b1 7b 45 5b  
iv= 84 31 9a a8 de 69 15 ca 1f 6b da 6b fb d8 c7 66  
z = 14 f1 c2 72 32 79 c4 19 4b 8e a4 1d 0c c8 08 63 d2 80 62 e1 e7 1d 3d da e3 c4 d1 58 a7  
f0 67 ac

Internal state at time -33  
15:2dc7ac66 14:22f89ac7 13:3dbc4dd8 12:58de26fb 11:0e9af16b 10:326bc4da 09:47af136b  
08:5acd781f  
07:5709afca 06:7ef13515 05:4135e269 04:355789de 03:74935ea8 02:25e26b9a 01:2626bc31  
00:1ec4d784  
R1:00000000 R2:00000000

Internal state at time -32  
15:3c7b93c0 14:2dc7ac66 13:22f89ac7 12:3dbc4dd8 11:58de26fb 10:0e9af16b 09:326bc4da  
08:47af136b  
07:5acd781f 06:5709afca 05:7ef13515 04:4135e269 03:355789de 02:74935ea8 01:25e26b9a  
00:2626bc31  
R1:9c62829f R2:5df00831

Internal state at time -31  
15:41901ee9 14:3c7b93c0 13:2dc7ac66 12:22f89ac7 11:3dbc4dd8 10:58de26fb 09:0e9af16b  
08:326bc4da  
07:47af136b 06:5acd781f 05:5709afca 04:7ef13515 03:4135e269 02:355789de 01:74935ea8  
00:25e26b9a  
R1:3d533f3a R2:80ff1faf

Internal state at time -30  
15:411efa99 14:41901ee9 13:3c7b93c0 12:2dc7ac66 11:22f89ac7 10:3dbc4dd8 09:58de26fb  
08:0e9af16b  
07:326bc4da 06:47af136b 05:5acd781f 04:5709afca 03:7ef13515 02:4135e269 01:355789de  
00:74935ea8  
R1:2ca57e9d R2:d1db72f9

Internal state at time -29  
15:24b3f49f 14:411efa99 13:41901ee9 12:3c7b93c0 11:2dc7ac66 10:22f89ac7 09:3dbc4dd8  
08:58de26fb  
07:0e9af16b 06:326bc4da 05:47af136b 04:5acd781f 03:5709afca 02:7ef13515 01:4135e269  
00:355789de  
R1:0e8dc40f R2:60921a4f

Internal state at time -28  
15:74265785 14:24b3f49f 13:411efa99 12:41901ee9 11:3c7b93c0 10:2dc7ac66 09:22f89ac7  
08:3dbc4dd8  
07:58de26fb 06:0e9af16b 05:326bc4da 04:47af136b 03:5acd781f 02:5709afca 01:7ef13515  
00:4135e269  
R1:16c81467 R2:da8e7d8a

Internal state at time -27  
15:481c5b9d 14:74265785 13:24b3f49f 12:411efa99 11:41901ee9 10:3c7b93c0 09:2dc7ac66  
08:22f89ac7  
07:3dbc4dd8 06:58de26fb 05:0e9af16b 04:326bc4da 03:47af136b 02:5acd781f 01:5709afca  
00:7ef13515  
R1:50c9eaa4 R2:3c3b2dfd

Internal state at time -26  
15:4b7f87ed 14:481c5b9d 13:74265785 12:24b3f49f 11:411efa99 10:41901ee9 09:3c7b93c0  
08:2dc7ac66  
07:22f89ac7 06:3dbc4dd8 05:58de26fb 04:0e9af16b 03:326bc4da 02:47af136b 01:5acd781f  
00:5709afca  
R1:59857b80 R2:be0fbdc1

Internal state at time -25  
15:0e633ce7 14:4b7f87ed 13:481c5b9d 12:74265785 11:24b3f49f 10:411efa99 09:41901ee9  
08:3c7b93c0  
07:2dc7ac66 06:22f89ac7 05:3dbc4dd8 04:58de26fb 03:0e9af16b 02:326bc4da 01:47af136b  
00:5acd781f  
R1:9528f8ea R2:bcc7f7eb

Internal state at time -24  
15:643ae5a6 14:0e633ce7 13:4b7f87ed 12:481c5b9d 11:74265785 10:24b3f49f 09:411efa99  
08:41901ee9  
07:3c7b93c0 06:2dc7ac66 05:22f89ac7 04:3dbc4dd8 03:58de26fb 02:0e9af16b 01:326bc4da  
00:47af136b  
R1:c59d2932 R2:e1098a64

Internal state at time -23  
15:625ac5d7 14:643ae5a6 13:0e633ce7 12:4b7f87ed 11:481c5b9d 10:74265785 09:24b3f49f  
08:411efa99  
07:41901ee9 06:3c7b93c0 05:2dc7ac66 04:22f89ac7 03:3dbc4dd8 02:58de26fb 01:0e9af16b  
00:326bc4da  
R1:755ebae8 R2:3f9e6e86

Internal state at time -22  
15:10e10abb 14:625ac5d7 13:643ae5a6 12:0e633ce7 11:4b7f87ed 10:481c5b9d 09:74265785  
08:24b3f49f  
07:411efa99 06:41901ee9 05:3c7b93c0 04:2dc7ac66 03:22f89ac7 02:3dbc4dd8 01:58de26fb  
00:0e9af16b  
R1:d643d938 R2:d799a5a3

Internal state at time -21  
15:1bdc9fab 14:10e10abb 13:625ac5d7 12:643ae5a6 11:0e633ce7 10:4b7f87ed 09:481c5b9d  
08:74265785  
07:24b3f49f 06:411efa99 05:41901ee9 04:3c7b93c0 03:2dc7ac66 02:22f89ac7 01:3dbc4dd8  
00:58de26fb  
R1:1798b822 R2:92245168

Internal state at time -20  
15:2567b94a 14:1bdc9fab 13:10e10abb 12:625ac5d7 11:643ae5a6 10:0e633ce7 09:4b7f87ed  
08:481c5b9d  
07:74265785 06:24b3f49f 05:411efa99 04:41901ee9 03:3c7b93c0 02:2dc7ac66 01:22f89ac7  
00:3dbc4dd8  
R1:4e29d84e R2:61b91f59

Internal state at time -19  
15:2af10db2 14:2567b94a 13:1bdc9fab 12:10e10abb 11:625ac5d7 10:643ae5a6 09:0e633ce7  
08:4b7f87ed  
07:481c5b9d 06:74265785 05:24b3f49f 04:411efa99 03:41901ee9 02:3c7b93c0 01:2dc7ac66  
00:22f89ac7  
R1:5b486570 R2:d97ebf32

Internal state at time -18

# BS ISO/IEC 18033-4:2011+A1:2020

## ISO/IEC 18033-4:2011+A1:2020(E)

15:3448fcc0 14:2af10db2 13:2567b94a 12:1bdc9fab 11:10e10abb 10:625ac5d7 09:643ae5a6  
08:0e633ce7  
07:4b7f87ed 06:481c5b9d 05:74265785 04:24b3f49f 03:411efa99 02:41901ee9 01:3c7b93c0  
00:2dc7ac66  
R1:421fbdfa R2:effe033f

Internal state at time -17  
15:789c639c 14:3448fcc0 13:2af10db2 12:2567b94a 11:1bdc9fab 10:10e10abb 09:625ac5d7  
08:643ae5a6  
07:0e633ce7 06:4b7f87ed 05:481c5b9d 04:74265785 03:24b3f49f 02:411efa99 01:41901ee9  
00:3c7b93c0  
R1:bff08d37 R2:28e1d53c

Internal state at time -16  
15:10da5941 14:789c639c 13:3448fcc0 12:2af10db2 11:2567b94a 10:1bdc9fab 09:10e10abb  
08:625ac5d7  
07:643ae5a6 06:0e633ce7 05:4b7f87ed 04:481c5b9d 03:74265785 02:24b3f49f 01:411efa99  
00:41901ee9  
R1:8d36a012 R2:bc320a23

Internal state at time -15  
15:5b6acbf6 14:10da5941 13:789c639c 12:3448fcc0 11:2af10db2 10:2567b94a 09:1bdc9fab  
08:10e10abb  
07:625ac5d7 06:643ae5a6 05:0e633ce7 04:4b7f87ed 03:481c5b9d 02:74265785 01:24b3f49f  
00:411efa99  
R1:92b7231b R2:9ec667b9

Internal state at time -14  
15:17060ce1 14:5b6acbf6 13:10da5941 12:789c639c 11:3448fcc0 10:2af10db2 09:2567b94a  
08:1bdc9fab  
07:10e10abb 06:625ac5d7 05:643ae5a6 04:0e633ce7 03:4b7f87ed 02:481c5b9d 01:74265785  
00:24b3f49f  
R1:538a936d R2:c036bc48

Internal state at time -13  
15:35368174 14:17060ce1 13:5b6acbf6 12:10da5941 11:789c639c 10:3448fcc0 09:2af10db2  
08:2567b94a  
07:1bdc9fab 06:10e10abb 05:625ac5d7 04:643ae5a6 03:0e633ce7 02:4b7f87ed 01:481c5b9d  
00:74265785  
R1:1a29f0af R2:7e65408e

Internal state at time -12  
15:5cf4385a 14:35368174 13:17060ce1 12:5b6acbf6 11:10da5941 10:789c639c 09:3448fcc0  
08:2af10db2  
07:2567b94a 06:1bdc9fab 05:10e10abb 04:625ac5d7 03:643ae5a6 02:0e633ce7 01:4b7f87ed  
00:481c5b9d  
R1:5e4350bb R2:d0826c98

Internal state at time -11  
15:479943df 14:5cf4385a 13:35368174 12:17060ce1 11:5b6acbf6 10:10da5941 09:789c639c  
08:3448fcc0  
07:2af10db2 06:2567b94a 05:1bdc9fab 04:10e10abb 03:625ac5d7 02:643ae5a6 01:0e633ce7  
00:4b7f87ed  
R1:5b20edbc R2:f327d61e

Internal state at time -10  
15:2753bab2 14:479943df 13:5cf4385a 12:35368174 11:17060ce1 10:5b6acbf6 09:10da5941  
08:789c639c  
07:3448fcc0 06:2af10db2 05:2567b94a 04:1bdc9fab 03:10e10abb 02:625ac5d7 01:643ae5a6  
00:0e633ce7  
R1:2d9c405e R2:11418c75

Internal state at time -9  
15:73775d6a 14:2753bab2 13:479943df 12:5cf4385a 11:35368174 10:17060ce1 09:5b6acbf6  
08:10da5941  
07:789c639c 06:3448fcc0 05:2af10db2 04:2567b94a 03:1bdc9fab 02:10e10abb 01:625ac5d7  
00:643ae5a6  
R1:972d2a15 R2:08a36a3b

Internal state at time -8  
15:43930a37 14:73775d6a 13:2753bab2 12:479943df 11:5cf4385a 10:35368174 09:17060ce1

08:5b6acbf6  
07:10da5941 06:789c639c 05:3448fcc0 04:2af10db2 03:2567b94a 02:1bdc9fab 01:10e10abb  
00:625ac5d7  
R1:8264eff1 R2:677f5747

Internal state at time -7  
15:77b4af31 14:43930a37 13:73775d6a 12:2753bab2 11:479943df 10:5cf4385a 09:35368174  
08:17060ce1  
07:5b6acbf6 06:10da5941 05:789c639c 04:3448fcc0 03:2af10db2 02:2567b94a 01:1bdc9fab  
00:10e10abb  
R1:0da66493 R2:9125bf61

Internal state at time -6  
15:15b2e89f 14:77b4af31 13:43930a37 12:73775d6a 11:2753bab2 10:479943df 09:5cf4385a  
08:35368174  
07:17060ce1 06:5b6acbf6 05:10da5941 04:789c639c 03:3448fcc0 02:2af10db2 01:2567b94a  
00:1bdc9fab  
R1:bbd3b2af R2:4b50ed23

Internal state at time -5  
15:24ff6e20 14:15b2e89f 13:77b4af31 12:43930a37 11:73775d6a 10:2753bab2 09:479943df  
08:5cf4385a  
07:35368174 06:17060ce1 05:5b6acbf6 04:10da5941 03:789c639c 02:3448fcc0 01:2af10db2  
00:2567b94a  
R1:8b5d75ba R2:0b92f50c

Internal state at time -4  
15:740c40b9 14:24ff6e20 13:15b2e89f 12:77b4af31 11:43930a37 10:73775d6a 09:2753bab2  
08:479943df  
07:5cf4385a 06:35368174 05:17060ce1 04:5b6acbf6 03:10da5941 02:789c639c 01:3448fcc0  
00:2af10db2  
R1:8ccae757 R2:ab3d746d

Internal state at time -3  
15:026a5503 14:740c40b9 13:24ff6e20 12:15b2e89f 11:77b4af31 10:43930a37 09:73775d6a  
08:2753bab2  
07:479943df 06:5cf4385a 05:35368174 04:17060ce1 03:5b6acbf6 02:10da5941 01:789c639c  
00:3448fcc0  
R1:f888aec9 R2:04223414

Internal state at time -2  
15:194b2a57 14:026a5503 13:740c40b9 12:24ff6e20 11:15b2e89f 10:77b4af31 09:43930a37  
08:73775d6a  
07:2753bab2 06:479943df 05:5cf4385a 04:35368174 03:17060ce1 02:5b6acbf6 01:10da5941  
00:789c639c  
R1:ee139bec R2:c6666f03

Internal state at time -1  
15:7a9alcff 14:194b2a57 13:026a5503 12:740c40b9 11:24ff6e20 10:15b2e89f 09:77b4af31  
08:43930a37  
07:73775d6a 06:2753bab2 05:479943df 04:5cf4385a 03:35368174 02:17060ce1 01:5b6acbf6  
00:10da5941  
R1:860a7dfa R2:bf0e0ffc

Internal state at time 0  
15:3d4aa9e7 14:7a9alcff 13:194b2a57 12:026a5503 11:740c40b9 10:24ff6e20 09:15b2e89f  
08:77b4af31  
07:43930a37 06:73775d6a 05:2753bab2 04:479943df 03:5cf4385a 02:35368174 01:17060ce1  
00:5b6acbf6  
R1:129d8b39 R2:2d7cdce1

Internal state at time 1  
15:71db1828 14:3d4aa9e7 13:7a9alcff 12:194b2a57 11:026a5503 10:740c40b9 09:24ff6e20  
08:15b2e89f  
07:77b4af31 06:43930a37 05:73775d6a 04:2753bab2 03:479943df 02:5cf4385a 01:35368174  
00:17060ce1  
R1:ab7cf688 R2:c1598aa6

Internal state at time 2  
15:258937da 14:71db1828 13:3d4aa9e7 12:7a9alcff 11:194b2a57 10:026a5503 09:740c40b9  
08:24ff6e20

# BS ISO/IEC 18033-4:2011+A1:2020

## ISO/IEC 18033-4:2011+A1:2020(E)

07:15b2e89f 06:77b4af31 05:43930a37 04:73775d6a 03:2753bab2 02:479943df 01:5cf4385a  
00:35368174  
R1:3cec1a4a R2:9053cc0e

Internal state at time 3  
15:52831a13 14:258937da 13:71db1828 12:3d4aa9e7 11:7a9a1cff 10:194b2a57 09:026a5503  
08:740c40b9  
07:24ff6e20 06:15b2e89f 05:77b4af31 04:43930a37 03:73775d6a 02:2753bab2 01:479943df  
00:5cf4385a  
R1:e5bf0e3b R2:75c177aa

Internal state at time 4  
15:7ce6f22f 14:52831a13 13:258937da 12:71db1828 11:3d4aa9e7 10:7a9a1cff 09:194b2a57  
08:026a5503  
07:740c40b9 06:24ff6e20 05:15b2e89f 04:77b4af31 03:43930a37 02:73775d6a 01:2753bab2  
00:479943df  
R1:e8f766fd R2:7eb070e5

Internal state at time 5  
15:5cc75397 14:7ce6f22f 13:52831a13 12:258937da 11:71db1828 10:3d4aa9e7 09:7a9a1cff  
08:194b2a57  
07:026a5503 06:740c40b9 05:24ff6e20 04:15b2e89f 03:77b4af31 02:43930a37 01:73775d6a  
00:2753bab2  
R1:b2829ddc R2:e21e038a

Internal state at time 6  
15:1fe39cdc 14:5cc75397 13:7ce6f22f 12:52831a13 11:258937da 10:71db1828 09:3d4aa9e7  
08:7a9a1cff  
07:194b2a57 06:026a5503 05:740c40b9 04:24ff6e20 03:15b2e89f 02:77b4af31 01:43930a37  
00:73775d6a  
R1:5f307d54 R2:ebfe08f3

Internal state at time 7  
15:04c3f285 14:1fe39cdc 13:5cc75397 12:7ce6f22f 11:52831a13 10:258937da 09:71db1828  
08:3d4aa9e7  
07:7a9a1cff 06:194b2a57 05:026a5503 04:740c40b9 03:24ff6e20 02:15b2e89f 01:77b4af31  
00:43930a37  
R1:4de35a6f R2:0b0b19d7

Internal state at time 8  
15:7df71d75 14:04c3f285 13:1fe39cdc 12:5cc75397 11:7ce6f22f 10:52831a13 09:258937da  
08:71db1828  
07:3d4aa9e7 06:7a9a1cff 05:194b2a57 04:026a5503 03:740c40b9 02:24ff6e20 01:15b2e89f  
00:77b4af31  
R1:7a55edeb R2:78f89fd1



## Annex D (informative)

### Security information

#### D.1 Security levels of stream ciphers

This annex lists the security levels of the stream ciphers specified in this part of ISO/IEC 18033.

**Table D.1 — Security levels of stream cipher modes**

Output function	Data confidentiality	Data integrity
Binary additive mode	An encryption mechanism based on these output functions is secure as long as the underlying keystream generator is secure. This is mathematically proven.	This mode does not provide security with respect to data integrity.
MULTI-S01 mode		The security with respect to data integrity is generally as secure as the keystream generator. However the security level is always upperbounded by the length of a forged message. A forged message whose length is $un$ bits can be accepted with the successful probability of $(u-2)2^{-n}$ , where $n$ is the security parameter of the mode.

**Table D.2 — Security levels of dedicated keystream generators**

Keystream generator	Security statement	Key length	Computational complexity of best known attack
CFB, OFB, CTR modes with $n$ -bit block cipher with $k$ -bit key length	Assuming that the underlying block cipher is secure, CFB, OFB, and CTR modes are generally indistinguishable from a random sequence. However if the processed number of blocks exceeds around $2^{n/2}$ then the probability of distinguishing between such keystream and a random sequence becomes substantial. The security of such a keystream is upperbounded by the security of the block cipher, such as the key length and the degree to which the cipher can be cryptanalysed.	$k$ bits	$2^{n/2}$
MUGI	There is no known cryptanalytic attack against MUGI which is faster than an exhaustive key search. MUGI has a 128-bit key length.	128 bits	$2^{128}$
SNOW 2.0	For 256-bit mode, distinguishing algorithms have been that distinguish the output of SNOW 2.0 from a random sequence. The best one is estimated to require about $2^{174}$ -bit known keystream. This barely violates the theoretical security level which can be deduced from the key length, 256 bits. However, any of attacks related to known distinguisher will not be threat for practical uses.	128 or 256 bits	$2^{174}$

Keystream generator	Security statement	Key length	Computational complexity of best known attack
Rabbit	There is no known distinguishing attack against Rabbit that is faster than exhaustive key search, as long as no more than $2^{64}$ keystream blocks are generated using one key.	128 bits	$2^{128}$
Decim <sup>v2</sup>	There is no known attack against Decim <sup>v2</sup> that is faster than exhaustive key search.	80 bits	$2^{80}$
KCipher-2 (K2)	There is no known attack against K2 which is faster than the key exhaustive search. K2 has a 128-bit key length.	128 bits	$2^{128}$
ZUC	There is no known attack against ZUC which is faster than the exhaustive key search. ZUC has a 128-bit key length.	128 bits	$2^{128}$
NOTE ZUC is also included in ETSI/SAGE specification document <sup>[21]</sup> , Chinese cryptographic standard GM/T 0001-2012 and Chinese national standard GB/T 33133.1-2016. During the evolution of ZUC, a flaw was found against its early public evaluation version, i.e. the differential attack in the paper <sup>[22]</sup> . The attack is not applicable to the final version of ZUC as stated in [22]. The version and evaluation history of ZUC is available in section 6 of [20]. ZUC refers to the final standard version. <span style="border: 1px solid black; padding: 0 2px;">A1</span>			

## D.2 Security-efficiency trade-off in MULTI-S01

Let  $n$  be the block size of MULTI-S01 output function. For  $n \cdot u$ -bit message, MULTI-S01 output function iterates a block process  $u+2$  times. In either of software and hardware implementation, the dominant calculation is the multiplication in the finite field  $GF(2^n)$ .

If two implementations, between which the block sizes  $n$  are different, are compared, the implementation with smaller  $n$  is generally faster than the other — although the implementation with the smaller value of  $n$  iterates more times, each multiplication can be calculated in less time and using smaller space. The factor by which the calculation speeds up depends on the algorithm or platform used to implement it.

Table D.3 shows some experimental results of MULTI-S01 implementations with  $n=64$ .



**Table D.3 — Experimental results of MULTI-S01 mode without keystream generation ( $n=64$ )**

Designation	Speed/size	Platform specification
High speed ASIC	5.1 Gbps@80MHz, 25.7 K Gates	ASIC using Hitachi cell library HG73C (0.35μm)
Small sized ASIC	2.0 K Gates, 100Mbps @ 100MHz	
Software	10 cycle/Byte (or equivalently 520 bps@650 MHz)	Intel Pentium III 650 MHz (Coppermine), Windows98 SE, RAM 64MB, Visual C++ Ver.6.0 Service Pack 3.

## D.3 Guidance on stream ciphers

This annex lists the features of the stream ciphers specified in this part of ISO/IEC 18033.







Table D.4 — Features of stream ciphers

Keystream generator	Property statement
CFB, OFB, CTR modes with $n$ -bit block cipher with $k$ -bit key length	The advantage of these three modes is that they share components with other modes of a block cipher. The modes can be used for any length of a plaintext on several architectures, and OFB and CTR are more suitable for encryption/decryption of $nx$ -bit length plaintexts. CFB is able to recover synchronization errors, and its performance depends on a length of keystream. CTR accepts random accesses to ciphertext.
MUGI	MUGI uses a 128-bit initial vector. It consists of 64-bit operations and generates 64-bit keystream for each execution. The cipher can be used for any length of a plaintext on several architectures, and it is more efficient for 64-bit CPU architecture and encryption/decryption of $64x$ -length plaintexts.
SNOW 2.0	SNOW 2.0 uses a 128-bit initial vector. It consists of 32-bit operations and generates 32-bit keystream for each execution. The cipher can be used for any length of a plaintext on several architectures, and it is more efficient for 32-bit CPU architecture and encryption/decryption of $32x$ -length plaintexts.
Rabbit	Rabbit uses a 64-bit initial vector. It consists of 32-bit operations and generates 128-bit keystream for each execution. The cipher can be used for any length of a plaintext on several architectures, and it is more efficient for 32-bit CPU architecture and encryption/decryption of $128x$ -length plaintexts.
Decim <sup>v2</sup>	Decim <sup>v2</sup> uses a 64-bit initial vector. It consists of bit-wise operations and generates 1-bit keystream for each execution. The cipher can be used for any length of a plaintext on several architectures, and it is more efficient for hardware implementation on resource-constrained devices.
KCipher-2 (K2)	Kcipher-2 uses a 128-bit initial vector. It consists of 32-bit operations and generates 64-bit keystream for each execution. The cipher can be used for any length of a plaintext on several architectures, and it is more efficient for 32-bit CPU architecture and encryption/decryption of $64x$ -length plaintexts.
 ZUC	ZUC uses a 128-bit initial vector. It consists of 32-bit operations and generates 32-bit keystream for each execution. The cipher can be used for any length of a plaintext on several architectures, and it is more efficient for 32-bit CPU architecture and encryption/decryption of $32x$ -length plaintexts. 



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