

# THE COMMUTATOR AS FIXED POINT OF A CLOSURE OPERATOR

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ABSTRACT. In this note we elaborate on the following remark of Keith Kearnes [Kea95, p. 930] giving an alternate description of the commutator (here  $\beta$  is a congruence of  $\mathbf{A}$  and  $\mathbf{A} \times_{\beta} \mathbf{A}$  denotes the subalgebra of  $\mathbf{A}^2$  with universe  $\beta$ ):

“Let  $\Delta_{\beta, \alpha}$  be the congruence on  $\mathbf{A} \times_{\beta} \mathbf{A}$  generated by

$$\{ \langle (x, x), (y, y) \rangle \mid (x, y) \in \alpha \}.$$

Call a subset  $G \subseteq A^2$   $\Delta$ -closed if

$$\Delta_{\beta, \alpha} \circ G \circ \Delta_{\beta, \alpha} \subseteq G.$$

When  $\alpha$  and  $\beta$  are reflexive, compatible relations, then  $[\alpha, \beta]$  is the smallest subset  $\gamma \subseteq A^2$  such that (i)  $\gamma$  is a congruence of  $\mathbf{A}$  and (ii)  $\gamma$  is  $\Delta$ -closed.”

## 1. INTRODUCTION

To be honest, the genesis of these notes was my futile attempt to interpret the expression  $\Delta_{\beta, \alpha} \circ G \circ \Delta_{\beta, \alpha}$ . Noting that  $\Delta_{\beta, \alpha} \subseteq (A \times A)^2$  while  $G \subseteq A^2$ , I wasn't aware of a standard means of composing relations of such different arities. In this note I try to reconcile this by giving an alternative “closure” operation that achieves the same end and verifies Keith's assertion that “there is a useful alternate description of  $[\alpha, \beta]$ .” Indeed, the commutator is a fixed point of some closure operator (i.e., closed set), and the closure operator involves the relation  $\Delta_{\beta, \alpha}$  defined above.

**1.1. Definitions.** For an algebra  $\mathbf{A}$  with congruence relations  $\alpha, \beta \in \text{Con } \mathbf{A}$ , let  $\underline{\beta}$  denote the subalgebra of  $\mathbf{A} \times \mathbf{A}$  with universe  $\beta$ . Let  $D = \{(a, a) \mid a \in A\}$  and  $D_{\alpha}^2 = \{((a, a), (b, b)) \in D^2 \mid a \alpha b\}$ . Finally, let  $\Delta_{\beta, \alpha} = \text{Cg}^{\underline{\beta}}(D_{\alpha}^2)$  be the congruence on  $\underline{\beta}$  generated by the set  $D_{\alpha}^2$ . The congruence class of  $\Delta_{\beta, \alpha}$  that contains  $(b, b')$  is denoted and defined as follows:

$$(b, b')/\Delta_{\beta, \alpha} = \{(a, a') \in \beta \mid (a, a') \Delta_{\beta, \alpha} (b, b')\}.$$

Let  $\Phi_{\beta, \alpha}: \mathcal{P}(\beta) \rightarrow \mathcal{P}(\beta)$  be the function that takes each  $B \subseteq \beta$  to

$$(1.1) \quad \Phi_{\beta, \alpha}(B) = \bigcup_{(b, b') \in B} (b, b')/\Delta_{\beta, \alpha}.$$

## 2. FIXED POINT LEMMA

**Lemma 2.1.** *Let  $\mathbf{A}$  be an algebra with  $\alpha, \beta \in \text{Con}(\mathbf{A})$ . If  $\Phi := \Phi_{\beta, \alpha}$  is defined as in (1.1), then*

- (i)  $\Phi$  is a closure operator on  $\mathcal{P}(\beta)$ ;
- (ii)  $[\alpha, \beta]$  is the least fixed point of  $\Phi$ .

*Proof.* (i) Fix  $B \subseteq \beta$ . We must prove the following: (a)  $B \subseteq \Phi(B)$ ; (b)  $B \subseteq C \Rightarrow \Phi(B) \subseteq \Phi(C)$ ; and (c)  $\Phi(\Phi(B)) = \Phi(B)$ . If  $(b, b') \in B$ , then  $(b, b') \in \Phi(B)$  since the operation (1.1) does not discard any of the pairs that were already in  $B$ . As for (b), if  $(a, a') \in \Phi(B)$ , then there exists  $(b, b') \in B \subseteq C$  such that  $(a, a') \Delta_{\beta, \alpha} (b, b')$ . Since  $(b, b')$  belongs to  $C$  we have  $(a, a') \in \Phi(C)$  as well. As for (c), it clearly follows from (a) and (b) that  $\Phi(B) \subseteq \Phi(\Phi(B))$ , so we prove the reverse inclusion. Let  $(d, d') \in \Phi(\Phi(B))$ . Then  $(c, c') \Delta_{\beta, \alpha} (d, d')$  for some  $(c, c') \in \Phi(B)$ , which implies  $(b, b') \Delta_{\beta, \alpha} (c, c')$  for some  $(b, b') \in B$ . By transitivity of  $\Delta_{\beta, \alpha}$  we conclude that  $(d, d') \in \Phi(B)$ , as desired.

(ii) Since  $[\alpha, \beta] \in \mathcal{P}(\beta)$  we have  $[\alpha, \beta] \subseteq \Phi([\alpha, \beta])$ , by (i). We prove the reverse inclusion. If  $(c, c') \in \Phi([\alpha, \beta])$ , then (1.1) implies there exists  $(b, b') \in [\alpha, \beta]$  such that

$$(2.1) \quad (b, b') \Delta_{\beta, \alpha} (c, c').$$

From the definition of  $\Delta_{\beta, \alpha}$  and Mal'tsev's theorem on congruence generation, (2.1) holds if and only if  $\exists z_i \beta z'_i$  ( $0 \leq i \leq n$ ),  $\exists x_i \alpha y_i$  ( $0 \leq i < n$ ),  $\exists f_i \in \text{Pol}_1(\mathbf{A} \times \mathbf{A})$  ( $0 \leq i < n$ ) such that  $(b, b') = (z_0, z'_0)$  and  $(z_n, z'_n) = (c, c')$ , and

$$(2.2) \quad \{(b, b'), (z_1, z'_1)\} = \{f_0(x_0, x_0), f_0(y_0, y_0)\}$$

$$(2.3) \quad \{(z_1, z'_1), (z_2, z'_2)\} = \{f_1(x_1, x_1), f_1(y_1, y_1)\}$$

⋮

$$\{(z_{n-1}, z'_{n-1}), (c, c')\} = \{f_{n-1}(x_{n-1}, x_{n-1}), f_{n-1}(y_{n-1}, y_{n-1})\}$$

For each  $(0 \leq i < n)$ ,  $f_i \in \text{Pol}_1(\mathbf{A} \times \mathbf{A})$ , which means

$$f_i(x, x') = g_i^\beta((x, x'), (a_1, a'_1), \dots, (a_k, a'_k)) = (g_i^\mathbf{A}(x, \mathbf{a}), g_i^\mathbf{A}(x, \mathbf{a}')),$$

for some  $k$ ,  $g_i \in \text{Clo}_{k+1}(\mathbf{A})$ , and constants tuples  $\mathbf{a} = (a_1, \dots, a_k)$  and  $\mathbf{a}' = (a'_1, \dots, a'_k)$  such that  $a_i \beta a'_i$  ( $1 \leq i \leq k$ ). By (2.2), either

$$(b, b') = (g_0(x_0, \mathbf{a}), g_0(x_0, \mathbf{a}')) \quad \text{and} \quad (z_1, z'_1) = (g_0(y_0, \mathbf{a}), g_0(y_0, \mathbf{a}')),$$

or vice-versa. Since  $x_0 \alpha y_0$  and  $a_i \beta a'_i$  ( $1 \leq i \leq k$ ), the  $\alpha, \beta$ -term condition entails

$$g_0(x_0, \mathbf{a}) [\alpha, \beta] g_0(x_0, \mathbf{a}') \iff g_0(y_0, \mathbf{a}) [\alpha, \beta] g_0(y_0, \mathbf{a}').$$

This and (2.2) yield  $(b, b') \in [\alpha, \beta]$  iff  $(z_1, z'_1) \in [\alpha, \beta]$ . Similarly (2.3) and  $x_1 \alpha y_1$  imply  $(z_1, z'_1) \in [\alpha, \beta]$  iff  $(z_2, z'_2) \in [\alpha, \beta]$ . Inductively, we arrive at  $(b, b') \in [\alpha, \beta]$  iff  $(c, c') \in [\alpha, \beta]$ , as desired.

We have thus proved  $[\alpha, \beta]$  is a fixed point of  $\Phi$ . In other words,  $[\alpha, \beta]$  is a “ $\Phi$ -closed” subset of  $\beta$ . (A set  $B \subseteq \beta$  is called  $\Phi$ -closed provided  $\Phi(B) \subseteq B$ .) Recall, if  $f$  is a monotone increasing function defined on a complete poset  $\langle P, \leq \rangle$ , then the least fixed point of  $f$  is  $\bigwedge \{p \in P \mid fp \leq p\}$ . Thus, Lemma 2.1 (ii) asserts that

$$(2.4) \quad [\alpha, \beta] = \bigwedge \{B \subseteq \beta \mid \Phi(B) \subseteq B\}.$$

We already proved  $[\alpha, \beta]$  is  $\Phi$ -closed, so it remains to check for every  $\Phi$ -closed subset  $B \subseteq \beta$  that  $[\alpha, \beta] \subseteq B$ . Fix a  $\Phi$ -closed subset  $B \subseteq \beta$ . It suffices to prove  $\mathbf{C}(\alpha, \beta; \Phi(B))$ , since this implies  $[\alpha, \beta] \subseteq \Phi(B) \subseteq B$ . Thus, our goal is to establish the  $\alpha, \beta$ -term condition.

Let  $p \in \text{Pol}_{k+1}(\mathbf{A})$  and  $a \alpha a'$  and  $c_i \beta c'_i$  ( $1 \leq i \leq k$ ); suppose  $p(a, \mathbf{c}) \Phi(B) p(a, \mathbf{c}')$ . We prove that these hypotheses entail the following relation:

$$(2.5) \quad p(a', \mathbf{c}) \Phi(B) p(a', \mathbf{c}').$$

By definition of  $\Phi$ , (2.5) is equivalent to the existence of some pair  $(b, b') \in B$  such that  $(b, b') \Delta_{\beta, \alpha} (p(a', \mathbf{c}), p(a', \mathbf{c}'))$ . Notice that the pair  $(p(a, \mathbf{c}), p(a, \mathbf{c}'))$  belongs to  $B$  since  $(p(a, \mathbf{c}), p(a, \mathbf{c}')) \in \Phi(B) \subseteq B$ . Also,  $c_i \beta c'_i$  ( $0 \leq i < k$ ) implies

$$\begin{aligned} ((a, a), (c_1, c'_1), (c_1, c'_1), \dots, (c_k, c'_k)) &\in \beta^{k+1} \quad \text{and} \\ ((a', a'), (c_1, c'_1), (c_1, c'_1), \dots, (c_k, c'_k)) &\in \beta^{k+1}. \end{aligned}$$

Therefore,

$$(2.6) \quad p^\beta((a, a), (c_1, c'_1), (c_1, c'_1), \dots, (c_k, c'_k)) = (p^{\mathbf{A}}(a, \mathbf{c}), p^{\mathbf{A}}(a, \mathbf{c}')) \in \beta \quad \text{and}$$

$$(2.7) \quad p^\beta((a', a'), (c_1, c'_1), (c_1, c'_1), \dots, (c_k, c'_k)) = (p^{\mathbf{A}}(a', \mathbf{c}), p^{\mathbf{A}}(a', \mathbf{c}')) \in \beta.$$

Finally,  $a \alpha a'$  implies  $p(a, \mathbf{c}) \alpha p(a', \mathbf{c})$ , and this—together with (2.6) and (2.7)—proves the pair  $((p(a, \mathbf{c}), p(a, \mathbf{c}')), (p(a', \mathbf{c}), p(a', \mathbf{c}')))$  belongs to  $\Delta_{\beta, \alpha}$ . Since  $(p(a, \mathbf{c}), p(a, \mathbf{c}')) \in B$ , this proves  $(p(a', \mathbf{c}), p(a', \mathbf{c}')) \in \Phi(B)$ , completing the proof.  $\square$

## REFERENCES

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