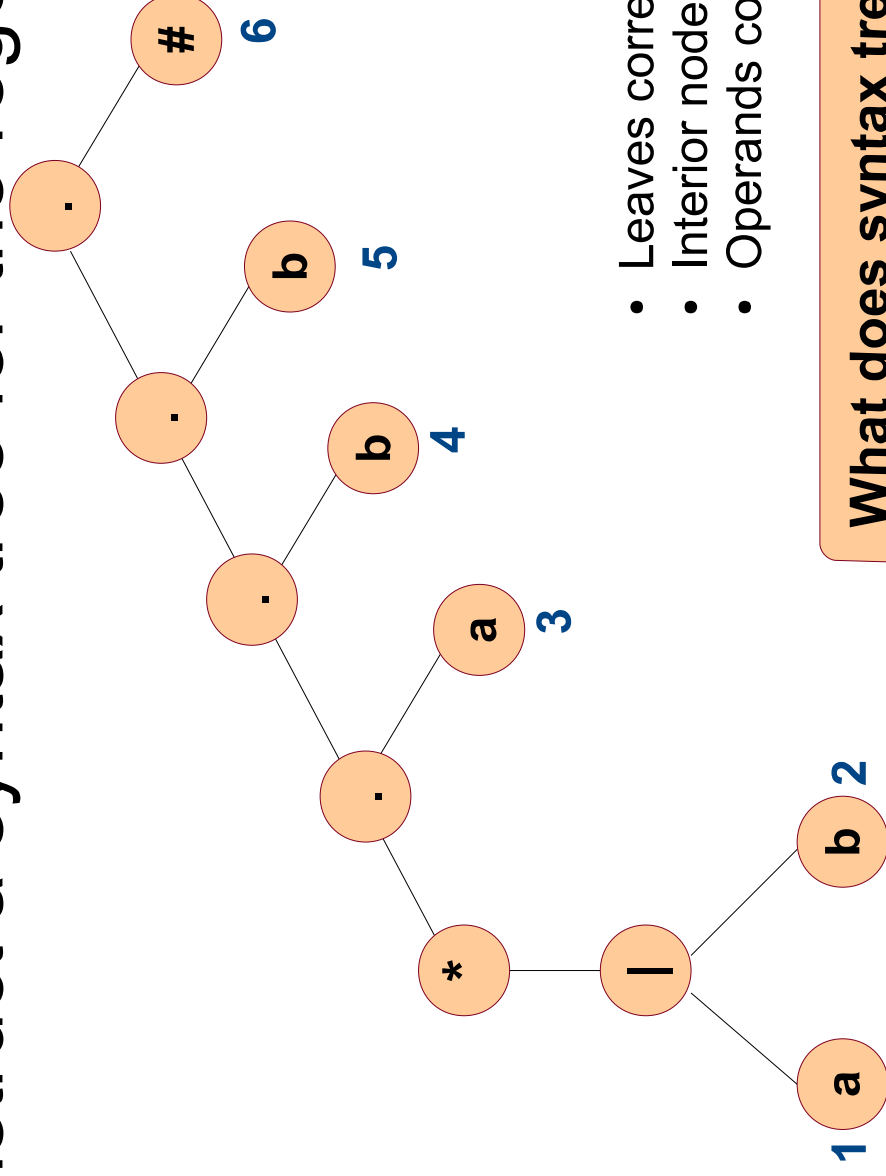


Regex → DFA

- Regex is $(a|b)^*abb\#$.
- Construct a syntax tree for the regex.



- Leaves correspond to operands.
- Interior nodes correspond to operators.
- Operands constitute strings.

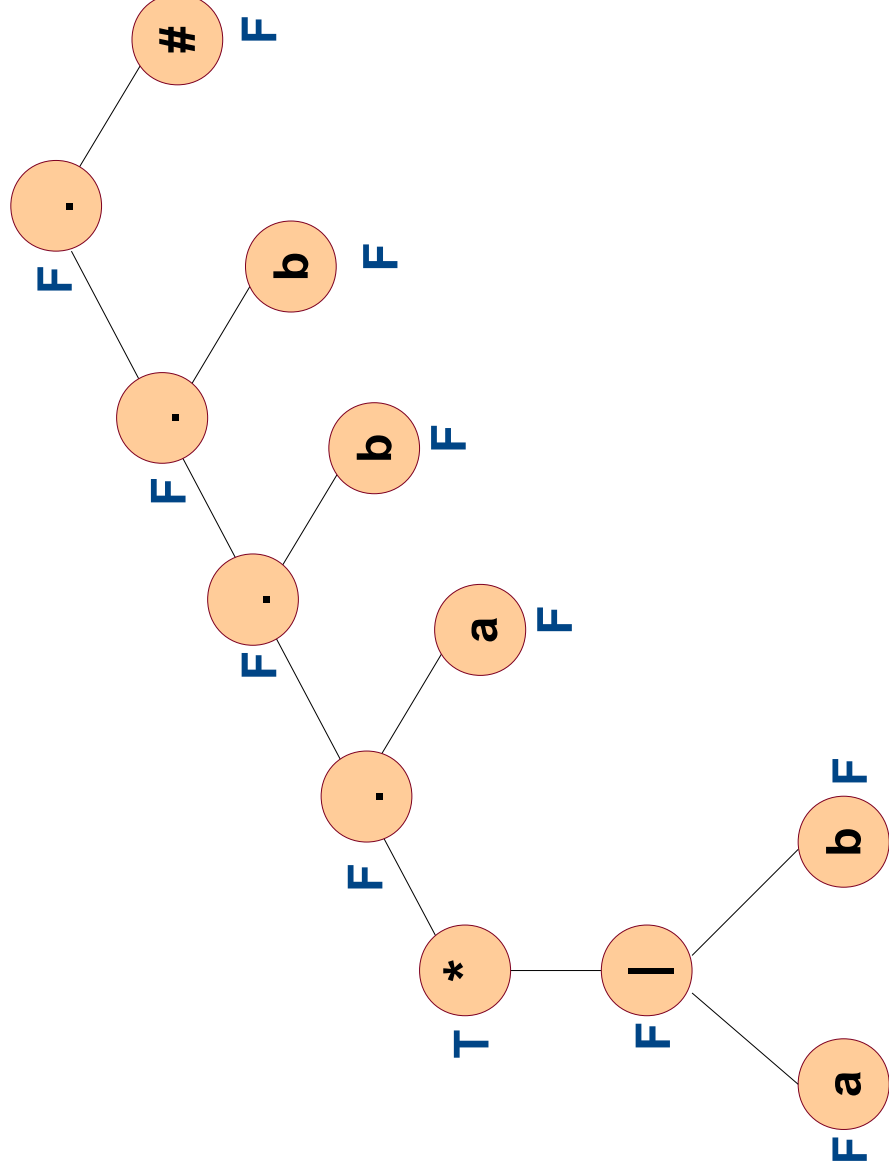
What does syntax tree for regex indicate?

Functions from Syntax Tree

- For a syntax tree node n
 - *nullable*(n): true if n represents ϵ .
 - *firstpos*(n): set of positions that correspond to the first symbol of strings in n 's subtree.
 - *lastpos*(n): set of positions that correspond to the last symbol of strings in n 's subtree.
 - *followpos*(n): set of next possible positions from n for valid strings.

nullable

- Regex is $(a|b)^*abb\#$.



nullable

Node n	$\text{nullable}(n)$
leaf labeled ϵ	true
leaf with position i	false
or-node $n = c1 \mid c2$	$\text{nullable}(c1)$ or $\text{nullable}(c2)$
cat-node $n = c1c2$	$\text{nullable}(c1)$ and $\text{nullable}(c2)$
star-node $n = c^*$	true

Classwork: Write down the rules for $\text{firstpos}(n)$.

firstpos

Node n	firstpos(n)
leaf labeled ϵ	$\{ \}$
leaf with position i	$\{i\}$
or-node $n = c1 \mid c2$	$\text{firstpos}(c1) \cup \text{firstpos}(c2)$
cat-node $n = c1c2$	
star-node $n = c^*$	$\text{firstpos}(c)$

firstpos

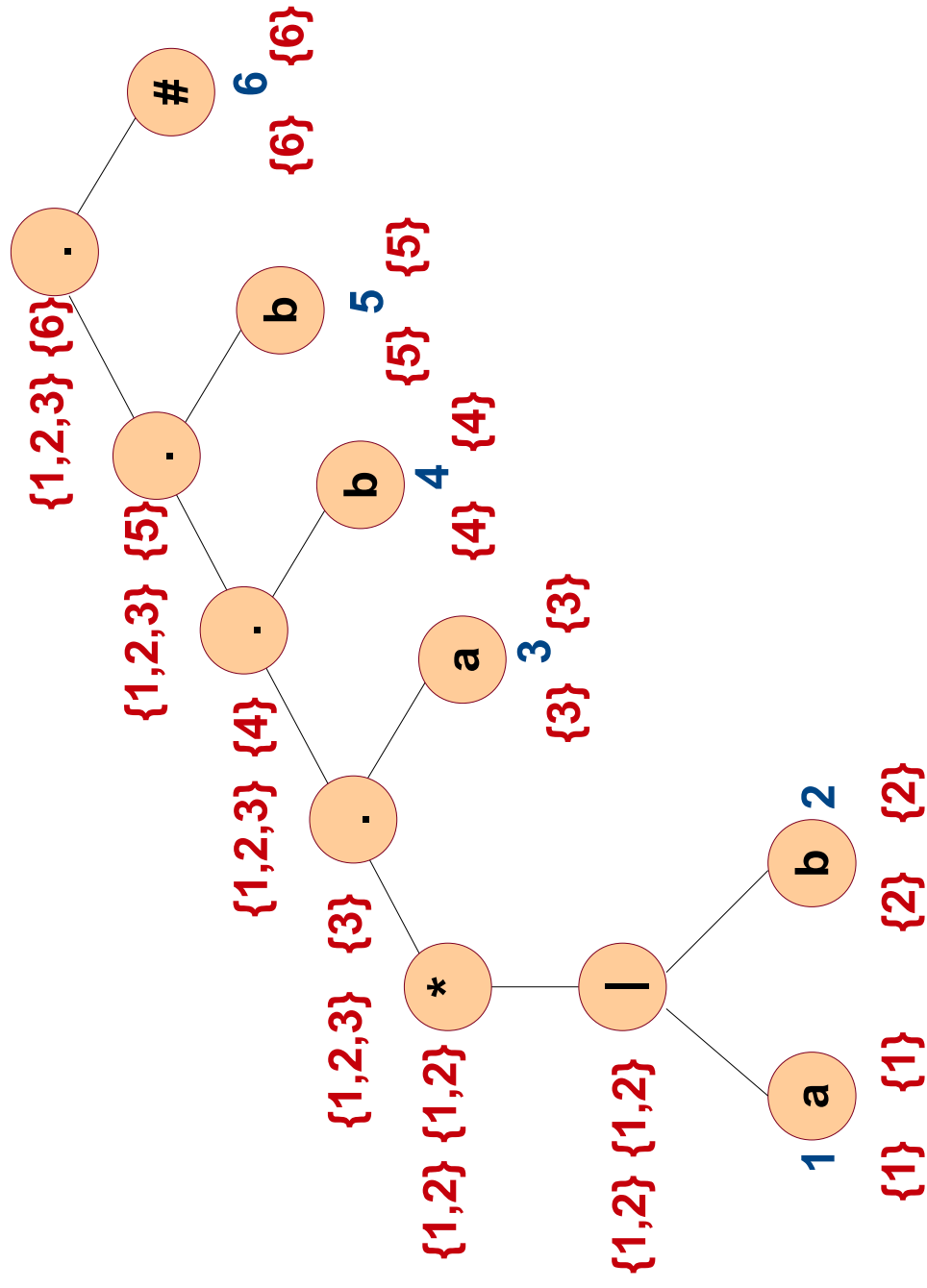
Node n	firstpos(n)
leaf labeled ϵ	$\{ \}$
leaf with position i	$\{i\}$
or-node $n = c1 \mid c2$	$\text{firstpos}(c1) \cup \text{firstpos}(c2)$
cat-node $n = c1c2$	if (nullable($c1$)) $\text{firstpos}(c1) \cup \text{firstpos}(c2)$ else $\text{firstpos}(c1)$
star-node $n = c^*$	$\text{firstpos}(c)$

Classwork: Write down the rules for lastpos(n).

lastpos

Node n	lastpos(n)
leaf labeled ϵ	$\{\}$
leaf with position i	$\{i\}$
or-node $n = c1 \mid c2$	$\text{lastpos}(c1) \cup \text{lastpos}(c2)$
cat-node $n = c1c2$	if (nullable(c2)) $\text{lastpos}(c1) \cup \text{lastpos}(c2)$ else $\text{lastpos}(c2)$
star-node $n = c^*$	$\text{lastpos}(c)$

firstpos lastpos

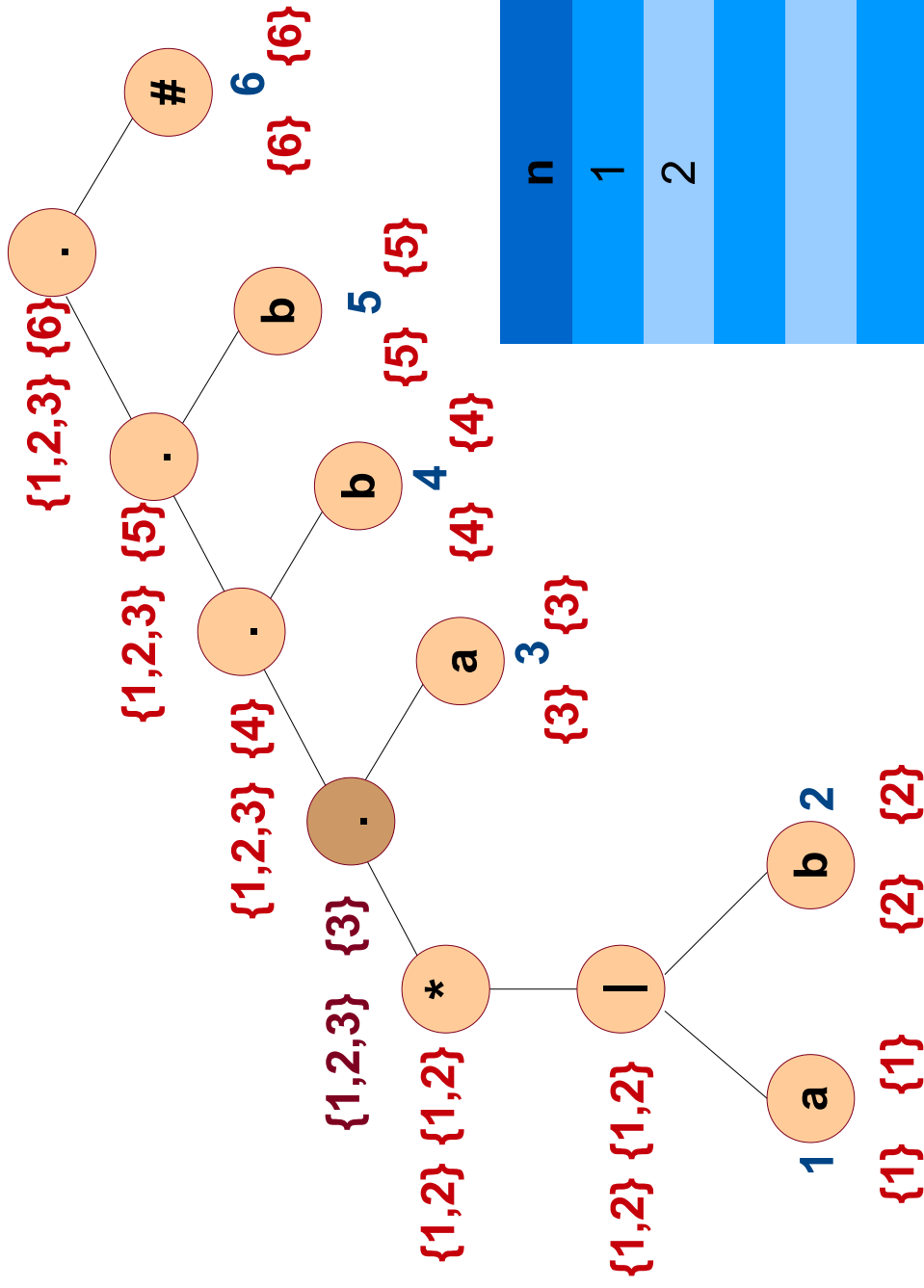


followpos

- *followpos*(n): set of next possible positions from n for valid strings.
 - If n is a **cat-node** with child nodes c1 and c2, then for each position in *lastpos*(c1), all positions in *firstpos*(c2) *follow*.
 - If n is a **star-node**, then for each position in *lastpos*(n), all positions in *firstpos*(n) *follow*.

followpos

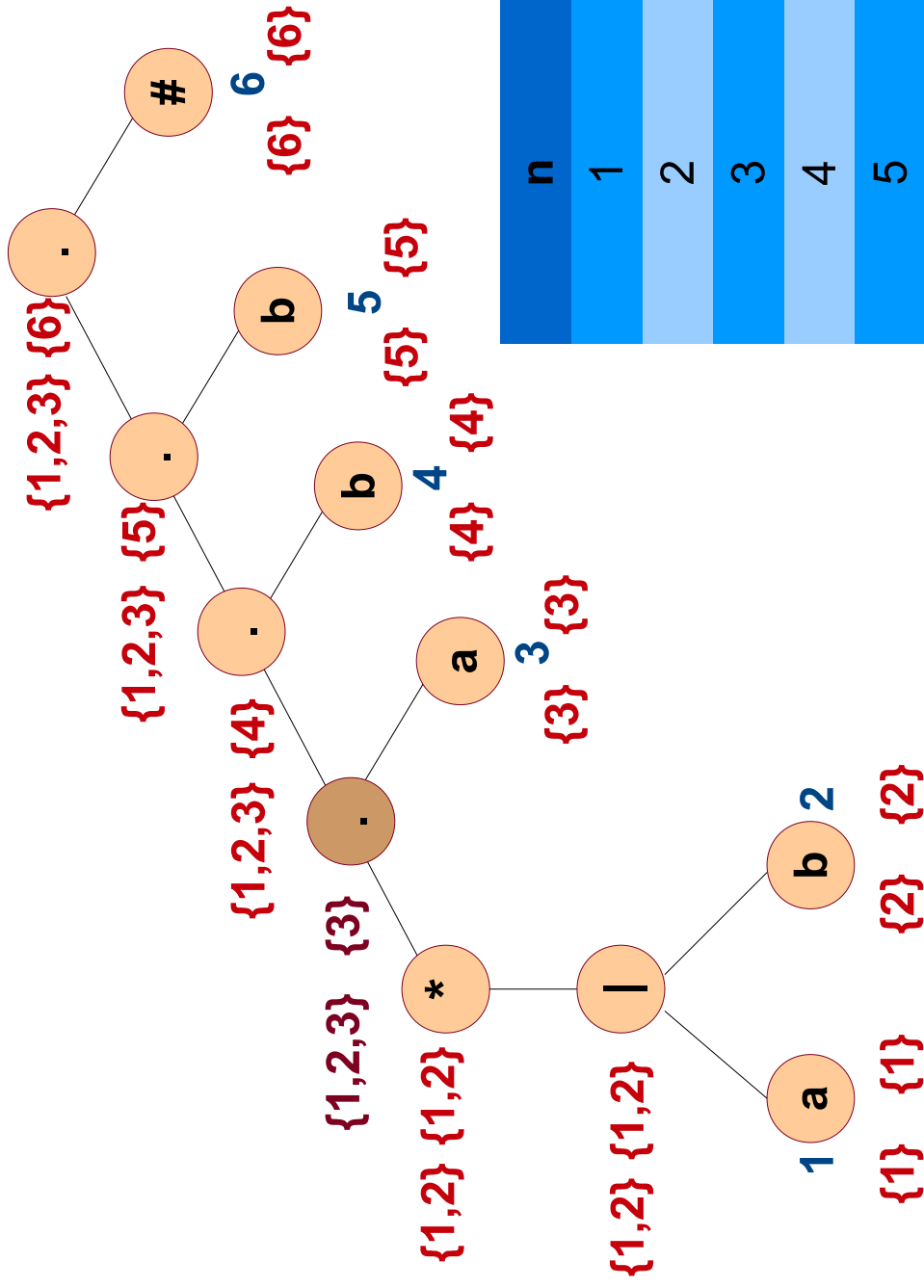
If n is a **cat-node** with child nodes $c1$ and $c2$, then for each position in $lastpos(c1)$, all positions in $firstpos(c2)$ follow.



n	followpos(n)
1	{3}
2	{3}

followpos

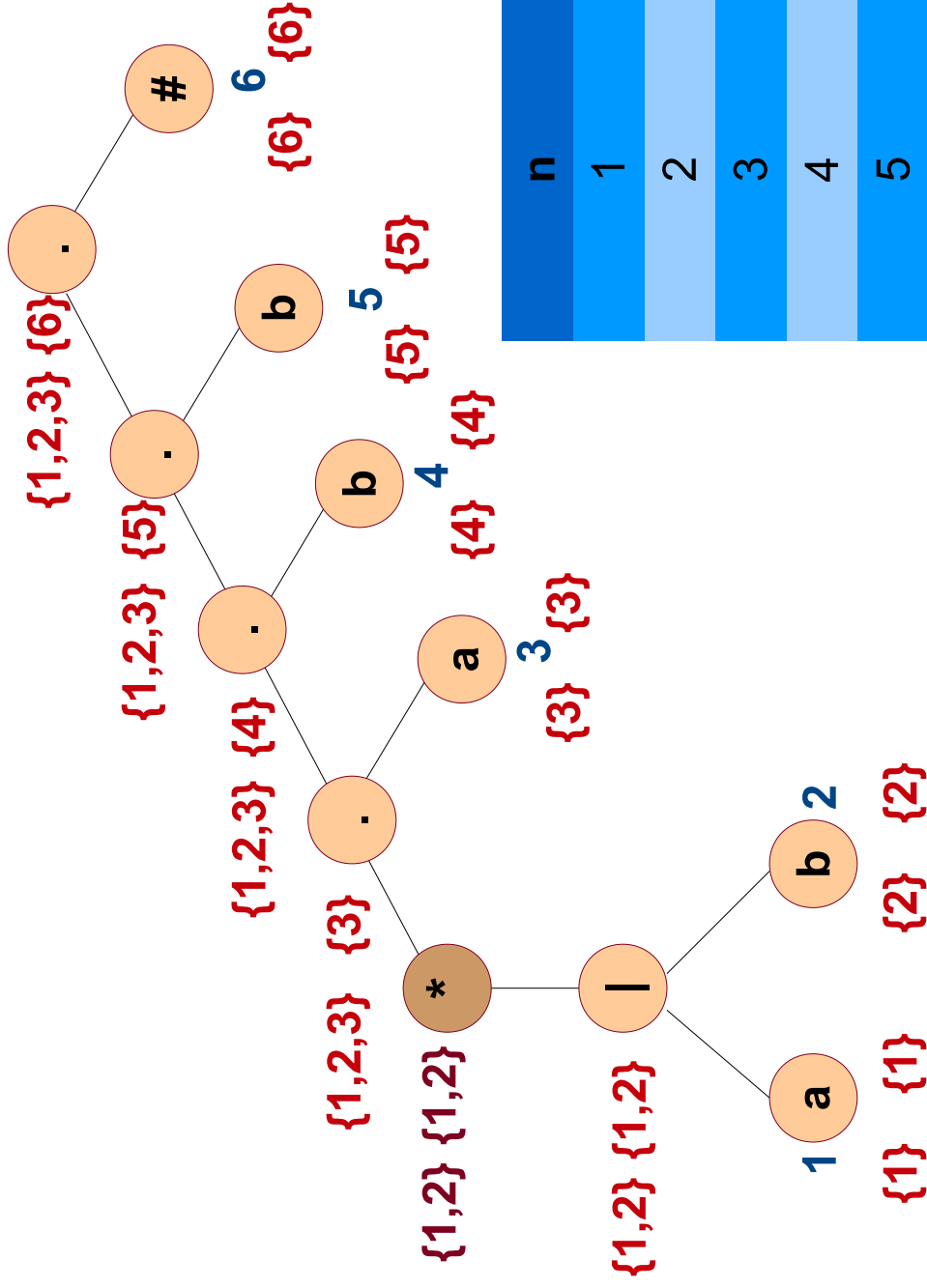
If n is a **cat-node** with child nodes $c1$ and $c2$, then for each position in $lastpos(c1)$, all positions in $firstpos(c2)$ follow.



n	followpos(n)
1	{3}
2	{3}
3	{4}
4	{5}
5	{6}
6	{ }

followpos

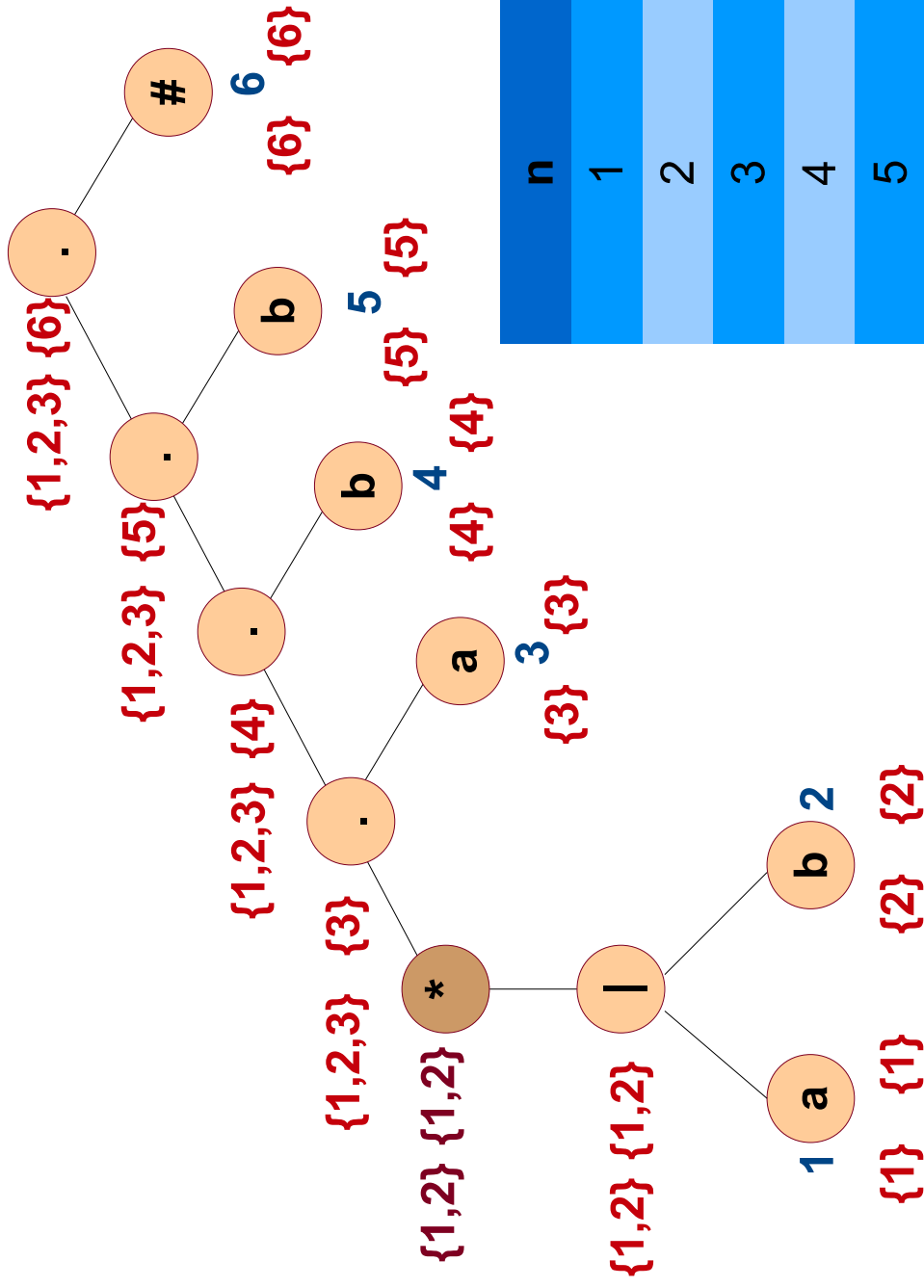
If n is a **star-node**, then for each position in $lastpos(n)$, all positions in $firstpos(n)$ *follow*.



n	$followpos(n)$
1	{3}
2	{3}
3	{4}
4	{5}
5	{6}
6	{ }

followpos

If n is a **star-node**, then for each position in $lastpos(n)$, all positions in $firstpos(n)$ *follow*.



n	$followpos(n)$
1	{3, 1, 2}
2	{3, 1, 2}
3	{4}
4	{5}
5	{6}
6	{ }

Regex \rightarrow DFA

1. Construct a syntax tree for regex#.
2. Compute *nullable*, *firstpos*, *lastpos*, *followpos*.
3. Construct DFA using transition function (*next slide*).
4. Mark *firstpos(root)* as start state.
5. Mark states that contain position of # as accepting states.

DFA Transitions

create unmarked state *firstpos*(*root*).

while there exists unmarked state *s* {

mark *s*

for each input symbol *a* {

uf = *U followpos*(*p*) where *p* is in *s* labeled *a*

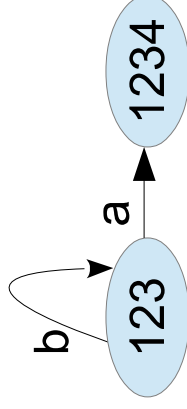
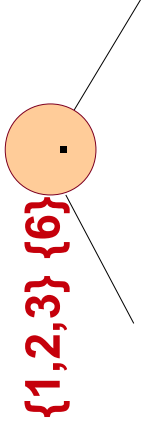
transition[*s*, *a*] = *uf*

if *uf* does not exist

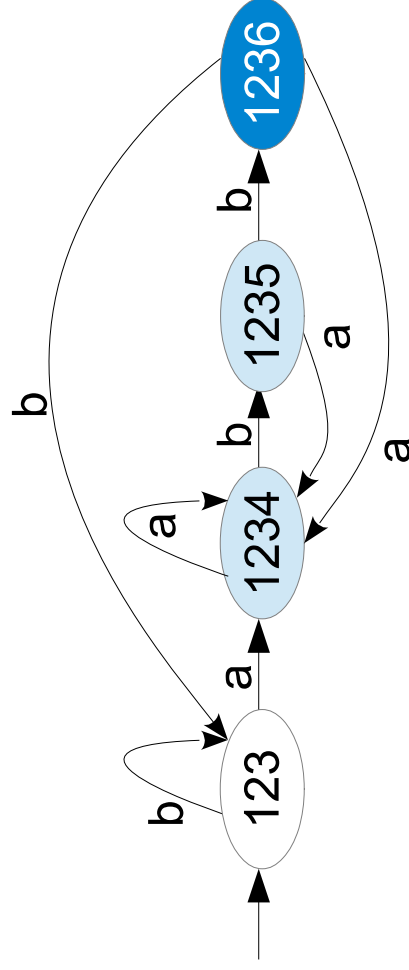
unmark *uf*

}

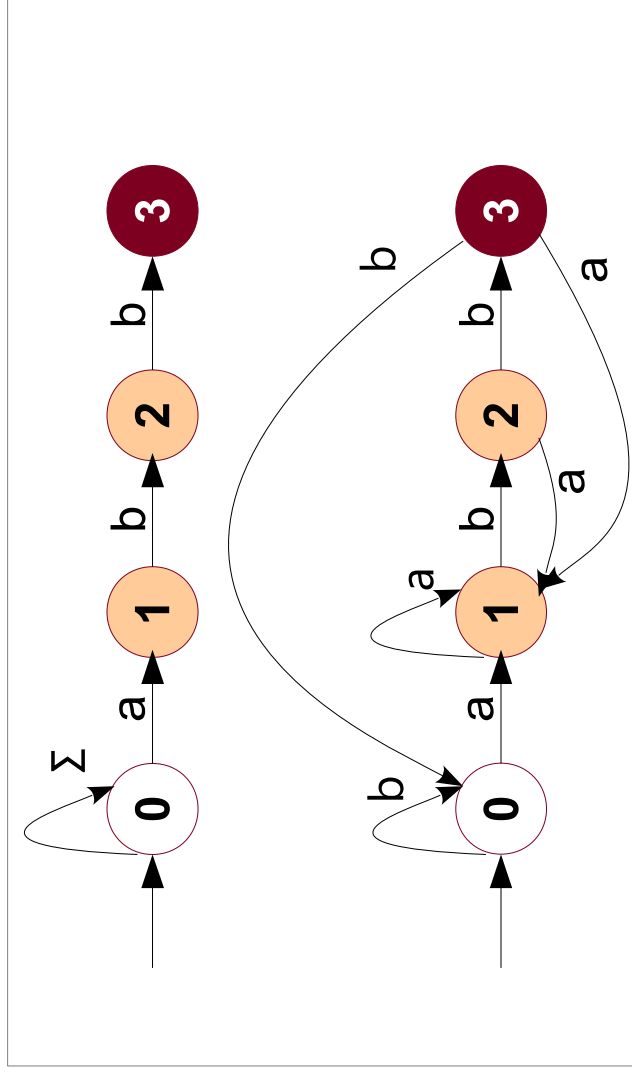
}



Final DFA



DFA



NFA

DFA

In case you are wondering...

- What to do with this DFA?
 - Recognize strings during lexical analysis.
 - Could be used in utilities such as *grep*.
 - Could be used in regex libraries as supported in php, python, perl, ... and Vipin's Ruby.