


Algorithmics	Student information	Date	Number of session
	UO: 300047	03/02/2025	1
	Surname: Rodríguez Torres	 Escuela de Ingeniería Informática Universidad de Oviedo	
	Name: Luisa Natalia		



## Activity 1. Divide and Conquer by subtraction

**Subtraction1.java:**

- **A = 1**
- **B = 1**
- **K = 0**
- **TOTAL COMPLEXITY:  $O(n)$**

**Subtraction2.java:**

- **A = 1**
- **B = 1**
- **K = 1**
- **TOTAL COMPLEXITY:  $O(n^2)$**

**Subtraction3.java:**

- **A = 2**
- **B = 1**
- **K = 0**
- **TOTAL COMPLEXITY:  $O(2^n)$**

After running the algorithms we can see that they all have their expected complexities.

For Subtraction1.java we stop computing at 16.384 due to a stackOverflow.

For Subtraction2. java we stop computing at  $n = 32.768$  due to a stackOverflow.

This happens due to how the recursions are programmed. These methods will use recursion  $n$  times, which at high numbers such as  $n = 16384$ , we will repeat the operation that amount of times. The Java stack cannot handle that much information.

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To calculate how long it would take for Substraction3 to compute at  $n=80$ , we can take a random measurement, in this example I will use 1788ms at  $n = 30$ . Since the complexity is  $O(2^n)$ , the time will be calculated as:  $1788 \times 2^{(n-30)}$ ms. In this case  $1788 \times 2^{(50)}$ ms. We get approximately  $2.01 \times 10^{18}$  milliseconds, or approximately 63,693,056 years.

Subtraction4.java:

- A = 1
- B = 1
- K = 2
- TOTAL COMPLEXITY:  $O(n^3)$
- 

For small values (100–800), the execution time is negligible, however after 1600, the time grows exponentially as expected, following a cubic pattern and quickly getting OoT.

n	Time (ms)
100	2
200	2
400	0
800	0
1600	33
3200	279
6400	2104
12800	15314
25000	OoT

Subtraction5.java:

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- **A = 2**
- **B = 2**
- **K = 1**
- **TOTAL COMPLEXITY:  $O(3^{(n/2)})$**

We can clearly see that the complexity of the problem translates to its real time. Every two sizes however, the time barely changes, due to perhaps, not having to divide the problem more from the previous iteration of the problem, so the same amount of invocations are needed.

n	Time (ms)
30	18
32	62
34	61
36	548
38	539
40	4819
42	4886
44	OoT

## Activity 2. Divide and Conquer by division

**Division1.java:**

- **A = 1**
- **B = 3**
- **K = 1**

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- **TOTAL COMPLEXITY:  $O(n)$**

#### Division2.java:

- **A = 2**
- **B = 2**
- **K = 1**
- **TOTAL COMPLEXITY:  $O(n \cdot \log n)$**

#### Division3.java:

- **A = 2**
- **B = 2**
- **K = 0**
- **TOTAL COMPLEXITY:  $O(n)$**

#### Division4.java:

- **A = 1**
- **B = 2**
- **K = 2**
- **TOTAL COMPLEXITY:  $O(n^2)$**

#### Division5.java:

- **A = 5**
- **B = 2**
- **K = 2**
- **TOTAL COMPLEXITY:  $O(n^2)$**

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## Activity 4. Two basic examples

Order	Fibonacci 1 time (ms)	Fibonacci 2 time (ms)	Fibonacci 3 time (ms)	Fibonacci 4 time (ms)
10	279	293	680	13652
20	337	635	1327	OoT
30	457	944	1919	OoT
40	598	1407	2608	OoT
50	579	2004	3233	OoT
60	563	2518	3802	OoT

Fibonacci 1 is an iterative solution with complexity  $O(n)$ . Fibonacci 2 is also an iterative solution with complexity  $O(n)$  but it utilizes a vector, making it slower when having to deal with bigger sizes. Fibonacci 3 is a recursive method with complexity  $O(n)$ , this method is slower even at smaller orders. Finally, Fibonacci 4 has exponential complexity  $O(1.6^n)$ , making it by far the worst version.

Size	VectorSum 1 time (ms)	VectorSum 2 time (ms)	VectorSum 3 time (ms)
3	2	3	2
24	1	2	3
192	5	13	37
1536	38	230	254
12288	309	2069	2448
98304	2392	4243	16467

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VectorSum 1 presents a complexity of  $O(n)$  due to iterating through the array once. VectorSum 2 presents also a complexity of  $O(n)$ , using recursion. VectorSum 3 also presents a complexity of  $O(n)$  by using divide and conquer. As we can observe, despite them all having the same complexity, VectorSum 1 is much faster than the others since it uses simpler operations. By using a simple for-loop with a sum it computes a very simple operation which reduces its overhead. Vector2 uses recursion, so it must allocate memory for every iteration. Vector3 also has the same problem as Vector 2, except with even more overhead, since we must calculate the middle point.

## Activity 2. Petanque championship organization

- The complexity of the algorithm I have implemented is:  $O(n \log(n))$ . It uses divide and conquer by division with the following parameters:  $A=2$ ,  $B=2$ ,  $K=1$ .

n	Time
2	1
4	0
8	0
16	0
32	0
64	0
128	0
256	0
512	1
1024	5
2048	46
4096	176
8192	608

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16384	3712
32768	OutOfMemory

We can see that this algorithm works very well for smaller numbers, and does indeed grow with a  $n \log(n)$  complexity. However, for bigger numbers, the amount of function calls becomes too much for Java, since it must allocate memory in each of them, so we get an OutOfMemory exception.