



ECE 281

Lesson 4 and 6 Notes

Objectives:

- Based on a problem description, design a solution using Boolean equations to solve the problem
- Utilize Boolean algebra to simplify Boolean equations or to derive a form that yields the most desirable hardware configuration
- Demonstrate the ability to describe a combinational digital system by a truth table, Boolean equation, Sigma/Pi notation, and schematic
- Given any one form, be able to produce the others

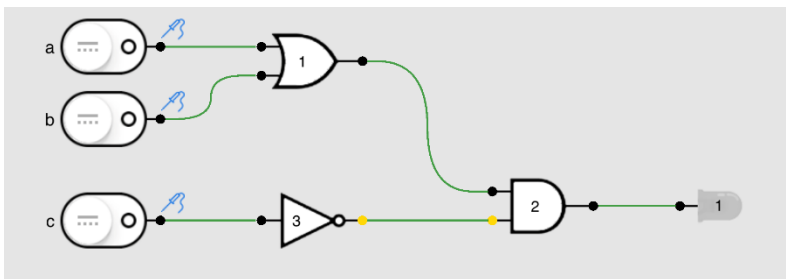
Review and Example #1: Lets create a logic circuit that will turn on an LED to see if you are eligible to apply for USAF Test Pilot School as a Flight Test Engineer. The key criteria are degree and time in service.

A – signals applicant has engineering degree

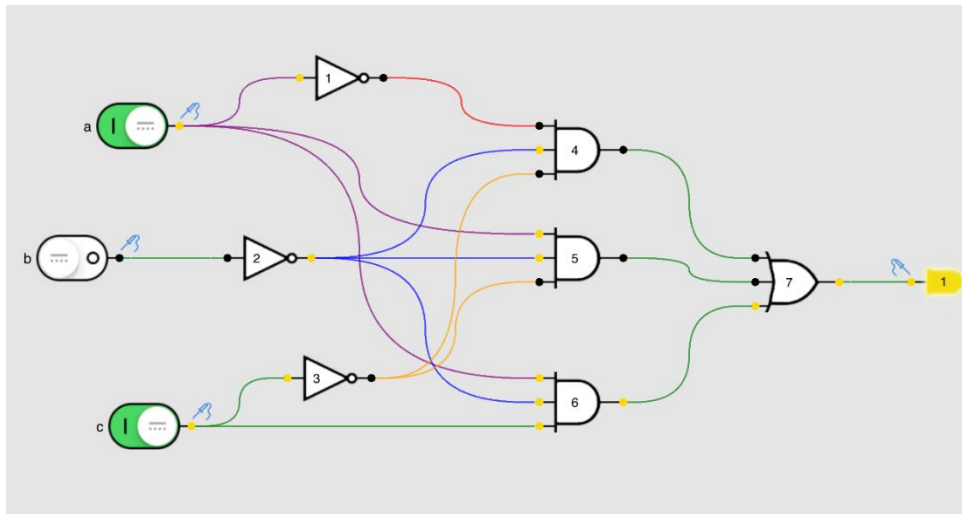
B – signals applicant has basic science degree

C – signals greater than 8 years time in service

$$\mathbf{Y} = (\mathbf{A} + \mathbf{B})\bar{\mathbf{C}}$$

[illegible]

Example #2) $Y = A'B'C' + AB'C' + AB'C$



Circuit Schematic Best Practices:

- A. Inputs on the left (or Top)
- B. Outputs on the right (or bottom)
- C. Gates flow from left to right
- D. Straight wires are best

Boolean Algebra Main Principles:

In Boolean Algebra, variables can take on either a true or false value. Using the rules we will discuss in this lesson, expressions can often be reduced to result in more efficient circuit implementation. While the second example above was already rather simple, we will prove an even simpler implementation through the course of today's lesson.

Terminology

Equivalent Representations: A and B \iff AB
 A or B \iff $A + B$
 Not A \iff A' \iff \overline{A}

Truth Table:

Minterm:

Maxterm:

Order of Operations:

Sum of Products:

| Input A | Input B | Output Y | Minterm (Form 1) | Minterm (Form 2) |
|---------|---------|----------|----------------------------|---------------------|
| 0 | 0 | | $\overline{A}\overline{B}$ | $A'B'$ |
| 0 | 1 | | $\overline{A}B$ | $A'B$ |
| 1 | 0 | | $A\overline{B}$ | AB' |
| 1 | 1 | | AB | AB |

Product of Sums:

| Input A | Input B | Output Y | Maxterm (Form 1) | Maxterm (Form 2) |
|---------|---------|----------|-------------------------------|---------------------|
| 0 | 0 | | $A + B$ | $A+B$ |
| 0 | 1 | | $A + \overline{B}$ | $A+B'$ |
| 1 | 0 | | $\overline{A} + B$ | $A'+B$ |
| 1 | 1 | | $\overline{A} + \overline{B}$ | $A'+B'$ |

NOTE: \overline{AB} is not the same as $\overline{A}\overline{B}$... more on this shortly

Commonly Used Boolean Algebra Properties (Review from zyBooks)

| Property | Name | Description |
|---|---|---|
| $A(B+C) = AB+AC$ $A+(BC) = (A+B)(A+C)$ | Distributive (AND) Distributive (OR) | (AND) Same as multiplication in regular algebra (OR) Not at all like regular algebra |
| $AB=BA$ $A+B=B+A$ | Commutative | Variable order does not matter. Good practice is to sort variables alphabetically |
| $(AB)C=A(BC)$ $(A+B)+C = A+(B+C)$ | Associative | Same as regular algebra |
| $AA' = 0$ $A+A' = 1$ | Complement (AND) Complement (OR) | (AND) Clearly one of A, A' must be 0. $1 \cdot 0 = 0 \cdot 1 = 0$ (OR) Clearly one of A, A' must be 1. $1+0 = 0+1 = 1$ |
| $A \cdot 1 = A$ $A+0 = A$ | Identity (AND) Identity (OR) | (AND) Result of $A \cdot 1$ is always A's value $0 \cdot 1 = 0$ $1 \cdot 1 = 1$ (OR) Result of $A+0$ is always A's value $0+0 = 0$ $1+0 = 1$ |
| $A \cdot 0 = 0$ $A+1 = 1$ | Null elements | Result doesn't depend on the value of A |
| $A \cdot A = A$ $A+A = A$ | Idempotent | Duplicate values can be removed |
| $(A')' = A$ | Involution | $(0')' = (1')' = 0$ $(1')' = (0')' = 1$ |
| $A(A+B) = A$ $A+(AB) = A$ | Covering | In either case, the value of B is irrelevant |
| $(AB)+(AB') = A$ $(A+B)(A+B') = A$ | Combining | This is an application of the complement and identity properties listed above |
| $AB+B'C = AC$ $(A+B)(B'+C) = (A+C)$ | Consensus | The conjunction of all unique terms |
| $(AB)' = A' + B'$ | DeMorgan's Law (for AND) | Each literal complemented ANDs become ORs |
| $(A+B)' = A'B'$ | DeMorgan's Law (for OR) | Each literal complemented, ORs become ANDs |

Visualization of Demorgan's Theorem:

$$Y = (AB)' =$$

$$Y = (A+B)'$$

Simplification of Example #2 Using Boolean Algebra:

Revisiting the concept that $\overline{A\overline{B}}$ is not the same as $\overline{A\overline{B}}$

$$\overline{A\overline{B}}$$

| A | B | Y |
|---|---|---|
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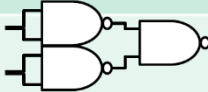
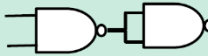

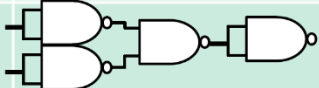
$$\overline{A\overline{B}}$$

| A | B | Y |
|---|---|---|
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Which logic gate represents $\overline{A\overline{B}}$?

Through application of DeMorgan's law, which single logic gate represents $\overline{A\overline{B}}$?

And for the final topic today, any logic operation can be represented solely with NAND gates.

| | Standard | NAND | Schematic |
|------|----------------------------|--|---|
| NOT | \overline{A} | $\overline{A\overline{A}}$ |  |
| OR | $A + B$ | $\overline{\overline{A}\overline{B}}$ |  |
| AND | AB | $\overline{\overline{A\overline{B}}}$ |  |
| NAND | $\overline{A\overline{B}}$ | $\overline{A\overline{B}}$ |  |
| NOR | $\overline{A + B}$ | $\overline{\overline{\overline{A}\overline{B}}}$ |  |