

Lecture 2

Process Concepts

Operating Systems

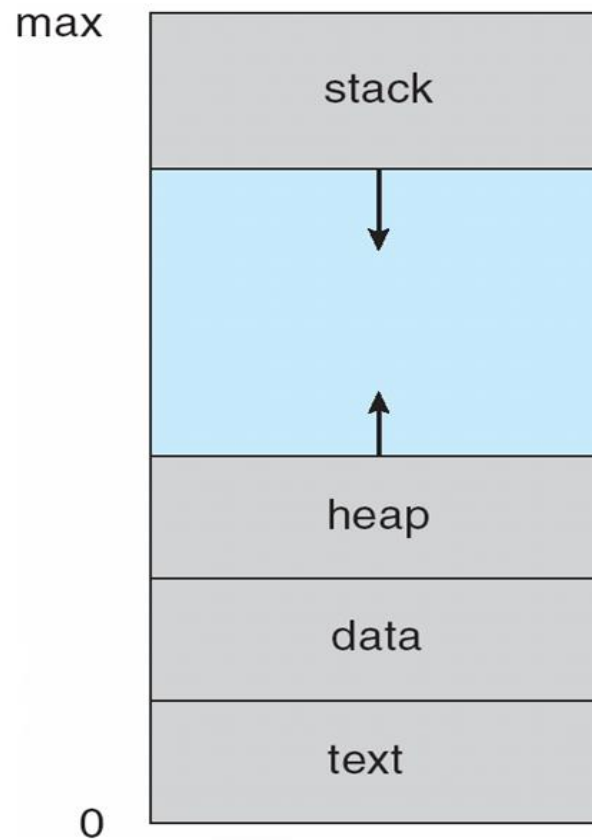
Process Concept

- An operating system executes a variety of programs:
 - Batch system – jobs
 - Time-shared systems – user programs or tasks
- The terms *job* and *process* used almost interchangeably
- Process – a program in execution; process execution must progress in sequential fashion
- A process includes:
 - program counter
 - stack
 - data section

The Process

- Multiple parts
 - The program code, also called **text section**
 - Current activity including **program counter**, processor registers
 - **Stack** containing temporary data
 - Function parameters, return addresses, local variables
 - **Data section** containing global variables
 - **Heap** containing memory dynamically allocated during run time
- Program is passive entity, process is active
 - Program becomes process when executable file loaded into memory
- Execution of program started via GUI mouse clicks, command line entry of its name, etc

Process in Memory



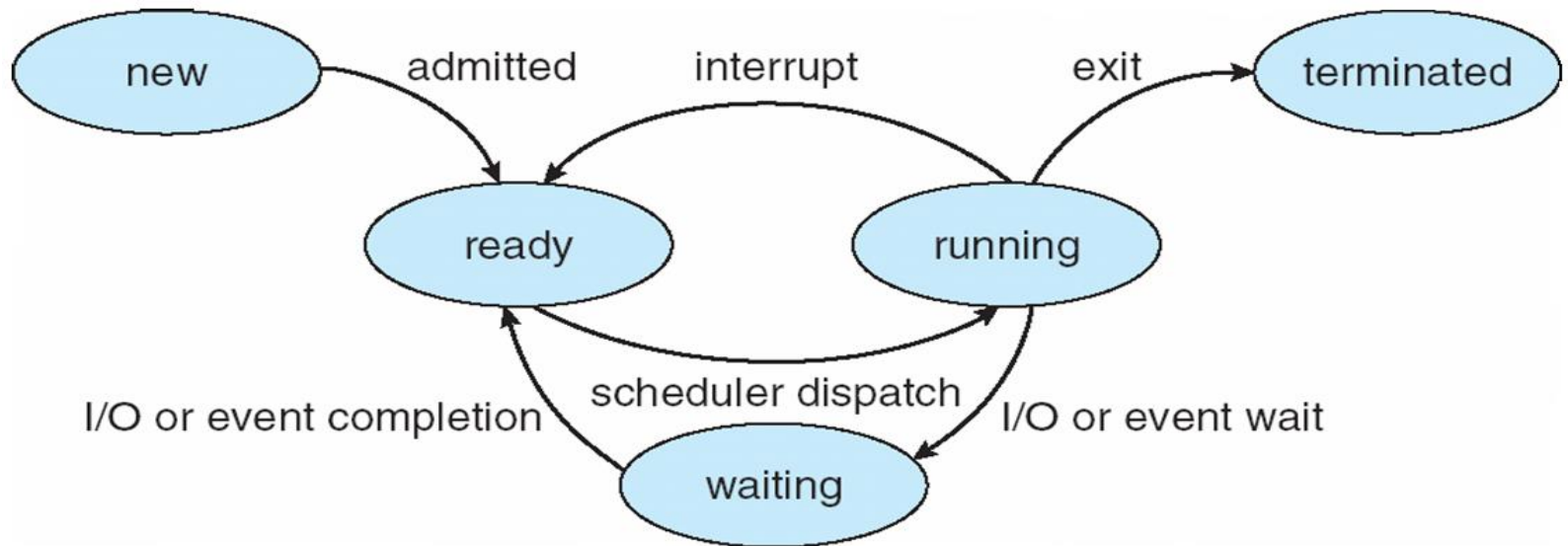
Process Address Space

- A list of memory locations from some min (usually 0) to some max that a process can read and write.
- Contains
 - the executable program
 - program's data
 - Stack?
 - Associated with a process is a set of registers e.g. PC, SP and other information to run the program.

Process State

- As a process executes, it changes *state*
 - **new**: The process is being created
 - **running**: Instructions are being executed
 - **waiting**: The process is waiting for some event to occur
 - **ready**: The process is waiting to be assigned to a processor
 - **terminated**: The process has finished execution

Diagram of Process State



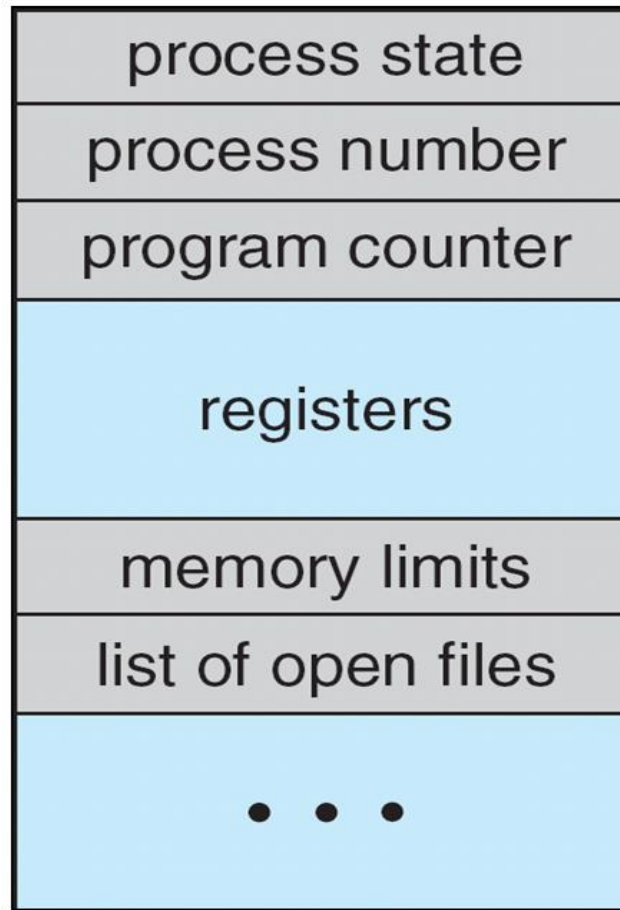


Process Control Block (PCB)

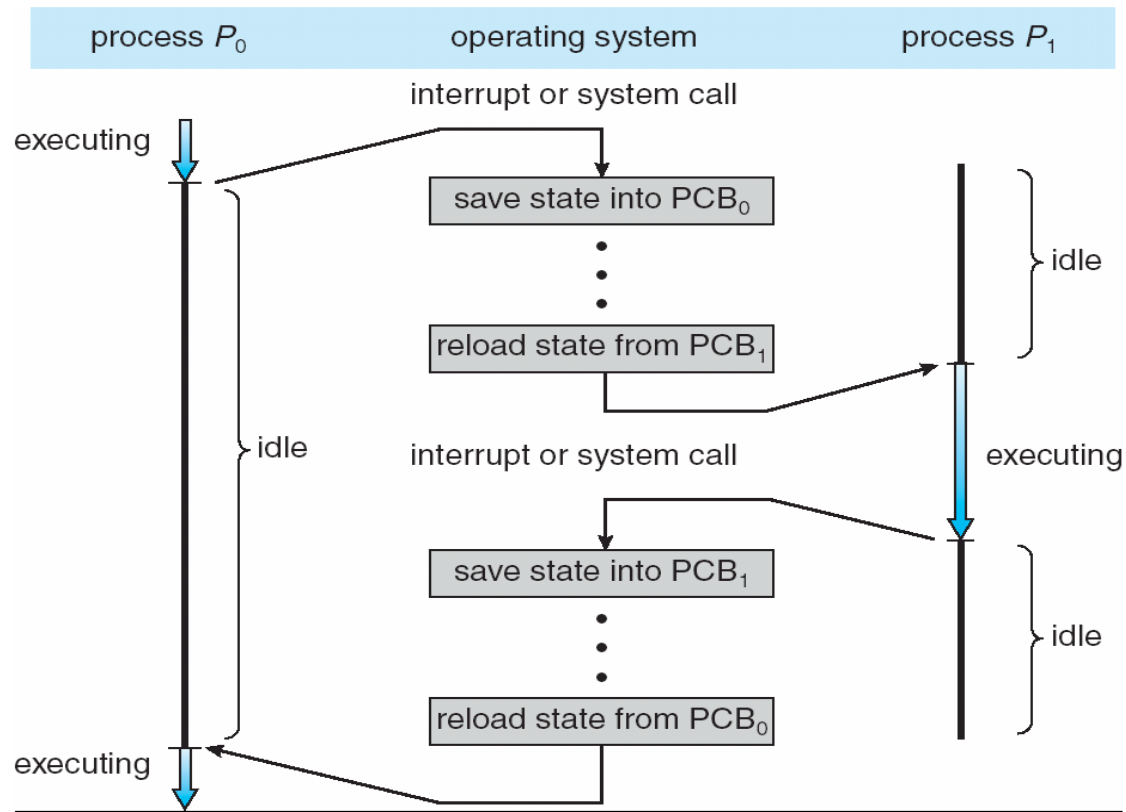
Information associated with each process

- Process state
- Program counter
- CPU registers
- CPU scheduling information
- Memory-management information
- Accounting information
- I/O status information

Process Control Block (PCB)



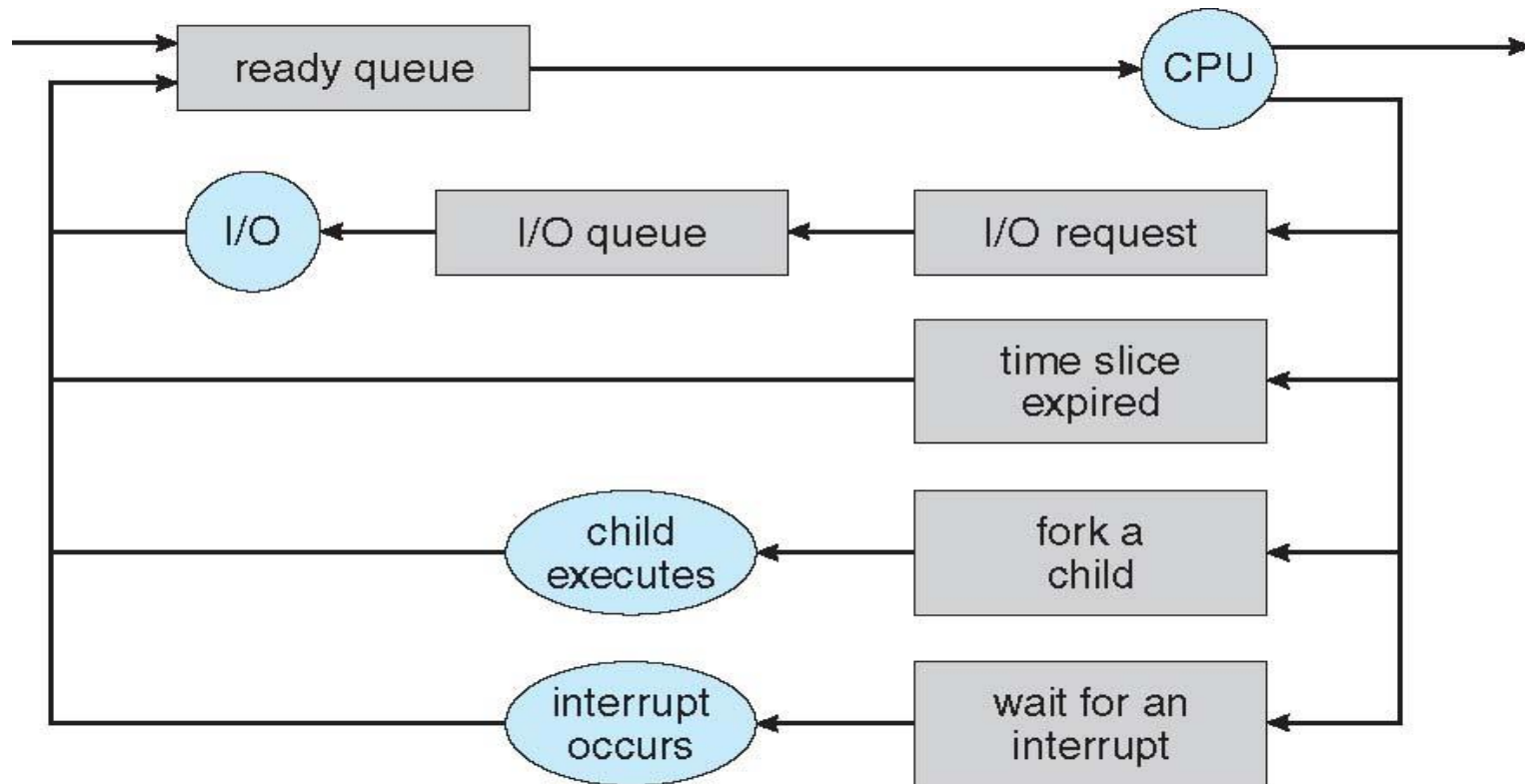
CPU Switch From Process to Process



Process Scheduling

- Maximize CPU use, quickly switch processes onto CPU for time sharing
- **Process scheduler** selects among available processes for next execution on CPU
- Maintains **scheduling queues** of processes
 - **Job queue** – set of all processes in the system
 - **Ready queue** – set of all processes residing in main memory, ready and waiting to execute
 - **Device queues** – set of processes waiting for an I/O device
 - Processes migrate among the various queues

Representation of Process Scheduling



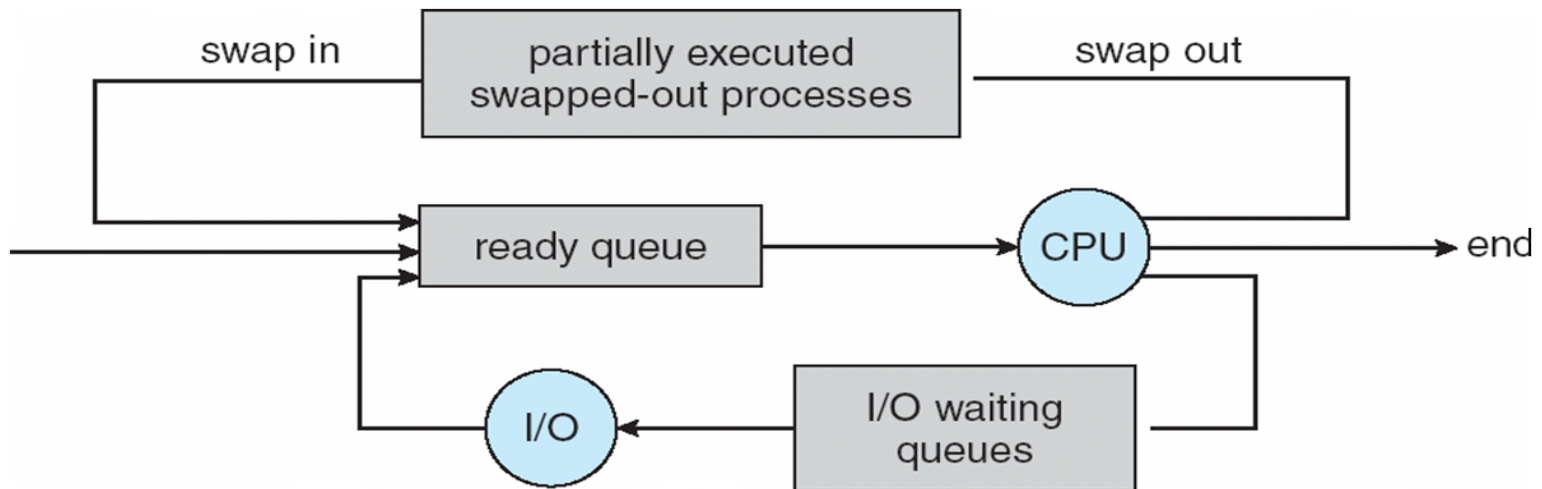
Schedulers

- **Long-term scheduler** (or job scheduler) – selects which processes should be brought into the ready queue
- **Short-term scheduler** (or CPU scheduler) – selects which process should be executed next and allocates CPU
 - Sometimes the only scheduler in a system

Schedulers (Cont.)

- Short-term scheduler is invoked very frequently (milliseconds) \Rightarrow (must be fast)
- Long-term scheduler is invoked very infrequently (seconds, minutes) \Rightarrow (may be slow)
- The long-term scheduler controls the *degree of multiprogramming*
- Processes can be described as either:
 - **I/O-bound process** – spends more time doing I/O than computations, many short CPU bursts
 - **CPU-bound process** – spends more time doing computations; few very long CPU bursts

Addition of Medium Term Scheduling



Context Switch

- When CPU switches to another process, the system must save the state of the old process and load the saved state for the new process via a **context switch**.
- **Context** of a process represented in the PCB
- Context-switch time is overhead; the system does no useful work while switching
 - The more complex the OS and the PCB -> longer the context switch
- Time dependent on hardware support
 - Some hardware provides multiple sets of registers per CPU -> multiple contexts loaded at once



Process Creation

- **Parent** process create **children** processes, which, in turn create other processes, forming a tree of processes
- Generally, process identified and managed via a **process identifier (pid)**
- Resource sharing
 - Parent and children share all resources
 - Children share subset of parent's resources
 - Parent and child share no resources
- Execution
 - Parent and children execute concurrently
 - Parent waits until children terminate

Process Creation (Cont.)

- Address space

- ☐ Child duplicate of parent
- ☐ Child has a program loaded into it

- UNIX examples

- ☐ **fork** system call creates new process
- ☐ **exec** system call used after a **fork** to replace the process' memory space with a new program

Process Creation

