# MEMORY SAFETY: CHAPEL, RUST, C/C++, PYTHON, MPI, OPENSHMEM

Slides based off blog post by Michael Ferguson

https://chapel-lang.org/blog/posts/memory-safety/

Michelle Strout put together slides for CSc 372 at UArizona April 17, 2025

#### **PLAN**

### Announcements

- Brad Chamberlain, Chapel co-creator, visiting Thursday April 24th
- LA3 is due on Friday April 25<sup>th</sup> (1 weeks left)
- Final projects are due Friday May 2<sup>nd</sup> (2 weeks left)

### Last time

- TopHat questions about Chapel data parallelism
- LA3 parallelism suggestions
- Heat diffusion in Chapel and implicit communication

# Today

- Recall extra credit opportunities
- TopHat questions about Chapel implicit communication
- Memory safety comparison between languages

#### **OUTLINE: MEMORY SAFETY**

- Visualizing
  - Use-after-free and memory leak bugs
  - How ownership types prevent them
  - How garbage collection prevents them
- Punchline from blog: comparison of languages
- Kinds of errors memory safety help prevent
  - Variable not initialized
  - Mishandling strings
  - Use-after-free
  - Out-of-bounds array access



chapel-lang.org/blog/posts/memory-safety/

# **Memory Safety in Chapel**

Posted on April 10, 2025.

Tags: Safety Language Comparison

By: Michael Ferguson

# VISUALIZING SOME MEMORY BUGS

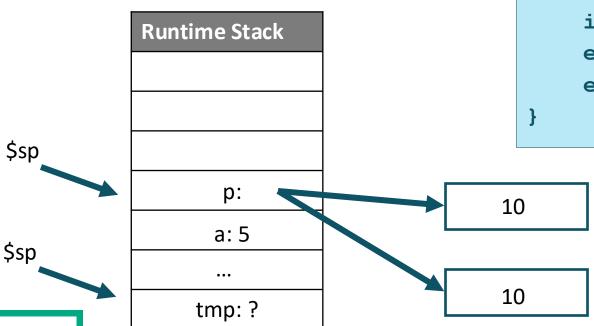
#### IN CLASS ACTIVITY: VISUALIZING SOME MEMORY BUGS

- Learning Objectives
  - Understand how memory is managed using stack frames and dynamic allocation
  - Identify common memory safety issues: use-after-free and memory leaks
  - See how different languages (Rust, Chapel, Java) handle memory safety
  - Role-play borrow checking to understand ownership models

#### RUNTIME STACK AND HEAP DIAGRAMS

### Memory allocation

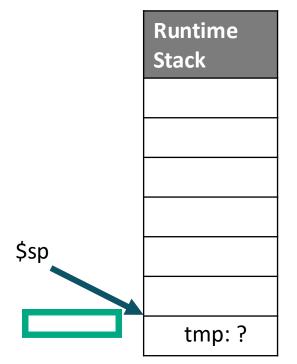
- Local variables are allocated on the runtime stack
- Malloc allocates memory in the heap
- Malloc returns a pointer to that memory in the heap that is typically then stored in the runtime stack

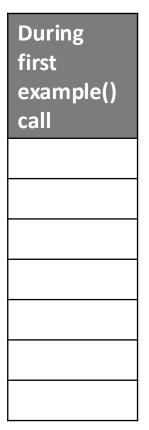


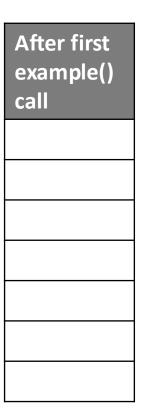
```
// example C code
void example() {
    int a = 5;
                          // local variable
    int* p = malloc(4); // heap allocation
    *p = 10;
} // p goes out of scope, but memory is still allocated unless freed
void main() {
                          // local variable
    int tmp;
    example();
    example();
```

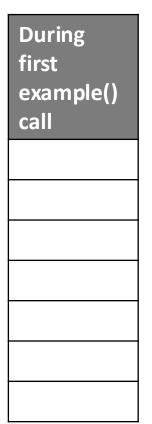
# ACTIVITY: SHOW THE STEPS DURING AND AFTER EACH EXAMPLE() CALL

```
void example() {
    int a = 5;
    int* p = malloc(4);
    *p = 10;
}
void main() {
    int tmp;
    example();
    example();
}
```









After first example() call

#### WHAT HAPPENS WITH THE FOLLOWING TWO EXAMPLES?

```
int* x = malloc(sizeof(int));
*x = 42;
free(x);
printf("%d\n", *x); // use-after-free
```

```
void leak() {
   int* y = malloc(sizeof(int));
   *y = 99;
   // no free(y);
} // memory leak
```

#### OWNERSHIP TYPES IN RUST AND CHAPEL

#### **Ownership types**

- Only one reference var owns the reference
- Can move ownership
- The compiler generates an error if trying to access through a reference that no longer owns something
- When that reference var goes out of scope, the heap allocation is freed

```
// example Rust code
struct MyBox {
    value: i32,
fn main() {
    let a = MyBox { value: 10 };
    let b = a; //ownership moves, a no longer valid
    println!("{}", a.value); //compile error
// example Chapel code
class MyClass {
  var x: int;
proc main() {
  var c = new owned MyClass();
  var d = c; // transfers ownership
  writeln(c); // error: c no longer owns
} // memory automatically freed
```

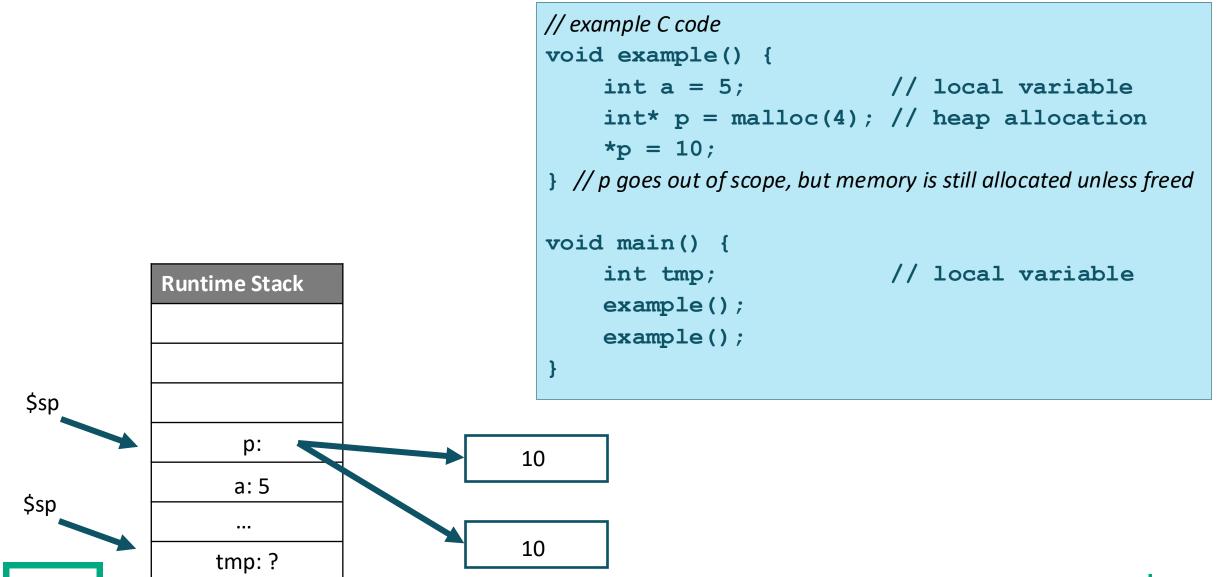
#### GARBAGE COLLECTION IN LANGUAGES LIKE JAVA AND PYTHON

#### **Garbage collection**

- Mark and sweep variants will traverse the active runtime stack and mark everything in the heap that is reachable
- During the sweep, anything in the heap that is not marked will be swept up and freed

```
// example Java code
void foo() {
    MyClass obj = new MyClass();
    obj.doSomething();
} // obj goes out of scope; object becomes
unreachable and is eligible for GC
```

#### HOW WOULD GC HELP WITH THIS EXAMPLE?



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- Kinds of errors memory safety help prevent
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  - Mishandling strings
  - Use-after-free
  - Out-of-bounds array access



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# **Memory Safety in Chapel**

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#### MEMORY SAFETY IN CHAPEL

#### **Memory Safety**

- Properties of programming languages that help reduce bugs
- Red means not much help, yellow means some help, green means protection from error

Error	С	C++	Rust	Python	Chapel
Variable Not Initialized	×	×	<b>▼</b>	<b>~</b>	<b>▽</b>
Mishandling Strings	×	<u> </u>	V	V	V
Use-After-Free	×	1	<u> </u>	$\overline{\checkmark}$	<u> </u>
Out-of-Bounds Array Access	×	×	V	V	<u> </u>

#### **Tradeoffs**

- Garbage collection makes it really easy for users to have memory safe code
- Unfortunately, it has performance impacts
- Thie second grid shows how various languages balance these tradeoffs

	C/C++	Rust	Python	MPI	OpenSHMEM	Chapel
Productivity	-	<b>√</b>	+	-	-	+
Performance	+	+	-	+	+	+
Scalability				+	+	+
Safety	-	+	+	-	-	✓

**Key:** +: great; √: good; -: drawback

#### VARIABLE NOT INITIALIZED

#### Languages that don't help

- C/C++ you can access junk left in stack from previous stack frames
- What if instead of junk it is confidential data?
- Leads to hard to find bugs and potential security issues

## Other approaches

- Rust will generate a compile-time error
- Chapel variables are initialized to a default value for their type
- How does Python work?

```
1  #include <stdio.h>
2  int main() {
3   int x;
4
5   printf("x is %i\n", x);
6   return 0;
7  }
```

```
$ gcc unset-variable.c
$ ./a.out
x is 32764
```

#### unset-variable.rs

```
1  fn main() {
2   let mut x: i64; // OOPS! forgot to initialize x
3   let y = x;
4   println!("y is {}", y);
5  }
```

#### unset-int-variable.chpl

```
proc main() {
   var x: int; // integer variables are set to 0 if not initialized
   writeln(x);
}
```

#### MISHANDLING STRINGS IN C

```
$ gcc string-greeting.c
                                string-greeting.c
                                                                $ ./a.out abcdefghijklmnopqrstuv
     #include <stdio.h>
                                                                Hello abcdefghijklmnopqrstuv
 2
     #include <string.h>
                                                                *** stack smashing detected ***: terminated
 3
                                                                Aborted (core dumped)
 4
     #define MAX_GREETING 16
 5
     int main(int argc, char** argv) {
 6
       char greeting[MAX_GREETING]; // C doesn't really have string support;
                                  // here we allocate an array to store the
 8
                                  // greeting
 9
10
                                  // OOPS! allocated array might not be big enough
11
12
       strcpy(greeting, "Hello "); // copy "Hello " into 'greeting'
13
14
       strcat(greeting, argv[1]); // append the passed name to the greeting
       printf("%s\n", greeting);
15
16
       return 0;
17
```

#### BETTER STRING HANDLING

#### C and C++

- In C, best practice is to use strncpy instead of strcpy and strlcat instead of strcat
- In C++, use the standard string type

### Rust, Python, and Chapel

 All of these languages have built in string types

#### ▼ string-greeting.cpp

#### ▼ string-greeting.chpl

```
config const who = ""; // enable command-line options like --who=world

var greeting = "Hello ";
greeting += who;
writeln(greeting);
```

#### **USE AFTER FREE**

#### C and C++

- In C, use-after-free is very easy to accidentally write
- In C++, unique\_ptr and shared\_ptr help to some extent, but use-after-free is still possible

### Rust, Python, and Chapel

- Rust, compile-time checking prevents a use-after-free in safe code, still possible in unsafe code
- Python, garbage collector avoids this issue
- Chapel, automatic memory management for most types means free is only needed when using unmanaged classes

#### use-after-free-scoped.chpl

```
class C { var x: int; }
      proc main() {
        // create a reference to a 'C' instance on the heap
        var b: borrowed C? = nil;
          var instance = new owned C(0);
          b = instance.borrow();
10
11
        b!.x = 42;
12
13
14
        // do other heap operations to make heap corruption more visible
15
          var x = new owned C(2);
16
17
          var y = new owned C(3);
18
19
        writeln(b);
20
21
```

#### ARRAY BOUNDS CHECKING

#### C and C++

- Out-of-bounds array accesses aren't checked unless running with a memorychecking tool like valgrind
- If you haven't learned valgrind, check it out, it is a life changing tool for debugging C/C++ code

#### Rust, Python, and Chapel

- Rust, out-of-bounds array accesses cause the program to halt with an out-ofbounds error
- Python, out-of-bounds array accesses raise and error
- Chapel, by default out-of-bounds will cause program to halt, compiling with -fast turns checking off

```
$ g++ out-of-bounds.cpp
$ ./a.out 123456789
zsh: segmentation fault ./a.out 123456789
```

```
$ python3 out-of-bounds.py 123456789
Traceback (most recent call last):
   File "/Users/mferguson/chapel-blog/content/
     main(sys.argv)
     ~~~^^^^^^^^^
File "/Users/mferguson/chapel-blog/content/
     x = array[idx]
     ~~~~^^^^^^^
IndexError: list index out of range
```

```
$ chpl out-of-bounds.chpl
$ ./out-of-bounds --idx=123456789
out-of-bounds.chpl:5: error: halt reached - array index out of bounds
note: index was 123456789 but array bounds are 0..9
```