

RECURSION.

Raj

c) Count Good no. : In digit string is good if the even index : Even no.
odd index : Odd no.
(zero may be at start)

Input: n = length of string
Output: good no. of length n ($0 \leq n \leq 10^{15}$)

* 8 eleven in modulo $10^9 + 7$

long, long helper (ll base, expo, mod)

$$12 \text{ cm} = 1$$

while ($\text{expo} > 0$)

{ if ($\text{expo}^{*102} == 0$) base = (base^2)% expo
 $\text{expo} = \text{expo}/2;$

else { ans = (ans * back) % mod;

$$r_{\text{exp}} = \text{exp}^{-1};$$

8

~~return ong~~

Int Count Good No. (long long n)

۱۷

$$\text{If even} = (n+1)2; \quad \text{odd} = n2$$

$$\text{mood} = \log + 7$$

8cten (inv) helper(5, even, mod)* helper(4, odd, mod);

7

Converted to int

$$= \frac{5^3 \times 4^2}{5 \times 4 \times 5}$$

$$\text{Pow}(5,3) * \text{Pow}(4,2)$$

Generalized

$$\text{Pow}(5, \frac{\text{no of even indices}}{4})^*$$

Q Pow funct (Fast calc.)

double Pow(double x, int n) {
if (n == 0) return 1

$$\left\{ x^y = (x^2)^{y/2} \right\}$$

long n = n;

if (N > 0) return Pow(x * x, N / 2);

else return x * Pow(x, N - 1);

↓
start 1 Base

divide in half part

Q Generate Parenthesis.

given n pairs of parenthesis, generate all well formed combination of parenthesis

eg: n = 3

output { "((()))", "(()())", "(())()", "()()()" }

* can never start with closing para. { close < open }

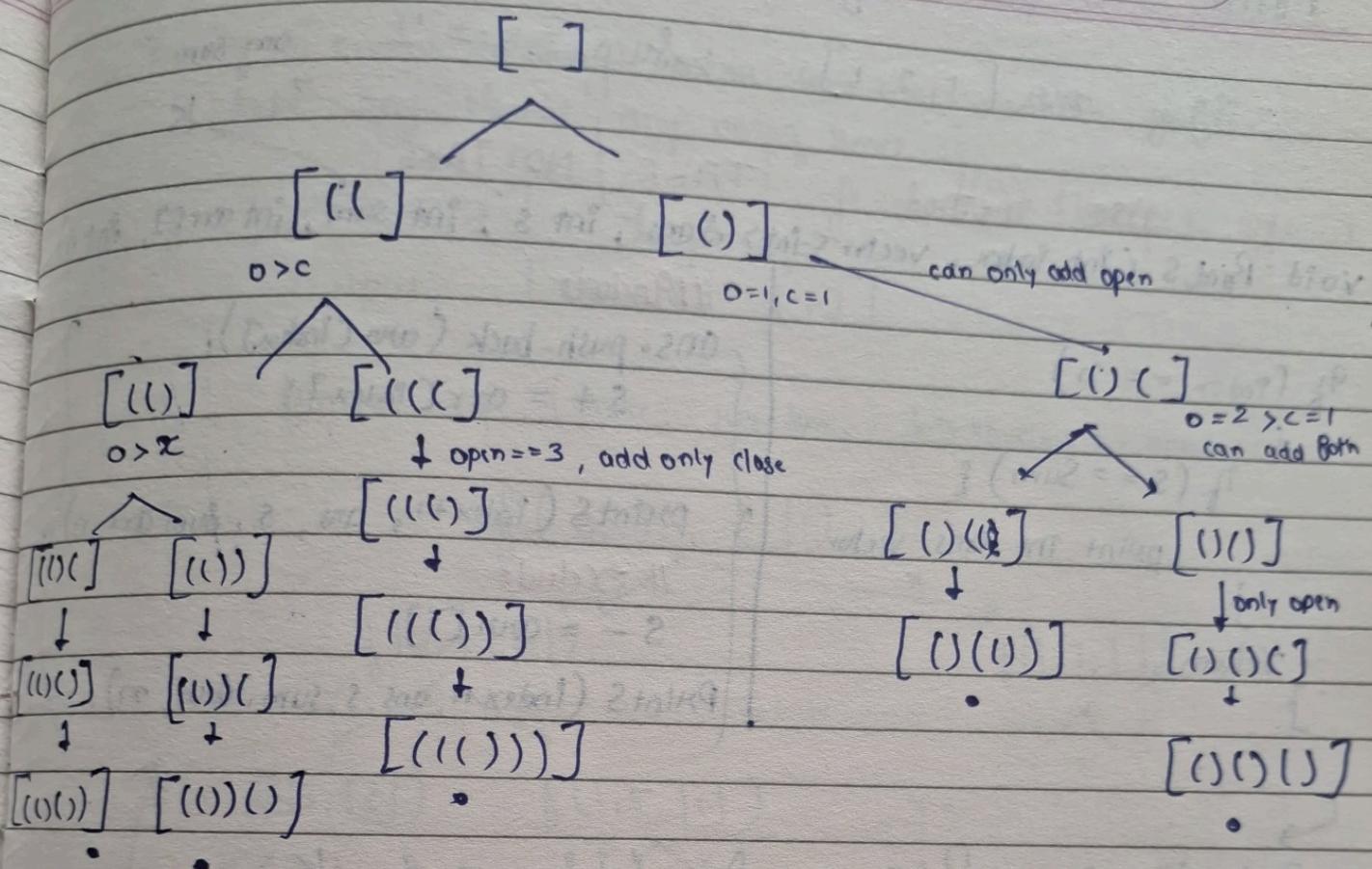
()

Open = 1, Close = 1

- * We have open limit of n, can't go more
- * open + close $\leq 2n$

for $n=3$

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void para (int n, close, open , vector<string>&v, strings)

{
 if (open < n) para (n, close, open + 1 , v, s + '(');

 if (close < open) para (n, close + 1, open , v, s + ')');

 if (open + close = 2 * n){

 v.push_back (s);

 return;

} got a completed answer

1) Put open

 ↓
(open < n)

2) Put Close

 ↓
(close < open)

vector<string> getparanthetis (int n)

{ Vector<string>ans = para (n, 0, 0, ans, "");

(unique elements)

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Print Subsequence of Sum 'K'

e.g.: arr [1, 2, 1]

sum = K = 2

arr sum

Use method of {TAKE | NOT TAKE}

void prints (int index, vector<int> &ans, int s, int sum, int arr[], int n)

if (index == n)

{

if (s == sum) {

print the ans vector

}

return;

}

// include

ans.push-back (arr [index]);

s+ = arr [index];

prints (index+1, ans, s, sum, arr, n);

// exclude.

s- = arr [i];

prints (index+1, ans, s, sum, arr, n);

Q Modify Q. : Print any One Subseq. of sum K.

• Make prints as a bool function

{ • At every recursive call, check if condition true.
Then return true }

if False: then only do recursive call.

ans.push-back (arr [index]); s+ = arr [index];

if (prints (index+1, ans, s, sum, arr, n) == true) return true;
s- = arr [i];

ans.pop-back();

if (prints (...)) = true) return true

return false

X
(Not

Q Combinational Sum

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Sum of all elements in array whose sum = k, where you can take any element any no. of times

eg: $[2, 3, 6, 7]$

target = 7

output: $[[2, 2, 3], [7]]$

Pick / Not Pick approach

function f (index, target sum, data structure for answer)

eg:

$f(0, 7, \boxed{1})$

$f(1, 7, \boxed{1})$

$f(0, 5, \boxed{1})$

$f(0, \boxed{3}, \boxed{2})$

$f(1, 5, \boxed{1})$

$f(0, 1, \boxed{2})$

$f(2, 2, \boxed{2})$

Conditions

element only added when $\text{arr}[i] < \text{target}$

easy code.

(Not possible)

$f(1, 1, \boxed{2})$

NULL

$f(2, 1, \boxed{2})$

NULL

$f(3, 1, \boxed{2})$

end

Comb. Sum II

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- There will be duplicate elements in array
- each element in given arr can be okay & used once
- final ans should not contain duplicates

• Sort given array

• Use a loop to skip duplicates

eg: $[10, 1, 2, 7, 6, 1, 5]$ target = 8

Output: $[[1, 1, 6], [1, 2, 5], [1, 7], [2, 6]]$

void comSum (int index, target, vector arr, vector<vector<int>> &ds)

```
{
    if (index == arr.size())
    {
        if (target == 0) ans.push_back(ds);
    }
}
```

II Skip duplicate, add, remove in some loop until
 $\star arr[i] > target$

```
for (int i = index; i < arr.size(); i++)
{
    if (i > index && arr[i] == arr[i-1]) continue;
    if (arr[i] > target) break;
}
```

```
ds.push_back(arr[i]);
comSum (index + 1, target - arr[i], arr, ans, ds);
```

Letter Combination

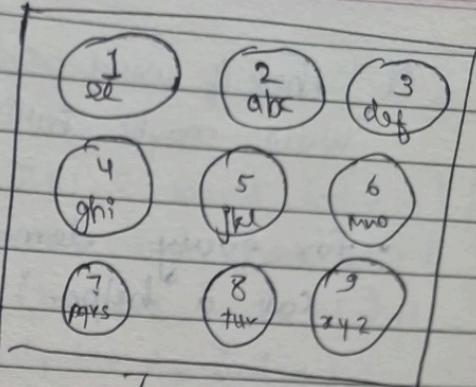
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String & digit (2-3) contain letters

return all possible letter combination
that the number could represent

eg: 2 3 = Digit (String)

[ad, ae, af, bd, be, bf, cd, ce, cf]



void letterCom (int index, String digits, String dial[], String curr, vector<String> ans)

{

if (index == digit.size())

{ ans.push_back (current);

return;

}

String letter = dial [digit [index] - '0'];

for (char c : letter)

{

letterCom (index + 1, digits, dial, curr + c, ans);

vector<String> letterCombination (String digits)

{

map string
to number

String dial [10] = {"", "", "abc", "def", "ghi", "jkl", "mno", "pqrs", "tuv", "wxyz"};

vector<String> ans

letterCom (0, digits, dial, "", ans);

} return ans

Word Search in matrix

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- A. Check if word exist in matrix or not eg word : SEE
Word maybe horizontal or vertical

- for every element in matrix call a "helper" function
- Recursive call for up | down | right | left

(check all boundaries at each call)

A	B	C	E
S	F	C	B
Z	Y	Z	S
A	D	E	E
H	E	L	L

bool dfs (vector<vector<^{char}string>& board, int r, c, index, string word)

```
{ if (r < 0 || c < 0 || r > board.size() || c > board[0].size() ||
      board[r][c] != word[index])
```

```
} return false; false → out of bounds →
      → char not matching
```

```
if (word.size() - 1 == index)
```

```
} return true;
```

```
// Mark current cell at '#' then go on
char temp = board[r][c]
board[r][c] = '#'
```

full
match

// Explore all four direction

```
bool iffound = dfs(board, r+1, c, word, index+1) ||
    dfs(board, r-1, c, word, index+1) ||
    dfs(board, r, c+1, word, index+1) ||
    dfs(board, r, c-1, word, index);
```

board[r][c] = temp; // back track and
change that char to original

return ifFound;

}

```
bool exist (vector<vector<char>> board, string word)
{
```

```
int rows = board.size();
int cols = board[0].size();
```

```
for (int i=0; i<rows; i++)
```

{

```
for (int j=0; j<cols; j++)
```

{

```
if (dfs(board, r, c, word, 0))
```

return true

}

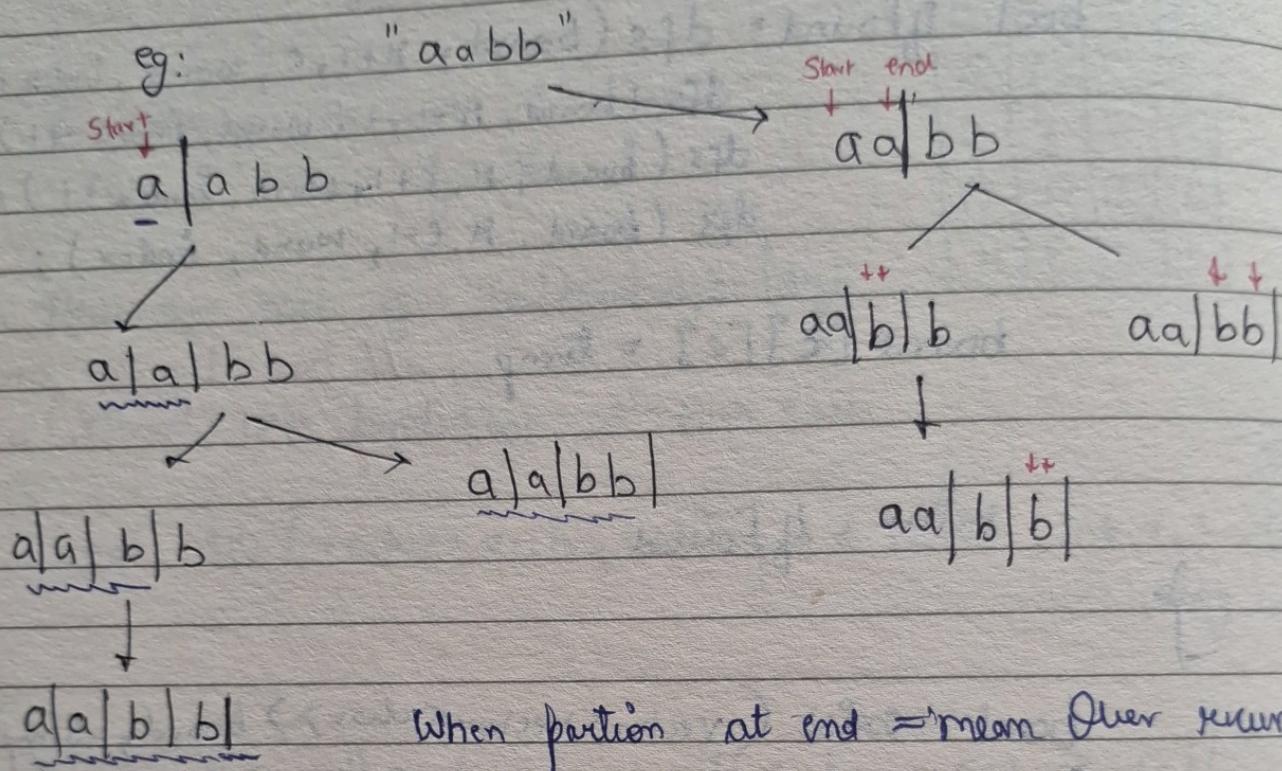
}

Palindrome Partition.

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Q Partition a string such that every substr is palindrome



- 1) Select a part for partition, check if it is palindrome
- 2) If Yes, Then only add it and explore further by recursion

```
vector<vector<string>> Partition (string s)
```

```
{ define vector and
  vec is vector. }
```

```
backtrack(0, s, path, result);
return result;
```

// function to check Palindrome.

bool isPalindrome (const string &s, int left, int right)

{

while (left < right) {

if (s[left] != s[right]) return false;

} left++; right--;

return true;

// Backtracking function

void backtrack (int start, const string &s, vector<string>& path, vector<vector<string>> &result)

{

if (start == s.size())

{ result.push_back (path);

} return;

Iterate until a palindrome is found, then only explore by Heuristic

for (int end = start; end < s.size(); ++end)

{

if (isPalindrome (s, start, end),

{

path.push_back (s.substr (start, end - start + 1))

→ Index, length

backtrack (end + 1, s, path, result);

path.pop();

}

}

choose

explore

(In-choose del)