ELEC0152 Software Development Report

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This report presents the development of "Fruit Catcher," a real-time embedded game implemented on the NUCLEO-F401RE board using the STM32F401RE microcontroller. The system integrates an SSD1306 OLED display and an analog joystick within a FreeRTOS architecture using the STM32 HAL. The software design employs a thread-safe, two-task model synchronized via mutexes: a high-priority task manages sensor acquisition using ADC with Direct Memory Access (DMA), while a normal-priority task handles game physics and I2C rendering. This project demonstrates the application of RTOS principles to ensure data integrity and timing determinism, validated through Segger SystemView analysis.

CCS Concepts: • **Computer systems organization** → **Embedded systems**; **Real-time operating systems**; • **Software and its engineering** → *Embedded software.*

Additional Key Words and Phrases: RTOS, STM32, I2C, HAL, Mutex, DMA, Segger SystemView, SSD1306

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1 Introduction

The objective of this project was to develop a real-time embedded game, "Fruit Catcher," on the STM32 NUCLEO-F401RE development board. The software architecture is implemented in C, utilizing the STM32 Hardware Abstraction Layer (HAL)[5] for low-level peripheral control and FreeRTOS for task scheduling and resource management.

Using a Real-Time Operating System (RTOS) offers significant advantages over a standard "superloop" architecture for this application. In a superloop, the game logic, input polling, and screen rendering would run sequentially. Since the I2C screen update is a blocking operation that can take approximately 30ms, a superloop architecture would suffer from input lag, as button presses occurring during the render phase would be missed. By leveraging FreeRTOS, this project decouples the high-frequency input sampling from the lower-frequency rendering loop. A high-priority task manages inputs to ensure responsiveness, while a lower-priority task handles the game physics and display, synchronized via mutexes to prevent data tearing.

1.1 Game Concept

"Fruit Catcher" is a skill-based arcade game designed for a monochrome 128x64 pixel environment. The core entities in the game are:

- **The Basket:** Controlled by the player, constrained to the bottom axis of the screen.
- **Fruits:** Falling objects represented as filled 6x6 pixel squares. Collecting these increments the score.
- **Bombs:** Falling hazards represented as "X" shapes. Collision with a bomb results in immediate game termination.

The game logic relies on randomized spawning algorithms to ensure replayability, with objects spawning at random X-coordinates and falling at varying speeds.

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1.2 Game Play

The gameplay experience focuses on responsiveness and state management. The user interacts with the system using an analog joystick module.

- **Controls:** The player moves the joystick along the X-axis to move the basket (the RECTANGLE in the bottom) left or right to catch fruits (the SQUARE shape) to gain scores. The input is handled via analog-to-digital conversion to allow for smooth movement logic. The integrated push-button serves as a state toggle.
- **Game States:** The system implements a Finite State Machine (FSM):
  (1) **Playing:** The active state where physics updates and collision checks occur (Fig. 2).
  (2) **Paused:** Triggered by pressing the joystick button. The game loop freezes, and a "PAUSED" overlay is rendered (Fig. 3).
  (3) **Game Over:** Triggered by colliding with a bomb (in CROSS "X" shape). The screen displays the final score. Pressing the button in this state resets the score and respawns the entities to restart the game (Fig. 4).

Fig. 1. Game Initiation

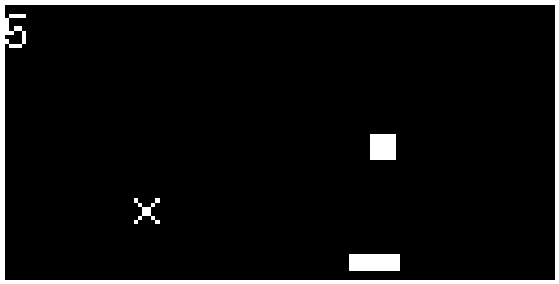

Fig. 2. Active Gameplay

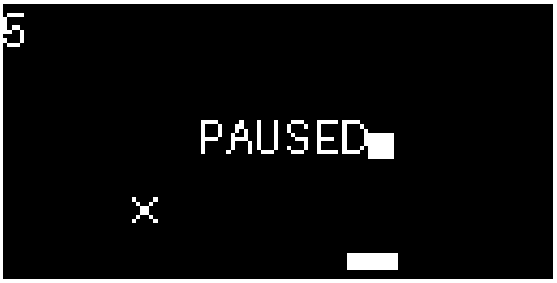

Fig. 3. Game Pause

Fig. 4. Game Over

2 Methodology

This section outlines the hardware integration and software architecture developed for the "Fruit Catcher" game. The system is designed to prioritize responsiveness and determinism by leveraging specific hardware peripherals on the STM32 microcontroller.

2.1 Hardware Setup

The physical system is built around the NUCLEO-F401RE development board, which houses the STM32F401RE microcontroller (ARM Cortex-M4)[6]. Two external modules are interfaced with the MCU: a SSD1306 OLED display for visual output and a generic dual-axis analog joystick for user input.

2.1.1 Display Interface. The SSD1306 OLED display (128x64 resolution) communicates via the I2C protocol. The STM32's I2C1 peripheral is configured in standard mode (100 kHz) to ensure stable data transmission. The display requires two data lines (SDA and SCL) along with power connections.

2.1.2 Input Interface. The joystick module provides three distinct signals: two analog voltages representing the X and Y axes, and one digital signal for the integrated push-button.

- **Analog Axes:** The VRX (Horizontal) and VRY (Vertical) pins are connected to the MCU's Analog-to-Digital Converter (ADC1). For this specific game implementation, primarily the X-axis is utilized for lateral basket movement.
- **Digital Button:** The switch (SW) pin is connected to a GPIO pin configured with an internal pull-up resistor. This ensures the pin reads a logical High (1) when released and Low (0) when pressed, eliminating the need for external resistors.

Table 1. Hardware Connections

(a) OLED Connections

Pin	STM32 Pin	Function
SCL	PB8	I2C1_SCL
SDA	PB9	I2C1_SDA
VCC	5V	Power
GND	GND	Ground

(b) Joystick Connections

Pin	STM32 Pin	Function
VRX	PA0	ADC1_IN0
VRY	PA1	ADC1_IN1
SW	PA9 (D8)	GPIO Input
+5V(VCC)	3.3V	Power
GND	GND	Ground

2.2 Software Development

The system architecture is built upon the FreeRTOS kernel, which manages the scheduling of concurrent tasks and creates a layer of abstraction between the application logic and the hardware[2]. Fig. 5 illustrates the data flow and synchronization mechanisms implemented in the design.

2.2.1 Architectural Diagram. The software architecture leverages FreeRTOS to implement a responsive producer-consumer model, effectively decoupling high-frequency sensor acquisition from the game rendering loop. As illustrated in Fig. 5, the system coordinates concurrent tasks through three specific synchronization mechanisms to maintain data integrity:

- A Mutex for global game state access.
- A Binary Semaphore (`ButtonSem`) for input signaling.
- A Message Queue (`LedQueue`) for collision events.

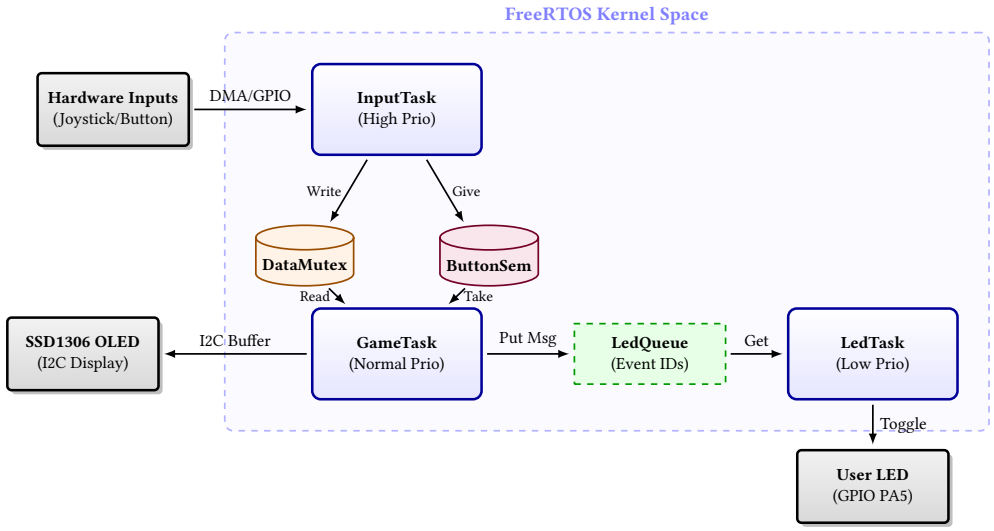

Fig. 5. Complete System Architecture. The InputTask acquires sensor data and signals the GameTask via Semaphore. The GameTask manages physics and offloads feedback events to the LedTask via Queue to prevent blocking.

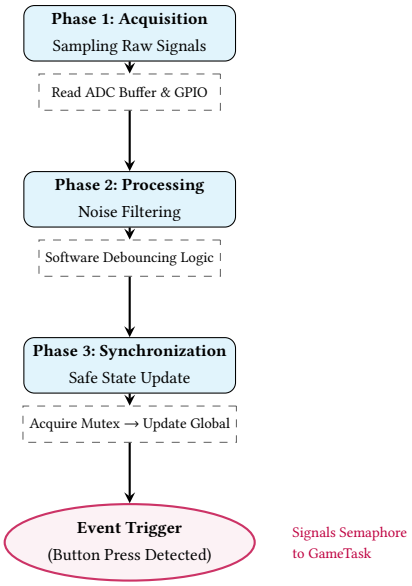

Fig. 6. Conceptual Logic of the InputTask. High-speed signal processing and noise filtering.

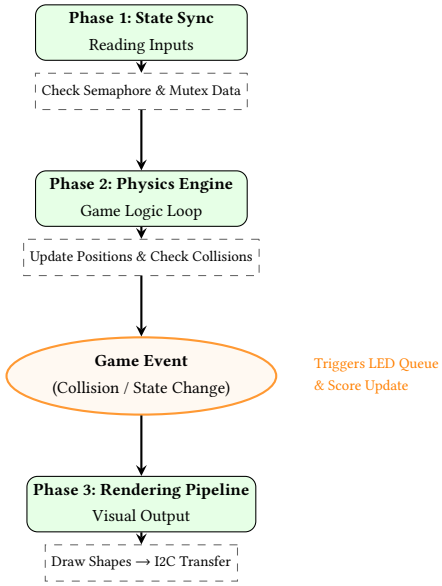

Fig. 7. Conceptual Logic of the GameTask. Cyclic engine for physics, events, and rendering.

2.2.2 Joystick Input Implementation. The joystick input subsystem is designed to provide low-latency, non-blocking control by offloading data acquisition to the microcontroller's hardware peripherals. The implementation manages two distinct signal types: analog position data (X/Y axes) and digital state data (push-button). The overall logic flow corresponds to the *Acquisition* and *Processing* phases illustrated in Fig. 6.

ADC and DMA Configuration. To handle the analog signals from the joystick's potentiometers without burdening the CPU, the STM32's Analog-to-Digital Converter (ADC1) is utilized in conjunction with Direct Memory Access (DMA). The peripheral was configured via the `.ioc` interface with the following parameters:

- **Scan Conversion Mode:** Enabled, allowing the ADC to automatically switch between Channel 0 (PA0/X-Axis) and Channel 1 (PA1/Y-Axis) within a single conversion sequence.
- **Direct Memory Access (DMA):** The DMA controller is configured in **Circular Mode** with a data width of *Half Word* (16-bit)[7].

This configuration establishes a continuous background data pipeline. As shown in the code snippet below, the DMA engine is activated once during system initialization in `main.c`. From that point forward, the hardware automatically writes the latest conversion results into the `adcValues` array, ensuring the `InputTask` always has access to the most recent data without needing to trigger conversions manually.

```
// Core/Src/main.c
// Start ADC1 in DMA mode to fill the buffer automatically
// adcValues[0] = X-Axis, adcValues[1] = Y-Axis
if (HAL_ADC_Start_DMA(&hadc1, (uint32_t*)adcValues, 2) != HAL_OK) {
    Error_Handler();
}
```

Digital Input Configuration. The joystick's integrated switch is connected to pin PA9. To minimize external component count, the pin is configured as a GPIO Input with the internal Pull-Up resistor enabled. This ensures the pin remains in a stable High state (Logic 1) when the button is open, transitioning to Low (Logic 0) only when physically pressed to ground.

Data Processing and Interpretation. The `InputTask` implements the signal processing logic shown in Phase 2 of Fig. 6. Instead of querying hardware registers, the task simply reads the global `adcValues` buffer populated by the DMA.

The raw 12-bit ADC values (ranging from 0 to 4095) are interpreted using a threshold-based approach to create a "Dead Zone" that filters out mechanical noise and prevents drift when the stick is centered.

```
// Core/Src/freertos.c - InputTask Logic
uint16_t raw_x = adcValues[0]; // Read from DMA buffer

// Threshold logic for digital-like movement
if (raw_x < 1000) {
    // Trigger Left Movement
}
else if (raw_x > 3000) {
    // Trigger Right Movement
}
// Values between 1000 and 3000 are ignored (Neutral)
```

This mapping ensures precise player control while accommodating the mechanical tolerances of the analog sensor. By decoupling the acquisition (DMA) from the interpretation (InputTask), the system achieves the high responsiveness depicted in the architectural diagram.

2.2.3 OLED Output Implementation. The visual output system is built upon a 128x64 pixel monochrome OLED display driven by the SSD1306 controller[4]. The implementation utilizes a "Frame Buffer" architecture to decouple the rendering logic from the physical data transmission, ensuring flicker-free animations and efficient I2C bus usage. The data flow corresponds to the *Rendering Pipeline* phase illustrated in Fig. 7.

I2C Peripheral Configuration. Communication with the display is established via the STM32's I2C1 peripheral. The configuration in the `.ioc` interface was selected to prioritize signal stability and compatibility with the specific OLED module:

- **Mode:** I2C Standard Mode (100 kHz). While Fast Mode (400 kHz) is possible, Standard Mode was chosen to ensure reliable data integrity over the jumper wire connections used in the prototype.
- **Pin Assignment:** PB8 is assigned to SCL (Serial Clock) and PB9 to SDA (Serial Data), utilizing the alternate function mappings of the microcontroller.
- **NVIC Settings:** I2C Event and Error interrupts are enabled to allow for future expansion into non-blocking DMA transfers, although the current implementation utilizes blocking mode for simplicity within the `GameTask`.

Frame Buffer and Rendering Strategy. Directly writing individual pixels to the OLED over I2C is inefficient due to protocol overhead. To mitigate this, the driver allocates a local RAM buffer of 1,024 bytes (128×64 bits/8).

All drawing operations—such as updating the basket position, drawing falling fruits, or rendering the score text—are performed on this local buffer. This allows the application to compose a complete frame in memory before committing it to the hardware.

```
// ssd1306.c - Simplified Drawing Logic
// Modifies the local RAM buffer, not the screen
void ssd1306_DrawPixel(uint8_t x, uint8_t y, SSD1306_COLOR color) {
    if (color == White) {
        SSD1306_Buffer[x + (y / 8) * SSD1306_WIDTH] |= (1 << (y % 8));
    } else {
        SSD1306_Buffer[x + (y / 8) * SSD1306_WIDTH] &= ~(1 << (y % 8));
    }
}
```

Screen Update and Data Transfer. The physical update of the display occurs at the end of the `GameTask` loop (Phase 3 in Fig. 7). The function `ssd1306_UpdateScreen()` initiates a bulk I2C transfer, flushing the entire contents of the RAM buffer to the OLED's Graphic Display Data RAM.

This "Draw-then-Flush" strategy prevents screen tearing (where half a frame is drawn over another) and ensures that the complex physics calculations do not interrupt the visual output stream.

```
// Core/Src/freertos.c - Rendering Sequence
ssd1306_Fill(Black);            // 1. Clear Buffer
Game_Draw(&player, &fruit);     // 2. Render Objects to Buffer
ssd1306_UpdateScreen();         // 3. Flush Buffer via I2C
```

2.2.4 FreeRTOS Tasks Implementation. The application logic is structured around the FreeRTOS real-time kernel, which orchestrates the execution of concurrent threads. This multitasking architecture allows the system to balance the

conflicting requirements of high-frequency input sampling (50 Hz) and computationally intensive graphical rendering (30 Hz). The system is decomposed into three distinct tasks, as summarized in Table 2.

Task Roles and Prioritization. The scheduler manages task execution based on a preemptive priority policy.

- **InputTask (High Priority):** This task is responsible for the critical path of user interaction. Assigned a priority of `osPriorityAboveNormal`, it can preempt the rendering loop at any moment. This ensures that button presses and joystick movements are captured deterministically every 20ms, eliminating input lag even when the I2C bus is saturated.
- **GameTask (Normal Priority):** Functioning as the main application engine, this task consumes the majority of CPU cycles. It manages the finite state machine, executes physics calculations, and drives the display updates. Its normal priority allows it to be interrupted by the `InputTask`, ensuring the system remains responsive.
- **LedTask (Low Priority):** This background worker handles visual feedback. By operating at low priority, it processes LED blink patterns without stalling the game physics.

Inter-Task Communication and Synchronization. To ensure data integrity and system stability, the tasks coordinate their actions through three specific synchronization mechanisms, as depicted in the architectural data flow (Fig. 5):

(1) **Shared State Protection (Mutex):** The global `JoystickData` structure acts as a shared memory resource between the Input and Game tasks. A Mutex (`DataMutex`) is employed to enforce mutual exclusion during read/write operations. This prevents "data tearing" anomalies where the game logic might inadvertently process an X-coordinate from the current frame combined with a Y-coordinate from the previous frame.

(2) **Event Signaling (Semaphore):** A Binary Semaphore (`ButtonSem`) is utilized to signal state transitions. When the `InputTask` detects a valid button press edge, it "gives" the semaphore. The `GameTask` checks for this semaphore at the start of each cycle to toggle between Playing and Paused states. This decoupled signaling mechanism avoids the need for complex global flags or polling within the game loop[1].

(3) **Asynchronous Processing (Message Queue):** To handle game events such as collisions, a Message Queue (`LedQueue`) is implemented. When a collision occurs, the `GameTask` posts an event ID to the queue and immediately continues execution. The `LedTask` consumes these messages asynchronously, performing blocking delays to blink the LED. This queue-based design isolates the real-time physics engine from the latency-heavy feedback logic.

Table 2. FreeRTOS Task and Synchronization Configuration

Task Name	Priority	Brief Description	Sync. Objects
`InputTask`	High	Polls buttons, reads DMA, updates global state	`DataMutex` (Write) `ButtonSem` (Give)
`GameTask`	Normal	Game loop, physics, collision, screen render	`DataMutex` (Read) `ButtonSem` (Take) `LedQueue` (Put)
`LedTask`	Low	Handles asynchronous LED blink events	`LedQueue` (Get)

3 Results and Analysis

The real-time behavior of the system was validated using SEGGER SystemView[3], which provides a cycle-accurate timeline of task execution and scheduler events. This analysis confirms that the system meets its timing requirements and demonstrates the correct operation of the synchronization mechanisms.

3.1 SEGGER SystemView and Timing Analysis

The real-time behaviour of the system was validated using SEGGER SystemView traces captured during active gameplay. The figures in this section provide direct, time-resolved evidence of task execution order, preemption, and the resulting CPU headroom.

3.1.1 Long-run scheduling stability and load distribution. Fig.8 presents a long-duration capture (tens of seconds), providing a macroscopic view of scheduler behaviour. The repeating execution pattern of `GameTask` (green) and `InputTask` (blue) remains consistent over the full interval, indicating that the system does not exhibit drift, runaway execution, or prolonged starvation of any task. The CPU load histogram in Fig. 8 is dominated by `GameTask`, while `InputTask` contributes a comparatively small fraction, consistent with the design goal that input sampling remains lightweight while rendering and game-state updates dominate runtime cost.

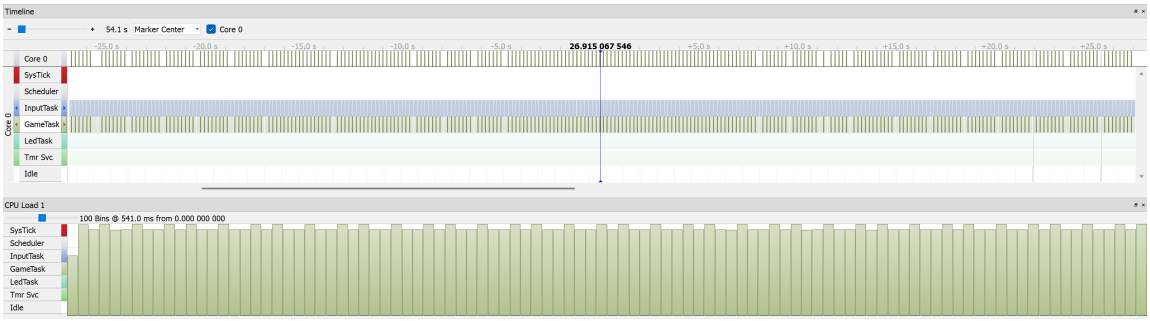

Fig. 8. Long-run SystemView trace and CPU load distribution. `GameTask` (green) dominates compute time over the capture window, while `InputTask` (blue) remains brief and periodic, indicating stable scheduling and consistent workload over time.

3.1.2 Frame-level execution profile and CPU headroom. A representative short-window overview is shown in Fig. 9. Within this interval, `GameTask` executes for approximately 19.1 ms, while `InputTask` executes for approximately 0.7692 ms plus a further 0.0739 ms burst. The remaining time is spent in `Idle` (approximately 20.1 ms). Summing the observed active execution time gives $\approx$ 19.943 ms out of $\approx$ 40.043 ms total, corresponding to an instantaneous utilisation of $\approx$ 49.8%, with $\approx$ 50.2% slack in `Idle`. This confirms that, during the captured gameplay segment, the CPU is not saturated and deadlines have substantial margin even when `GameTask` performs its longest operations.

3.1.3 Preemptive responsiveness and worst-case input latency. Fig. 10 provides direct evidence of correct preemptive scheduling. `GameTask` occupies a long execution region (labelled $\approx$ 19.9 ms), and `InputTask` becomes ready during this interval, executing immediately as a short burst ($\approx$ 13.6 μs followed by $\approx$ 73.9 μs). The execution order indicates that `InputTask` is scheduled without waiting for `GameTask` to complete, i.e., it preempts ongoing work rather than being delayed until the end of the frame. From a responsiveness perspective, this behaviour bounds the additional

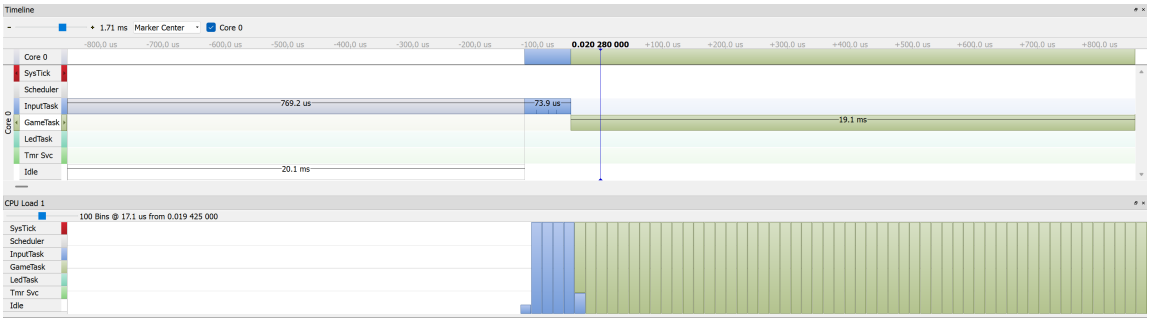

Fig. 9. Short-window overview of task execution. GameTask runs for ≈ 19.1 ms while InputTask executes in shorter bursts (≈ 0.769 ms and ≈ 0.074 ms), with the remaining time spent in Idle (≈ 20.1 ms), indicating significant CPU headroom.

software-induced latency on input processing to the scheduler/context-switch overhead plus the current instruction boundary, rather than the full duration of the render/update phase. In practical terms, user input sampling is protected from the long GameTask execution time, preventing perceptible control lag even when the frame update is expensive.

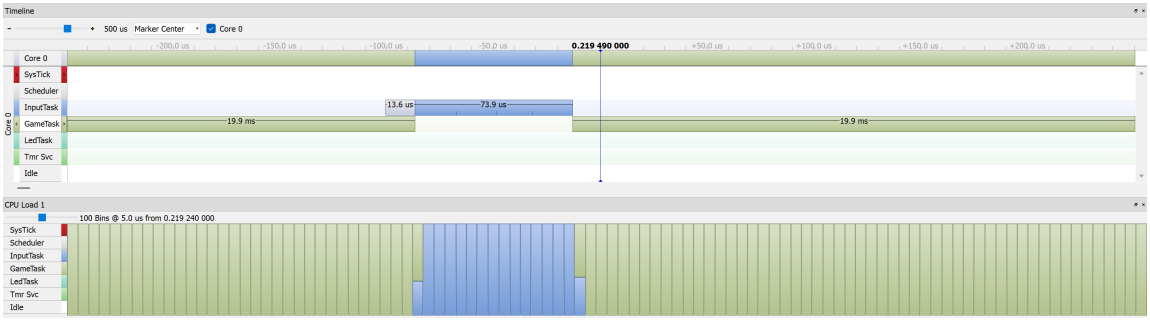

Fig. 10. Preemption behaviour during a long GameTask execution window (≈ 19.9 ms). InputTask executes immediately upon wake-up as short bursts (≈ 13.6 μs and ≈ 73.9 μs), demonstrating low-latency preemption.

3.1.4 InputTask execution time and computational efficiency. A zoomed capture focusing on InputTask is shown in Fig. 11. The trace annotates two consecutive segments of approximately 13.9 μs and 79.2 μs, giving a combined observed execution time of ≈ 93.1 μs for the input-processing burst in this window. Relative to the millisecond-scale execution of GameTask, this confirms that InputTask is computationally lightweight and well suited to high-frequency scheduling. The surrounding Idle regions (labelled ≈ 19.9 ms before the burst and ≈ 7.9 ms after) further indicate that input handling does not materially reduce available slack time, supporting the design objective of maintaining responsive controls with minimal CPU overhead.

3.1.5 Runtime variability across frames. Fig. 12 shows multiple consecutive GameTask executions across a wider time axis. Several GameTask blocks are labelled close to ≈ 19.9 ms, while other instances are shorter (e.g., ≈ 14.9 ms and ≈ 11.9 ms). This demonstrates that the per-frame workload is not constant, which is expected in interactive systems where rendering content and game-state complexity vary (e.g., different numbers of objects to draw or conditional logic paths). Critically, despite this variability, the pattern remains well structured (repeated execution blocks with

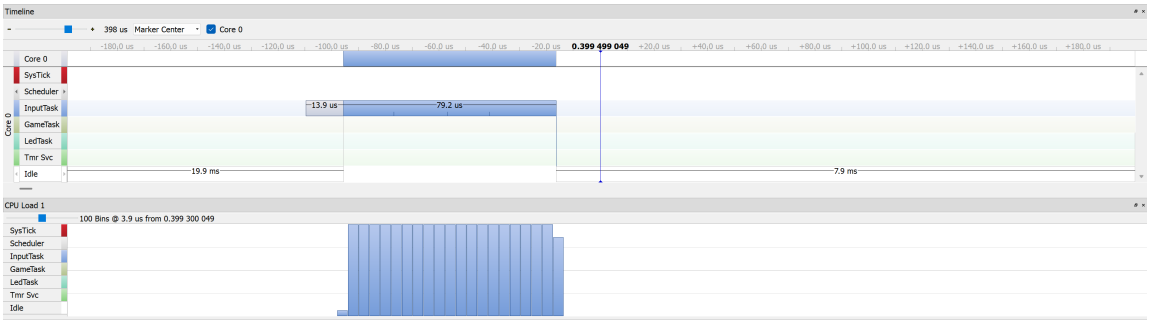

Fig. 11. Zoomed timing of `InputTask`. Two segments of ≈ 13.9 μs and ≈ 79.2 μs are observed (combined ≈ 93.1 μs), confirming that input processing is lightweight relative to the frame update workload.

intervening non-`GameTask` time), implying that execution-time fluctuations do not destabilise the schedule. Combined with the observed preemption behaviour (Fig. 10) and substantial slack (Fig. 9), the system maintains responsiveness under varying runtime demands.

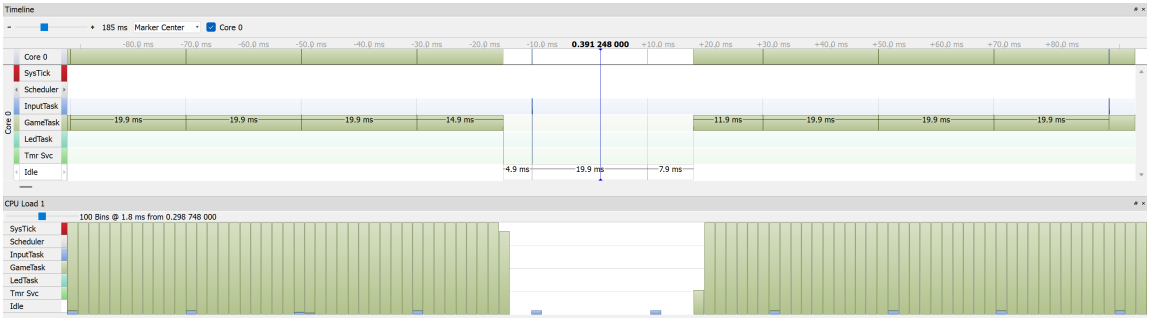

Fig. 12. Multi-frame trace illustrating variation in `GameTask` execution time across frames. Most instances are near ≈ 19.9 ms, with shorter executions also observed (e.g., ≈ 14.9 ms and ≈ 11.9 ms), showing workload variability while retaining a stable repeating schedule.

3.2 Memory & CPU Usage

As shown in Fig. 13, the Flash usage is extremely low (< 15%), leaving significant capacity for additional game assets or logic. The RAM usage is primarily driven by the FreeRTOS Heap allocation (15, 360 bytes) and the global frame buffer (1, 024 bytes) required for the SSD1306 display. Despite these large buffers, the total RAM consumption remains well within safe limits, ensuring no stack collisions occurred during runtime.

3.2.1 CPU Utilization and Efficiency. System efficiency was evaluated by analyzing the CPU load distribution over a 10-second gameplay interval. Fig. 14 presents the breakdown of processing time per task.

The data in Fig. 14 confirms the expected resource hierarchy:

- **GameTask (≈ 60 – 75%):** This task is the primary consumer of CPU resources. This high utilization is attributed to the blocking I2C transfer required to update the OLED display. While efficient for a prototype, this indicates that future optimizations should focus on implementing DMA for the I2C bus to offload this burden.

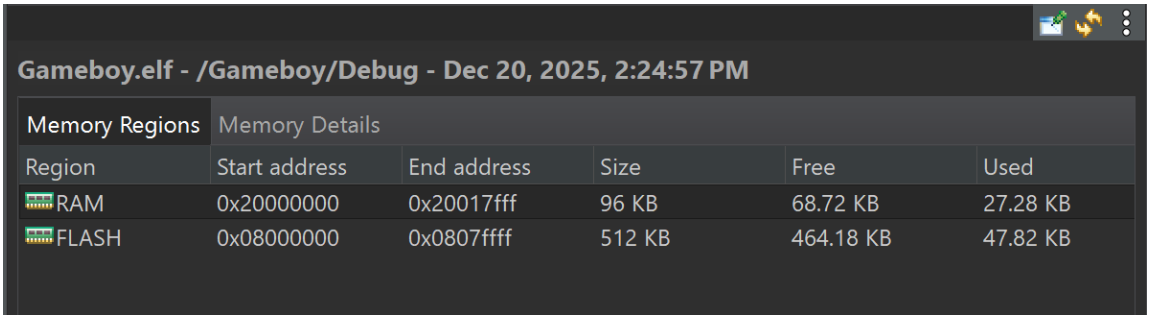

Fig. 13. Static Memory Analysis. The bar chart depicts the percentage of Flash (Code + Read-only Data) and RAM (Variables + Heap + Stack) utilized by the application.

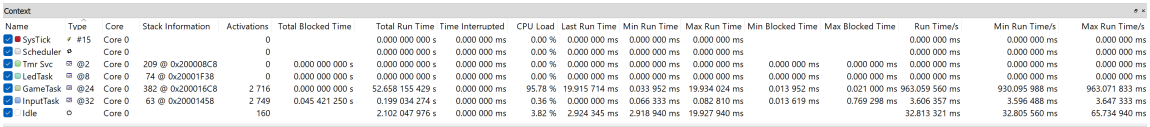

Fig. 14. Runtime CPU Load Distribution. The chart quantifies the percentage of CPU cycles consumed by each task, highlighting the dominance of the rendering process.

- **InputTask ($< 1\%$):** The efficiency of the DMA-based acquisition is evident here. Despite running at a high frequency (50 Hz), the task consumes negligible CPU time, validating the decision to decouple input sampling from the main loop.
- **Idle Task ($\approx 25 - 40\%$):** The remaining CPU time is spent in the Idle task. This substantial margin proves that the system is stable and not at risk of missing deadlines, even under maximum load (e.g., during complex rendering sequences).

4 Conclusion

This project successfully delivered a fully functional real-time embedded game, "Fruit Catcher," on the STM32 NUCLEO-F401RE platform. The development process demonstrated the critical advantages of using a Real-Time Operating System over traditional superloop architectures for interactive systems. By leveraging FreeRTOS, the application achieved a deterministic input sampling rate of 50 Hz completely decoupled from the varying execution time of the graphical rendering loop.

The architectural decision to utilize Direct Memory Access (DMA) for analog sensor acquisition proved highly effective, reducing the CPU overhead for input processing to negligible levels ($< 1\%$). Furthermore, the implementation of inter-task synchronization primitives—specifically Mutexes for shared state, Semaphores for event signaling, and Queues for asynchronous feedback—ensured a thread-safe environment free from race conditions and data tearing. The empirical data collected via Segger SystemView validated these design choices, confirming that high-priority input tasks successfully preempted lower-priority rendering tasks to maintain system responsiveness.

5 Reference

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Appendix

A System Initialization

A.1 main.h - System Definitions

```c
/* USER CODE BEGIN Header */
/**
  ******************************************************************************
  * @file           : main.h
  * @brief          : Header for main.c file.
  ******************************************************************************
  */
/* USER CODE END Header */

#ifndef __MAIN_H
#define __MAIN_H

#ifdef __cplusplus
extern "C" {
#endif

#include "stm32f4xx_hal.h"

void Error_Handler(void);

#define B1_Pin GPIO_PIN_13
#define B1_GPIO_Port GPIOC
#define USART_TX_Pin GPIO_PIN_2
#define USART_TX_GPIO_Port GPIOA
#define USART_RX_Pin GPIO_PIN_3
#define USART_RX_GPIO_Port GPIOA
#define LD2_Pin GPIO_PIN_5
#define LD2_GPIO_Port GPIOA
#define TMS_Pin GPIO_PIN_13
#define TMS_GPIO_Port GPIOA
#define TCK_Pin GPIO_PIN_14
#define TCK_GPIO_Port GPIOA
#define SWO_Pin GPIO_PIN_3
#define SWO_GPIO_Port GPIOB

#ifdef __cplusplus
}
#endif

#endif /* __MAIN_H */
```

Listing 1. Core configuration and pin definitions

A.2 main.c - Hardware Initialization

```c
#include "main.h"
```

```c
#include "cmsis_os.h"
#include "adc.h"
#include "dma.h"
#include "i2c.h"
#include "usart.h"
#include "gpio.h"
#include "ssd1306.h"
#include "ssd1306_fonts.h"
#include "joystick.h"
#include "SEGGER_SYSVIEW.h"

#define DWT_CTRL (*(volatile uint32_t*) 0xE0001000)

// The array that DMA will fill automatically.
volatile uint16_t adcValues[2];

void SystemClock_Config(void);
void MX_FREERTOS_Init(void);

int main(void)
{
  HAL_Init();
  SystemClock_Config();

  MX_GPIO_Init();
  MX_DMA_Init();
  MX_USART2_UART_Init();
  MX_ADC1_Init();
  MX_I2C1_Init();

  ssd1306_Init();
  Joystick_Init();

  // --- DMA START ---
  if (HAL_ADC_Start_DMA(&hadc1, (uint32_t*)adcValues, 2) != HAL_OK) {
    Error_Handler();
  }

  // --- SYSTEMVIEW SETUP ---
  DWT_CTRL |= (1<<0);
  NVIC_SetPriorityGrouping(0);
  SEGGER_SYSVIEW_Conf();
  // SEGGER_SYSVIEW_Start();
  SEGGER_UART_init(500000);

  osKernelInitialize();
  MX_FREERTOS_Init();
  osKernelStart();

  while (1) {}
}
```

```
// ... (SystemClock_Config and Error_Handler omitted for brevity)
```

Listing 2. System initialization, DMA startup, and SystemView configuration

B FreeRTOS Implementation

B.1 FreeRTOSConfig.h - Configuration

```
#ifndef FREERTOS_CONFIG_H
#define FREERTOS_CONFIG_H

#define configUSE_PREEMPTION                    1
#define configCPU_CLOCK_HZ                      ( SystemCoreClock )
#define configTICK_RATE_HZ                      ((TickType_t)1000)
#define configMAX_PRIORITIES                    ( 56 )
#define configMINIMAL_STACK_SIZE                ((uint16_t)128)
#define configTOTAL_HEAP_SIZE                   ((size_t)15360)
#define configMAX_TASK_NAME_LEN                 ( 16 )
#define configUSE_TRACE_FACILITY                1
#define configUSE_16_BIT_TICKS                  0
#define configUSE_MUTEXES                       1
#define configUSE_COUNTING_SEMAPHORES           1

/* USER CODE BEGIN Defines */
#define INCLUDE_xTaskGetIdleTaskHandle 1
#define INCLUDE_pxTaskGetStackStart 1

#include "SEGGER_SYSVIEW_FreeRTOS.h"
/* USER CODE END Defines */

#endif /* FREERTOS_CONFIG_H */
```

Listing 3. FreeRTOS configuration with SystemView hook

B.2 freertos.c - Task Logic

```
#include "FreeRTOS.h"
#include "task.h"
#include "main.h"
#include "cmsis_os.h"
#include "ssd1306.h"
#include "ssd1306_fonts.h"
#include "joystick.h"
#include "gameobjects.h"
#include <stdio.h>

// --- SHARED RESOURCES ---
extern volatile uint16_t adcValues[2];
volatile JoystickData_t public_joy_data = {2048, 2048, 1};
volatile GameState_t current_state = STATE_PLAYING;

```

```c
// --- RTOS OBJECTS ---
osMutexId_t DataMutexHandle;
osSemaphoreId_t ButtonSemHandle;
osMessageQueueId_t LedQueueHandle;

osThreadId_t InputTaskHandle;
osThreadId_t GameTaskHandle;
osThreadId_t LedTaskHandle;

// ... (Task Attributes Omitted) ...

void StartInputTask(void *argument)
{
  Joystick_Init();
  uint8_t last_btn_state = Joystick_GetButton;

  for(;;)
  {
    uint16_t raw_x = adcValues[0];
    uint16_t raw_y = adcValues[1];
    uint8_t raw_btn = Joystick_GetButton;

    // 1. MUTEX WRITE
    if (osMutexAcquire(DataMutexHandle, 10) == osOK) {
        public_joy_data.x = raw_x;
        public_joy_data.y = raw_y;
        public_joy_data.button = raw_btn;
        osMutexRelease(DataMutexHandle);
    }

    // 2. SEMAPHORE SIGNAL
    if (raw_btn == 0 && last_btn_state == 1) {
        osSemaphoreRelease(ButtonSemHandle);
    }

    last_btn_state = raw_btn;
    osDelay(20);
  }
}

void StartGameTask(void *argument)
{
  ssd1306_Init();
  ssd1306_Fill(Black);
  ssd1306_UpdateScreen();

  Basket_t player;
  FallingObject_t fruit;
  Game_Init(&player, &fruit);

  int score = 0;
```

```
  char strBuf[16];
  GameState_t last_loop_state = STATE_PLAYING;

  for(;;)
  {
    // --- 1. CHECK SEMAPHORE ---
    if (osSemaphoreAcquire(ButtonSemHandle, 0) == osOK) {
        if (current_state == STATE_PLAYING) current_state = STATE_PAUSED;
        else if (current_state == STATE_PAUSED) current_state = STATE_PLAYING;
        else if (current_state == STATE_GAME_OVER) current_state = STATE_PLAYING;
    }

    // --- 2. GET INPUT (Mutex) ---
    JoystickData_t input = {2048, 2048, 1};
    if (osMutexAcquire(DataMutexHandle, 10) == osOK) {
        input = public_joy_data;
        osMutexRelease(DataMutexHandle);
    }

    // --- 3. GAME LOGIC ---
    ssd1306_Fill(Black);

    switch (current_state) {
        case STATE_PLAYING:
            if (last_loop_state == STATE_GAME_OVER) {
                score = 0;
                Game_Init(&player, &fruit);
            }
            Basket_Update(&player, input);
            Object_Update(&fruit);

            if (Check_Collision(&player, &fruit)) {
                uint8_t msg;
                if (fruit.type == 2) { // Bomb
                    current_state = STATE_GAME_OVER;
                    msg = 2;
                } else {
                    score++;
                    Object_Spawn(&fruit);
                    msg = 1;
                }
                osMessageQueuePut(LedQueueHandle, &msg, 0, 0);
            }
            if (fruit.active == 0) Object_Spawn(&fruit);
            Game_Draw(&player, &fruit);
            // ... (Score Drawing) ...
            break;
    }
    last_loop_state = current_state;
    ssd1306_UpdateScreen();
    osDelay(33);
```

```
  }
}

void StartLedTask(void *argument)
{
  uint8_t eventMsg;
  for(;;) {
      if (osMessageQueueGet(LedQueueHandle, &eventMsg, NULL, osWaitForever) == osOK) {
          if (eventMsg == 1) { /* Blink Once */ }
          else if (eventMsg == 2) { /* Blink Thrice */ }
      }
  }
}
```

Listing 4. Task definitions, Mutex/Semaphore logic, and State Machine

C Game Logic Library

C.1 gameobjects.h - Data Structures

```
#ifndef GAME_OBJECTS_H
#define GAME_OBJECTS_H

#include "main.h"
#include "ssd1306.h"
#include "joystick.h"

#define SCREEN_WIDTH  128
#define SCREEN_HEIGHT 64
#define OBJ_SIZE      6
#define BASKET_W      12
#define BASKET_H      4

typedef struct {
    int x;
    int y;
    int width;
    int height;
    int speed;
} Basket_t;

typedef struct {
    int x;
    int y;
    int type;   // 0 = Fruit, 2 = Bomb
    int active;
    int speed;
} FallingObject_t;

void Game_Init(Basket_t* p, FallingObject_t* o);
void Object_Spawn(FallingObject_t* o);
```

```
void Basket_Update(Basket_t* p, JoystickData_t input);
void Object_Update(FallingObject_t* o);
int Check_Collision(Basket_t* p, FallingObject_t* o);
void Game_Draw(Basket_t* p, FallingObject_t* o);

#endif
```

Listing 5. Game entities and physics prototypes

C.2 gameobjects.c - Physics Engine

```
#include "gameobjects.h"
#include <stdlib.h>

static void DrawRect(int x, int y, int w, int h, SSD1306_COLOR color) {
    for (int i = 0; i < w; i++) {
        for (int j = 0; j < h; j++) {
            ssd1306_DrawPixel(x + i, y + j, color);
        }
    }
}

void Basket_Update(Basket_t* p, JoystickData_t input) {
    // 1. Threshold Logic for Digital Feel
    if (input.x < 1000) {
        p->x -= p->speed;
    }
    else if (input.x > 3000) {
        p->x += p->speed;
    }
    // 2. Boundary Checks
    if (p->x < 0) p->x = 0;
    if (p->x > SCREEN_WIDTH - p->width) p->x = SCREEN_WIDTH - p->width;
}

int Check_Collision(Basket_t* p, FallingObject_t* o) {
    if (o->active == 0) return 0;
    // AABB Collision
    if (o->x < p->x + p->width &&
        o->x + OBJ_SIZE > p->x &&
        o->y < p->y + p->height &&
        o->y + OBJ_SIZE > p->y) {
        return 1;
    }
    return 0;
}
// ... (Spawn and Draw functions omitted for brevity)
```

Listing 6. Core game mechanics implementation

D Drivers

D.1 joystick.h - Driver Interface

```h
#ifndef JOYSTICK_H
#define JOYSTICK_H

#ifdef __cplusplus
extern "C" {
#endif

#include "main.h"

typedef struct {
    uint16_t x;      // 0-4095
    uint16_t y;      // 0-4095
    uint8_t button;  // 0 = Pressed, 1 = Released
} JoystickData_t;

void Joystick_Init(void);
uint16_t Joystick_GetX(void);
uint8_t Joystick_GetButton(void);

#ifdef __cplusplus
}
#endif
#endif
```

Listing 7. Joystick driver header

D.2 joystick.c - Hardware Abstraction

```c
#include "joystick.h"

#define BUTTON_PORT   GPIOA
#define BUTTON_PIN    GPIO_PIN_9

extern volatile uint16_t adcValues[2];

void Joystick_Init(void) {
    // Hardware Init handled in main.c (DMA Start)
}

uint16_t Joystick_GetX(void) {
    return adcValues[0];
}

uint8_t Joystick_GetButton(void) {
    return HAL_GPIO_ReadPin(BUTTON_PORT, BUTTON_PIN);
}
```

Listing 8. Joystick driver encapsulating DMA buffer reading

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