

$M = 10011010$

$K_1 = \{11101001\}$

$K_2 = \{10100111\}$

Apply initial permutation on Plaintext (M)

Plaintext = 10011010

| | | | | | | | | |
|----|---|---|---|---|---|---|---|---|
| IP | 2 | 6 | 3 | 1 | 4 | 8 | 5 | 7 |
|----|---|---|---|---|---|---|---|---|

| Bit # | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
|-------|---|---|---|---|---|---|---|---|
| P | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 |
| IP(P) | 0 | 0 | 0 | 1 | 1 | 0 | 1 | 1 |

Now divide it two half

Left half = 0001

Right half = 1011

| |
|-----------------|
| EIP |
| 4 1 2 3 2 3 4 1 |

Steps

| Bits | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
|---------------------|---|---|---|---|---|---|---|---|
| R | 1 | 0 | 1 | 1 | | | | |
| EIP(R) | 1 | 1 | 0 | 1 | 0 | 1 | 1 | 1 |
| K_1 | 1 | 1 | 1 | 0 | 1 | 0 | 0 | 1 |
| $EIP(R) \oplus K_1$ | 0 | 0 | 1 | 1 | 1 | 1 | 1 | 0 |

S Boxes ($EIP(R) \oplus k_1$)

| | | | | | | |
|---|---|---|---|---|---|---|
| 1 | 0 | 0 | 0 | 1 | 1 | 1 |
|---|---|---|---|---|---|---|

$P_4(S \text{ Boxes } (EIP(R) \oplus k_1))$

| | | | |
|---|---|---|---|
| 0 | 0 | 0 | 1 |
|---|---|---|---|

$$S_0 = \begin{matrix} & \begin{matrix} 0 & 1 & 2 & 3 \end{matrix} \\ \begin{matrix} 0 \\ 1 \\ 2 \\ 3 \end{matrix} & \begin{bmatrix} 1 & 0 & 3 & 2 \\ 3 & 2 & 1 & 0 \\ 0 & 2 & 1 & 3 \\ 3 & 1 & 3 & 2 \end{bmatrix} \end{matrix}$$

| | | | |
|---|---|---|---|
| 0 | 0 | 1 | 1 |
|---|---|---|---|

Rearrange P_4

| | | | | |
|-------|---|---|---|---|
| P_4 | 2 | 4 | 3 | 1 |
|-------|---|---|---|---|

Now

Calculate $f(k_1)(L, R)$
 $= (0001 \oplus 0001, 1011)$
 $= (0000, 1011)$

Now we apply swap function

$L = 0000$

$R = 0000$

After swapping value

$L = 1011$

$R = 0000$

Row =

| | |
|---|---|
| 0 | 1 |
|---|---|

 Col =

| | |
|---|---|
| 0 | 1 |
|---|---|

 $\rightarrow 2$

$$S_1 = \begin{matrix} & \begin{matrix} 0 & 1 & 2 & 3 \end{matrix} \\ \begin{matrix} 0 \\ 1 \\ 2 \\ 3 \end{matrix} & \begin{bmatrix} 0 & 1 & 2 & 3 \\ 2 & 0 & 1 & 3 \\ 3 & 0 & 1 & 0 \\ 2 & 1 & 0 & 3 \end{bmatrix} \end{matrix}$$

| | | | |
|---|---|---|---|
| 1 | 1 | 1 | 0 |
|---|---|---|---|

Row =

| | |
|---|---|
| 1 | 0 |
|---|---|

 $= 2$
 Col =

| | |
|---|---|
| 1 | 1 |
|---|---|

 $= 3$ $\rightarrow 0$

Use $K_2 = 10100111$

Again apply initial permutation on plaintext

| Bits | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
|-------|---|---|---|---|---|---|---|---|
| P | 1 | 0 | 0 | 1 | 1 | 0 | 1 | 0 |
| IP(P) | 0 | 0 | 0 | 1 | 1 | 0 | 1 | 1 |

After swapping left or right value

Right = 0000

Left = 1011

Steps

| | | | | | | | | | |
|---|---|---|---|---|---|---|---|---|---|
| E | P | 4 | 1 | 2 | 3 | 2 | 3 | 4 | 1 |
|---|---|---|---|---|---|---|---|---|---|

| Bits | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
|---------------------|---|---|---|---|---|---|---|---|
| R | 0 | 0 | 0 | | | | | |
| E/P(R) | 1 | 0 | 0 | 0 | 0 | 0 | 1 | 0 |
| K_2 | 1 | 0 | 1 | 0 | 0 | 1 | 1 | 1 |
| $E/P(R) \oplus K_2$ | 0 | 0 | 1 | 0 | 0 | 1 | 0 | 1 |

Left value in S Box.

001010

$$\left. \begin{array}{l} \text{Row} = 010 = 2 \\ \text{Col} = 01 = 1 \end{array} \right\} 2 \text{ (two bit value is 00)}$$

Right value in S Box 1

0101

$$\left. \begin{array}{l} \text{Row} = 01 = 1 \\ \text{Col} = 110 = 3 \end{array} \right\} = 3 \text{ (two bit value is 01)}$$

| | 1 | 2 | 3 | 4 |
|------------------------------|----|---|----|---|
| SBox ($E/P(R) \oplus k_2$) | 01 | 0 | 10 | 1 |

Apply P_4

P_4

| | | | |
|---|---|---|---|
| 2 | 4 | 3 | 1 |
|---|---|---|---|

| | | | | |
|--------------------------------------|---|---|----|----|
| P_4 (SBox ($E/P(R) \oplus k_2$)) | 0 | 1 | 10 | 01 |
|--------------------------------------|---|---|----|----|

→ Now calculate $f(k_2)$ (L, R)
 $(1011 \oplus 0111, 0000)$
 $1100, 0000$

→ Now we apply the IP^{-1}

IP^{-1}

| | | | | | | | |
|---|---|---|---|---|---|---|---|
| 4 | 1 | 3 | 5 | 7 | 2 | 8 | 6 |
|---|---|---|---|---|---|---|---|

| Bit | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 |
|----------------|---|---|---|---|---|---|---|---|
| R.L | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 |
| $IP^{-1}(R.L)$ | 0 | 1 | 0 | 0 | 0 | 1 | 0 | 0 |

→ The encryption text is

Plaintext = 10011010

$\epsilon(\text{text}) = 01000100$