

OCT 2023 CHAPEL TUTORIAL UPDATED FOR USE IN SPRING 2025 CS372 CLASS

Chapel Team, edited by Michelle Strout March 27 and April 1, 2025

PLAN

Announcements

- MT2 and LA2 grades have been posted
- SA6 is due Monday, survey for final project preferences
- Some reading assignment for Chapel have been posted

• Last time

• Midterm 2

Today

- Chapel introduction
- Please start the docker pull for Chapel (see next slide)

Example codes for Chapel tutorial slides

• https://github.com/UofA-CSc-372-Spring-2025/CSc372Spring2025-CourseMaterials/tree/main/Sandboxes/ChapelTutorialExamples

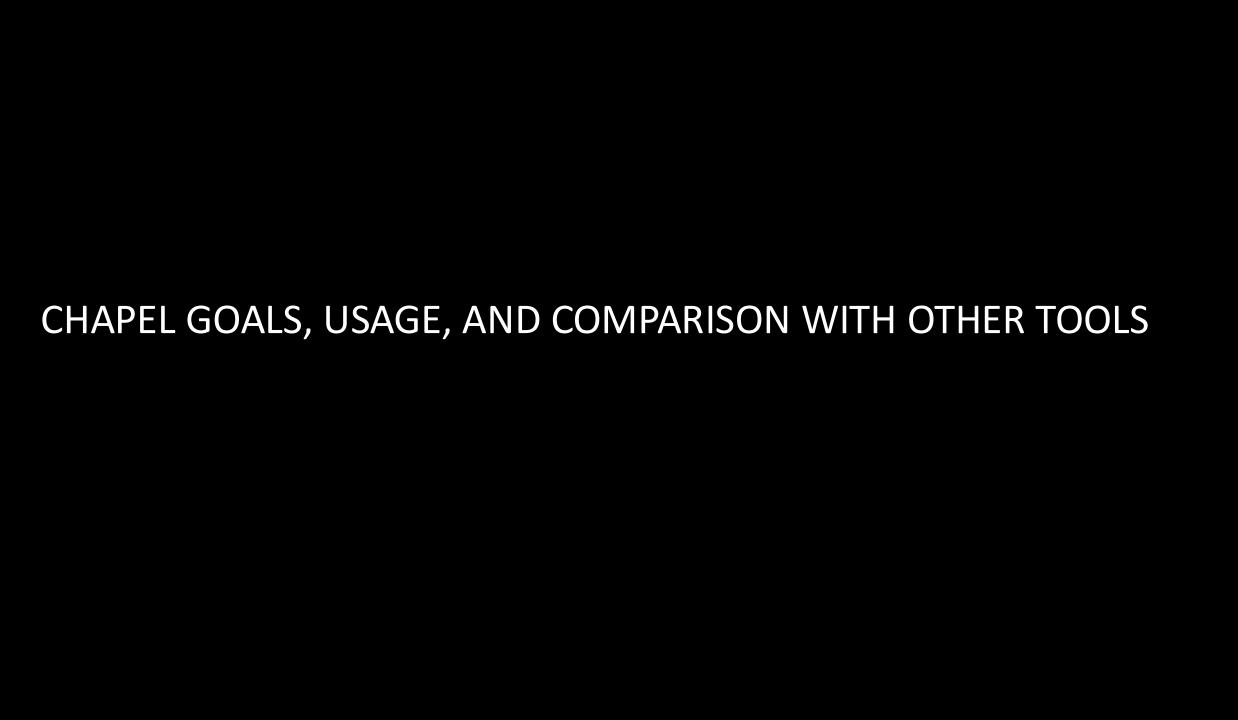
Using a container on your laptop

- First, install docker for your machine and start it up (see the README.md for more info)
- Then, use the chapel-gasnet docker container

```
docker pull docker.io/chapel/chapel-gasnet  # takes about 5 minutes
cd CSc372Spring2025-CourseMaterials/Sandboxes/ChapelTutorialExamples/
docker run --rm -it -v "$PWD":/workspace chapel/chapel-gasnet
root@589405d07f6a:/opt/chapel# cd /workspace
root@xxxxxxxxx:/myapp# chpl 01-hello.chpl
root@xxxxxxxxx:/myapp# ./01-hello -nl 1
```

OUTLINE: OVERVIEW OF PROGRAMMING IN CHAPEL

- Chapel Goals, Usage, and Comparison with other Tools
- Hello World (Hands On)
- Chapel Execution Model and Parallel Hello World (Hands On)
- kmer counting using file IO, config consts, strings, maps (Hands On)
- Parallelizing a program that processes files (Hands On)
- GPU programming support
- Learning goals for rest of tutorial



CHAPEL PROGRAMMING LANGUAGE

Chapel is a general-purpose programming language that provides ease of parallel programming, high performance, and portability.

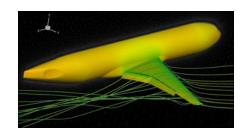
And is being used in applications in various ways:

refactoring existing codes,
developing new codes,
serving high performance to Python codes (Chapel server with Python client), and
providing distributed and shared memory parallelism for existing codes.



APPLICATIONS OF CHAPEL: LINKS TO USERS' TALKS (SLIDES + VIDEO)

CHIUW 2023



CHAMPS: 3D Unstructured CFD

Python3 Client
Socket

Dispatcher

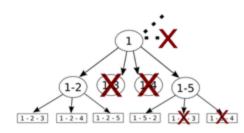
Code Modules

Distributed
Object Store
Platform

MPP, SMP, Cluster, Laptop, etc.

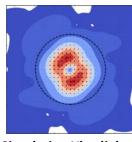
Arkouda: Interactive Data Science at Massive Scale

CHIUW 2020



ChOp: Chapel-based Optimization

CHIUW 2021 CHIUW 2023



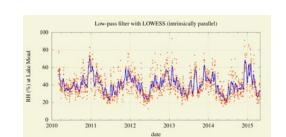
ChplUltra: Simulating Ultralight Dark Matter

ChOp: Chaper-based Op

CHIUW 2020

020 CHIUW 2022



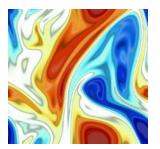


Lattice-Symmetries: a Quantum Many-Body Toolbox Desk dot chpl: Utilities for Environmental Eng.



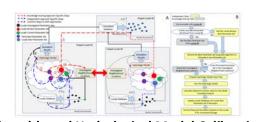
RapidQ: Mapping Coral Biodiversity

CHIUW 2023



ChapQG: Layered Quasigeostrophic CFD

CHIUW 2022



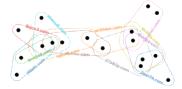
Chapel-based Hydrological Model Calibration CHIUW 2023

CHIUW 2022

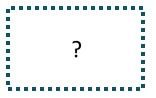


CrayAl HyperParameter Optimization (HPO)

CHIUW 2021



CHGL: Chapel Hypergraph Library CHIUW 2020



Your Application Here?



HIGHLIGHTS OF CHAPEL USAGE

CHAMPS: Computational Fluid Dynamics framework for airplane simulation

- Professor Eric Laurendeau's team at Polytechnique Montreal
- Performance: achieves competitive results w.r.t. established, world-class frameworks from Stanford, MIT, etc.
- Programmability: "We ask students at the master's degree to do stuff that would take 2 years and they do it in 3 months."

Arkouda: data analytics framework (https://github.com/Bears-R-Us/arkouda)

- Mike Merrill, Bill Reus, et al., US DOD
- Python front end client, Chapel server that processes dozens of terabytes in seconds
- April 2023: 1200 GiB/s for argsort on an HPE EX system



Other recent users

- Marjan Asgari et al, "Development of a knowledge-sharing parallel computing approach for calibrating distributed watershed hydrologic models", Environmental Modeling and Software.
- Scott Bachman has written some coral reef image analysis applications in Chapel.



CHAPEL IS HIGHLY PERFORMANT AND SCALABLE

HPE Apollo (May 2021)



- HDR-100 Infiniband network (100 Gb/s)
- 576 compute nodes
- 72 TiB of 8-byte values
- ~480 GiB/s (~150 seconds)

HPE Cray EX (April 2023)

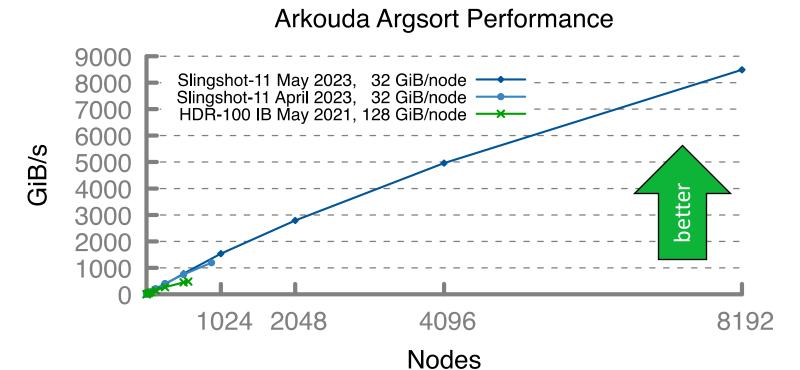


- Slingshot-11 network (200 Gb/s)
- 896 compute nodes
- 28 TiB of 8-byte values
- ~1200 GiB/s (~24 seconds)

HPE Cray EX (May 2023)



- Slingshot-11 network (200 Gb/s)
- 8192 compute nodes
- 256 TiB of 8-byte values
- ~8500 GiB/s (~31 seconds)



A notable performance achievement in ~100 lines of Chapel



COMPARE WITH OTHER PARALLEL PROGRAMMING MODELS

- Shared-memory parallelism
 - Pthreads: low-level library for creating and managing thread
 - OpenMP: pragmas added before loops and other statements
 - Rust, Julia: programming languages with some threaded para
 - RAJA, Kokkos: C++ libraries that use template metaprogramn
- Distributed-memory parallelism and shared-memory para
 - MPI+X:
 - -MPI stands for message passing interface
 - MPI is a library for sending and receiving messages between proce
 - All processes allocate their own memory and run the same progra
 - -There are many options for X: OpenMP, Pthreads, Python, Julia, R,
 - OpenSHMEM: library for implementing a partitioned global a
 - Spark: Python, Scala, and Java accessible library for especially
 - Regent and Legion: programming language and runtime that implements implicit task parallelism
 - Kokkos Remote Spaces: extends Kokkos C++ template views to distributed views

Chapel:

- shared memory parallelism,
- distributed-memory parallelism,
- data parallelism,
- task parallelism,
- map-reduce parallelism,
- vector parallelism,
- GPU parallelism, ...

All can be expressed in the same programming language.

COMPARE WITH OTHER PARALLEL PROGRAMMING MODELS (W/OUT CHAPEL BOX)

- Shared-memory parallelism
 - Pthreads: low-level library for creating and managing threads of execution that share memory
 - OpenMP: pragmas added before loops and other statements in C/C++/Fortran programs
 - Rust, Julia: programming languages with some threaded parallelism constructs
 - RAJA, Kokkos: C++ libraries that use template metaprogramming to express parallel policies
- Distributed-memory parallelism and shared-memory parallelism
 - MPI+X:
 - -MPI stands for message passing interface
 - -MPI is a library for sending and receiving messages between processes
 - -All processes allocate their own memory and run the same program, SPMD: Single Program Multiple Data
 - -There are many options for X: OpenMP, Pthreads, Python, Julia, RAJA, Kokkos, Chapel, ...
 - OpenSHMEM: library for implementing a partitioned global address space
 - Spark: Python, Scala, and Java accessible library for especially the map-reduce parallelism
 - Regent and Legion: programming language and runtime that implements implicit task parallelism
 - Kokkos Remote Spaces: extends Kokkos C++ template views to distributed views

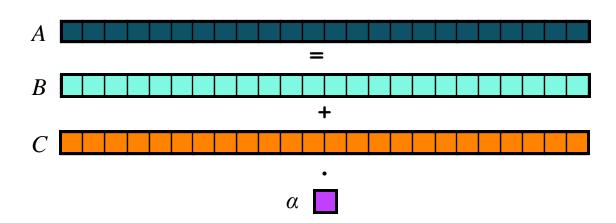


LET'S COMPARE WITH STREAM TRIAD: A PARALLEL COMPUTATION

Given: *m*-element vectors *A*, *B*, *C*

Compute: $\forall i \in 1..m, A_i = B_i + \alpha \cdot C_i$

In pictures:



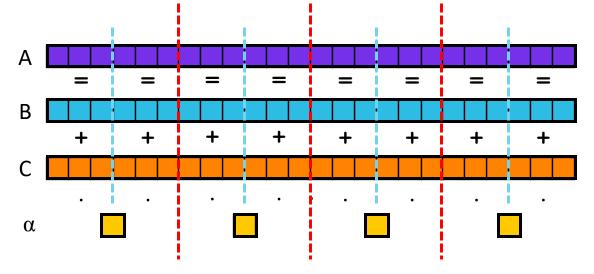
LET'S COMPARE WITH STREAM TRIAD: A PARALLEL COMPUTATION

Given: n-element vectors A, B, C

Compute: $\forall i \in 1..n, A_i = B_i + \alpha \cdot C_i$

In pictures, in parallel (distributed memory multicore, global-

view):



STREAM TRIAD: IN MPI+OPENMP

```
MPI + OpenMP
#include <hpcc.h>
                                                        if (!a | | !b | | !c) {
                                                                                         #define N 2000000
                                                                                                                CUDA
#ifdef OPENMP
                                                          if (c) HPCC free(c);
#include <omp.h>
                                                          if (b) HPCC free(b);
                                                                                         int main() {
                                                                                           float *d a, *d b, *d c;
#endif
                                                          if (a) HPCC free(a);
                                                                                           float scalar;
                                                          if (doIO) {
static int VectorSize;
                                                            fprintf( outFile, "Failed to
                                                                                           cudaMalloc((void**)&d a, sizeof(float)*N);
static double *a, *b, *c;
                                                              allocate memory (%d).\n",
                                                                                           cudaMalloc((void**)&d b, sizeof(float)*N);
                                                              VectorSize );
                                                                                           cudaMalloc((void**)&d c, sizeof(float)*N);
int HPCC StarStream(HPCC Params *params) {
                                                            fclose( outFile );
 int myRank, commSize;
                                                                                           dim3 dimBlock(128);
```

HPC suffers from too many distinct notations for expressing parallelism and locality. This tends to be a result of bottom-up language design.

```
MPI_Reduce( &rv, &errCount, 1, MPI_INT, MPI_SUM,
    0, comm );

return errCount;
}

int HPCC_Stream(HPCC_Params *params, int doIo) {
    register int j;
    double scalar;

VectorSize = HPCC_LocalVectorSize( params, 3,
        sizeof(double), 0);

a = HPCC_XMALLOC( double, VectorSize);
    b = HPCC_XMALLOC( double, VectorSize);
    c = HPCC_XMALLOC( double, VectorSize);
```

```
for (j=0; j<VectorSize; j++) {
   b[j] = 2.0;
   c[j] = 1.0;
}
scalar = 3.0;
#ifdef _OPENMP
#pragma omp parallel for
#endif
for (j=0; j<VectorSize; j++)
   a[j] = b[j]+scalar*c[j];

HPCC_free(c);
HPCC_free(b);
HPCC_free(a);
return 0; }</pre>
```

WHY SO MANY PROGRAMMING MODELS?

HPC tends to approach programming models bottom-up:

Given a system and its core capabilities...

...provide features that permit users to access the available performance.

Type of HW Parallelism	Programming Model	Unit of Parallelism
Inter-node	MPI	executable
Intra-node/multicore	OpenMP / pthreads	iteration/task
Instruction-level vectors/threads	pragmas	iteration
GPU/accelerator	CUDA / Open[MP CL ACC]	SIMD function/task

benefits: lots of control; decent generality; easy to implement

downsides: lots of user-managed detail; brittle to changes



STREAM TRIAD: IN CHAPEL

```
#include <omp.h>
static double *a, *b, *c;
int HPCC StarStream(HPCC Params *pa
 int rv, errCount;
                  sizeof(double),
 a = HPCC XMALLOC(
 b = HPCC XMALLOC (
                 and HPC expert to each focus on their strengths.
```

c = HPCC XMALLOC(

#include <hpcc.h>

```
use BlockDist;
                        config const m = 1000,
                                      alpha = 3.0;
                        const ProblemSpace = blockDist.createDomain({1..m})
                        var A, B, C: [ProblemSpace] real;
                        B = 2.0;
                        C = 1.0;
rv = HPCC_Stream( params, 0 == my A = B + alpha * C;
```

The special sauce: How should this index set and any arrays and computations over it—be mapped to the system?

```
(float) *N);
(float) *N);
```

```
Philosophy: Top-down language design can tease system-specific implementation
details away from an algorithm, permitting the compiler, runtime, applied scientist,
```

```
*b, float *c,
```

alue, int len) {

HELLO WORLD (HANDS ON)

Example codes for Chapel tutorial slides

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Using a container on your laptop

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```
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cd CSc372Spring2025-CourseMaterials/Sandboxes/ChapelTutorialExamples/
docker run --rm -it -v "$PWD":/workspace chapel/chapel-gasnet
root@589405d07f6a:/opt/chapel# cd /workspace
root@xxxxxxxxx:/myapp# chpl 01-hello.chpl
root@xxxxxxxxx:/myapp# ./01-hello -nl 1
```

"HELLO WORLD" IN CHAPEL: TWO VERSIONS

Fast prototyping

```
writeln("Hello, world!");
```

"Production-grade"

```
module Hello {
   proc main() {
     writeln("Hello, world!");
   }
}
```





"HELLO WORLD" IN CHAPEL: TWO VERSIONS

Fast prototyping (configurable)

```
config const audience = "world";
writeln("Hello, ", audience, "!");
```

```
01-hello-configurable.chpl
```

"Production-grade" (configurable)

```
module Hello {
  config const audience = "world";

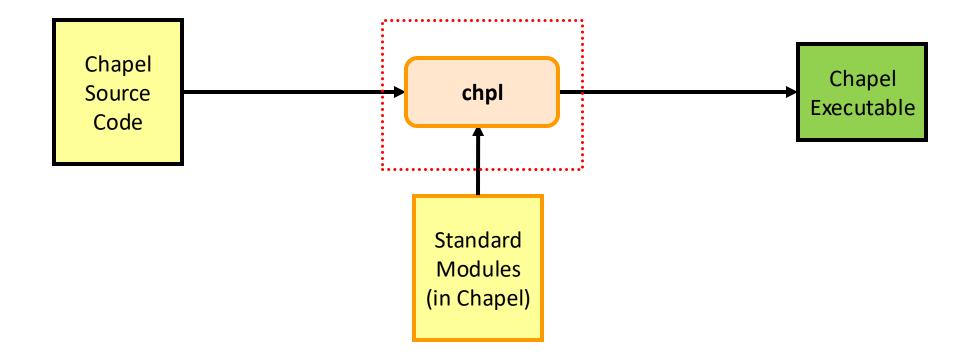
  proc main() {
    writeln("Hello, ", audience, "!");
  }
}
```

01-hello-production-configurable.chpl

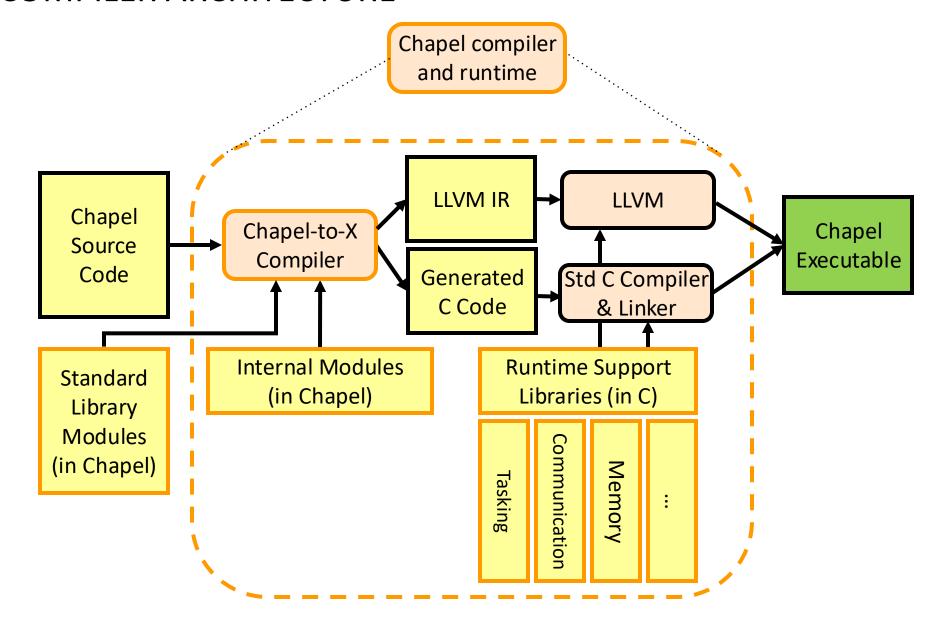
• To change 'audience' for a given run:

./01-hello-configurable -nl 1 --audience="y'all"

COMPILING CHAPEL



CHAPEL COMPILER ARCHITECTURE

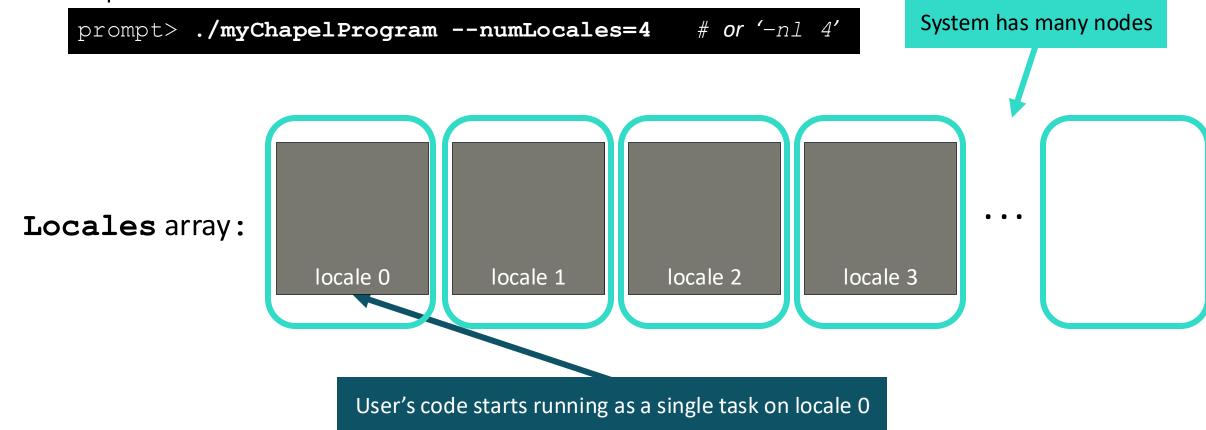


CHAPEL EXECUTION MODEL AND PARALLEL HELLO WORLD (HANDS ON)

CHAPEL EXECUTION MODEL AND TERMINOLOGY: LOCALES

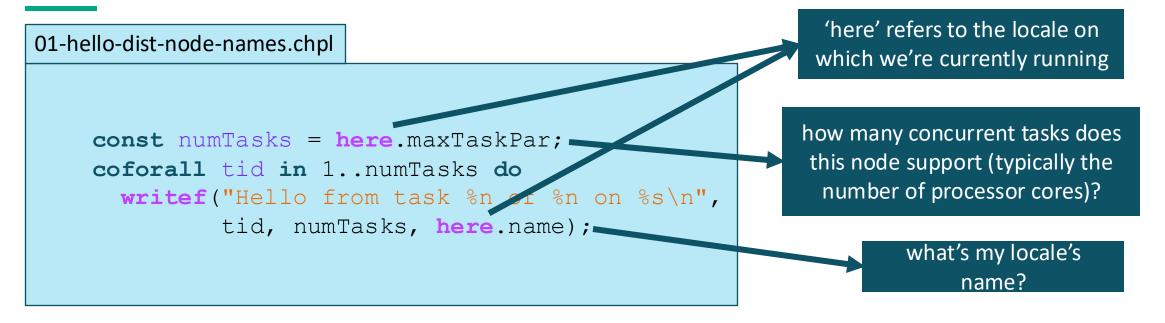
Locales can run tasks and store variables

- Each locale executes on a "compute node" on a parallel system
- User specifies number of locales on executable's command-line



01-hello-dist-node-names.chpl

```
const numTasks = here.maxTaskPar;
coforall tid in 1..numTasks do
   writef("Hello from task %n of %n on %s\n",
        tid, numTasks, here.name);
```



01-hello-dist-node-names.chpl

```
const numTasks = here.maxTaskPar;
coforall tid in 1..numTasks do
   writef("Hello from task %n of %n on %s\n",
        tid, numTasks, here.name);
```

a 'coforall' loop executes each iteration as an independent task

```
> chpl 01-hello-dist-node-names.chpl
> ./01-hello-dist-node-names -nl 1
Hello from task 1 of 4 on n1032
Hello from task 4 of 4 on n1032
Hello from task 3 of 4 on n1032
Hello from task 2 of 4 on n1032
```

01-hello-dist-node-names.chpl

```
const numTasks = here.maxTaskPar;
coforall tid in 1..numTasks do
   writef("Hello from task %n of %n on %s\n",
        tid, numTasks, here.name);
```

```
> chpl 01-hello-dist-node-names.chpl
> ./01-hello-dist-node-names -nl 1
Hello from task 1 of 4 on n1032
Hello from task 4 of 4 on n1032
Hello from task 3 of 4 on n1032
Hello from task 2 of 4 on n1032
```

So far, this is a shared-memory program

Nothing refers to remote locales, explicitly or implicitly

TASK-PARALLEL "HELLO WORLD" (DISTRIBUTED VERSION)



TASK-PARALLEL "HELLO WORLD" (DISTRIBUTED VERSION)

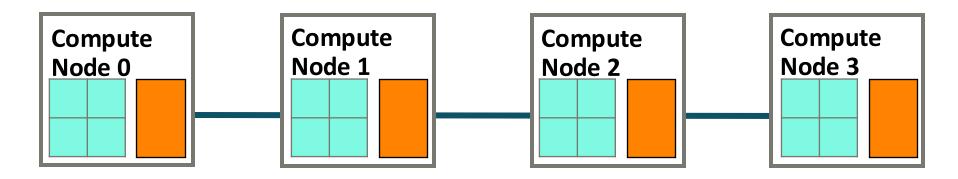
```
create a task per locale
01-hello-dist-node-names.chpl
                                                              on which the program is running
coforall loc in Locales {
  on loc
                                                                 have each task run 'on' its
    const numTasks = here.maxTaskPar;
                                                                         locale
    coforall tid in 1..numTasks do
                                                                 then print a message per
       writef("Hello from task %n of %n on %s\n",
                                                                         core,
               tid, numTasks, here.name);
                                                                       as hefore
                                                           > chpl 01-hello-dist-node-names.chpl
                                                           > ./01-hello-dist-node-names -nl=4
                                                           Hello from task 1 of 4 on n1032
                                                           Hello from task 4 of 4 on n1032
                                                           Hello from task 1 of 4 on n1034
                                                           Hello from task 2 of 4 on n1032
                                                           Hello from task 1 of 4 on n1033
                                                           Hello from task 3 of 4 on n1034
                                                           Hello from task 1 of 4 on n1035
```

LOCALES AND EXECUTION MODEL IN CHAPEL

In Chapel, a locale refers to a compute resource with...

- processors, so it can run tasks
- memory, so it can store variables

For now, think of each compute node as having one locale run on it



Processor Core

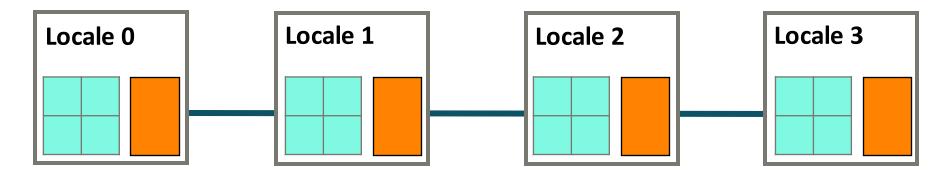
Memory

LOCALES AND EXECUTION MODEL IN CHAPEL

Two key built-in variables for referring to locales in Chapel programs:

• **Locales**: An array of locale values representing the system resources on which the program is running

• here: The locale on which the current task is executing



Processor Core

Memory

GETTING STARTED WITH LOCALES

Users specify # of locales when running Chapel programs

```
% a.out --numLocales=8 % a.out -nl 8
```

Chapel provides built-in locale variables

```
config const numLocales: int = ...;
const Locales: [0..#numLocales] locale = ...;
```

• User's main() begins executing on locale #0, i.e. 'Locales[0]'

LOCALE OPERATIONS

Locale methods support queries about the target system:

```
proc locale.physicalMemory(...) { ... }
proc locale.maxTaskPar { ... }
proc locale.id { ... }
proc locale.name { ... }
```

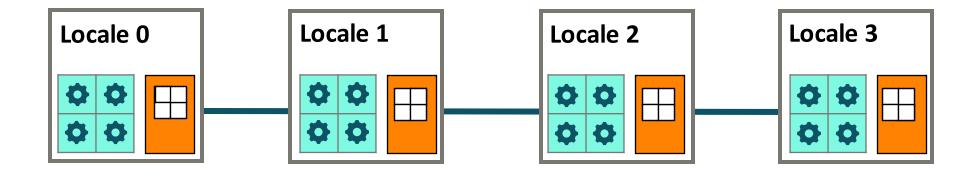
• *On-clauses* support placement of computations:

```
writeln("on locale 0");
on Locales[1] do
   writeln("now on locale 1");
writeln("on locale 0 again");
```

```
on A[i,j] do
  bigComputation(A);
on node.left do
  search(node.left);
```

KEY CONCERNS FOR SCALABLE PARALLEL COMPUTING

- 1. parallelism: Which tasks should run simultaneously?
- 2. locality: Where should tasks run? Where should data be allocated?



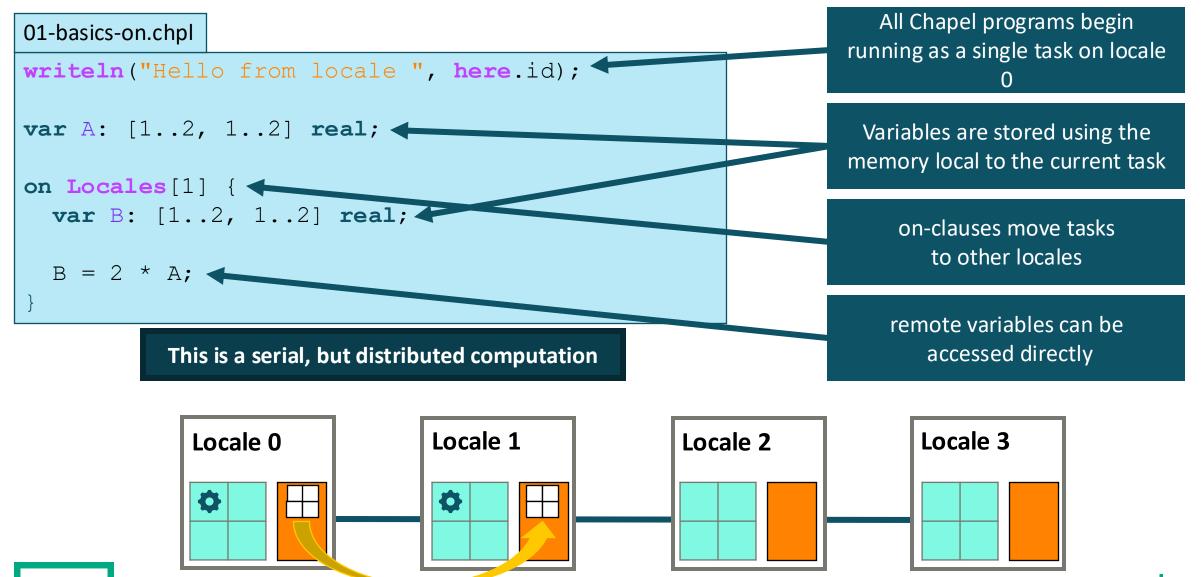
Processor Core

Memory

BASIC FEATURES FOR LOCALITY



01-basics-on.chpl



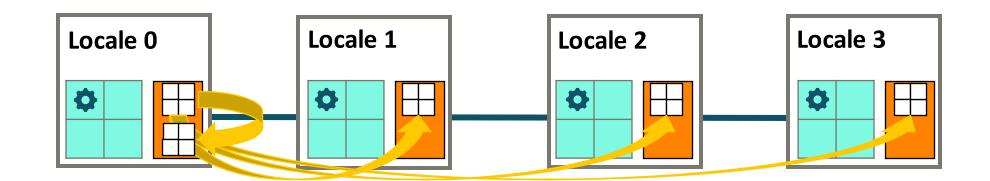
```
01-basics-for.chpl

writeln("Hello from locale ", here.id);

var A: [1..2, 1..2] real;

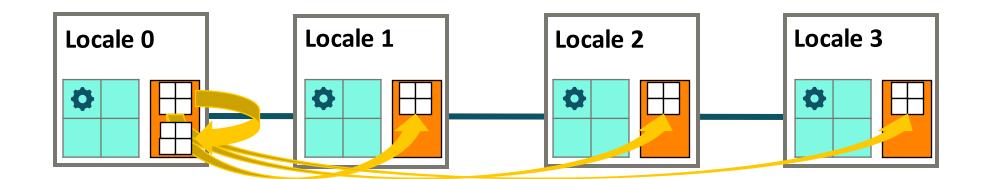
for loc in Locales {
    on loc {
      var B = A;
    }
}
This loop will serially iterate over the program's locales
```

This is also a serial, but distributed computation



01-basics-coforall.chpl writeln("Hello from locale ", here.id); var A: [1..2, 1..2] real; coforall loc in Locales { var B = A; } } The coforall loop creates a parallel task per iteration

This results in a parallel distributed computation



ARRAY-BASED PARALLELISM AND LOCALITY

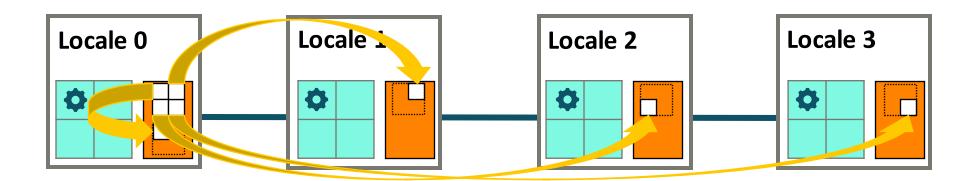


```
01-basics-distarr.chpl
writeln("Hello from locale ", here.id);
var A: [1..2, 1..2] real;
use BlockDist;

var D = blockDist.createDomain({1..2, 1..2}); <
var B: [D] real;
B = A;</pre>
```

Chapel also supports distributed domains (index sets) and arrays

They also result in parallel distributed computation



HANDS ON: PARALLELISM ACROSS AND WITHIN LOCALES



Parallel hello world

• 01-hellopar.chpl

Things to try

```
chpl 01-hellopar.chpl
./01-hellopar -nl 1 --tasksPerLocale=3
./01-hellopar -nl 2 --tasksPerLocale=3
```

Key concepts

- 'coforall' over the `Locales` array with an `on` statement
- 'coforall' creating some number of tasks per locale
- configuration constants, 'config const'
- range expression, '0..<tasksPerLocale'
- 'writeln'
- inline comments start with '//'

```
// can be set on the command line with --
tasksPerI ocale=2
config const tasksPerLocale = 1;
// parallel loops over nodes and then over threads
coforall loc in Locales do on loc {
  coforall tid in 0..<tasksPerLocale {</pre>
    writeln ("Hello world! ",
              "(from task ", tid,
              " of ", tasksPerLocale,
              " on locale ", here.id,
              " of ", numLocales, ")" );
```

PARALLELISM AND LOCALITY ARE ORTHOGONAL IN CHAPEL

1 01-parallelism-and-locality.chpl

• This is a parallel, but local program:

```
coforall i in 1..msgs do
  writeln("Hello from task ", i);
```

• This is a distributed, but serial program:

```
writeln("Hello from locale 0!");
on Locales[1] do writeln("Hello from locale 1!");
on Locales[2] {
   writeln("Hello from locale 2!");
   on Locales[0] do writeln("Hello from locale 0!");
}
writeln("Back on locale 0");
```

This is a distributed parallel program:

```
coforall i in 1..msgs do
  on Locales[i%numLocales] do
  writeln("Hello from task ", i, " running on locale ", here.id);
```

HANDS ON: PARALLELISM AND LOCALITY IN CHAPEL

Goals

- Compile and run some of the examples from the last section
- Experiment some with '01-basics-distarr.chpl'

Compile and run some of the other examples from the last section

```
chpl 01-parallelism-and-locality.chpl
./01-parallelism-and-locality -nl 1
./01-parallelism-and-locality -nl 4
```

Experiment some with '01-basics-distarr.chpl'

- 1. what happens when you add a 'writeln(D)' to write out the domain 'D'?
- 2. what happens when you change 'D's initial value to '{0..3,0..3}'?
- 3. where does the computation on locales other than locale 0 happen?