



- Security Goals
- Threats & Defenses
- Encryption
- Secure connections
- Secure servers

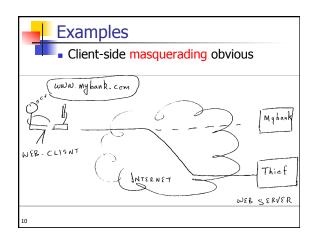
Security Goals

Lots of different dimensions to Web security
Confidentiality/Privacy
Data Integrity
Service Integrity (availability)
Authenticity
Non-repudiation

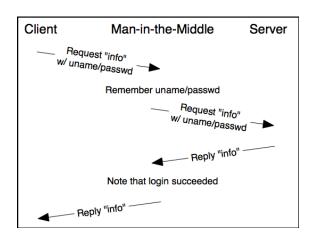
# 4

# Threads/evil techniques

- Masquerading: pretend to be someone else
  - Address spoofing: pretend to be somewhere else
  - Eavesdropping: listen in on comm
  - Man-in-the-middle: manipulate comm
  - Replay: record and use comm data later
  - Physical: steal the post-it with password









## **Examples**

- Replay attack
  - Record good conversation, then replay to masquerade as one party
  - E.g., record the uname/passwd, then replay it to authenticate in the future
- How vulnerable is TCP to replay-attack hijacking?
- Man-in-the-middle?



#### Defenses

- Authentication
  - A party establishes identity to others
  - Usually requires credentials
    - ID badge
    - Passport
    - Password
  - In real world, assymmetric relationships mean one-sided authentication
  - On Web, no social clues to indicate identity; all parties must authenticate
  - Makes masquerading impossible

#### **Defenses**

- Non-repudiation
  - A party to transaction cannot later deny it
  - In real-world, we use signatures
  - On Web, can you repudiate a login?
- Authorization
  - Granting privileges to authenticated parties
  - A policy is a spec of authorization rules
- A mechanism is the system by which a security policy is implemented
  - Most fundamental is encryption
  - Encryption != security

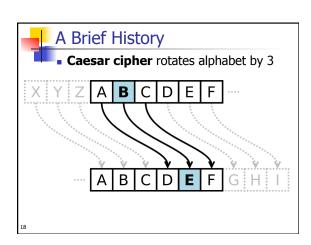
# **Encryption**

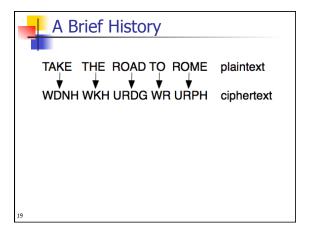
- Cryptography vs cryptanalysis
- Encryption applies a reversible fn to some piece of data, yielding something unreadable
- Decryption recovers the original data from the unreadable encryption-output
- The encryption/decryption algorithm assumed known; the key is secret



# Encryption (2)

- Plaintext string s
- Encryption key K<sub>enc</sub>
- Decryption key K<sub>dec</sub>
- Encrypt s with K<sub>enc</sub> to obtain ciphertext
- Decrypt K<sub>enc</sub>(s) with decryption key K<sub>dec</sub> to reobtain s
- $K_{dec}(K_{enc}(s)) = s$

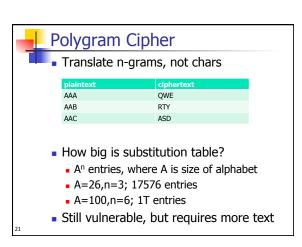






## Substitution Ciphers

- No need to shift 3 chars
  - You could do 2! Or even 4!
- You also don't have to shift the alphabet at all. Just arbitrary 1:1 mapping of alphabet chars, using a substitution table
- All of these are vulnerable to frequency analysis
  - Letter
  - Word
  - Common phrases





### **Substitution Rules**

- Don't store table explicitly; derive table rows using substitution rule
  - E.g., **s XOR k**, where k is *key*
  - Remember: security level depends on size of key
  - Key of len  $b => 2^b$  possible keys



#### **Substitution Rules**

- XOR "flips a bit" for input bits that correspond to key's "1"
  - Correspond to a 0? No change

0000000001010101 plaintext 1011010010011100 key 1011010011001001 XOR

- Encrypted string should ideally show no pattern for frequency analysis attack
- Use key long enough to make ciphertext appear random



## Substitution Rules (2)

- What's the right size key?
  - Who is trying to break the scheme?
  - 3GHz CPU => 300 inst for possible key test
  - 1 sec, 10M keys
  - 1 day, 1T keys
  - 60-bit key takes 100 CPUs 3 years
- Is that good enough?
- Also, use statistical techniques to determine ideal key length

24



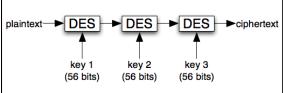
# **Data Encryption Standard**

- DES is a block cipher with 56-bit key
  - 64-bits at a time
  - Perform 16 rounds of encryption, w/std. permutations of keys and data
  - DES is not secure
- Data xmitted in 64-bit blocks, each may be coded independently

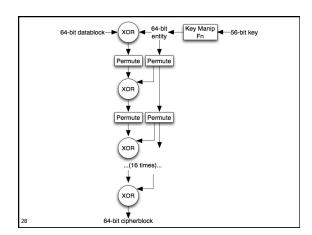
DES in 64-bit blocks TAKE THE ROAD TO ROME TAKE THE ROAD TO ROME DES DES DES a594b3cdg802318 687e39824a6b987c 243ace1976358bd6 a594b3cdg802318243ace1976358bd6687e39824a6b987c



- Triple-DES is 168 bits
- Break the key into 3 parts



DES' bit-logic techniques make it fast





# **Key Management**

- Encryption depends on everyone having the same key
- Key distribution is the weak link
  - Hard to distribute
  - Vulnerable to key theft
- What we've been discussing is best called symmetric encryption
  - Only kind from 5000BC to 1976
- Assymmetric, or public-key encryption, uses two keys
  - One of the greatest achievements of CS