Practical 3 Compiler Construction

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Aim:

Write a program to find first(), and follow() set for each non-terminal of given grammar.

Input Production Rules:

```
E -> BA
A -> aB A| #
B -> D C
C -> +D C | #
D -> bEc | i
```

C code:

```
#include <stdio.h>
#include <ctype.h>
#include <string.h>
void followfirst(char, int, int);
void follow(char c);
void findfirst(char, int, int);
int count, n = 0;
char calc_first[10][100];
char calc_follow[10][100];
int m = 0;
char production[10][10];
char f[10], first[10];
int k;
char ck;
int e;
int main(int argc, char **argv)
   int jm = 0;
    int km = 0;
    int i, choice;
    count = 8;
   // The Input grammar
    strcpy(production[0], "E=BA");
```

```
strcpy(production[1], "A=aBA");
strcpy(production[2], "A=#");
strcpy(production[3], "B=DC");
strcpy(production[4], "C=+DC");
strcpy(production[5], "C=#");
strcpy(production[6], "D=bEc");
strcpy(production[7], "D=i");
int kay;
char done[count];
int ptr = -1;
// Initializing the calc_first array
for (k = 0; k < count; k++)</pre>
    for (kay = 0; kay < 100; kay++)
        calc_first[k][kay] = '!';
    }
int point1 = 0, point2, xxx;
for (k = 0; k < count; k++)
    c = production[k][0];
    point2 = 0;
    xxx = 0;
    // Checking if First of c has
    for (kay = 0; kay <= ptr; kay++)</pre>
        if (c == done[kay])
            xxx = 1;
    if (xxx == 1)
        continue;
    // Function call
    findfirst(c, 0, 0);
    ptr += 1;
    // Adding c to the calculated list
    done[ptr] = c;
    printf("\n First(%c) = { ", c);
    calc_first[point1][point2++] = c;
    // Printing the First Sets of the grammar
    for (i = 0 + jm; i < n; i++)
```

```
int lark = 0, chk = 0;
       for (lark = 0; lark < point2; lark++)</pre>
        {
           if (first[i] == calc_first[point1][lark])
               chk = 1;
               break;
            }
       if (chk == 0)
           printf("%c, ", first[i]);
           calc_first[point1][point2++] = first[i];
   printf("}\n");
   jm = n;
   point1++;
printf("\n");
printf("----\n\n");
char donee[count];
ptr = -1;
// Initializing the calc_follow array
for (k = 0; k < count; k++)</pre>
   for (kay = 0; kay < 100; kay++)
       calc_follow[k][kay] = '!';
    }
point1 = 0;
int land = 0;
for (e = 0; e < count; e++)
   ck = production[e][0];
   point2 = 0;
   xxx = 0;
   // Checking if Follow of ck
   for (kay = 0; kay <= ptr; kay++)</pre>
       if (ck == donee[kay])
           xxx = 1;
```

```
if (xxx == 1)
            continue;
        land += 1;
        follow(ck);
        ptr += 1;
        // Adding ck to the calculated list
        donee[ptr] = ck;
        printf(" Follow(%c) = { ", ck);
        calc_follow[point1][point2++] = ck;
        // Printing the Follow Sets of the grammar
        for (i = 0 + km; i < m; i++)
            int lark = 0, chk = 0;
            for (lark = 0; lark < point2; lark++)</pre>
                if (f[i] == calc_follow[point1][lark])
                    chk = 1;
                    break;
            if (chk == 0)
                printf("%c, ", f[i]);
                calc_follow[point1][point2++] = f[i];
        }
        printf(" }\n\n");
        km = m;
        point1++;
    }
void follow(char c)
    int i, j;
    // Adding "$" to the follow
    // set of the start symbol
    if (production[0][0] == c)
        f[m++] = '$';
```

```
for (i = 0; i < 10; i++)
        for (j = 2; j < 10; j++)
            if (production[i][j] == c)
                if (production[i][j + 1] != '\0')
                    followfirst(production[i][j + 1], i, (j + 2));
                }
                if (production[i][j + 1] == '\0' \&\& c != production[i][0])
                    // Calculate the follow of the Non-Terminal
                    follow(production[i][0]);
void findfirst(char c, int q1, int q2)
    int j;
    // The case where we encounter a Terminal
    if (!(isupper(c)))
        first[n++] = c;
    for (j = 0; j < count; j++)</pre>
        if (production[j][0] == c)
            if (production[j][2] == '#')
                if (production[q1][q2] == '\0')
                    first[n++] = '#';
                else if (production[q1][q2] != '\0' && (q1 != 0 || q2 != 0))
                {
                    // Non-Terminal we encounter after epsilon
                    findfirst(production[q1][q2], q1, (q2 + 1));
                else
```

```
first[n++] = '#';
            else if (!isupper(production[j][2]))
                first[n++] = production[j][2];
                // encounter at the beginning
                findfirst(production[j][2], j, 3);
void followfirst(char c, int c1, int c2)
    int k;
   // The case where we encounter
    if (!(isupper(c)))
       f[m++] = c;
    else
        int i = 0, j = 1;
        for (i = 0; i < count; i++)
            if (calc_first[i][0] == c)
                break;
        }
        // Including the First set of the Non-Terminal in the Follow of
        // the original query
        while (calc_first[i][j] != '!')
            if (calc_first[i][j] != '#')
                f[m++] = calc_first[i][j];
            else
                if (production[c1][c2] == '\0')
                    follow(production[c1][0]);
```

Output:

```
First(E) = { b, i, }

First(A) = { a, #, }

First(B) = { b, i, }

First(C) = { +, #, }

First(D) = { b, i, }

Follow(E) = { $, c, }

Follow(A) = { $, c, }

Follow(B) = { a, $, c, }

Follow(C) = { a, $, c, }

Follow(D) = { +, a, $, c, }

PS C:\Users\unnat\OneDrive\Desktop\CC>
```

Conclusion:

From this practical, we learnt about first and follow of grammar.