Depth-First Search (DFS) Algorithms with Time Complexity Analysis

Recursive DFS Algorithm

```
Algorithm 1 Recursive DFS(G, u)
 1: Input: Graph G = (V, E), starting vertex u
 2: Output: Visited vertices in depth-first order
 3: function DFS-RECURSIVE(G, u, visited)
        visited[u] \leftarrow \text{TRUE}
                                                                                        \triangleright Mark vertex u as visited
        for all v \in G.Adj[u] do
                                                                                  \triangleright Iterate over all neighbors of u
 5:
            if visited[v] = FALSE then
                                                                                                \triangleright If v is not visited
 6:
 7:
               DFS-RECURSIVE(G, v, visited)
                                                                                              \triangleright Recursive call for v
            end if
 8:
        end for
10: end function
```

Time Complexity Analysis of Recursive DFS

- **Initialization**: Marking all vertices as unvisited takes $\mathcal{O}(V)$.
- **Traversal**:
 - Each vertex is visited exactly once: $\mathcal{O}(V)$.
 - Each edge is explored exactly once in the adjacency list: $\mathcal{O}(E)$.
- **Overall Time Complexity**: $\mathcal{O}(V+E)$.

Iterative DFS Algorithm

```
Algorithm 2 Iterative DFS(G, s)
```

```
1: Input: Graph G = (V, E), starting vertex s
 2: Output: Visited vertices in depth-first order
 3: function DFS-ITERATIVE(G, s)
        Initialize:
        visited[v] \leftarrow \text{FALSE for all } v \in V
                                                                                      ▷ Mark all vertices as unvisited
 5:
        stack \leftarrow \text{Empty Stack}
 6:
 7:
        Push(stack, s)
                                                                          ▶ Push the starting vertex onto the stack
        while stack \neq \text{Empty do}
 8:
                                                                                     ▶ Process stack until it is empty
            u \leftarrow \text{Pop}(stack)
                                                                               ▶ Remove and process the top vertex
 9:
            if visited[u] = FALSE then
                                                                                                    \triangleright If u is not visited
10:
                visited[u] \leftarrow \text{TRUE}
11:
                                                                                                    \triangleright Mark u as visited
                for all v \in G.Adj[u] do
                                                                                     \triangleright Iterate over all neighbors of u
12:
                     if visited[v] = FALSE then
                                                                                                    \triangleright If v is not visited
13:
                        Push(stack, v)
                                                                                                  \triangleright Add v to the stack
14:
                     end if
15:
                end for
16:
17:
            end if
        end while
19: end function
```

Time Complexity Analysis of Iterative DFS

• **Initialization**: Marking all vertices as unvisited takes $\mathcal{O}(V)$.

- **Traversal**:
 - Each vertex is pushed and popped from the stack exactly once: $\mathcal{O}(V)$.
 - Each edge is explored exactly once in the adjacency list: $\mathcal{O}(E)$.
- **Overall Time Complexity**: $\mathcal{O}(V+E)$.

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Key Differences Between Recursive and Iterative DFS

- **Recursive DFS**:
 - Uses function call stack for maintaining the DFS traversal order.
 - Simpler to implement but limited by system recursion depth.
- **Iterative DFS**:
 - Uses an explicit stack to simulate the recursion.
 - More memory-efficient for very deep graphs, as it avoids stack overflow.