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Program Structures and Algorithms
Spring 2021(SEC 05)

Task:

- For weighted quick union, store the depth rather than the size;
- For weighted quick union with path compression, do two loops, so that all intermediate nodes point to the root, not just the alternates.

For both of these, code the alternative and benchmark it against the implementation in the repository.

Conclusion:

If we have a tree of height 1, it must have at least 2 nodes i.e., 2^h nodes, where h = height of the tree.

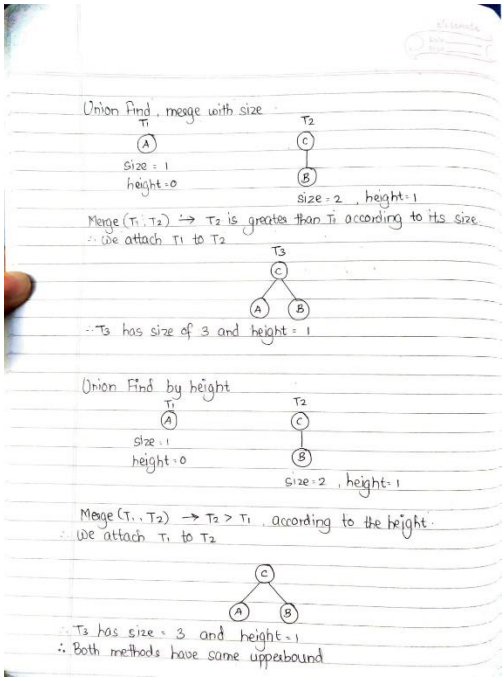
Union Find by size

So, for height = 0, minimum number of nodes will be 1 ∴ size of tree is 1
Similarly, for height = 1, minimum number of nodes will be 2 ∴ size of tree is 2
If we increase the height of the tree, it will always increment by 1
So, we get h = log(n)

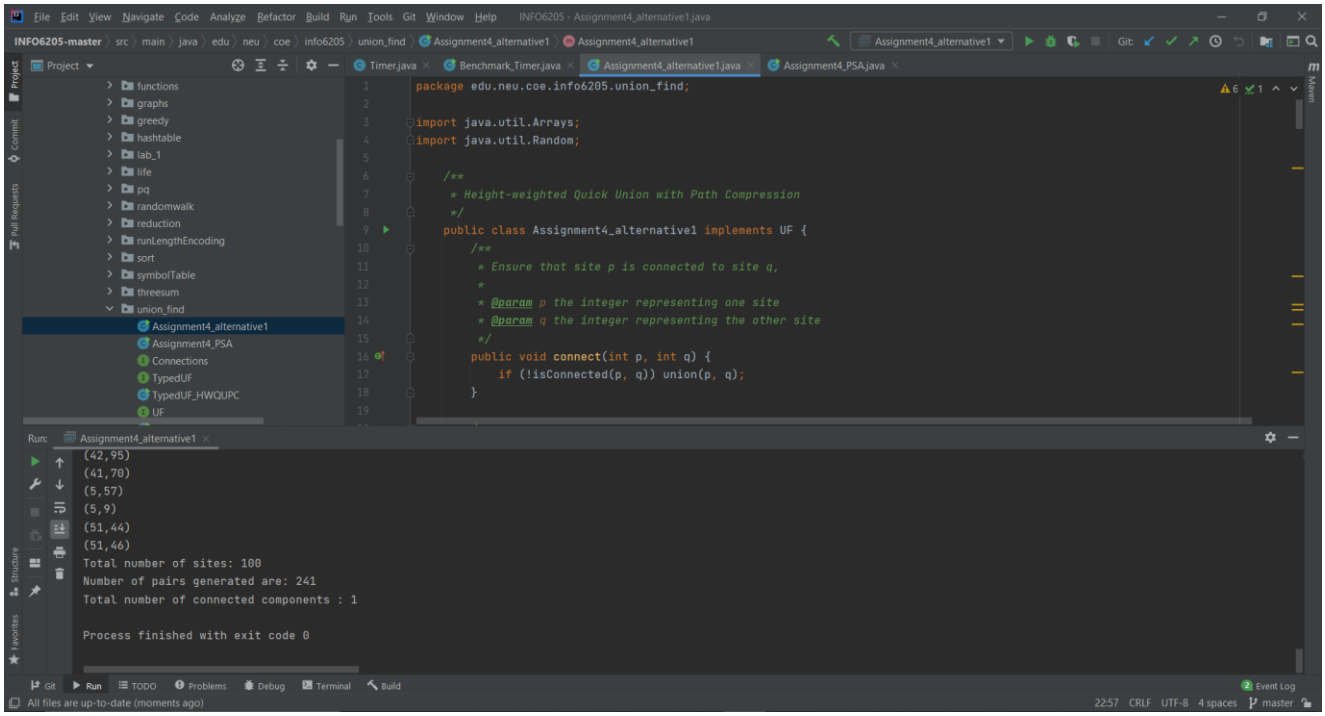
Union Find by height

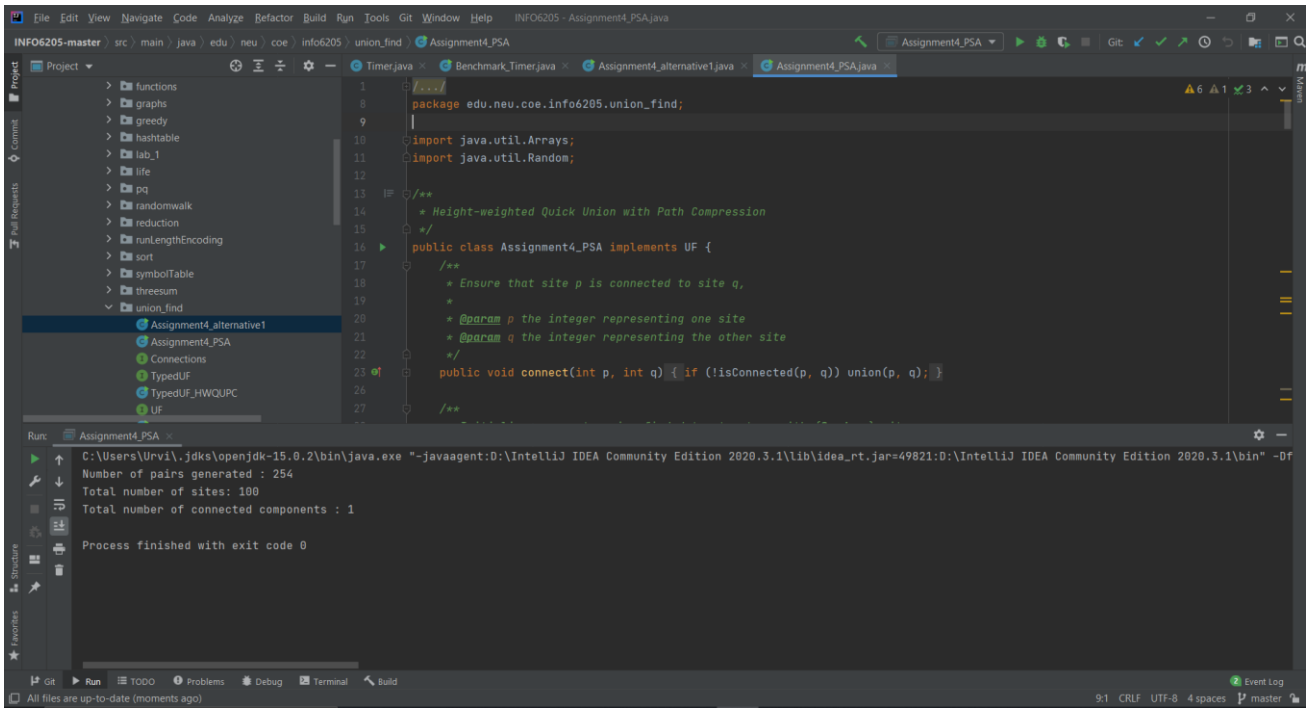
For height = 0, minimum number of nodes will be 1
Similarly, for height = 1, minimum number of nodes will be 2
Here, we get h = log(n)

Evidence to support that conclusion:



Output Screenshots





Benchmarking implementation