

VERY IMPRESSIVE TEXT

Collection of important mathematical bullshit and definitions

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1 VfdA

The abbreviation *VfdA* stands for the German expression *Voll für den Ar**** and can be used in pseudo-academic papers and documents like the current one. It is generally used before a long-winded and utterly useless mathematical proof, which stands in no connection to the rest of the paper. It is only used to impress possible readers and to boast about the non-existent knowledge of the author about mathematical subjects. A perfect example for this abbreviation is the following one:

$$\begin{aligned}
 \text{VfdA : } \quad \mathcal{F}(f)(t) &= \frac{1}{(2\pi)^{\frac{n}{2}}} \int_{\mathbb{R}^n} f(x) e^{-it*x} dx & \int_{-\infty}^{\infty} |f(t)| dx < \infty \\
 f_m &= \sum_{k=0}^{n-1} x_{2k} e^{-\frac{2\pi i}{2n} m(2k)} + \sum_{k=0}^{n-1} x_{2k+1} e^{-\frac{2\pi i}{2n} m(2k+1)} \\
 &= \sum_{k=0}^{n-1} x'_k e^{-\frac{2\pi i}{2n} mk} + e^{-\frac{\pi i}{n} m} \sum_{k=0}^{n-1} x''_k e^{-\frac{2\pi i}{n} mk} \\
 &= \begin{cases} f'_m + e^{-\frac{\pi i}{n} m} f''_m & \text{if } m < n \\ f'_{m-n} - e^{-\frac{\pi i}{n} (m-n)} f''_{m-n} & \text{if } m \geq n \end{cases}
 \end{aligned}$$

One shall note, that *VfdA* automatically implicates *OBdA* (German: *Ohne Beschränkung der Allgemeinheit*, English: *Without loss of generality*) to compensate the so-called *LoC* (English: *Loss of context*) when used.

2 Plustorial

The *plustorial* (German: *Die Plusultät*) of a number is defined as follows:

$$n? := \sum_{i=1}^n i \quad (n \in \mathbb{Z}) \qquad n? = \sum_{i=1}^n i = \frac{n(n-1)}{2}$$

3 Closed Interval

The alternate notation for a closed interval over a set $K \subseteq \mathbb{K}$, which has the comparison operator \leq defined for every elements $k, l \in K$, can be written as follows:

$$\begin{aligned}
 \langle k, l \rangle &:= \begin{cases} [k, l], & \text{if } k \leq l \\ [l, k], & \text{otherwise} \end{cases} & l, k \in K \subseteq \mathbb{K} \\
 \langle \pm k \rangle &:= [-k, k]
 \end{aligned}$$

4 Set with a finite amount of elements

Let K be the subset of the field \mathbb{K} and let $f : K \rightarrow \mathbb{B}$ be a function, which defines for every element $k \in \mathbb{K}$, whether it is also an element of the subset K .

$$\forall k \in \mathbb{K} : f(k) \Leftrightarrow k \in K$$

Based on the equation above, the subset K can now be re-defined as follows:

$$K = \{k \in \mathbb{K} \mid f(k)\} \subset \mathbb{K}$$

The following notation can be used to indicate, that the subset $K \subset \mathbb{K}$ has only a finite amount of elements k :

$$\{k \in \mathbb{K} \mid f(k)\}_{\infty}^< : \Leftrightarrow |\{k \in \mathbb{K} \mid f(k)\}| < \infty$$

5 Assembler command "ABK"

The i386 assembler command ABK triggers a quadruple-fault, when loaded into the instruction cache during execution and simultaneously short-circuits the machine's DC voltage regulator with the CPU power inlet, causing the CPU to be grilled with with the given DC voltage (usually 240V in Europe). Have Fun! Example usage:

```
mov    dword ptr [ebp-18h], esp
push   1
call   dword ptr ds:[404090h]
add    esp, 4
mov    dword ptr ds:[403030h], 0FFFFFFFFh
mov    ecx, dword ptr ds:[403020h]
call   dword ptr ds:[404088h]
mov    edx, dword ptr ds:[403028h]
mov    dword ptr [eax], edx
mov    dword ptr ds:[403038h], ecx
mov    eax, [403010]
call   dword ptr ds:[404080h]
add    esp, 4
call   401C60
push   403008h
add    esp, 8
mov    edx, dword ptr ds:[403024h]
mov    dword ptr [ebp-28h], edx
push   eax
mov    ecx, dword ptr ds:[403020h]
lea    ecx, dword ptr [ebp-10h]
abk    // initiate quadruple-faulting
```

6 ε -Potato

The so-called *Epsilon-Kartoffel* (German expression for *epsilon-potato*) is a special form of an open convex topological ε -sphere or ε -neighbourhood. It is a subset of the topological space \mathbb{K}^n , which is grouped around the element $m \in \mathbb{K}^n$. The following rules apply for a subset $K_\varepsilon(m) \subset \mathbb{K}^n$ being qualified as a *epsilon-potato* around the point m :

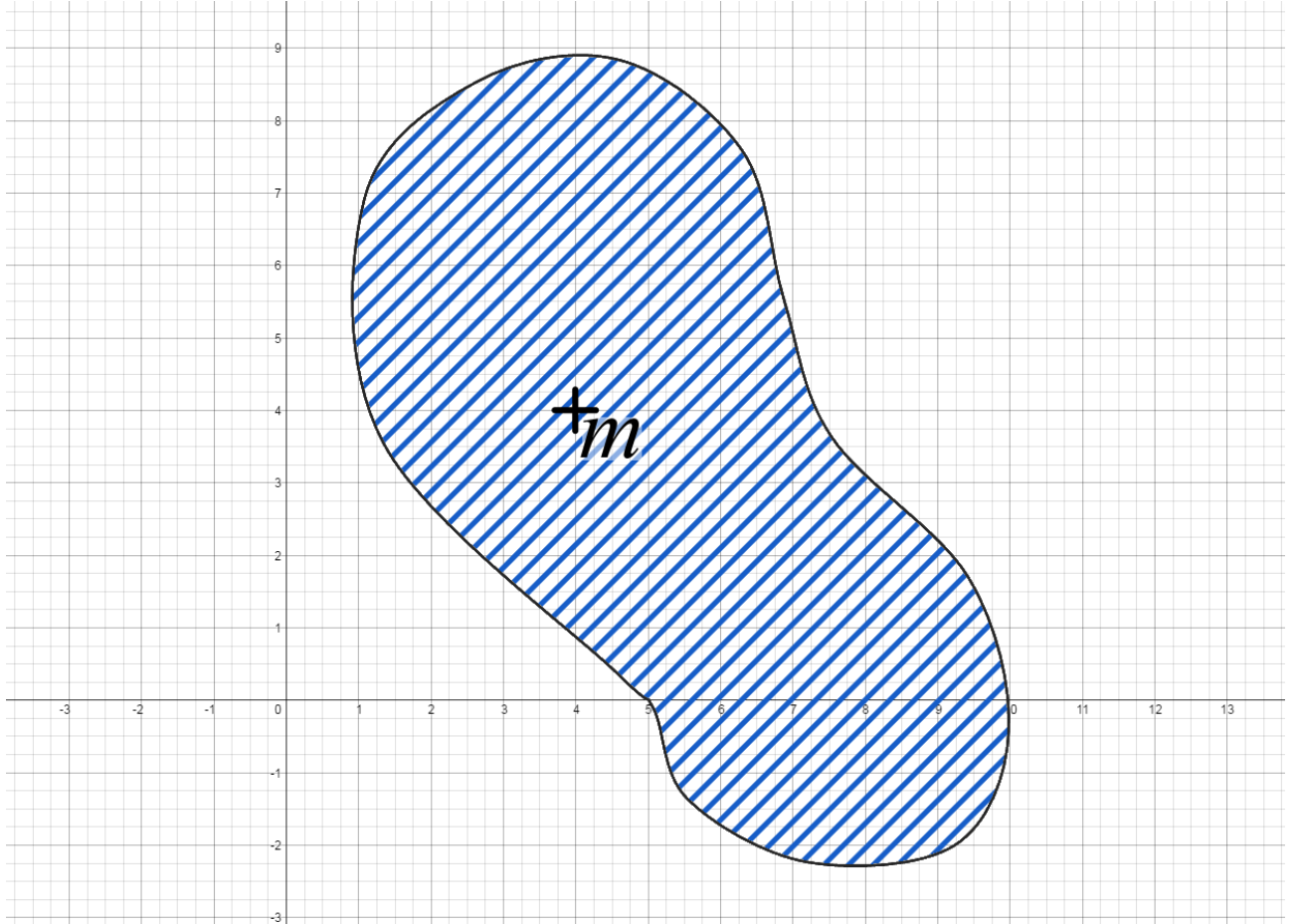
- (1) $m \in \mathbb{K}^n$, $m \in K_\varepsilon(m)$
- (2) $K_\varepsilon(m)_{\infty} <$
- (3) $K_\varepsilon(m), \mathfrak{S}(K_\varepsilon(m)) \subseteq \mathcal{C}^\infty(\mathbb{K}^n)$
- (4) $\frac{\sup \{\|p - m\| : p \in K_\varepsilon(m)\}}{\inf \{\|q - m\| : q \in \mathbb{K}^n \setminus K_\varepsilon(m)\}} < \infty$
- (5) $\forall p \in K_\varepsilon(m) : \|p - m\| < \infty$

As the requirement (3) states, the surface $\mathfrak{S}(K_\varepsilon(m))$ must be an absolute continuously one. It can also be represented by the following function \mathcal{S} :

$$\mathcal{S} : \mathbb{K}^n \rightarrow \mathfrak{S}(K_\varepsilon(m)) \quad \mathcal{S} \in \mathcal{C}^\infty(\mathbb{K}^n)$$

which is an absolute continuous one over the field \mathbb{K}^n and represents each point on the potato's surface $\mathfrak{S}(K_\varepsilon(m))$ based on the given n -dimensional rotation angle $\varphi \in \mathbb{K}^n$.

The point m is also defined as the *physical center of mass* of the ε -potato $K_\varepsilon(m)$.



7 The *for-probably-all*-quantifier

The *for-probably-all*-quantifier (written as a W combined with an vertically flipped A) is a quantifier, which states that the given condition is *probably* valid for all elements. The number of elements, for which the given condition is *probably* valid must be a finite number.

The following example demonstrates the use of the *for-probably-all*-quantifier by assuming that one could divide by *probably* every number in the field of rational numbers \mathbb{Q} :

$$\forall p, q \in \mathbb{Q} : pq^{-1} \in \mathbb{Q}$$

As we very well know, one cannot divide a number by zero, thus the use of a *for-all*-quantifier would be illegal. The *for-probably-all*-quantifier, however, solves this problem, as the amount of elements whom which a number cannot be divided by (only zero) is finite:

$$\{q \in \mathbb{Q} \mid \text{number cannot be divided by } q\}_{\infty}^<$$

The formal definition for the *for-probably-all*-quantifier is therefore:

$$\text{let } M \subseteq \mathbb{K}, \quad \text{let } f : M \rightarrow \mathbb{B}$$

$$\forall k \in M : f(k) \quad :\Longleftrightarrow \quad \{k \in M \mid \neg f(k)\}_{\infty}^< =: T \subsetneq \mathbb{K} \quad \wedge \quad \forall k \in M \setminus T : f(k)$$

The *for-probably-all*-quantifier can be seen as an implication of the *for-all*-quantifier:

$$\forall A \quad \Longrightarrow \quad \forall \forall A \quad (A \in \mathbb{B})$$

The probability, that the given statement $A \in \mathbb{B}$ is probably true inside a *for-all*-quantifier, when it has been proven true inside a *for-probably-all*-quantifier, is called the *magnitude of the for-probably-all-quantifier* and is usually written as $\mathfrak{p} \in (0.5, 1]$.

Example:

$$\text{let } M = [-10, 10] \subseteq \mathbb{K}, \quad \text{let } f : M \rightarrow \mathbb{B} : k \mapsto \exists^1 l \in M : k = -l$$

$$\forall k \in M : f(k) \quad \Longrightarrow \quad \mathfrak{p}(\forall k \in M : f(k)) > 99\%$$

In the example above, the *for-probably-all*-quantifier assumes, that with a probability of over 99%, each element $k \in M$ has an additive inverse element $l \in M$, for which the equation $k = -l$ is true.

Only with the use of highly sophisticated tools like <http://www.wolframalpha.com/>, one can decide whether the \forall - or \exists^+ -quantifier shall be used instead of the $\forall \forall$ -quantifier (hence also the symbol of the W and A , which can be interpreted as the initials of the name *WolframAlpha*).

8 The *troll*-quantifier

The so-called *troll*-quantifier is an variation of the *exists*-quantifier, which states that the the given element usually exists *but* one shall note, that it may also not exist:

$$\text{let } M \subseteq \mathbb{K}, \quad \text{let } f : M \rightarrow \mathbb{B}$$

$$\perp k \in M : f(k) \quad :\Longleftrightarrow \quad \exists k \in M : f(k) \quad (\vee \nexists k \in M : f(k))$$

The probability that the *exists*-quantifier is true, when the *troll*-quantifier has been proven true is called the *magnitude of the troll-quantifier* and is usually written as $t \in (0.5, 1]$. The magnitudes t and p are equal as shown as follows:

$$\forall A \implies \perp A \qquad p(\forall A) = t(\exists A) \qquad (A \in \mathbb{B})$$

The quantifiers can be ordered as follows (according to their implicity):

$$\forall A \implies \forall \forall A \implies \exists^+ A \implies \exists A \implies \perp A \qquad (A \in \mathbb{B})$$

9 Polynomial final solution

A polynomial $K \in \mathbb{K}[X]$ has a *final solution* $\mathcal{F}_x(K)$, if a graphic calculator yields the result 10 (or sometimes 9.5) after the polynomial has been typed. The procedure to find the *final solution* $\mathcal{F}(K)$ is unknown to mankind, yet some gifted mathematicians have found the way to calculate the *final solution* $\mathcal{F}_x(K)$ using a graphic calculator from the TEXAS INSTRUMENTS TI-83-series. A possible procedure of solution could look as follows:

$$\begin{aligned} \text{let } K \in \mathbb{K}[X], \quad \text{let } S \in \mathbb{K} \\ \exists \mathcal{F}_x(K) = \langle 9.5 \dots 10 \rangle \subsetneq \mathbb{K} \quad :\iff \quad \exists S \in \text{img alt="TI-83 calculator icon" data-bbox="528 433 562 480"} (K) : ??? \end{aligned}$$

10 The *blah blah blah*-operator

The so-called *blah blah blah-operator* (also sometimes known as the *dot-notation*) is a mathematical operator, which indicates the continuation of a previously given expression (modified by given contextual parameters). The following equations are examples for the usage of the *blah blah blah-operator*:

$$\begin{aligned} a_1 + a_2 + a_3 + \dots \quad :\iff \quad \sum_{k=1}^{\infty} b_k \\ f(x) = \langle \sin, \cos, -\sin, -\cos, \dots \rangle(x) \quad :\iff \quad f(x) = \langle (\sin(x))^{(n)} \mid n \in \mathbb{N}_0 \rangle \\ \vdots \qquad \qquad \qquad \vdots \end{aligned}$$

The *blah blah blah-operator* may also be used inside a *set-notation* as follows:

$$\begin{aligned} [a \dots b] \quad :\iff \quad \{x \in \mathbb{K} \mid a \leq x \leq b : a, b \in \mathbb{K}\} \\ [a_k, a_{k+1}, a_{k+2}, \dots] \quad :\iff \quad \{a_i \in \mathbb{K} \mid i \geq k : i, k \in \mathbb{N}_0\} \\ \vdots \qquad \qquad \qquad \vdots \end{aligned}$$

11 Unsort-sort

The *unsort-sort* algorithm is a sorting algorithm, which uses a rather philosophical approach to array/list sorting:

The sorting algorithm is initialized with a ordered set $S_0 \subseteq \mathbb{K}$, which will be defined as ORDERED list (independent from the 'real' value order of the list's elements). Any further given set $S_n \subseteq \mathbb{K}$ will be sorted based on the given definition of ORDER.

The *unsort-sort* algorithm may be visualized by the following example:

let $\{3, 1, 5, 4, 2\} =: S_0 \subseteq \langle 1, 5 \rangle =: K \subsetneq \mathbb{N}$ a ordered set and let $f : \mathcal{P}(\mathbb{K}) \mapsto \mathcal{P}(\mathbb{K})$ be the sorting algorithm

TODO

12 Interval constraints

A *interval constraint* or *set constraint* is usually applied to an interval- or range-notation and consists itself of a set, field or range. The constraint indicates that the target set is constructed only using the elements of the given constraint. e.g.:

$$[1..5]_{\mathbb{N}} := \{1, 2, 3, 4, 5\} \neq \{1, 2, 0, 1, 2\} =: [1..5]_{\mathbb{F}_3}$$

The formal definition of an *interval/range/set/field constraint* is:

$$\begin{aligned} &\text{let } \mathbb{K}_0 \text{ a field and } \mathcal{P}(\mathbb{K}_0) \ni K \subseteq \mathbb{K}_0 \text{ a subset of the field } \mathbb{K}_0 \\ &[a..b]_K = [a, b]_K := \{k \in K \mid a \leq k \leq b\} \quad (a, b \in K, a \leq b) \end{aligned}$$

$$\begin{aligned} &\text{let } \mathbb{K}_1 \subseteq \mathbb{K}_0 \text{ be another field and } \mathcal{P}(\mathbb{K}_1) \ni L \subseteq \mathbb{K}_1 \text{ a subset of the field } \mathbb{K}_1 \\ &K_L := \{k \in K \mid k \in L\} = K \cap L \subseteq \mathbb{K}_0 \end{aligned}$$

The the target set, range, interval or field must be a subset of the *constraint*, rendering the following statement mathematically incorrect:

$$(-6, -1]_{\mathbb{N}^+}$$

13 <to be defined>

<to be defined>