CHAPEL BASICS, PART II

Chapel Team, edited by Michelle Strout April 8, 2025

Example codes for Chapel tutorial slides

• https://github.com/UofA-CSc-372-Spring-2025/CSc372Spring2025-CourseMaterials/tree/main/Sandboxes/ChapelTutorialExamples

Using a container on your laptop

- First, install docker for your machine and start it up (see the README.md for more info)
- Then, use the chapel-gasnet docker container

```
docker pull docker.io/chapel/chapel-gasnet # takes about 5 minutes
cd CSc372Spring2025-CourseMaterials/Sandboxes/ChapelTutorialExamples/
docker run --rm -it -v "$PWD":/workspace chapel/chapel-gasnet
root@589405d07f6a:/opt/chapel# cd /workspace
root@xxxxxxxxx:/myapp# chpl 01-hello.chpl
root@xxxxxxxxx:/myapp# ./01-hello -nl 1
```

PLAN

Announcements

- Everyone should have their Final Project assignments (go over final project expectations)
- SA7 is due Friday April 11th

• Last time

- TopHat Questions
- ICA10 Quiz
- Chapel Programming Basics in the context of an Nbody simulation, part I

Today

• Chapel programming basics in the context of an Nbody simulation, part II

SA7 EDITS FROM YESTERDAY

• See piazza post, screen shot below for reference

Actions •

SA7 Chapel edits

The SA7 assignment writeup, https://github.com/UofA-CSc-372-Spring-2025/CSc372Spring2025-CourseMaterials/blob/main/SmallAssignmentWriteups/sa7-chapel.md, and initial files, https://github.com/UofA-CSc-372-Spring-2025/sa7-chapel-start, have undergone some modifications.

- The array of files is being tested in sa7-student-tests.chpl using a set instead of just comparing the array values since sometimes the order didn't match.
- The rgb to gray conversion and the sobel edge detector descriptions have been clarified in the writeup.
 They gray value needs to be put in all of the rgb channels, before and after edge detection.
- There is are some tests for the values in the histogram now.
- The Gradescope autograder has been set up as of Monday April 7th at 8:15pm. Already submitted
 assignments have been regraded. The grading tests are somewhat different that the tests provided in sa7student-tests.chpl.

OUTLINE: OVERVIEW OF PROGRAMMING IN CHAPEL

- Main() Procedure
- Ranges and basic control flow
- Procedures and iterators
- Where might we parallelize the n-body computation? (Hands On)

MAIN() PROCEDURE

```
proc main() {
  initSun();
  writef("%.9r\n", energy());
  for 1..numsteps do
    advance(0.01);
  writef("%.9r\n", energy());
```

```
Procedure Definition
proc main() {
  initSun();
  writef("%.9r\n", energy());
  for 1..numsteps do
    advance(0.01);
  writef("%.9r\n", energy());
```

```
proc main() {
                                   Procedure Call
  initSun();
  writef("%.9r\n", energy());
  for 1..numsteps do
    advance(0.01);
  writef("%.9r\n", energy());
```

```
Activity: Using the table at <a href="https://chapel-">https://chapel-</a>
                       lang.org/docs/modules/standard/IO/FormattedIO.html, format
                       the energy values in three different ways.
proc main() {
   initSun();
  writef("%.9r\n", energy());
                                                            Formatted I/O
   for 1...numsteps do
     advance (0.01);
  writef("%.9r\n", energy());
```

```
proc main() {
  initSun();
  writef("%.9r\n", energy());
  for 1..numsteps do
    advance (0.01);
  writef "%.9r\n", energy());
                   Range Value
```

RANGES: INTEGER SEQUENCES

RANGE VALUES: INTEGER SEQUENCES

Syntax

```
range-expr:
[low] .. [high]
```

Definition

Regular sequence of integers

low <= high: low, low+1, low+2, ..., high</pre>

low > high: degenerate (an empty range)

low or high unspecified: unbounded in that direction

Examples

```
1..6 // 1, 2, 3, 4, 5, 6
6..1 // empty
3.. // 3, 4, 5, 6, 7, ...
```

RANGE OPERATORS

```
const r = 1..10;
printVals(r);
printVals(r # 3);
printVals(r by 2);
printVals(r by -2);
printVals(r by 2 # 3);
printVals(r # 3 by 2);
printVals(0.. #n);
proc printVals(r) {
  for i in r do
    write(i, " ");
  writeln();
```

```
1 2 3 4 5 6 7 8 9 10
1 2 3
1 3 5 7 9
10 8 6 4 2
1 3 5
1 3
0 1 2 3 4 ... n-1
```

Activity: Experiment with 02-range-operators.chpl. What are the two different ways the pound sign (#) are being used?

```
proc main() {
  initSun();
 writef("%.9r\n", energy());
  for 1...numsteps do
    advance (0.01);
  writef("%.9r\n", energy());
                                             Serial for loop
```

BASIC SERIAL CONTROL FLOW

Syntax

```
for-loop:
   for [index-expr in] iteratable-expr { stmt-list }
```

Meaning

- Executes loop body serially, once per loop iteration
- Declares new variables for identifiers in *index-expr*
 - -type and const-ness determined by *iteratable-expr*
 - -iteratable-expr could be a range, array, iterator, iterable object, ...

Examples

```
var A: [1..3] string = [" DO", " RE", " MI"];

for i in 1..3 { write(A[i]); } // DO RE MI
for a in A { a += "LA"; } write(A); // DOLA RELA MILA
```

CONTROL FLOW: OTHER FORMS

Conditional statements

```
if cond { computeA(); } else { computeB(); }
```

While loops

```
while cond {
  compute();
}
```

For loops

```
for indices in iteratable-expr {
  compute();
}
```

Select statements

```
select key {
  when value1 { compute1(); }
  when value2 { compute2(); }
  otherwise { compute3(); }
}
```

CONTROL FLOW: BRACES VS. KEYWORDS

Control flow statements specify bodies using curly brackets (compound statements)

Conditional statements

```
if cond { computeA(); } else { computeB(); }
```

While loops

```
while cond {
  compute();
}
```

For loops

```
for indices in iteratable-expr {
  compute();
}
```

Select statements

```
select key {
  when value1 { compute1(); }
  when value2 { compute2(); }
  otherwise { compute3(); }
}
```

CONTROL FLOW: BRACES VS. KEYWORDS

They also support keyword-based forms for single-statement cases

Conditional statements

```
if cond then computeA(); else computeB();
```

While loops

```
while cond do
  compute();
```

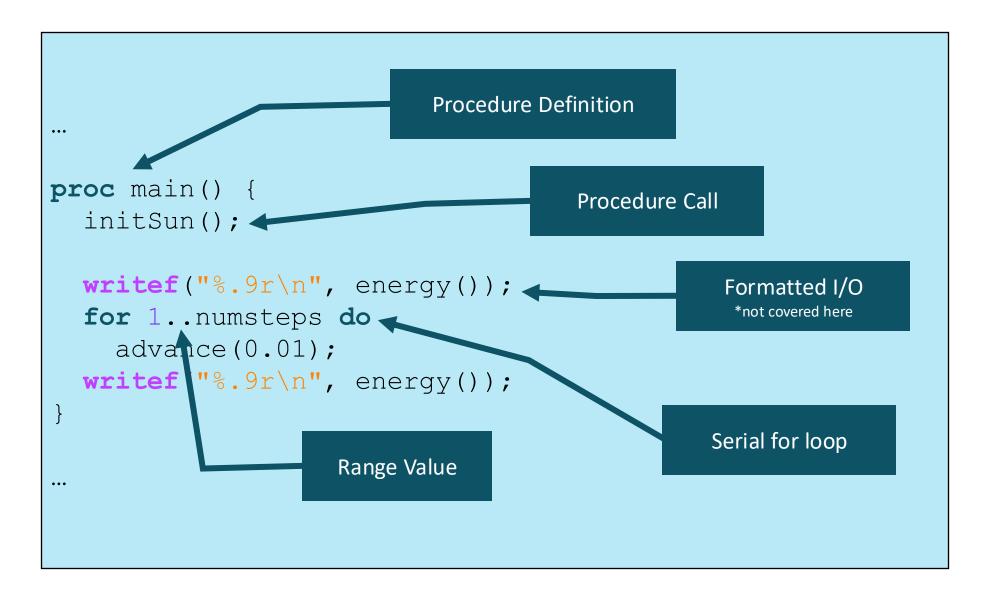
For loops

```
for indices in iteratable-expr do
  compute();
```

Select statements

```
select key {
  when value1 do compute1();
  when value2 do compute2();
  otherwise do compute3();
}
```

PROCEDURES AND ITERATORS



```
advance (0.01);
proc advance(dt) {
  for i in 1..numbodies {
    for j in i+1..numbodies {
      const dpos = bodies[i].pos - bodies[j].pos,
            mag = dt / sqrt(sumOfSquares(dpos))**3;
      bodies[i].v -= dpos * bodies[j].mass * mag;
      bodies[j].v += dpos * bodies[i].mass * mag;
  for b in bodies do
   b.pos += dt * b.v;
```

```
m_1 \mathbf{a}_1 = \frac{Gm_1 m_2}{r_{12}^3} (\mathbf{r}_2 - \mathbf{r}_1) Sun-Earth
advance (0.01);
                                    m_2 \mathbf{a}_2 = \frac{Gm_1 m_2}{r_{21}^3} (\mathbf{r}_1 - \mathbf{r}_2) Earth-Sun
proc advance(dt) {
  for i in 1..numbodies {
     for j in i+1..numbodies {
        const dpos = bodies[i].pos - bodies[j].pos,
                 mag = dt / sqrt(sumOfSquares(dpos))**3;
       bodies[i].v -= dpos * bodies[j].mass * mag;
       bodies[j].v += dpos * bodies[i].mass * mag;
  for b in bodies do
     b.pos += dt * b.v;
```

```
advance(0.01); ←
                                      Procedure call
proc advance(dt) { <---</pre>
  for i in 1..numbodies {
                                       Procedure definition
    for j in i+1..numbodies {
      const dpos = bodies[i].pos - bodies[j].pos,
             mag = dt / sqrt(sumOfSquares(dpos))**3;
      bodies[i].v -= dpos * bodies[j].mass * mag;
      bodies[j].v += dpos * bodies[i].mass * mag;
  for b in bodies do
    b.pos += dt * b.v;
```

• Example to compute the area of a circle

```
proc area(radius: real): real {
  return 3.14 * radius**2;
}
writeln(area(2.0)); // 12.56
```

```
proc area(radius) {
  return 3.14 * radius**2;
}

Argument and return
  types can be omitted
```

• Example of argument default values, naming

ARGUMENT INTENTS

Question: Why might it be useful for the compiler to have this information?

Arguments can optionally be given intents

- (blank): varies with type; follows principle of least surprise
 - -most types: const in or const ref
 - -sync/single vars, atomics: ref
- ref: formal is a reference back to the actual
- const [ref | in]: disallows modification of the formal
- param/type: actual must be a param/type
- in: initializes formal using actual; permits formal to be modified
- out: copies formal into actual at procedure return
- inout: does both of the above

• For some types, argument intents are needed so as to avoid inadvertent races

- Arguments can optionally be given intents.
- 'ref' intent means the actual being passed in will be modified

• Can't pass a 'const' to a 'ref' intent

- Can pass a 'const' to a 'const ref' intent
- However, can't write to a formal coming in as 'const' intent

Question: Why would we want to pass something by reference if we didn't plan on modifying it?

Can't pass 'const' and 'var' into 'param' intents

```
proc foo(param x: real, type t) {
    ...
    ...
}

const r: real,
    A: [1..3] real;

// foo(r, A); // illegal: can't pass vars and consts to params and types
writeln((r, A)); // writes(0.0, [0.0, 0.0, 0.0])
```

• Can pass a literal, param, or a type into 'param' intent

```
proc foo(param x: real, type t) {
    ...
    ...
}

const r: real,
    A: [1..3] real;

foo(1.2, r.type); // OK: passing a literal/param and a type
writeln((r, A)); // writes(0.0, [0.0, 0.0, 0.0])
```

• 'in' intents cause the formal variable to get its own value of the actual argument

• 'out' intents cause the formal value to be copied into actual argument upon return from procedure

• 'inout' intent is a combination of 'in' and 'out' intent

```
proc advance(dt) {
  for i in 1..numbodies {
    for j in i+1..numbodies {
      const dpos = bodies[i].pos - bodies[j].pos,
             mag = dt / sqrt(sumOfSquares(dpos))**3;
      bodies[i].v -= dpos * bodies[j].mass * mag;
      bodies[j].v += dpos * bodies[i].mass * mag;
  for b in bodies do
   b.pos += dt * b.v;
```

Use of iterator proc advance(dt) { for (i,j) in triangle(numbodies) { const dpos = bodies[i].pos - bodies[j].pos, mag = dt / sqrt(sumOfSquares(dpos)) **3; iter triangle(n) { Definition of iterator for i in 1..n do for j in i+1...n do yield (i,j);

```
proc advance(dt) {
  for (i, j) in triangle(numbodies) {
    const dpos = bodies[i].pos - bodies[j].pos,
          mag = dt / sqrt(sumOfSquares(dpos))**3;
    bodies[i].v -= dpos * bodies[j].mass * mag;
    bodies[j].v += dpos * bodies[i].mass * mag;
  for b in bodies do
   b.pos += dt * b.v;
```

HANDS ON: WHERE MIGHT WE CONSIDER PARALLELIZING N-BODY

Look at 'nbody.chpl' and identify...

nbody.chpl

- 'for' loops that can be parallelized
- 'for' loops that need to stay serial to keep meaning

'for' loops that are "mostly" parallel but have something like +=

```
Can be parallelized
Inherently serial loop
```

Can be parallelized but have to avoid races when adding into velocity field

```
for b in bodies do
 b.pos += dt * b.v;
for 1...numsteps do
  advance (0.01);
for i in 1..numbodies {
  for j in i+1..numbodies {
    const dpos = bodies[i].pos - bodies[j].pos,
        mag = dt / sqrt(sumOfSquares(dpos))**3;
   bodies[i].v -= dpos * bodies[j].mass * mag;
   bodies[j].v += dpos * bodies[i].mass * mag;
```

See https://chapel-

lang.org/docs/technotes/reduceIntents.html

Chapel homepage: https://chapel-lang.org

• (points to all other resources)

Social Media:

• Twitter: <a>@ChapelLanguage

• Facebook: <a>@ChapelLanguage

• YouTube: http://www.youtube.com/c/ChapelParallelProgrammingLanguage

Community Discussion / Support:

• Discord: https://discord.com/invite/xu2xg45yqH

• Stack Overflow: https://stackoverflow.com/questions/tagged/chapel

• GitHub Issues: https://github.com/chapel-lang/chapel/issues



DWNLOAD DOCS - LEARN RESOURCES - COMMUNITY BLOG

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