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Metacircular Virtualization & Practical Nock Interpretation

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Abstract

This paper presents a detailed examination of meta-circular virtualization in the Nock combinator calculus, focusing on the practical implementation of a virtualized Nock interpreter within the Urbit ecosystem. We explore the design and functionality of key components, including the ++mink interpreter, and their roles in enabling robust Nock execution environments. Extensions of Nock using fake opcodes, particularly opcode 12, are discussed, and a minimal working example of a Nock-virtualized Nock formula is provided.

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1 Introduction

The Nock combinator calculus is used as a target instruction set architecture (ISA) for the Hoon and Jock languages. However, Nock itself is a “crash-only” paradigm which has no error handling mechanism within the terms of the ISA itself. Furthermore, sometimes information is necessary to complete a calculation which may not be present in the apparent subject. Practical Nock interpreters handle this by running Nock within Nock, a process of metacircular virtualization. This article explores the design of `++mock`, the basic virtualized Nock interpreter, and its role within a Nock-based kernel and standard library.

2 Background

In 1958, John McCarthy began his lifelong work on the Lisp programming language (McCarthy1996, McCarthy1996). While the form would evolve over the years, the fundamental components included S-expressions, or parenthesized lists of homoiconic code–data; `'` or quote, a deferred expression; and

eval, a function to evaluate quoted code.¹ Together these components allowed McCarthy and colleague Steve Russell to implement a Lisp interpreter in Lisp itself, the first metacircular interpreter (**McCarthy1978, McCarthy1978**). Such was a radical departure from the imperative programming paradigm of the time, exemplified in Fortran (1954–55); Lisp’s reliance on recursion and higher-order functions was both a precursor to the functional program paradigm and a rejuvenation of the Church and Curry schools of fundamental computing logic.

Metacircularity was considered sufficiently important that several programming languages implemented it as a first-class feature, including Forth, TeX, and Prolog. They have generally been used either for tightly interpreted features like debuggers and code introspection tools, or for language extensions like macros.

Amin and Rompf (**Amin2017**) analyzed how metacircular interpreters build on top of each other, each adding a new layer of abstraction. While primarily examining a Lisp dialect, they showed practically how interpreted language extensions (even in cascading towers) can be collapsed to bottom-level interpretation, similar to Nock’s own `++mink` function. They concisely epitomized a line of research on reflective languages and self-interpreted languages, such as quasiquotation in Lisp (**Bawden1999, Bawden1999**) and Template Meta-Haskell (**Sheard2002, Sheard2002**). The ultimate point of this enquiry was to show that metacircularity was a powerful and well-grounded tool for building interpreters, and that it could be used to in fact build practical interpreters.

Thus, before Nock was even named, one of its earliest design criteria was the ability to straightforwardly implement a metacircular interpreter (`~sorreg-namtyv; ~sorreg-namtyv`, 2006; 2010). While the original motivation was to implement a hyper-Turing operator over a referentially transparent bound namespace, Nock’s support for virtualization also soon proved useful for managing error traces and crashes. Yarvin wrote in an early document:

¹Paul Graham, in his 2001 short essay “What Made Lisp Different”, grouped these concepts into the phrase “The whole language always available” (**Graham2001, Graham2001**).

Metacircularity in most languages is a curiosity and corner case. Programmers are routinely warned not to use it, and for good reason. Practical metacircularity is a particular strength of Nock, or any functional assembly language. Dynamic loading and/or virtualization works, at least logically, as in an OS.

We hypothesize that a metacircular functional dissemination protocol can serve as a complete, general-purpose system software environment—there is no problem it cannot solve, in theory or in practice. Such a protocol can be described as an “operating function” or *OF*—the functional analog of an imperative operating system. (*~sorreg-namtyv*, 2010)

Metacircularity thus closely ties to the idea of a functional operating system. This is a theme that has been explored in other contexts, such as the Lisp machines and the Haskell GHC runtime system. For Nock, a virtualized interpreter is no mere corner case but a first-class component of the system, both for practical execution and for resource discovery.

3 Virtualized Nock

Given an argument for metacircularity, we can now examine how Nock is actually virtualized in practice. The Urbit standard library provides a metacircular Nock interpreter, `++mink`, along with several other components for building virtualized Nock environments. Some essential components provided by vanilla Hoon for virtualizing Nock include:

- `++mink`, compute bottom-level virtual Nock.
- `++mock`, compute formula on subject with hint. Wraps `++mink`; common endpoint.
- `++mule`, kick a trap (deferred computation). Wraps `++mock`; common endpoint.

- `++soft`, wrap a typecast as a unit (i.e., fail with `~` rather than crash).

Nock execution environments virtualize Nock for two primary reasons:

1. Crash handling (provisional execution).
2. Language extension (additional opcodes).

3.1 Crash handling

While virtualization necessarily incurs some overhead relative to more direct execution, it gives the Nock programmer a way to bail (crash), trace (log with debugging information), and bind (time-limit) a computation. In fact, stack traces are computed through virtualization, and so are guaranteed to be as correct as their implementation.

A Nock virtualization environment needs to distinguish at least three kinds of results:

1. A successful result of computation, along with a valid Nock noun.
2. A block, which indicates that the computation is not yet complete and should be continued.
3. A halt or deterministic error, which indicates that the computation has failed and should not be continued.

These correspond to standard scry results (see “Opcode 12” below) in this order:

1. [`~ ~ noun`], a successful result.
2. `~`, a result is not currently available (block).
3. [`~ ~`], a result will never become available (halt).

Because Nock is untyped and we need to distinguish a “zero” from a “block”, the structure of each scry result fundamentally differs from the other options.

Nondeterministic errors—crashes in the runtime, such as an out-of-memory error—are not handled by or visible to the Nock interpreter as a solid-state computer.

3.2 Language extension

Virtualization enables the extension of Nock using “fake opcodes”. At the current time, only opcode 12 has been added;² this opcode provides a generalized way to supply the OS context to the scope in a way that appears like a GOTO or remote call but actually preserves Nock semantics including strict subject scoping. There is a sense in which the scry operation and the bound (referentially transparent) scry namespace are identical to the metacircularity of the Nock interpreter itself. The operating function itself can be treated as a first-class member of its own ecosystem.

Hoon defines `+$nock` already suited for virtualized Nock execution including fake opcode 12 (Listing 1).

However, while `++mink` is compiled to and evaluated in Nock, it is written in Hoon and accordingly uses Hoon’s affordances, particularly vases. Vases are a pair of type and data intended to store relevant type information for correctly evaluating raw untyped Nock nouns.³ Accordingly, it is Hoon’s higher-level introspective tools that facilitate `++mink`’s operation.

The remainder of this section analyzes how each `++mink` implements each Nock opcode, including particularly the fake opcode 12. The following code samples are taken from `/sys/hoon` in the Urbit codebase.

`++mink` produces a head atom marking the kind of result the interpreter has raised: if a `%0`, the result is a valid Nock noun; if a `%1`, a block is indicated (i.e. the halting problem continues); if a `%2`, then a deterministic error has occurred with an error trace.

3.3 Opcode 0

Opcode 0 is the identity function, which returns its argument unchanged. This is implemented in `++mink` as follows:

²Compare, however, the approach of **Atman2025** (**Atman2025**), pp. XX–XX in this volume, who presented an explicitly constructive Nock-like paradigm in Ax.

³See **Davis2025** (**Davis2025**), pp. XX–XX in this volume, for a discussion of vases and their role in practical Nock compilation.

Listing 1: Hoon's +\$nock specification

```

::          ::::: virtual knock
+$  knock  $^  :: autocons
          [p=nock q=nock]
          $%    :: constant
5         [%1 p=★]
          :: compose
          [%2 p=nock q=nock]
          :: cell test
          [%3 p=nock]
10        :: increment
          [%4 p=nock]
          :: equality test
          [%5 p=nock q=nock]
          :: if, then, else
15        [%6 p=nock q=nock r=nock]
          :: serial compose
          [%7 p=nock q=nock]
          :: push onto subject
          [%8 p=nock q=nock]
20        :: select arm and fire
          [%9 p=@ q=nock]
          :: edit
          [%10 p=[p=@ q=nock] q=nock]
          :: hint
25        [%11 p=$@(@ [p=@ q=nock]) q=nock]
          :: grab data from sky
          [%12 p=nock q=nock]
          :: axis select
          [%0 p=@]
30        ==

```

```

187
188 :: *[a 0 b]      /[b a]
189   [%0 axis=]
190 =/ part (frag axis.formula subject)
191 ?~ part [%2 trace]
192 [%0 u.part]
193

```

194 In essence, an axis is provisionally extracted from the subject;
 195 if this is ~ then an error is raised with the current execution
 196 trace, else the resulting value from that axis is produced.

197 3.4 Opcode 1

198 Opcode 1 is the constant function, which returns a constant
 199 value extricated from the formula. This is implemented in ++mink
 200 as follows:

```

201
202 :: *[a 1 b]      b
203   [%1 constant=]
204 [%0 constant.formula]
205

```

206 3.5 Opcode 2

207 Opcode 2 is the evaluation function, which evaluates the for-
 208 mula on the subject. This is implemented in ++mink as follows:

```

209
210 :: *[a 2 b c]    *[*[a b] *[a c]]
211   [%2 subject=* formula=]
212 =/ subject $(formula subject.formula)
213 ?. ?=(%0 -.subject) subject
214 =/ formula $(formula formula.formula)
215 ?. ?=(%0 -.formula) formula
216 %= $
217   subject product.subject
218   formula product.formula
219 ==
220

```

221 Errors are passed directly through, whereas the resulting for-
 222 mula is invoked on the resulting subject.

223 3.6 Opcode 3

224 Opcode 3 is the cell test function, which tests whether the sub-
 225 ject is a cell. This is implemented in ++mink as follows:

```
226 :: *[a 3 b]      ?*[a b]
227               [%3 subject=*]
228 =/ argument $(formula argument.formula)
229 ?. ?=(%0 -.argument) argument
230 [%0 .?(product.argument)]
231
```

233 This reduces to an evaluation of a base-level opcode 3 after a
 234 preliminary check that the formula evaluates correctly.

235 3.7 Opcode 4

236 Opcode 4 is the increment function, which increments the sub-
 237 ject. This is implemented in ++mink as follows:

```
238 :: *[a 4 b]      **[a b]
239               [%4 subject=*]
240 =/ argument $(formula argument.formula)
241 ?. ?=(%0 -.argument) argument
242 ?^ product.argument [%2 trace]
243 [%0 .+(product.argument)]
244
```

246 Note that like opcode 3, this reduces to an evaluation of a base-
 247 level opcode 4.

248 3.8 Opcode 5

249 Opcode 5 is the equality test function, which tests whether the
 250 subject and formula are equal. This is implemented in ++mink
 251 as follows:

```
252 :: *[a 5 b c]    =[*[a b] *[a c]]
253               [%5 subject=* formula=*]
254 =/ a $(formula a.formula)
255 ?. ?=(%0 -.a) a
256 =/ b $(formula b.formula)
257 ?. ?=(%0 -.b) b
258 [%0 =(product.a product.b)]
259
```

261 Like opcodes 3 and 4, this reduces to an evaluation of a base-
 262 level opcode 5 after checking the left-hand and right-hand for-
 263 mulas.

264 3.9 Opcode 6

265 Opcode 6 is the if-then-else or conditional branch function,
 266 which evaluates the formula on the subject and then evaluates
 267 either the first or second formula depending on the result. This
 268 is implemented in ++mink as follows:

```

269 :: *[a 6 b c d]   *[a *[[c d] 0 *[[2 3] 0 *[a 4 4 b]]]]
270   [%6 test=* yes=* no=*]
271   /=  result  $(formula test.formula)
272   ?.  ?=(%0 -.result)  result
273   ?+  product.result
274       [%2 trace]
275       %&  $(formula yes.formula)
276       %|  $(formula no.formula)
277   ==
278
279
```

280 As with the direct opcode 6, there is no short-circuit evaluation
 281 in the conditional statement.

282 3.10 Opcode 7

283 Opcode 7 is the composition function, which evaluates the for-
 284 mula on the subject and then evaluates the next formula on the
 285 result. This is implemented in ++mink as follows:

```

286 :: *[a 7 b c]           *[[a b] c]
287   [%7 subject=* next=*]
288   /=  subject  $(formula subject.formula)
289   ?.  ?=(%0 -.subject)  subject
290   %=  $
291       subject  product.subject
292       formula  next.formula
293   ==
294
295
```

296 This implementation simply drops the evaluation of `subject`
 297 through as the explicit subject for the next formula.

3.11 Opcode 8

Opcode 8 is the variable push function, which pushes a value onto the subject. This is implemented in ++mink as follows:

```

301
302 :: *[a 8 b c]          *[[*[a b] a] c]
303   [%8 head=* next=*]
304   =/ head $(formula head.formula)
305   ?. ?=(%0 -.head) head
306   %= $
307   subject [product.head subject]
308   formula next.formula
309   ==
310
```

Here the conceptual differences between opcode 7 and opcode 8 can be sharply compared: how 7 requires an evaluation against the current subject, leading to something of the nature of $c(b(a))$ in terms of function composition, versus how 8 simply augments the subject.

3.12 Opcode 9

Opcode 9 is the arm select function, which selects an arm of the subject and evaluates it. This is implemented in ++mink as follows:

```

320
321 :: *[a 9 b c]          *[[*[a c] 2 [0 1] 0 b]
322   [%9 axis=@ core=*]
323   =/ core $(formula core.formula)
324   ?. ?=(%0 -.core) core
325   =/ arm (frag axis.formula product.core)
326   ?~ arm [%2 trace]
327   %= $
328   subject product.core
329   formula u.arm
330   ==
331
```

For all the mechanics involved in opcode 9, the implementation is perhaps surprisingly simple. Extract the arm at the axis and evaluate it structurally using an opcode 2. (As an aside, Hoon which compiles to Nock aggressively utilizes lambdas to define

functions at the point of use, and frequently employs a gate-building gate pattern. This latter resolves to two opcode 9s in a series.)

3.13 Opcode 10

Opcode 10 is the edit function, which edits a value in the subject at the given axis. This is implemented in ++mink as follows:

```

342 :: *[a 10 [b c] d] #[b *[a c] *[a d]]
343    [%10 [axis=@ value=*] target=*]
344    ?= (0 axis.formula) [%2 trace]
345    /= target $(formula target.formula)
346    ?= (%0 -.target) target
347    /= value $(formula value.formula)
348    ?= (%0 -.value) value
349    /= mutant=(unit *)
350    (edit axis.formula product.target product.value)
351    ?~ mutant [%2 trace]
352    [%0 u.mutant]
353
```

++edit traverses the noun to attempt to access a given axis; if it succeeds, it returns a new noun with the original value at that axis replaced by a new target value. If it fails, it returns an empty unit, indicating without a crash that the edit failed.

There's more plumbing here because of the mutation of the noun, but the essential concept of locating and replacing a valid subtree is straightforward.

3.14 Opcode 11

Opcode 11 is the hint function, which provides a hint to the interpreter about how to interpret the argument. This is the only way to produce side effects from Nock as a pure function, and ++mink respects and passes through hints raised by the virtualized Nock formula. This is implemented in ++mink as follows:

```

369 :: *[a 11 b c]          *[a c]
370 :: *[a 11 [b c] d]    *[[*[a c] *[a d]] 0 3]
371    [%11 tag=@ next=*
```

```

373 =/ next $(formula next.formula)
374 ?. ?=(%0 -.next) next
375 :- %0
376 .* subject
377 [11 tag.formula 1 product.next]
378 ::
379 [%11 [tag=@ clue=*] next=*]
380 =/ clue $(formula clue.formula)
381 ?. ?=(%0 -.clue) clue
382 =/ next
383 =? trace
384 ?=(?(%hunk %hand %lose %mean %spot) tag.formula)
385 [[tag.formula product.clue] trace]
386 $(formula next.formula)
387 ?. ?=(%0 -.next) next
388 :- %0
389 .* subject
390 [11 [tag.formula 1 product.clue] 1 product.next]
391

```

Two options match for opcode 11, a static and a dynamic hint. Even though this virtualized code “sandboxes” execution, the Nock formula including the `clue` is in fact evaluated. The runtime listens for a `clue`, which is compared to a whitelisted set of supported hints, as in normal (nonvirtualized) Nock execution.

Note that the argument cannot simply be discarded by the interpreter, because it may result in a crash. Thus, as with direct Nock execution, the hint must be evaluated.⁴

⁴Here, this should be read “evaluated safely”, with the caveat that some hints may result in a runtime crash. For instance, an ill-formed `%slog` hint (`printf`) will crash the interpreter if the `clue` is not structured correctly. This is a consequence of the fact that Nock is untyped and so the interpreter cannot know whether the hint is well-formed or not. The `++mink` interpreter does not attempt to validate hints, but merely passes them through to the runtime for evaluation. (One could imagine a Nock virtualizer which itself handled certain hints as if they were “side effects” to the evaluation.)

401 3.15 Opcode 12

402 Opcode 12 is the `scry` function, which retrieves a value from the
 403 sky namespace environment.⁵ This is implemented in `++mink`
 404 as follows:

```

405     [%12 ref=* path=*]
406   /= ref $(formula ref.formula)
407   ?. ?=(%0 -.ref) ref
408   /= path $(formula path.formula)
409   ?. ?=(%0 -.path) path
410   /= result (scry product.ref product.path)
411   ?~ result
412   [%1 product.path]
413   ?~ u.result
414   [%2 [%hunk product.ref product.path] trace]
415   [%0 u.u.result]
416
417
```

418 Opcode 12 does not have a direct Nock expression since it “steps
 419 outside” the subject–formula pair via the expedient of the `scry`
 420 gate.

421 In order to interpret opcode 12, `++mink` takes as an
 422 argument, in addition to the subject and formula,
 423 a gate which will be invoked with the results of
 424 the subformulas of 12 as a sample. The result of
 425 executing this gate (not as interpreted by `++mink`
 426 but by whichever interpreter is running `++mink`)
 427 are then treated as the result of the Nock 12 oper-
 428 ation. (“4n `++mink`”, Urbit documentation)

429 Much of consequence is entailed by the use of a `scry` gate,
 430 which is a Hoon function that retrieves a value from an asso-
 431 ciated sky namespace (nominally, the bound `scry` namespace,
 432 although this is not necessarily entailed by Urbit in contempo-
 433 rary practice). The `scry` function bears the signature

```

434 $-(^ (unit (unit)))
```

⁵There is a sense in which opcodes 6–11 are also virtualized extensions to Nock, acting as they do as macros.

The result of this operation is either a noun or a block, and it is returned as a (unit (unit noun)). These correspond to standard scry results in this order:

1. [`~ ~ noun`], a successful result.
2. `~`, a result is not currently available (block).
3. [`~ ~`], a result will never become available (halt).

For instance, the reference `roof` function for Arvo, the Urbit kernel, is schematically defined as:

```

443 ++ roof
444 |$ [vase]
445 $- $: lvc=(unit (set ship))      :: leakset
446      pov=path                    :: provenance
447      [vis=view bem=beam]         :: perspective
448      ==                          ::
449      %- unit                     :: ~      unknown
450      %- unit                     :: ~ ~    invalid
451      (pair mark vase)            :: envased result
452 +$ view $@ (term [way=term car=term]) :: perspective
453
454
```

At this point, an excursus into the Arvo core is warranted. The Arvo core is the Nock-based kernel of Urbit, and it is responsible for managing the state of the system and providing a consistent interface for agents to interact with that state. The Arvo core is built on top of the Nock interpreter, and it uses the `++mink` function to interpret Nock formulas. Each vane in the Arvo core receives a `+$roof` in its function call, which it can call when it needs to scry with opcode 12. This gate provides access to the vane's state, and it is used to retrieve values from the vane's state in a way that is consistent with the Nock semantics.

`++look` is a gate that is used to convert the `++mink`-compatible version of the scry handler into a version that can be used by the Arvo core. This gate acts as a bridge between the `.^` dotket and the scry handler (`++look`) with access to the roof that came from the vane. This pattern imposes constraints on interpreters, but provides a way to mediate access to the state

of the system in a way that is consistent with Nock semantics. Thus `.^` dotket gives mediated access into the sky without exposing the state of the entire system in the subject (untyped permissions-free access).

3.16 Et cetera

`++mink` also features some machinery to correctly implement the cell distribution rule.

```
[ ^ *]
=/ head $(formula -.formula)
?. ?=(%0 -.head) head
=/ tail $(formula +.formula)
?. ?=(%0 -.tail) tail
[%0 product.head product.tail]
```

There is little to be said of this machinery, but particularly notice that the error is propagated from the member of the cell which crashes.

`++mink` is naturally jetted for efficiency, ultimately by `u3m_soft_run()` in Urbit’s Vere (C) and by `interpret()` in NockApp’s NockVM (née) Sword (Rust), both the virtualization context for their respective runtimes. Some internal affordances for the runtime are in place to keep execution consistent; e.g., scry gates are stored on a stack to prevent re-entry if it scries.

Nock is a crash-only language, which means that when something goes wrong, the exception is an exception for the runtime or virtual machine to handle as it will. Typically, for a runtime this means an injection of the error trace back into the Nock kernel (cf. section “`u3n: nock execution`” in `u3.md` in the Urbit documentation⁶). (For `++mink`, errors are attempted to be front-run by the interpreter and thus bad Nock should not be executed as such.)

Harkening back to **Amin2017** (**Amin2017**)’s collapsed towers of interpreters, the affordances of `++mock` as a metacircular interpreter permit us to build a tower of interpreters on

⁶This section devoted a fair amount of attention to opcode 11 now 12 (it being written for Nock 5K) and hint evaluation.

top of it without incurring excessive performance overhead. As the author of `u3.md` concluded, “we [even] simply treat Nock proper as a special case of `mock`.”

What would such an interpreter look like, and how would it relate to actual operation on a runtime in practice? Let us extend `++mink` to implement a logical loobean operator. (Call the modified gate `++pink`.) This operation will be invoked with a new fake opcode 13. Under the hood, opcode 13 will simply supply logically negate a value (via an opcode 6), much as how opcode 12 supplies a consistent Nock formula for the `scry` operation. (Similar to `[0 0]`, we will be able to use `[13 1 2]` to initiate a crash if desired.)

```

520
521 :: *[a 13 b]          NOT *[a b]
522   [%13 subject=*]
523 =/  argument $(formula argument.formula)
524 ?= (%0 -. argument) argument
525 ?^  product.argument [%2 trace]
526 ?= (%1 product.argument) [%0 %.y]
527 ?= (%0 product.argument) [%0 %.n]
528 [%2 trace]
529

```

This is invoked as:

```

531
532 > (pink [0 13 1 0] *$-(^ (unit (unit))))
533 [%0 product=0]
534 > (pink [0 13 1 1] *$-(^ (unit (unit))))
535 [%0 product=1]
536 > (pink [0 13 1 2] *$-(^ (unit (unit))))
537 [%2 trace=~]
538

```

Finally, it is instructive to consider what the minimal required amount of code machinery would be for one to successfully implement a Nock metacircular interpreter in Nock itself. An examination of `++mink` shows modest reliance on Hoon’s higher-level features; tree math, for instance, in `++cap` for distinguishing head vs. tail and `++mas` for finding the relative address of an axis. (Neither of these are particularly complicated, but they entail some basic arithmetic operations as well.) By omitting `++scry` and opcode 12, as well as reducing to a structural match with a `+$tone` instead of an explicit Hoon type

549 match, one could produce a reasonably sized Hoon program
 550 compiling to something close to the minimal viable Nock vir-
 551 tualized interpreter. No doubt intrepid code golfers will find it
 552 very much an upper bound rather than a lower.

553 This Nock formula produces a working virtualized Nock
 554 interpreter derived from Hoon code with some hand-tuned
 555 design. The input function signature is [subject=*
 556 formula=*] with output of the form [⌵ product=*]; if the
 557 head is %0, then the product is a valid Nock noun, while a head
 558 of %bail indicates a crash in the evaluation. Opcode 9 calls
 559 may be addressed into the core as subject, and the formula it-
 560 self is complete with an empty subject. (Note as well that the
 561 %bail return may itself be the natural result of a valid Nock
 562 computation, which would be avoidable if we switched to a
 563 unit return instead.)

```

564 [[1
565   [:: +mul multiplication of two atoms
566     8 [1 1 1]
567       [1 8 [1 0]
568         8 [1 6 [5 [1 0] 0 60]
569           [0 6]
570             9 2 10 [60 8 [9 686 0 31]
571               9 2 10 [6 0 124] 0 2]
572                 10 [6 8 [9 20 0 31]
573                   9 2 10 [6 [0 125] 0 14]
574                     0 2] 0 1]
575             9 2 0 1] 0 1]
576   [:: +add addition of two atoms
577     8 [1 0 0]
578       [1 6 [5 [1 0] 0 12]
579         [0 13]
580           9 2 10
581             [6 [8 [9 686 0 7] 9 2 10 [6 0 28] 0 2]
582               4 0 13]
583                 0 1] 0 1]
584   [[8
585     :: +list type definition
586       [8 [1 0]
587         [1 8 [0 6] 8 [5 [0 14] 0 2] 0 6]
588           0 1]
589     ]

```

```

590     [1 8 [1 0]
591       [1 8
592         [6 [3 0 6]
593           [[6 [5 [1 0] 0 12] [1 0] 0 0]
594             8 [0 30] 9 2 10 [6 0 29] 0 2]
595           6 [5 [1 0] 0 6] [1 0] 0 0]
596         8 [5 [0 14] 0 2] 0 6]
597       0 1] 0 1]
598 [:: +dvr division with remainder
599 8 [1 1 1]
600   [1 6 [5 [1 0] 0 13]
601     [0 0]
602     8 [1 0]
603     8 [1 6 [8 [9 687 0 31]
604       9 2 10 [6 [0 124] 0 125] 0 2]
605       [[0 6] 0 60]
606       9 2 10
607       [60 8 [9 23 0 31]
608         9 2 10 [6 [0 124] 0 125] 0 2]
609       10 [6 4 0 6] 0 1]
610     9 2 0 1] 0 1]
611 [:: +cap:
612 8 [1 0]
613   [1 6 [5 [1 2] 0 6]
614     [1 2]
615     6 [5 [1 3] 0 6]
616     [1 3]
617     6 [6 [5 [1 1] 0 6] [1 0] 5 [1 0] 0 6]
618     [0 0]
619     9 2 10
620     [6 7 [8 [9 170 0 7]
621       9 2 10
622       [6 [0 14] 7 [0 3] 1 2]
623     0 2] 0 2]
624   0 1] 0 1]
625 [:: +dec decrement an atom
626 8 [1 0]
627   [1 6 [5 [1 0] 0 6]
628     [0 0]
629     8 [1 0]
630     8 [1 6 [5 [0 30] 4 0 6]
631       [0 6]

```

```

632             9 2 10 [6 4 0 6] 0 1]
633         9 2 0 1] 0 1]
634     :: +lth less-than test
635     8 [1 0 0]
636         [1 6 [6 [5 [0 12] 0 13] [1 1] 1 0]
637             [6 [8
638                 [1 6 [5 [1 0] 0 28]
639                     [1 0]
640                         6 [6 [6 [5 [1 0] 0 29]
641                             [1 1]
642                                 1 0]
643                                     [6 [9 2 10
644                                         [14 [8 [9 686 0 15]
645                                             9 2 10 [6 0 60]
646                                                 0 2]
647                                                     8 [9 686 0 15]
648                                                         9 2 10 [6 0 61] 0 2]
649                                                             0 1]
650                                                                 [1 0]
651                                                                     1 1]
652                                                                         1 1]
653                                                         [1 0] 1 1] 9 2 0 1]
654         [1 0] 1 1] 1 1] 0 1]
655     [:: +edit edit an atom at a given axis
656     8 [1 0 0 0]
657         [1 6 [5 [1 1] 0 12]
658             [[1 0] 0 27]
659                 6 [6 [3 0 26] [1 1] 1 0]
660                     [1 0]
661                         8 [9 2 10 [26
662                             8 [8 [9 342 0 63]
663                                 9 2 10 [6 0 28] 0 2]
664                                     6 [5 [1 2] 0 2]
665                                         [0 116]
666                                             6 [5 [1 3] 0 2]
667                                                 [0 117]
668                                                     1 1.818.845.538 0]
669                                                         10 [12 8 [9 87 0 63]
670                                                             9 2 10 [6 0 28] 0 2]
671                         0 1]
672                 6 [5 [1 0] 0 2]
673                     [1 0]

```

```
674             8 [8 [9 342 0 15]
675               9 2 10 [6 0 60] 0 2]
676             6 [5 [1 2] 0 2]
677               [7 [0 3] [1 0] [0 5] 0 117]
678             6 [5 [1 3] 0 2]
679               [7 [0 3] [1 0] [0 116] 0 5]
680             1 1.818.845.538 0]
681 0 1]
682 [:: +]ink, raw virtual Nock
683 8 [1 0 0]
684   [1 8
685     6 :: Cell distribution
686     [6 [3 0 13] [3 0 26] 1 1]
687     [8 [9 2 10 [13 0 26] 0 1]
688       6 [5 [1 0] 0 4]
689       [8 [9 2 10 [13 0 59] 0 3]
690       6 [5 [1 0] 0 4]
691       [[1 0] [0 13] 0 5]
692     0 2]
693   0 2]
694 6 :: Opcode 0
695 [6 [3 0 13]
696   [6 [5 [1 0] 0 26]
697     [6 [3 0 27] [1 1] 1 0]
698     1 1]
699   1 1]
700 [8 [8 [9 22 0 7]
701   9 2 10 [6 [0 59] 0 28] 0 2]
702   6 [5 [1 0] 0 2]
703   [1 1.818.845.538 0]
704   [1 0]
705   0 5]
706 6 :: Opcode 1
707 [6 [3 0 13]
708   [5 [1 1] 0 26] 1 1]
709 [[1 0] 0 27]
710 6 :: Opcode 2
711 [6 [3 0 13]
712   [6 [5 [1 2] 0 26] [3 0 27] 1 1] 1 1]
713 [8 [9 2 10 [13 0 54] 0 1]
714   6 [5 [1 0] 0 4]
715   [8 [9 2 10 [13 0 119] 0 3]
```

```

716             6 [5 [1 0] 0 4]
717             [9 2 10 [6 [0 13] 0 5] 0 7]
718             0 2]
719         0 2]
720     6 :: Opcode 3
721     [6 [3 0 13]
722       [5 [1 3] 0 26] 1 1]
723     [8 [9 2 10 [13 0 27] 0 1]
724       6 [5 [1 0] 0 4]
725       [[1 0] 3 0 5]
726       0 2]
727     6 :: Opcode 4
728     [6 [3 0 13]
729       [5 [1 4] 0 26]
730       1 1]
731     [8 [9 2 10 [13 0 27] 0 1]
732       6 [5 [1 0] 0 4]
733       [6 [6 [3 0 5] [1 1] 1 0]
734         [[1 0] 4 0 5]
735         1 1.818.845.538 0]
736       0 2]
737     6 :: Opcode 5
738     [6 [3 0 13]
739       [6 [5 [1 5] 0 26] [3 0 27] 1 1] 1 1]
740     [8 [9 11 10 [29 0 118] 0 1]
741       6 [5 [1 0] 0 4]
742       [8 [9 2 10 [13 0 119] 0 3]
743       6 [5 [1 0] 0 4]
744       [[1 0] 5 [0 13] 0 5]
745       0 2]
746     0 2]
747     6 :: Opcode 6
748     [6 [3 0 13]
749       [6 [5 [1 6] 0 26]
750         [6 [3 0 27] [3 0 55] 1 1] 1 1] 1 1]
751     [8 [9 11 10 [29 0 118] 0 1]
752       6 [5 [1 0] 0 4]
753       [6 [5 [1 0] 0 5]
754         [9 2 10 [13 0 238] 0 3]
755       6 [5 [1 1] 0 5]
756       [9 2 10 [13 0 239] 0 3]
757       1 1.818.845.538 0]

```

```

758         0 2]
759 6 :: Opcode 7
760 [6 [3 0 13]
761     [6 [5 [1 7] 0 26] [3 0 27] 1 1] 1 1]
762 [8 [9 2 10 [13 0 54] 0 1]
763     6 [5 [1 0] 0 4]
764         [9 2 10 [6 [0 5] 0 119] 0 3]
765         0 2]
766 6 :: Opcode 8
767 [6 [3 0 13]
768     [6 [5 [1 8] 0 26] [3 0 27] 1 1] 1 1]
769 [8 [9 2 10 [13 0 54] 0 1]
770     6 [5 [1 0] 0 4]
771         [9 2 10 [6 [[0 5] 0 28] 0 119] 0 3]
772         0 2]
773 6 :: Opcode 9
774 [6 [3 0 13]
775     [6 [5 [1 9] 0 26]
776         [6 [3 0 27]
777             [6 [3 0 54] [1 1] 1 0]
778                 1 1]
779                 1 1]
780     1 1]
781 [8 [9 2 10 [13 0 55] 0 1]
782     6 [5 [1 0] 0 4]
783         [8 [8 [9 22 0 15]
784             9 2 10 [6 [0 246] 0 13] 0 2]
785         6 [5 [1 0] 0 2]
786             [1 1.818.845.538 0]
787         9 2 10 [6 [0 13] 0 5]
788         0 7]
789     0 2]
790 6 :: Opcode 10
791 [6 [3 0 13]
792     [6 [5 [1 10] 0 26]
793         [6 [3 0 27]
794             [6 [3 0 54]
795                 [6 [3 0 108] [1 1] 1 0]
796                     1 1]
797                     1 1]
798                 1 1]
799                 1 1]

```

```

800      [6 [5 [1 0] 0 108]
801        [1 1.818.845.538 0]
802        8 [9 2 10 [13 0 55] 0 1]
803        6 [5 [1 0] 0 4]
804        [8 [9 2 10 [13 0 237] 0 3]
805        6 [5 [1 0] 0 4]
806        [8 [8 [9 86 0 31]
807          9 2 10
808          [6 [0 1.004] [0 29] 0 13]
809          0 2]
810          6 [5 [1 0] 0 2]
811          [1 1.818.845.538 0]
812          [1 0]
813          0 5]
814          0 2]
815          0 2]
816      6 :: Opcode 11
817      [6 [3 0 13]
818        [6 [5 [1 11] 0 26]
819        [6 [3 0 27]
820        [6 [3 0 54]
821        [1 1]
822        1 0]
823        1 1]
824        1 1]
825        1 1]
826      [8 [9 2 10 [13 0 55] 0 1]
827      6 [5 [1 0] 0 4]
828      [[1 0]
829      2 [0 28] [1 11] [0 118] [1 1] 0 5]
830      0 2]
831      1 1.818.845.538 0]
832      0 1]
833      :: +mas
834      8 [1 0]
835      [1 6 [6 [5 [1 3] 0 6] [1 0] 5 [1 2] 0 6]
836      [1 1]
837      6 [6 [5 [1 1] 0 6] [1 0] 5 [1 0] 0 6]
838      [0 0]
839      8 [9 20 0 7]
840      9 2 10
841      [6 [7 [0 3] 7 [8 [9 170 0 7]

```



```

842                                     9 2 10
843                                     [6 [0 14] 7 [0 3] 1 2]
844                                     0 2] 0 3]
845             7 [0 3] 8 [9 4 0 7]
846             9 2 10 [6
847             [7 [0 3]
848             9 2 10
849             [6 7 [8 [9 170 0 7]
850             9 2 10 [6 [0 14] 7 [0 3] 1 2]
851             0 2] 0 2]
852             0 1]
853             7 [0 3] 1 2]
854             0 2]
855             0 2]
856             0 1]
857 [:: +frag fragmentary subtree of noun
858 8 [1 0 0]
859 [1 6 [5 [1 0] 0 12]
860 [1 0]
861 8 [1 6 [5 [1 1] 0 28]
862 [[1 0] 0 29]
863 6 [6 [3 0 29] [1 1] 1 0]
864 [1 0]
865 9 2 10
866 [14 [8 [9 175 0 15]
867 9 2 10 [6 0 60] 0 2]
868 8 [8 [9 342 0 15]
869 9 2 10 [6 0 60] 0 2]
870 6 [5 [1 2] 0 2]
871 [0 122]
872 6 [5 [1 3] 0 2]
873 [0 123]
874 1 1.818.845.538 0]
875 0 1]
876 9 2 0 1]
877 0 1]
878 [:: +unit, type definition of unit/Maybe
879 8
880 [8 [1 0] [1 8 [0 6] 8 [5 [0 14] 0 2] 0 6] 0 1]
881 [1 8 [1 0]
882 [1 8 [6 [3 0 6]
883 [[8 [0 30] 9 2 10 [6 0 28] 0 2]

```

```

884           8 [7 [0 7] 8 [9 46 0 7]
885             9 2 10 [6 0 14] 0 2]
886           9 2 10 [6 0 29] 0 2]
887         6 [5 [1 0] 0 6]
888         [1 0]
889         0 0]
890       8 [5 [0 14] 0 2]
891     0 6]
892   0 1]
893   0 1]
894   :: +sub subtraction of two atoms
895   8 [1 0 0]
896     [1 6 [5 [1 0] 0 13]
897       [0 12]
898       9 2 10
899       [6 [8 [9 686 0 7] 9 2 10 [6 0 28] 0 2]
900         8 [9 686 0 7] 9 2 10 [6 0 29] 0 2]
901     0 1]
902   0 1]
903   0 1]
904

```

4 Conclusion

Metacircular virtualization is a powerful tool for building interpreters, and Nock is not only no exception, it exemplifies the practice. The design of `++mock` and its underlying `++mink` function provides a practical way to interpret Nock in a way that is both efficient and extensible. One could imagine, for instance, a Nock interpreter cousin to `++mock` that is instrumented for partial evaluation of formulas (as in a debugger or syntax highlighter robust to syntax errors). Nock supports a rich capacity to build metacircular interpreters as a first-class feature, a potential which remains as yet relatively unexplored in the Nock ecosystem.

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