





### 1. Players Relation

Given FDs:

$\text{player\_id} \rightarrow \{\text{batting\_style}, \text{bowling\_style}, \text{name}, \text{role}, \text{team\_Id}\}$

Compute the Closure of Player\_ID :

$\text{player\_id}^+ = \{ \text{player\_id}, \text{batting\_style}, \text{bowling\_style}, \text{name}, \text{role}, \text{team\_id} \}$

Since  $\text{player\_id}^+$  includes all attributes of the relation,  $\text{player\_id}$  is a Key.

For all functional dependencies in this relation, the left side is  $\text{player\_id}$ , which is a Key.

Hence, the Players relation satisfies BCNF.

### 2. Team Relation

Given FDs:

$\text{team\_id} \rightarrow \{\text{tname}, \text{caption\_id}, \text{home\_ground}\}$

Compute the Closure of team\_id :

$\text{team\_id}^+ = \{ \text{team\_id}, \text{tname}, \text{caption\_id}, \text{home\_ground} \}$

Since  $\text{team\_id}^+$  includes all attributes of the relation,  $\text{team\_id}$  is a Key.

For all functional dependencies in this relation, the left side is  $\text{team\_id}$ , which is a Key.

Hence, the Team relation satisfies BCNF.

### 3. Venue Relation

Given FDs:

$\text{venue\_id} \rightarrow \{\text{city}, \text{state}, \text{grd\_name}\}$

Compute the Closure of Venue\_id :

$\text{venue\_id}^+ = \{ \text{venue\_id}, \text{city}, \text{state}, \text{grd\_name} \}$

Since  $\text{venue\_id}^+$  includes all attributes of the relation,  $\text{venue\_id}$  is a Key.

For all functional dependencies in this relation, the left side is  $\text{venue\_id}$ , which is a Key.

Hence, the Venue relation satisfies BCNF.

### 4. Match\_details Relation

Given FDs:

$\text{match\_id} \rightarrow \{\text{match\_date}, \text{toss\_decision}, \text{status}, \text{team1}, \text{team2}, \text{toss\_winner}, \text{venue}, \text{match\_time}\}$

Compute the Closure of Match\_id :

match\_id+ = { match\_id, match\_date, toss\_decision, status, team1, team2, toss\_winner, venue, match\_time }

Since match\_id+ includes all attributes of the relation, match\_id is a Key.

For all functional dependencies in this relation, the left side is match\_id, which is a Key.

Hence, the Match\_details relation satisfies BCNF.

## 5. Playing\_11 Relation

Given FDs:

{ match\_id, player\_id } → team\_id  
{ match\_id, player\_id } → is\_substitute  
{ match\_id, player\_id } → substituted\_out

Compute the Closure of { match\_id, player\_id } :

{ match\_id, player\_id }+ = { match\_id, player\_id, team\_id, is\_substitute, substituted\_out }

Since { match\_id, player\_id }+ includes all attributes of the relation, { match\_id, player\_id } is a Key.

For all functional dependencies in this relation, the left side is { match\_id, player\_id }, which is a Key.

Hence, the Playing\_11 relation satisfies BCNF.

## 6. Match\_officials Relation

Given FDs:

match\_id → {on\_field1, on\_field2, third\_ump}

Compute the Closure of match\_id :

match\_id+ = { match\_id, on\_field1, on\_field2, third\_ump }

Since match\_id+ includes all attributes of the relation, match\_id is a Key.

For all functional dependencies in this relation, the left side is match\_id, which is a Key.

Hence, the Match\_officials relation satisfies BCNF.

## 7. Partnership Relation

Given FDs:

partnership\_id → {player1, player2, match\_id, runs, partnership\_no, inning\_no}

Compute the Closure of partnership\_id :  
partnership\_id+ = { partnership\_id, player1, player2, match\_id, runs,  
partnership\_no, inning\_no }

Since partnership\_id+ includes all attributes of the relation,  
partnership\_id is a Key.

For all functional dependencies in this relation, the left side is  
partnership\_id, which is a Key.  
Hence, the Partnership relation satisfies BCNF.

#### 8. Per\_ball\_data Relation

Given FDs:

ball\_id → {over\_no, ball\_no, of\_match, inning, extras, wicket\_no,  
total\_score, runs\_scored, on\_strike, off\_strike, bowled\_by,  
striker\_batting\_position, type\_of\_extra, boundary}

Compute the Closure of ball\_id :  
ball\_id+ = { ball\_id, over\_no, ball\_no, of\_match, inning, extras,  
wicket\_no, total\_score, runs\_scored, on\_strike, off\_strike, bowled\_by,  
striker\_batting\_position, type\_of\_extra, boundary }

Since ball\_id+ includes all attributes of the relation, ball\_id is a Key.

For all functional dependencies in this relation, the left side is  
ball\_id, which is a Key.  
Hence, the Per\_ball\_data relation satisfies BCNF.

#### 9. Wickets Relation

Given FDs:

wicket\_id → {wicket\_type, player\_got\_out, caught\_by, run\_out\_by, on\_ball,  
stumped\_by}

Compute the Closure of Wicket\_id :  
wicket\_id+ = { wicket\_id, wicket\_type, player\_got\_out, caught\_by,  
run\_out\_by, on\_ball, stumped\_by }

Since wicket\_id+ includes all attributes of the relation, wicket\_id is a  
Key.

For all functional dependencies in this relation, the left side is  
wicket\_id, which is a Key.  
Hence, the Wickets relation satisfies BCNF.

#### 10. Final\_result Relation

Given FDs:

match\_id → {winner, loser, score\_of\_winner, score\_of\_loser,  
win\_run\_margin, win\_wicket\_margin, player\_of\_the\_match}

Compute the Closure of Match\_id :

match\_id+ = { match\_id, winner, loser, score\_of\_winner, score\_of\_loser, win\_run\_margin, win\_wicket\_margin, player\_of\_the\_match }

Since match\_id+ includes all attributes of the relation, match\_id is a Key.

For all functional dependencies in this relation, the left side is match\_id, which is a Key.

Hence, the Final\_result relation satisfies BCNF.

#### 11. Player\_performance\_in\_a\_match Relation

Given FDs:

{ in\_match, player } → {runs, balls\_played, wickets\_taken, overs\_bowled, runs\_conceded, fours, sixes, fours\_conceded, sixes\_conceded}

Compute the Closure of { in\_match, player } :

{ in\_match, player }+ = { in\_match, player, runs, balls\_played, wickets\_taken, overs\_bowled, runs\_conceded, fours, sixes, fours\_conceded, sixes\_conceded }

Since { in\_match, player }+ includes all attributes of the relation, { in\_match, player } is a Key.

For all functional dependencies in this relation, the left side is { in\_match, player }, which is a Key.

Hence, the Player\_performance\_in\_a\_match relation satisfies BCNF.

#### 12. Points\_table Relation

Given FDs:

team\_id → {matches\_played, losses, wins, no\_result, points, NRR}

Compute the Closure of Team\_id :

team\_id+ = { team\_id, matches\_played, losses, wins, no\_result, points, NRR }

Since team\_id+ includes all attributes of the relation, team\_id is a Key.

For all functional dependencies in this relation, the left side is team\_id, which is a Key.

Hence, the Points\_table relation satisfies BCNF.

#### 13. Users Relation

Given FDs:

user\_id → {password, last\_access, date\_of\_sign\_in}

Compute the Closure of User\_id :

$user\_id^+ = \{ user\_id, password, last\_access, date\_of\_sign\_in \}$

Since  $user\_id^+$  includes all attributes of the relation,  $user\_id$  is a Key.

For all functional dependencies in this relation, the left side is  $user\_id$ , which is a Key.

Hence, the User relation satisfies BCNF.

Hence, All relations are in Boyes Code Normal Form.