MATLAB for Optimization

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Outline

- Introduction of mathematical optimization
- How to use optimization to model and solve engineering problems
- How to solve optimization models using MATLAB and CVX
- Applications
 - Production plan
 - Smart electric vehicle charging
 - Vehicle-to-grid
 - Cast some machine learning to optimization
 - Stochastic optimization
- Solving optimization problems

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- Applications
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- Solving optimization problems
- ► The slides and all the codes are available at my website: http://uregina.ca/~wang233z

What is an optimization model?

- ➤ A mathematical approach for seeking a "best" "decision/action" from a "set of alternatives"
- ► An objective function that is to be maximized or minimized
- A set of constraints (possibly empty) that must be satisfied

Formal Formulation

Mathematical programs are problems of the form

minimize
$$f(x)$$

subject to $q_i(x) < 0, \forall i = 1, ..., m$ (1)

- $ightharpoonup x \in \mathbb{R}^n$ is the optimization variable
- $f: \mathbb{R}^n \to \mathbb{R}$ is objective function
- $g_i: \mathbb{R}^n o \mathbb{R}$ are (inequality) constraint functions
- ► Feasible region: $C = x : g_i(x) \le 0, \forall i = 1, ..., m$
- $\mathbf{x}^{\star} \in \mathbb{R}^n$ is an optimal solution if $x^{\star} \in \mathcal{C}$, and $f(x^{\star}) \leq f(x), \ \forall x \in \mathcal{C}$

Minimization vs Maximization

- ▶ Without loss of generality, it is sufficient to consider a minimization objective since maximize $\{f(x):x\in\mathcal{C}\}\equiv -\text{minimize}\ \{f(x):x\in\mathcal{C}\}$
- ► Thus to develop the theory we will only consider minimization problems. When actually solving problems we can use the actual min or max objective

Program vs Optimization Problem

- A "program" or "mathematical program" is an optimization problem with a finite number of variables and constraints written out using explicit mathematical (algebraic) expressions
- ► The word "program", "programming" means "plan", "planning"
- Early applications of optimization arose in planning resource allocations and gave rise to "programming" to mean optimization (predates computer programming)
- ► We will use "program", "programming" and "optimization problem", "optimization" interchangeably

Functions	Variables	Problem type	Difficulty

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All linear	Continuous variables	Linear Program (LP) or Linear Optimization problem	easy

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Some nonlinear	Continuous variables	Nonlinear Program (NLP)or	Easy or Difficult	
Some nonlinear	Continuous variables	Nonlinear Optimization Problem	Easy or Dillicuit	
Linear/nonlinear	C d'	Integer Program (IP) or	Difficult	
Linear/ nonlinear	Some discrete	Discrete optimization problem	Dillicuit	

Develop optimization models

- Many problems of real importance can be formulated as an optimization problem
- There are two important aspects
 - Mathematical modeling of engineering problems
 - Problem solving
- Reducing a seemingly new problem to an instance of a well-known problem allows one to use pre-existing methods for solving them

Formulation Steps

- Encode decisions/actions as decision variables whose values we are seeking
- Identify the relevant problem data
- Express constraints on the values of the decision variables as mathematical relationships (inequalities) between the variables and problem data
- Express the *objective function* as a function of the decision variables and the problem data

Ex1: Production plan

- ▶ A company produces two types of drinks: Drink 1 and Drink 2 and they yield different profit.
- ➤ The capacity of the barrels to contain Drink 1 and Drink 2 are 500 L and 400 L.
- Each type drink requires some amount of ingredient A and B.
- ▶ There are 30 L ingredient A and 44 L ingredient B in stock.
- ▶ The manager is developing a plan to maximize the profit.

Drink (L)	A (L)	B (L)	Profit (\$)
Drink 1	0.03	0.08	1.00
Drink 2	0.06	0.04	1.25

► Decision variables:

▶ Decision variables: amount of drinks

i.e., x_1 : Drink 1; x_2 : Drink 2

▶ Decision variables: amount of drinks i.e., x₁: Drink 1; x₂: Drink 2

► Data:

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- Data: required ingredient to produce drinks, amount of ingredient in stock, barrel capacity, profit of each drink

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 - Ingredient constraints:

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$$0.08x_1 + 0.04x_2 \le 44$$

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- ▶ Barrel constraints: $x_1 \le 500, x_2 \le 400$
- Production constraints: $x_1 \ge 0, x_2 \ge 0$
- ▶ Objective function: $1x_1 + 1.25x_2$

minimize
$$-(1x_1+1.25x_2)$$

subject to $0.03x_1+0.06x_2 \leq 30$
 $0.08x_1+0.04x_2 \leq 44$
 $x_1 \leq 500$
 $x_2 \leq 400$
 $x_1, x_2 \geq 0$ (2)

minimize
$$-(1x_1 + 1.25x_2)$$

subject to $0.03x_1 + 0.06x_2 \le 30$
 $0.08x_1 + 0.04x_2 \le 44$
 $x_1 \le 500$
 $x_2 \le 400$
 $x_1, x_2 > 0$ (2)

Linear programming

minimize
$$-(1x_1 + 1.25x_2)$$

subject to $0.03x_1 + 0.06x_2 \le 30$
 $0.08x_1 + 0.04x_2 \le 44$
 $x_1 \le 500$
 $x_2 \le 400$
 $x_1, x_2 > 0$ (2)

- Linear programming
- can be solve by linprog using MATLAB

MATLAB's Optimization Toolbox

▶ Not a complete list

Problem type	Solvers
Linear Programming	linprog
Mixed Linear Programming	intlinprog
Linear Least Squares	lsqlin, lsqnonneg
Quadratic Programming	quadprog

Linear programming solver

- ightharpoonup linprog(c, A, b, Aeq, beq, lb, ub, options)
- Finds the minimum of a problem specified by

minimize
$$c^T x$$

subject to $Ax \leq b$
 $A_{eq}x = b_{eq}$
 $lb \leq x \leq ub$ (3)

where c,x,b,beq,lb, and ub are vectors, and A and Aeq are matrices.

$\underset{x}{minimize}$	$1x_1 + 1.25x_2$	$\underset{x}{minimize}$	$c^T x$
subject to	$0.03x_1 + 0.06x_2 \le 30$	subject to	$Ax \leq b$
	$0.08x_1 + 0.04x_2 \le 44$		$A_{eq}x = b_{eq}$
	$x_1 \le 500$		$lb \le x \le ub$
	$x_2 \le 400$		
	$x_1, x_2 \ge 0$		

ightharpoonup What are A,b,c,x?

 \blacktriangleright What are A, b, c, x?

$$A = \begin{bmatrix} 0.03 & 0.06 \\ 0.08 & 0.04 \\ 1 & 0 \\ 0 & 1 \\ -1 & 0 \\ 0 & -1 \end{bmatrix} b = \begin{bmatrix} 30 \\ 44 \\ 500 \\ 400 \\ 0 \\ 0 \end{bmatrix} c = \begin{bmatrix} 1 \\ 1.25 \end{bmatrix}$$

 \blacktriangleright What are A, b, c, x?

$$A = \begin{bmatrix} 0.03 & 0.06 \\ 0.08 & 0.04 \\ 1 & 0 \\ 0 & 1 \\ -1 & 0 \\ 0 & -1 \end{bmatrix} b = \begin{bmatrix} 30 \\ 44 \\ 500 \\ 400 \\ 0 \\ 0 \end{bmatrix} c = \begin{bmatrix} 1 \\ 1.25 \end{bmatrix}$$

Matlab code: x = linprog(c, A, b) here:

$$x = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$$

Model reformulation (alternative)

ightharpoonup x = linprog(c, A, b, [], [], lb, ub)

$$A = \begin{bmatrix} 0.03 & 0.06 \\ 0.08 & 0.04 \end{bmatrix} \ b = \begin{bmatrix} 30 \\ 44 \end{bmatrix} \ c = \begin{bmatrix} 1 \\ 1.25 \end{bmatrix}$$

$$lb = \begin{bmatrix} 0 \\ 0 \end{bmatrix} \ ub = \begin{bmatrix} 500 \\ 400 \end{bmatrix} \ x = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$$

Quadratic programming solver

- ightharpoonup quadprog(H, c, A, b, Aeq, beq, lb, ub)
- Finds the minimum of a problem specified by

minimize
$$\frac{1}{2}x^T H x + c^T x$$
 subject to
$$Ax \leq b$$

$$A_{eq}x = b_{eq}$$

$$lb \leq x \leq ub$$
 (4)

where c, x, b, beq, lb, and ub are vectors, and H, A and Aeq are matrices.

▶ It is harder to reformulate a problem to standard quadratic program.

CVX

- CVX is a Matlab-based modeling system for convex optimization.
- CVX turns Matlab into a modeling language, allowing constraints and objectives to be specified using standard Matlab expression syntax.

CVX code

```
\begin{array}{ll} \underset{x}{\text{minimize}} & 1x_1 + 1.25x_2 \\ \text{subject to} & 0.03x_1 + 0.06x_2 \leq 30 \\ & 0.08x_1 + 0.04x_2 \leq 44 \\ & x_1 \leq 500 \\ & x_2 \leq 400 \\ & x_1, x_2 \geq 0 \end{array}
```

CVX code

```
minimize 1x_1 + 1.25x_2
                                    cvx_begin
                                      variable x(2);
subject to 0.03x_1 + 0.06x_2 \le 30
                                      minimize -(1*x(1)+
           0.08x_1 + 0.04x_2 < 44
                                         1.25*x(2)
           x_1 \le 500
                                      subject to
                                      0.03*x(1)+0.06*x(2) <= 30:
           x_2 \le 400
                                      0.08*x(1)+0.04*x(2) <=44:
           x_1, x_2 > 0
                                      x(1) \le 500:
                                      x(2) < =400;
                                      x(1) >= 0;
                                      x(2) >= 0;
                                    cvx_end
```

CVX Python code

```
import numpy as np
import cvxpy as cp
n = 2
x = cp.Variable(n)
objective = cp.Minimize(-(100*x[0]+125*x[1]))
constraints = [3*x[0]+6*x[1] <= 30.
                8*x[0]+4*x[1]<=44.
                x[0] \le 5.
                x[1] \le 4.
                x[0] >= 0.
                x[1] >= 0
prob = cp.Problem(objective, constraints)
result = prob.solve()
print(x.value)
```

Electric vehicles (EV)

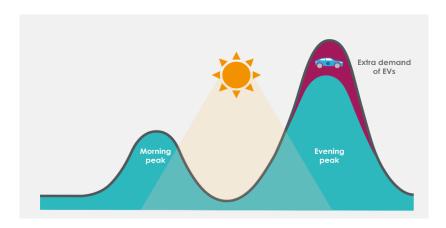
- Lower maintenance costs
- Lower taxes
- ► Cheaper fuel (electricity)
- ► Government subsidy



Basic EV charging

- ▶ Typical battery capacity: 30kWh 100kWh
- ▶ Onboard charger: 1.9 22 kW
- ▶ DC offboard charger: 50 350 kW
- Charging level
 - ► Level 1: 0 10 kW
 - ► Level 2: 10 50 kW
 - ► Level 3: 50 350 kW

Impact on power system



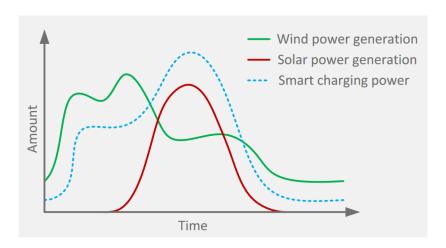
Smart charging and vehicle-to-grid (V2G)

- Using the electric vehicle battery to feed power back to the grid using a bidirectional EV charger
- Advantages: storage for renewables, reduce peak demand, ancillary services, etc
- Challenges: optimal control, bidirectional charger, battery degradation, standardization, regulatory framework, etc

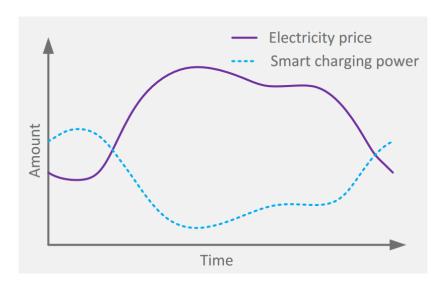
V2G applications

- Local load balancing
 - Adjust charging time/power according to load
 - Balance multiple charge points with priority
- Renewable energy utilization
- Price based charging/discharging
- Peak shaving
- Grid back up

Renewable energy availability



Price based charging

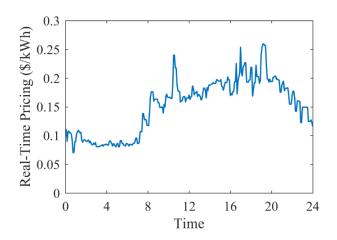


Ref: Electric Cars, edx.org

Ex2: Smart charging

- You have an EV and can be charged at home.
- ► The electricity price is real-time pricing (RTP).
- ▶ You need 20kW to charge up your EV.
- ▶ The rated power of the charger is 2kW.
- Assume that the EV can be charged anytime in the day.
- Now, you need to determine what time to charge to minimize the electricity bill in the day.

Real-time-pricing



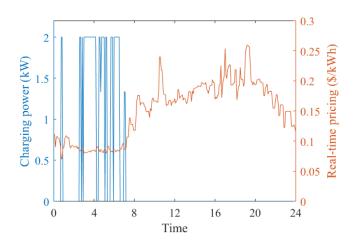
Smart charging model

$$\begin{array}{ll} \text{minimize} & \sum_{t \in \mathcal{T}} \pi_t p_t \\ \\ \text{subject to} & \sum_{t \in \mathcal{T}} p_t = 20 \\ \\ & 0 \leq p_t \leq 2 \end{array}$$

where π_t is price, p_t is charging power.

CVX code

```
load('As3_data')
price = price';
cvx_begin
    variable p(288)
    minimize price * p
    subject to
        sum(p)==20*60/15; %kW15min
        0<=p<=2;
cvx_end</pre>
```



▶ Is this too aggressive?

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 - home arriving, leaving and driving time?
 - state of charge?

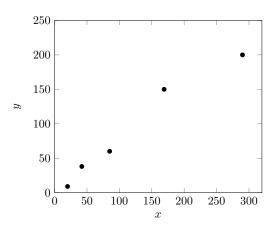
- Is this too aggressive?
- How about battery degradation cost?
- ► How about
 - home arriving, leaving and driving time?
 - state of charge?
 - multiple vehicles?
 - convenience and privacy?
 - power system reliability?

Ex3: V2G

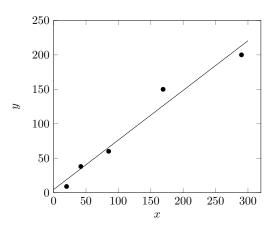
$$\begin{split} & \underset{P_t^{EV_t}}{\text{minimize}} & \text{var} \left(P_t^{EV_t} - P_{t \in [t_1, t_2] \cup [t_3, t_4]}^{EV_t} + P_t^{EV_{-i}} + P_{base, t} \right) \\ & + \upsilon \sum_{t \in T} DC(DOD_t) \left| P_t^{EV_t} \right| + \kappa \sum_{t \in T} RTP_t \ P_t^{EV_t} \\ & \text{subject to:} \\ & - P_{rated} \leq P_t^{EV_t} \leq P_{rated} \ , \forall \ t \in [t_2, t_3] \cup [t_4, t_1] \\ & \sum_{t} P_t^{EV_t} = \beta E_0, \ \forall \ t \in [t_1, t_2] \cup [t_3, t_4] \\ & SOC_t = SOC_0 + \sum_{t=0}^{t} P_t^{EV_t} \\ & SOC_{min} \leq SOC_t \leq SOC_{max} \ , \forall \ t \in T \\ & SOC_{t_2} \geq SOC_{acc} \\ & SOC_{t_4} \geq SOC_{acc} \end{split}$$

Ref: K. Ginigeme and Z. Wang, "Distributed Dobtimal Vehicles To C_t , Approaches with Consideration of Battery Degradation C_t , "Cost under Teacher Pricing", IEEE Access, 2020.

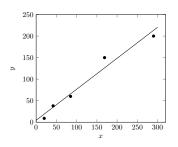
Machine learning: data fitting



Machine learning: data fitting



Machine learning: data fitting



• Given $x_i, y_i, i = 1, ..., m$, find $f(\theta) = \theta_1 * x + \theta_2$ that optimizes

$$\underset{\theta}{\mathsf{minimize}} \sum_{i=1}^{m} (\theta_1 x_i + \theta_2 - y_i)^2$$

where θ_1 is slope, θ_2 is intercept.

Formal problem setting

- ▶ Input: $x_i \in \mathbb{R}^n, i = 1, ..., m$
- ▶ Output: $y_i \in \mathbb{R}, i = 1, ..., m$
- lacktriangle Model parameters: $heta \in \mathbb{R}^k$
- ▶ Predicted output: $\hat{y} \in \mathbb{R}$

Formal problem setting

- ▶ Input: $x_i \in \mathbb{R}^n, i = 1, ..., m$
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- ▶ Model parameters: $\theta \in \mathbb{R}^k$
- ▶ Predicted output: $\hat{y} \in \mathbb{R}$
- Let's define a function that maps inputs to feature vectors

$$\phi: \mathbb{R}^n \to \mathbb{R}^k$$

Then, we can write

$$\hat{y}_i = \sum_{j=1}^k \theta_j \cdot \phi_j(x_i) \equiv \theta^T \phi(x_i)$$

Formal problem setting

$$\hat{y}_i = \sum_{i=1}^J \theta_j \cdot \phi_j(x_i) \equiv \theta^T \phi(x_i)$$

Let's write it in a more compact way

$$\Phi \in \mathbb{R}^{m \times k} = \begin{bmatrix} \phi(x_1)^T \\ \phi(x_2)^T \\ \vdots \\ \phi(x_m)^T \end{bmatrix}, \ y \in \mathbb{R}^m = \begin{bmatrix} y_1 \\ y_2 \\ \vdots \\ y_m \end{bmatrix}$$

Then,

$$\hat{y} = \Phi \theta$$

Loss function

► The loss function is defined as

$$\ell: \mathbb{R} \times \mathbb{R} \to \mathbb{R}_+$$

Square loss

$$\ell(\hat{y}, y) = (\hat{y} - y)^2$$

Absolute loss

$$\ell(\hat{y}, y) = |\hat{y} - y|$$

Deadband loss

$$\ell(\hat{y}, y) = max\{0, |\hat{y} - y| - \varepsilon\}, \varepsilon \in \mathbb{R}_+$$

Data fitting with square Loss

Square loss

$$\ell(\hat{y}, y) = (\hat{y} - y)^2$$

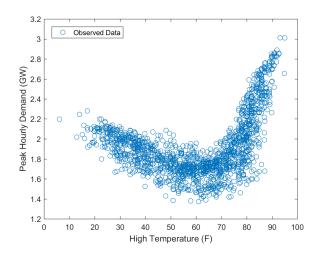
► Called least-squares objective function

$$\underset{\theta}{\operatorname{minimize}} \; (\hat{y} - y)^2 = \underset{\theta}{\operatorname{minimize}} \; \|\hat{y} - y\|_2^2$$

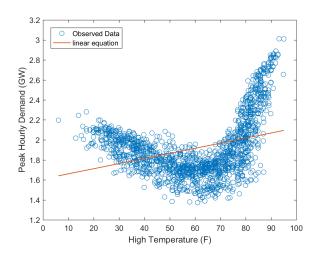
Has analytical solution

$$\theta^{\star} = (\Phi^T \Phi)^{-1} \Phi^T y$$

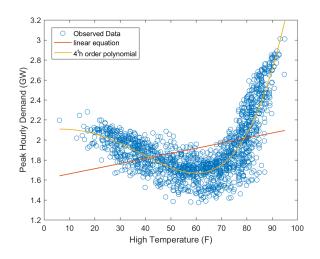
Ex4: Electricity peak demand forecasting



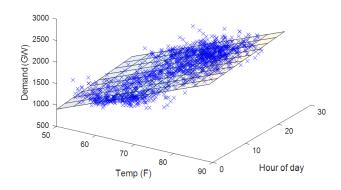
Ex4: Electricity peak demand forecasting



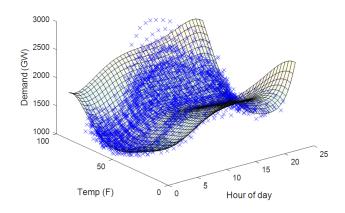
Ex4: Electricity peak demand forecasting



Linear regression 2d vector



"Nonlinear regression" 2d vector



Matlab code

```
X = load('max_temp_b.txt');
y = load('max_demand_b.txt');
[m,n] = size(X);
%% linear
Phi = [X \text{ ones}(m,1)];
% theta = inv(Phi' * Phi) * Phi' * y;
theta = Phi \setminus y;
%% 4th order polynomial
Phi = [X.^4 X.^3 X.^2 X ones(m,1)];
theta = Phi \ y;
```

Ex5: Newsvendor problem

- A company produces winter coats.
- ▶ The company must commit to specific production quantity x before knowing the exact demand d, 3 months before the winter season.

Ex5: Newsvendor problem

- A company produces winter coats.
- ▶ The company must commit to specific production quantity x before knowing the exact demand d, 3 months before the winter season.
- After seeing demand d, the company decides the amount y_r to sell in regular price π_r , and the amount y_s to sell at a salvage/discounted price π_s .
- ► This is called decision making under *uncertainty*, because decision *x* is made under uncertain demand *d*.

Two stage stochastic programming

- Decision variables:
 - Here-and-Now decision: production quantity x
 - lackbox Wait-and-See decision: regular price quantity y_r , discounted price quantity y_s
- Objective: minimize production cost and expected future cost

Two stage stochastic programming

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Two stage stochastic programming

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- Objective: minimize production cost and expected future cost
- Stochastic program

$$\begin{aligned} & \underset{x}{\text{minimize}} & & f(x) + \mathbb{E}_d[Q(x,d)] \\ & \text{subject to} & & 0 \leq x \leq \hat{x} \\ \\ Q(x,d) &= \underset{y_r(d),y_s(d)}{\text{minimize}} & & - \left(\pi_r \; y_r(d) + \pi_s \; y_s(d)\right) \\ & \text{subject to} & & y_r(d) \leq d, \forall d \in \mathcal{D} \\ & & & y_r(d) + y_s(d) \leq x, \forall d \in \mathcal{D} \\ & & & y_r(d), y_s(d) \geq 0, \forall d \in \mathcal{D} \end{aligned}$$

► Combine the two stages

$$\begin{aligned} & \underset{x, \ y_r(d), \ y_s(d)}{\text{minimize}} & f(x) + \mathbb{E}_d[-(\pi_r \ y_r(d) + \pi_s \ y_s(d))] \\ & \text{subject to} & 0 \leq x \leq \hat{x} \\ & y_r(d) \leq d, \forall d \in \mathcal{D} \\ & y_r(d) + y_s(d) \leq x, \forall d \in \mathcal{D} \\ & y_r(d), y_s(d) \geq 0, \forall d \in \mathcal{D} \end{aligned}$$

Combine the two stages

$$\label{eq:subject_to_subject_to} \begin{split} & \underset{x, \, y_r(d), \, y_s(d)}{\text{minimize}} & f(x) + \mathbb{E}_d[-(\pi_r \, y_r(d) + \pi_s \, y_s(d))] \\ & \text{subject to} & 0 \leq x \leq \hat{x} \\ & y_r(d) \leq d, \forall d \in \mathcal{D} \\ & y_r(d) + y_s(d) \leq x, \forall d \in \mathcal{D} \\ & y_r(d), y_s(d) \geq 0, \forall d \in \mathcal{D} \end{split}$$

- Suppose demand d is a discrete random variable with S scenarios $(d_1, ..., d_s)$, and each scenario d_i with a probability p_i .
- ▶ Correspondingly, the sell quantities have $y_{r,i}$ and $y_{s,i}$ for each scenario d_i .

$$\begin{split} & \underset{x, \ y_r(d), \ y_s(d)}{\text{minimize}} & f(x) + \mathbb{E}_d[-(\pi_r \ y_r(d) + \pi_s \ y_s(d)) \\ & \text{subject to} & 0 \leq x \leq \hat{x} \\ & y_r(d) \leq d, \forall d \in \mathcal{D} \\ & y_r(d) + y_s(d) \leq x, \forall d \in \mathcal{D} \\ & y_r(d), y_s(d) \geq 0, \forall d \in \mathcal{D} \end{split}$$

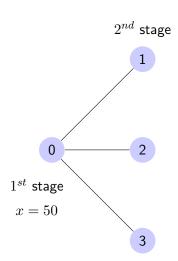
A concrete example

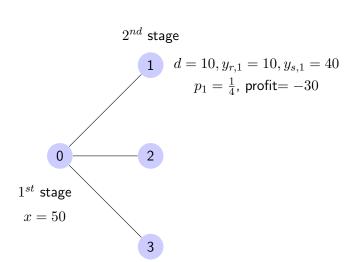
- Suppose there are 3 scenarios, $d_1 = 10$ with probability of $\frac{1}{4}$; $d_2 = 30$ with probability of $\frac{5}{12}$; $d_3 = 50$ with probability of $\frac{1}{3}$.
- ▶ Unit cost to produce coats: c=5, regular price: $\pi_r=10$, discounted price: $\pi_s=3$.
- ▶ Production capacity: $\hat{x} = 70$.

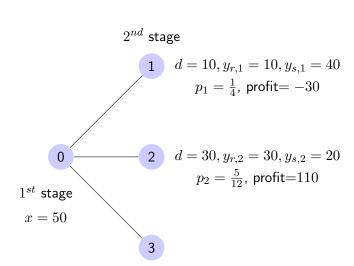
A concrete example

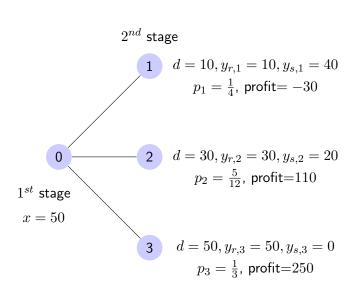
- Suppose there are 3 scenarios, $d_1=10$ with probability of $\frac{1}{4}$; $d_2=30$ with probability of $\frac{5}{12}$; $d_3=50$ with probability of $\frac{1}{3}$.
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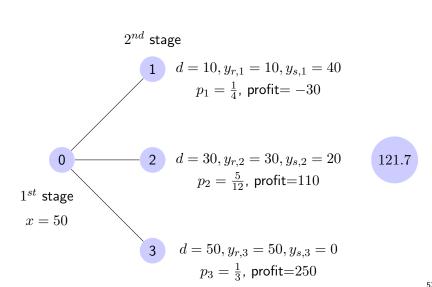
$$\begin{split} \underset{x, \, y_{r,i}, \, y_{s,i}}{\text{minimize}} & \quad 5x - [\frac{1}{4}(10y_{r,1} + 3y_{s,1}) + \frac{5}{12}(10y_{r,2} + 3y_{s,2}) \\ & \quad + \frac{1}{3}(10y_{r,3} + 3y_{s,3})] \\ \text{subject to} & \quad 0 \leq x \leq 70 \\ & \quad y_{r,1} \leq 10, y_{r,2} \leq 30, y_{r,3} \leq 50 \\ & \quad y_{r,i} + y_{s,i} \leq x, \forall i = 1, ..., S \\ & \quad y_{r,i}, y_{s,i} \geq 0, \forall i = 1, ..., S \end{split}$$











Dealing with the General Case

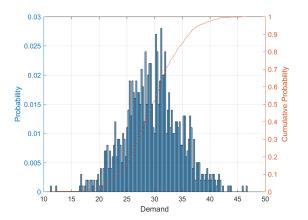
▶ What if it is hard to define scenarios?

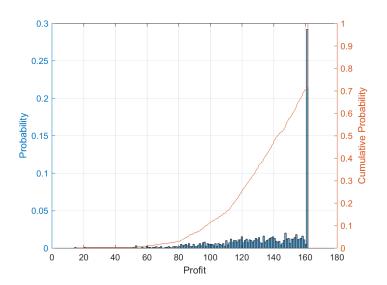
Dealing with the General Case

- ▶ What if it is hard to define scenarios?
- ▶ What if our distribution is not discrete?

Dealing with the General Case

- What if it is hard to define scenarios?
- What if our distribution is not discrete?
- Sampling is enough under most general conditions





Electricity market-clearing problem

 A two-settlement market clearing problem with day ahead (DA) and real time (RT) stages

$$\begin{split} & \underset{p_g^{DA}, p_{g,\omega}^{RT}, p_{d,\omega}^{shed}}{\text{minimize}} & cost^{DA}(p_g^{DA}) + \mathbb{E}_{\omega}[cost^{RT}(p_{g,\omega}^{RT}, p_{d,\omega}^{shed})] \\ & \text{subject to} & f(p_g^{DA}) \leq 0 \\ & g(p_{g,\omega}^{RT}, p_{d,\omega}^{shed}) \leq 0, \ \forall \omega \\ & g(p_g^{DA}, p_{g,\omega}^{RT}, p_{d,\omega}^{shed}) \leq 0, \ \forall \omega \end{split}$$

How is the optimization problem actually solved?

 Unconstrained problem. Set the derivative (one dimensional) or gradient (high dimensional)

$$\nabla_x f(x) = 0$$

- ▶ Then how to find $\nabla_x f(x) = 0$
 - Direct solution: Analytically compute
 - Gradient descent

Repeat:
$$x \leftarrow x - \alpha \nabla_x f(x)$$

Newton's method

Repeat:
$$x \leftarrow x - (\nabla_x^2 f(x))^{-1} \nabla_x f(x)$$

Constrained optimization

Barrier method: Approximate problem via unconstrained optimization

$$\underset{x}{\operatorname{minimize}} f(x) - t \sum_{i=1}^{m} log(-g_i(x))$$

as $t \to 0$, this approaches original problem

Maximize Lagrangian dual problem

$$\underset{\lambda}{\operatorname{maximize}} \; \{ \underset{x}{\operatorname{minimize}} \; f(x) + \sum_{i=1}^{m} \lambda_i g_i(x) \}$$

Practically solving optimization problems

- The good news, for many classes of optimization problems, people have already done all the "hard work" of developing numerical algorithms
- ► A wide range of tools that can take optimization problems in "natural" forms and compute a solution
- Some well-known libraries: CVX (MATLAB), CVXPY (Python), YALMIP(MATLAB), AMPL (custom language), GAMS (custom language), Gurobi (custom language)

Brief history of convex optimization

- ► Theory (convex analysis): ca1900–1970
- Algorithms
 - ▶ 1947: simplex algorithm for linear programming (Dantzig)
 - ▶ 1960s: early interior-point methods (Fiacco & McCormick, Dikin, . . .)
 - ▶ 1970s: ellipsoid method and other subgradient methods
 - ▶ 1980s: polynomial-time interior-point methods for linear programming (Karmarkar 1984)
 - late 1980s-now: polynomial-time interior-point methods for nonlinear convex optimization (Nesterov & Nemirovski 1994)
- Applications
 - before 1990: mostly in operations research; few in engineering
 - since 1990: many new applications in engineering (control, signal
 - processing, communications, circuit design, . . .); new problem classes (semidefinite and second-order cone programming, robust optimization)

MATLAB history

- Invented by Prof. Cleve Moler, University of New Mexico in late 1970s
- ► The MathWorks, Inc. was formed in 1984 by Moler and Jack little. One product: MATLAB
- ► Today: 100 products; over 1 million users worldwide.

Take home messages

- Many problems of real importance can be formulated as an optimization problem.
- Particularly, convex optimization problems that can be solved efficiently and which still find a huge number of applications
- Reducing a seemingly new problem to an instance of a well-known problem allows one to use pre-existing methods for solving them

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Thank you!