## **HOMEWORK 4**

For this week, please answer the following questions from the text. I've copied the problem itself below and the question numbers for your convenience.

- (1) (2.3) Let g be a primitive root for  $\mathbb{F}_p$ .
  - (a) Suppose that x = a and x = b are both integer solutions to the congruence  $gx = h \mod p$ . Prove that  $a = b \mod (p-1)$ . Explain why this implies the map (2.1) on page 65 is well-defined.
  - (b) Prove that

$$\log_g(h_1h_2) = \log_g(h_1) + \log_g(h_2)$$

for all  $h_1, h_2 \in \mathbb{F}_p$ .

(c) Prove that

$$\log_q(h^n) = n \log_q(h)$$

for all  $h \in \mathbb{F}_p$  and  $n \in \mathbb{Z}$ .

- (2) (2.4) Compute the following discrete logarithms:
  - (a)  $\log_2(13)$  for the prime 23, i.e., p = 23, g = 2, and you must solve the the congruence  $2^x = 13 \mod 23$ .
  - (b)  $\log_{10}(22)$  for the prime p = 47.
  - (c)  $\log_6 27(608)$  for the prime p = 941. (Hint: Look in the second column of Table 2.1 on page 66.)
- (3) (2.16) Verify the following assertions from Exmaple 2.16.
  - (a)  $x^2 + \sqrt{x} = \mathcal{O}(x^2)$
  - (b)  $5 + 6x^2 37x^5 = \mathcal{O}(x^5)$
  - $(c) k^{300} = \mathcal{O}(2^k)$
  - (d)  $(\ln k)^{375} = \mathcal{O}(k^{0.001})$
  - (e)  $k^2 2^k = \mathcal{O}(e^{2k})$
  - (f)  $N^{10}2^N = \mathcal{O}(e^N)$
- (4) (1.44) Consider the Hill cipher defined by (1.11)

$$e_k(m) = k_1 \cdot m + k_2 \mod p \text{ and } d_k(m) = k_1^{-1} \cdot (c - k_2) \mod p$$

where m, c, and  $k_2$  are column vectors of dimension n, and  $k_1$  is an  $n \times n$ -matrix.

- (a) We use the Hill cipher with p = 7 and the key  $k_1 = \begin{pmatrix} 1 & 3 \\ 2 & 2 \end{pmatrix}$  and  $k_2 = \begin{pmatrix} 5 \\ 4 \end{pmatrix}$ .
  - (i) Encrypt the message  $m_1 = \begin{pmatrix} 2 \\ 1 \end{pmatrix}$ .
  - (ii) What is the matrix  $k_1^{-1}$  used for decryption?
  - (iii) Decrypt the message  $c = \begin{pmatrix} 3 \\ 5 \end{pmatrix}$ .
- (b) Explain why the Hill cipher is vulnerable to a known plaintext attack.

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(c) The following plaintext/ciphertext pairs were generated using a Hill cipher with the prime p = 11. Find the keys  $k_1$  and  $k_2$ .

$$m_1 = \begin{pmatrix} 5 \\ 4 \end{pmatrix}, c_1 = \begin{pmatrix} 1 \\ 8 \end{pmatrix}, m_2 = \begin{pmatrix} 8 \\ 10 \end{pmatrix}, c_2 = \begin{pmatrix} 8 \\ 5 \end{pmatrix}, m_3 = \begin{pmatrix} 7 \\ 1 \end{pmatrix}, c_3 = \begin{pmatrix} 8 \\ 7 \end{pmatrix}$$

- (d) Explain how any simple substitution cipher that involves a permutation of the alphabet can be thought of as a special case of a Hill cipher.
- (5) (1.48) Explain why the cipher

$$e_k(m) = k \oplus m$$
 and  $d_k(c) = k \oplus c$ 

defined by XOR of bit strings is not secure against a known plaintext attack. Demonstrate your attack by finding the private key used to encrypt the 16-bit ciphertext c = 1001010001010111 if you know the corresponding plaintext is m = 0010010000101100.