## Homework 2

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All code used in this assignment, the markdown files used to build this PDF, and the PDF building script can be found in GitHub here.

## **Problems**

#### Problem 1a

Pseudo-code implementation. In this,  $\mathtt{i}$  is the element to insert at base of stack and  $\mathtt{s}$  is the input stack:

```
FUNCTION [i: ELEMENT, s: STACK] -> STACK
START
    # stack s is the input stack and i is the element to place at bottom
    LET s_0 = INITIALIZE EMPTY STACK
WHILE s IS NOT EMPTY:
        s_0.push(s.pop) # pop from one stack and immediately move to second stack
s.push(i)
WHILE s_0 IS NOT EMPTY:
        s.push(s.pop)
RETURN s
END
```

## Problem 1b

Pseudo-code implementation. In this,  ${\tt i}$  is the element to insert in the third position of the stack and  ${\tt s}$  is the input stack:

```
FUNCTION [i: ELEMENT, s: STACK] -> STACK
START
    # stack s is the input stack and i is the element to place in the third position
    LET s_0 = INITIALIZE EMPTY STACK
    WHILE s IS NOT EMPTY:
        s_0.push(s.pop) # pop from one stack and immediately move to second stack
    s.push(s_0.pop)
    s.push(s_0.pop)
    s.push(i)
    WHILE s_0 IS NOT EMPTY:
        s.push(s_0.pop)
    RETURN s
END
```

# Problem 2a

The below table represents each character as it's being iterated over, what the stack looks like, and any stack operation that is being applied:

Character	Stack	Stack Operation
{	{	Push '{'
[	[{	Push '['
A	[{	None
+	[{	None
В	[{	None
]	{	Pop '['
-	{	None
[	[{	Push '['
(	([{	Push '('
$\mathbf{C}$	([{	None
-	([{	None
D	([{	None
)	[{	Pop '('
]	{	Pop '['
END	{	Check empty

After iterating through all characters, stack is non-empty so delimiters are not properly matching.

# Problem 2b

The below table represents each character as it's being iterated over, what the stack looks like, and any stack operation that is being applied:

Character	Stack	Stack Operation
(	(	Push '('
(	((	Push '('
H	((	None
)	(	Pop '('
*	(	None
{	{(	Push '{'
(	({(	None
[	[({(	Push '['
J	[({(	None
+	[({(	None
K	[({(	None
]	({(	Pop '['
)	{(	Pop '('
}	(	Pop '{'
)		Pop '('
END		Check empty

After iterating through all characters, stack is empty so delimiters are properly matching.

Pseudo-code implementation checking mirrored strings. In this function, w is the input string in form xCy that will be tested for being mirrored:

```
FUNCTION [w: STRING] -> BOOLEAN
START
  # string w is input string of format xCy
 LET stack = INITIALIZE EMPTY STACK
 LET found_c = FALSE
 LET output = TRUE
 FOR char IN w:
    IF NOT found_c AND char != 'C':
      stack.push(char)
   ELSE IF char == 'C':
      LET found_c = TRUE
    ELSE:
      IF stack.pop != char:
       LET output = FALSE
  RETURN output
END
```

A Python implementation of the above algorithm with some test cases is as follows. As these problems become more complex, I feel it more necessary to convert the pseudo-code into real code that I can use to test various test cases.

```
class Stack(list):
    # This is just a wrapper class around Python lists to so that we can interact
    # with the lists in a way that idiomatic of a stack with methods to push and check
    # if the stack is empty.
    def __init__(self):
        super().__init__()
    def push(self, item):
        self.append(item)
    def is_empty(self):
        return len(self) == 0
def check_mirrored(s: str) -> bool:
    stack = Stack()
    found_c = False
    output = True
    for char in s:
        if (not found_c) and (char != 'C'):
            stack.push(char)
        elif char == 'C':
            found_c = True
        else:
            if stack.pop() != char:
```

```
output = False
return output and stack.is_empty() # using 'and' here to validate stack is empty

if __name__ == "__main__":
    print("Running test cases...")
    assert check_mirrored('xCx') == True
    assert check_mirrored('xyCyx') == True
    assert check_mirrored('xyCy') == False
    assert check_mirrored('xyCxy') == False
    assert check_mirrored('xyyzaCazyyx') == True
    assert check_mirrored('xyyzaCazyyu') == False
```

Within the below pseudo-code, the function named check\_each\_section takes an input string and returns true if it matches the pattern and false if there is a non-symmetric xCy substring.

```
FUNCTION check_mirrored [input_stack: STACK] -> BOOLEAN
START check_mirrored:
  # stack inpt_stack is input stack representation with format xCy
  LET stack = INITIALIZE EMPTY STACK
  LET found_c = FALSE
  WHILE NOT s.is_empty:
   LET char = s.pop
    IF NOT found_c AND char != 'C':
      stack.push(char)
   ELSE IF char == 'C':
      LET found_c = TRUE
      IF stack.pop != char:
        RETURN FALSE
 RETURN stack.is_empty
END check_mirrored
FUNCTION check_each_section [w: STRING] -> BOOLEAN
START check_each_section:
  LET section_stack = INITIALIZE EMPTY STACK
  FOR char IN w:
    IF char != 'D':
      section_stack.push(char)
   ELSE:
      IF NOT check_mirrored(section_stack):
        RETURN FALSE
      LET section_stack = INITIALIZE EMPTY STACK
  RETURN check_mirrored(section_stack)
END check_each_section
```

The below code is a slight modification and extension of the above Python code. For this implementation, we use two stacks. One stack collects everything inbetween the 'D's (the 'D-segments') and the second stack is used within the function to test if the extracted D-segment sequence stack represents a valid xCy mirrored string. Again, test strings are asserted at the end.

```
class Stack(list):
    def __init__(self):
        super().__init__()

def push(self, item):
        self.append(item)

def is_empty(self):
        return len(self) == 0
```

```
def check_mirrored(s: Stack) -> bool:
    stack = Stack()
   found_c = False
    for char in s:
        if (not found_c) and (char != 'C'):
            stack.push(char)
        elif char == 'C':
            found_c = True
        else:
            if stack.pop() != char:
                return False
    return stack.is_empty() # Need to ensure stack is empty after popping all matching elements
def check_each_section(s: str) -> bool:
    section_stack = Stack()
    for char in s:
        if char != 'D':
            section_stack.push(char)
        else:
            if not check_mirrored(section_stack):
                return False
            section_stack = Stack()
    return check_mirrored(section_stack) # Need to check last D-segment
if __name__ == "__main__":
   print("Running test cases...")
   assert check_each_section('xCxDxCx') == True
   assert check_each_section('DxCxDxCx') == True
   assert check_each_section('xCxDxCxD') == True
   assert check_each_section('xyCyxDxyCyx') == True
   assert check_each_section('xyCxyDxyCyx') == False
    assert check_each_section('xyyzaCazyyxDxyCyxDuwqCqwu') == True
    assert check_each_section('xyyzaCazyyxDxyyzaCazyyu') == False
```

Within the pseudo-code implementation below, the function insert shows how a second stack can be used to insert a new element into the primary stack and the function read will read the element from the stack.

For both these functions, the input parameters i are the index to insert/read at, s is the stack to do the reading/insertion in, and e is the element to insert (for the insertion function).

In this implementation, the bottom of the stack will be indexed at 0 and will increase by increments of +1 for each additional element added to the stack. Once the stack has been popped a number of times equivalent to the index, the element to insert is pushed onto the stack before moving everything from the secondary stack back onto the primary stack. The **read** function works similarly to the **insert** function, but it simply peeks once it gets to the appropriate index instead of pushing a new element onto the stack.

```
FUNCTION insert [e: ELEMENT, i: INDEX, s: STACK] -> STACK
# function which uses a second stack to insert into another stack
START insert
  LET secondary stack = INITIALIZE EMPTY STACK
  FOR _ FROM O to i:
    secondary_stack.push(s.pop)
  s.push(e)
  FOR _ FROM O to i:
    s.push(secondary_stack.pop)
  RETURN s
END insert
FUNCTION read [i: INDEX, s: STACK] -> STACK
# This function works almost the same as 'insert', just
# without doing the actual insertion
START read
  LET secondary_stack = INITIALIZE EMPTY STACK
  FOR _ FROM O to i:
    secondary_stack.push(s.pop)
  LET output = s.read
  FOR _ FROM O to i:
    s.push(secondary_stack.pop)
  RETURN output
END read
Python implementation with test cases are shown below:
FUNCTION insert [e: ELEMENT, i: INDEX, s: STACK] -> STACK
# function which uses a second stack to insert into another stack
START insert
  LET secondary_stack = INITIALIZE EMPTY STACK
  FOR _ FROM 0 to i:
    secondary_stack.push(s.pop)
  s.push(e)
```

```
FOR \_ FROM 0 to i:
    s.push(secondary_stack.pop)
  RETURN s
END insert
FUNCTION read [i: INDEX, s: STACK] -> STACK
\# This function works almost the same as 'insert', just
# without doing the actual insertion
START read
 LET secondary_stack = INITIALIZE EMPTY STACK
 FOR _ FROM 0 to i:
   secondary_stack.push(s.pop)
 LET output = s.read
 FOR \_ FROM 0 to i:
   s.push(secondary_stack.pop)
  RETURN output
END read
```

Need to come back to complete.

- # Design a method for keeping two stacks within a single linear array
- # s[SPACESIZE] so that neither stack overflows until all of memory is used and an
- # entire stack is never shifted to a different location within the array. Write
- # methods push1, push2, pop1, and pop2 to manipulate the two stacks. (Hint:
- # the two stacks grow toward each other.)

TODO!

Problem 7a

We can parse the expression (A+B)\*(C\$(D-E)+F)-G for postfix in the following steps. The stack starts with '(' inside.

Character	Stack	Current expression
(	((	NONE
À	(	A
+	((+	A
В	((+	AB
)	(	AB+
*	(*	AB+
(	(*(	AB+
$^{\mathrm{C}}$	(*(	AB+C
\\$	(*(	AB+C\$
(	' ''	AB+C\$
D	( ( (	AB+C\$D
-		AB+C\$D
$\mathbf{E}$		AB+C\$DE
)	( (	AB+C\$DE-
+		AB+C\$DE-
$\mathbf{F}$		AB+CDE-F
)	(*	AB+C\$DE-F+
-	(-	AB+C\$DE-F+*
G	(-	AB+C\$DE-F+*G
)	NONE	AB+C\$DE-F+*G-

So postfix notation for the expression is AB+C\$DE-F+\*G-

For prefix notation, we can parse as follows:

Character	Stack	Current expression
G	(	G
-	(-	G
(	(-(	G
F	(-(	$\operatorname{GF}$
+	(-(+	$\operatorname{GF}$
(	(-(+(	$\operatorname{GF}$
$\mathbf{E}$	(-(+(	GFE
-	(-(+(-	GFE
D	(-(+(-	GFED
)	(-(+	GFED-
\$.g	(-(+	GFED-\$
$\mathbf{C}$	(-(+	GFED-\$.gC
)	(-	GFED-\$.gC+
*	(-*	GFED-\$.gC+
(	(-*(	GFED-\$.gC+
В	(-*(	GFED-\$.gC+B
+	(-*(+	GFED-\$.gC+B
A	(-*(+	GFED-\$.gC+BA
)	(-*	GFED-\$.gC+BA+
)	NONE	GFED-\$.gC+BA+*-

So the prefix notation for that expression is  ${\tt GFED-\$C+BA+*-}.$ 

Problem 7b

We can parse the expression A+(((B-C)\*(D-E)+F)/G)\$(H-J) in postfix notation using the following steps:

Character	Stack	Current expression
A	(	A
+	(+	A
(	(+(	A
+ (	(+((	A
(	(+(((	A
В	(+(((	AB
-	(+(((-	AB
C ) *	(+(((-	ABC
)	(+((	ABC-
*	(+((*	ABC-
(	(+((*(	ABC-
D	(+((*(	ABC-D
-	(+((*(-	ABC-D
E	(+((*(-	ABC-DE
)	(+((*	
+ F	(+((+	
F	(+((+	ABC-DE-*F
)	(+(	ABC-DE-*F+
/	(+(/	ABC-DE-*F+
G	(+(/	ABC-DE-*F+G
/ G ) \$	(+	ABC-DE-*F+G/
\$	(+	ABC-DE-*F+G/\$
(	(+(	ABC-DE-*F+G/\$
H	(+(	ABC-DE-*F+G/\$H
-	(+(-	
J	(+(-	ABC-DE-*F+G/\$HJ
)	(+	ABC-DE-*F+G/\$HJ-
)	NONE	ABC-DE-*F+G/\$HJ-+

So the postfix notation for the expression is ABC-DE-\*F+G/\$HJ-+.

To parse the expression into prefix notation, we complete the following steps:

Character	Stack	Current expression
(	((	NONE
Ĵ	((	J
-	((-	J
H	((-	JH
)	(	JH-
\$	(	JH-\$
(	((	JH-\$
( G / ( F	((	$_{ m JH-\$G}$
/	((/	$_{ m JH-\$G}$
(	((/(	$_{ m JH-\$G}$
	((/(	$_{ m JH-\$GF}$
+	((/(+	$_{ m JH-\$GF}$
+ (	(() ( . (	$_{ m JH-\$GF}$
$\mathbf{E}$	((/(+(	$_{ m JH-\$GFE}$
-	((/(+(-	$_{ m JH-\$GFE}$
D		JH-\$GFED
)		JH-\$GFED-
*		JH-\$GFED-
(	((/(+*(	JH-\$GFED-
C	((/(+*(	m JH-\$GFED-C
-	((/(+*(-	JH-\$GFED-C
В	((/(+*(-	JH-\$GFED-CB
)		JH-\$GFED-CB-
) ) )	((/	JH-\$GFED-CB-*+
)	(	JH-\$GFED-CB-*+/
+ A	(+	JH-\$GFED-CB-*+/
	(+	JH-\$GFED-CB-*+/A
)	NONE	JH-\$GFED-CB-*+/A+

So the prefix notation for the expression is  $\mathtt{JH-\$GFED-CB-*+/A+}$ 

# Problem 8a

Out of time.