

A simplicial homology algorithm for Lipschitz optimisation

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Abstract The simplicial homology global optimisation (SHGO) algorithm is a general purpose global optimisation algorithm based on applications of simplicial integral homology and combinatorial topology. SHGO approximates the homology groups of a complex built on a hypersurface homeomorphic to a complex on the objective function. This provides both approximations of locally convex subdomains in the search space through Sperner's lemma and a useful visual tool for characterising and efficiently solving higher dimensional black and grey box optimisation problems. This complex is built up using sampling points within the feasible search space as vertices. The algorithm is specialised in finding all the local minima of an objective function with expensive function evaluations efficiently which is especially suitable to applications such as energy landscape exploration. SHGO was initially developed as an improvement on the topographical global optimisation (TGO) method. It is proven that the SHGO algorithm will always outperform TGO on function evaluations if the objective function is Lipschitz smooth. In this paper SHGO is applied to non-convex problems with linear and box constraints with bounds placed on the variables. Numerical experiments on linearly constrained test problems show that SHGO gives competitive results compared to TGO and the recently developed Lc-DISIMPL algorithm as well as the PSwarm, LGO and DIRECT-L1 algorithms. Furthermore SHGO is compared with the TGO, basinhopping (BH) and differential evolution (DE) global optimisation algorithms over a large selection of

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black-box problems with bounds placed on the variables from the SciPy benchmarking test suite. A Python implementation of the SHGO and TGO algorithms published under a MIT license can be found from <https://bitbucket.org/upiamcompthermo/shgo/>.

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1 Introduction

1.1 Objective function statement and nomenclature

Consider a general optimisation problem of the form

$$\begin{aligned} \min_{\mathbf{x}} f(\mathbf{x}) \\ \text{s.t. } \mathbf{g}(\mathbf{x}) \geq 0 \end{aligned} \quad (1)$$

The continuous real objective function $f(\mathbf{x})$ maps a vector of dimension n to a scalar value. It can be either smooth or non-smooth depending on the local minimisation method used. The variables \mathbf{x} are assumed to be bounded. In this publication we mainly consider real, smooth, but not necessarily convex functions with linear constraint functions. In addition we will assume that the objective function has a finite number of local minima

$$f : \mathbb{R}^n \rightarrow \mathbb{R} \quad (2)$$

\mathbf{g} maps the set of linear constraints

$$\mathbf{g} : [\mathbf{l}, \mathbf{u}]^n \rightarrow \mathbb{R}^m \quad (3)$$

For example, if lower and upper bounds l_i and u_i are implemented for each variable then we have an initially defined hyperrectangle

$$\mathbf{x} \in \Omega \subseteq [\mathbf{l}, \mathbf{u}]^n = [l_1, u_1] \times [l_2, u_2] \times \dots \times [l_n, u_n] \subseteq \mathbb{R}^n \quad (4)$$

where Ω is the limited feasible subset excluding points outside the bounds and constraints:

$$\Omega = \{\mathbf{x} \in [\mathbf{l}, \mathbf{u}]^n \mid \mathbf{g}_i(\mathbf{x}) \geq 0, \forall i = 1, \dots, m\} \quad (5)$$

Since the constraints in \mathbf{g} are linear the set Ω is always a compact space.

In the development of SHGO several concepts from algebraic and combinatorial topology [18] are required. The following definition was adapted from Hatcher [15, p. 9].

Definition 1 A k -simplex is a set of $n + 1$ vertices in a convex polyhedron of dimension n . Formally if the $n + 1$ points are the $n + 1$ standard $n + 1$ basis vectors for $\mathbb{R}^{(n+1)}$. Then the n -dimensional k -simplex is the set

$$S^n = \left\{ (v_1, \dots, v_{n+1}) \in \mathbb{R}^{n+1} \mid \sum_{i=1}^{n+1} v_i = 1, v_i \geq 0 \right\}$$

For example, a 2-simplex is a triangle and a 3-simplex is a tetrahedron. We will use the following combinatorial definition of a simplicial complex [15, p. 107].

Definition 2 A *simplicial complex* \mathcal{H} is a set \mathcal{H}^0 of vertices together with sets \mathcal{H}^n of n -simplices, which are $(n + 1)$ -element subsets of \mathcal{H}^0 . The only requirement is that each $(k + 1)$ -elements subset of the vertices of an n -simplex in \mathcal{H}^n is a k -simplex, in \mathcal{H}^k .

Thus each n -simplex has $n + 1$ distinct vertices, and no other n -simplex has this same set of vertices.

In this publication the \mathcal{H} symbol will be used to represent a (finite) simplicial complex rather than the more standard Δ to avoid confusion with the difference and Laplacian operators common in optimisation. The superscript \mathcal{H}^k represents the subset of k -dimensional simplices where for an n dimensional problem the highest dimensional k -simplex contains $n + 1$ vertices. Finally we define a k -chain [18].

Definition 3 A k -chain is a union of simplices.

For example a 0-chain is a set of vertices, a 1-chain is a set of edges and a 2-chain is a set of triangles. $C(\mathcal{H}^k)$ denotes a k -chain of k -simplices. A vertex in \mathcal{H}^0 is denoted by v_i . If v_i and v_j are two endpoints of a directed edge in \mathcal{H}^1 from v_i to v_j then the symbol $\overline{v_i v_j}$ represents the edge so that it is bounded by the 0-chain $\partial(\overline{v_i v_j}) = v_j - v_i$ and similarly for an edge directed from v_j to v_i , we have, $\partial(\overline{v_j v_i}) = \partial(-\overline{v_i v_j}) = v_i - v_j$. Higher dimensional simplices can be represented and directed in a similar manner, for example a triangle consisting of three vertices v_i, v_j and v_k directed as $\overline{v_i v_j v_k}$ has the boundary of directed edges $\partial(\overline{v_i v_j v_k}) = \overline{v_i v_j} + \overline{v_j v_k} + \overline{v_k v_i}$.

1.2 Multimodal objective functions and local minima mapping

Non-convex problems are commonly solved using global optimisation methods. One such example is the topographical global optimisation (TGO) method [17,44–46] which is a clustering algorithm that finds several local minima from which the (approximate) global minimum is found. It is often desirable to find all the local minima of the objective function. For example in applications such as energy landscape exploration of potential models wherein mapping the local minima of the potential functions can provide valuable insights into the system. Algorithms such as the basin-hopping global optimisation algorithm are typically used to find these points [49].

The graph extracted from the topographical global optimisation (TGO) [17,44–46] topograph (as described in Sect. 2) is unsatisfactory in some ways. Primarily because several starting points in the same locally convex domain can be generated even when enough information from the objective function sampling is known to prevent this from occurring. This leads to superfluous function evaluations in the local minimisation step of the algorithm. Contrary to intuition, this problem is exacerbated by increasing the number of initial sampling points used in the algorithm as demonstrated in Sect. 2. This can lead to a very large number of function evaluations required to solve the problem. In particular in multimodal energy surfaces where the local minima can often be located in short distances relative to the search space [52] and thus requires a large number of initial sampling to locate all these domains. Some shortcomings in using the TGO method to map local minima are:

- Geometric information available from the sampling points is being disregarded by the graphs built up using only the Euclidean distance metric.
- Knowledge of the number and location of local minimisers in a given sampling set is not being used to the full extent.
- More than one minimiser might be produced in the same locally convex domain and there is no guarantee that a minimiser set produced by TGO will be in the locally convex

domains of all local minima even if the number of local minima is known and a minimiser set of this cardinality is produced.

By constructing a directed simplicial complex we show that the simplicial homology global optimisation (SHGO) algorithm does not produce superfluous starting points for the class of all Lipschitz smooth functions resulting in more efficient performance for these problems compared to TGO. The directed complex is also used to approximate the homology group of the objective function hypersurface which, using integral homology version of the Invariance Theorem [18], allows for efficient mapping of optimisation problems where the number of local minima is known a-priori.

1.3 Derivative-free methods for Lipschitz optimisation problems

Both the SHGO and TGO algorithms only make use of function evaluations without requiring the derivatives of objective functions. This makes them applicable to black-box global optimisation problems. A recent review and experimental comparison of 22 derivative-free optimisation algorithms by Rios and Sahinidis [38] concluded that global optimisation solvers such as TOMLAB/MULTI-MIN, TOMLAB/GLCCLUSTER, MCS and TOMLAB/LGO perform better, on average, than other derivative-free solvers in terms of solution quality within 2500 function evaluations. Both the TOMLAB/GLCCLUSTER and MCS [20] implementations are based on the well-known DIRECT (DIviding RECTangle) algorithm [23].

The DISIMPL (DIviding SIMPLices) algorithm was recently proposed by Paulavičius and Žilinskas [33]. The experimental investigation in [33] shows that the proposed simplicial algorithm gives very competitive results compared to the DIRECT algorithm. DISIMPL has been extended in [34,36]. The Gb-DISIMPL (Globally-biased DISIMPL) was compared in [36] to the DIRECT and DIRECTI methods in extensive numerical experiments on 800 multidimensional multiextremal test functions.

In a recent adaption of DISIMPL for linearly constrained optimisation problems, Lc-DISIMPL [35] showed extremely competitive results compared to the PSwarm [47] and DIRECT-L1 algorithms [13]. In particular the Lc-DISIMPL-v algorithm was shown to solve the problems in a fewer number of function evaluations on average and was the only algorithm to converge on all of the test problems. In this publication both the SHGO and TGO algorithms were tested on the same problem set and the results are compared to the data from [35] which also contains results on the PSwarm [47] and DIRECT-L1 algorithms [13].

The DISIMPL algorithm is the most similar to SHGO in the sense that both make use of a simplicial complex. DISIMPL uses a simplicial complex in a spatial partitioning of the initial search space. Since the geometric structure of the two algorithms are related, it is reasonable to expect some theoretical relation of its properties. In particular the graph structure in the DISIMPL-v algorithm [35] can be used to build the directed simplicial complex used by SHGO. In Sect. 5 we also show how some of the same principles developed for SHGO can also be applied in the DISIMPL-v algorithm since the same information is readily available to the algorithm.

1.4 Overview of this publication

The TGO method is briefly reviewed in Sect. 2 closely following the formalism developed by Henderson et al. [17]. In Sect. 3 we provide numerical examples of TGO which is then used as an informal experimental motivation for extending the algorithm. These two sections are important for continuity and understanding of the improved features of SHGO, in particular

Definition 9 which will be used as a performance criterion. In Sect. 4 we present the most immediately apparent extension of TGO and illustrate the shortcomings of that approach. The new SHGO method is then formally presented in Sect. 5. In Sect. 6 we provide experimental results of linearly constrained problems comparing the SHGO, TGO, Lc-DISIMPL [35], PSwarm [47] and DIRECT-L1 [13] algorithms. Furthermore SHGO is compared with the TGO, basinhopping (BH) and differential evolution (DE) global optimisation algorithms over a large selection of black-box problems from the SciPy [22] global optimisation benchmarking test suite. We conclude with various recommendations for possible further improvements of SHGO.

2 Topographical global optimisation (TGO)

The Topographical Global Optimisation (TGO) was originally conceived by Törn [45] and Henderson et al. [16, 17] introduced new formalisms and empirical methods to determine hyperparameters described in this section. Henderson et al. [17] also presents the algorithm in an introductory fashion. It is in essence an iterative clustering algorithm that maps the hypersurface of the objective function into a topography matrix (called a t -matrix) and then finds a certain number of starting points referred to as local minimisers. A local search using the local minimisers as starting points is then used to find each minimum from which the global minimum is finally calculated. Henderson et al. [17] used the feasible direction interior-point method proposed by Herskovits [19] in this step. The feasible direction interior-point method allows for minimisation of problems with linear and/or nonlinear equality constraints; an extension by Henderson et al. [17] of the original applications of Törn [45]. The TGO method consists of three steps:

1. Uniform random sampling generation of N points in the search space.
2. Construction of the topograph, which is a directed graph with the sampled points as vertices on a k -nearest neighbours basis with the direction of the arc directed towards a point with a larger function value.
3. Local minimisation of topograph minimisers.

2.1 Step 1: Random sampling point generation

In order to generate the uniform sampling points within Ω the deterministic Sobol sequence is used in this publication [17, 40]. Other possible low discrepancy sequences such as the Halton and Van der Corput sequences [26] can also be used in this step. An efficient Gray code implementation was proposed by Antonov and Saleev [2] wherein a single XOR operation for each dimension can be used to find the next sampling point in the sequence $x_{n,i} = x_{n-1,i} \oplus v_{k,i}$. An adaptation of this method is available in the open source Python library UQToolbox [5]. The Sobol sequenced points are generated within the n -dimensional hypercube $[0, 1]^n \in \mathbb{R}^n$, providing a uniform distribution on the hypersurface within this space. In the current implementation this set of points is stretched across the lower and upper bounds to form the hyperrectangle $[\mathbf{l}, \mathbf{u}]^n = [l_1, u_1] \times [l_2, u_2] \times \cdots \times [l_n, u_n] \subseteq \mathbb{R}^n$. The subset of feasible points contained in Ω is found by discarding any points lying outside the constraints $\mathbf{g}(\mathbf{x}) > 0$.

2.2 Step 2: Construction of the topograph

The topograph is constructed from the generated sampling points within Ω . From the topograph several global minimisers in f are found using the definitions developed in this section

which are then used as starting points for local minimisation routines. First N points are selected from the uniformly generated sequence of points within the feasible domain of $\Omega \subset \mathbb{R}^n$. Points generated by the sequence that lie outside the constraints are excluded. The points are denoted by $\mathbf{p}_i, i = 1, 2, 3 \dots N$. Next for each point \mathbf{p}_i a reference list is constructed by ordering the other $N - 1$ points from their nearest to farthest Euclidean distances. These ordered lists make up the rows of the topography matrix (or topograph). Furthermore, for some point $\mathbf{p}_j \in \{1, 2, 3 \dots (N - 1)\}$ in the row with the first entry \mathbf{p}_i , a sign is assigned as follows:

$$\text{sign}(\mathbf{p}_j) = \begin{cases} f(\mathbf{p}_j) \geq f(\mathbf{p}_i) & \rightarrow + \\ f(\mathbf{p}_j) < f(\mathbf{p}_i) & \rightarrow - \end{cases}$$

In order to demonstrate this construction we will define this ordered list in such a way that the increasing indices represent an ordered list of the nearest points to \mathbf{p}_i , that is $\|\mathbf{p}_i - \mathbf{p}_{i+1}\| \leq \|\mathbf{p}_{i+1} - \mathbf{p}_{i+2}\| \forall i$. Suppose for example that $f(\mathbf{p}_2) \geq f(\mathbf{p}_1)$, $f(\mathbf{p}_3) < f(\mathbf{p}_1)$ and $f(\mathbf{p}_N) \geq f(\mathbf{p}_1)$, the resulting topograph with the first row known is:

$$t\text{-matrix} = \left(\begin{array}{c|cccc} \mathbf{p}_1 & +\mathbf{p}_2 & -\mathbf{p}_3 & \dots & +\mathbf{p}_N \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ \mathbf{p}_N & \mathbf{p}_j & \dots & \mathbf{p}_j & \mathbf{p}_j \end{array} \right) \quad (6)$$

Note that the remaining rows (represented by unknown points and signs \mathbf{p}_j) are constructed similarly to the first row for every \mathbf{p}_i row. The topography matrix can be interpreted as a directed graph, where the signs represent the directed arcs on the graph. It should also be noted that if \mathbf{g} contains non-linear constraints then the graphs produced by the topograph may be connected across disconnected and/or non-convex subspaces of Ω . Finally, it should further be noted that these signs represent direction of the graph structure only, they are not the usual operation of a scalar acting on a vector. Example 1 in Sect. 3 demonstrates the construction of the topograph numerically.

Given an integer $1 \leq k \leq (N - 1)$, the $N \times k$ submatrix obtained by considering only the k -nearest neighbours is called the k - t -matrix. For example for $k = 1$:

$$1\text{-}t\text{-matrix} = \left(\begin{array}{c|c} \mathbf{p}_1 & +\mathbf{p}_2 \\ \vdots & \vdots \\ \mathbf{p}_N & \mathbf{p}_j \end{array} \right) \quad (7)$$

for $k = 2$:

$$2\text{-}t\text{-matrix} = \left(\begin{array}{c|cc} \mathbf{p}_1 & +\mathbf{p}_2 & -\mathbf{p}_3 \\ \vdots & \vdots & \vdots \\ \mathbf{p}_N & \mathbf{p}_j & \mathbf{p}_j \end{array} \right) \quad (8)$$

and so forth. The k - t -matrix is a representation of its k^+ -topograph where every row forms a directed subgraph.

The following definitions adapted from Henderson et al. [17] are used to find the global minimisers of the objective function.

Definition 4 Given an integer $1 \leq k \leq (N - 1)$, the i th row of the k - t -matrix is said to be a positive row, if all its elements have a plus sign. That is iff $f(\mathbf{p}_j) \geq f(\mathbf{p}_i) \forall j$.

Definition 5 Given an integer $1 \leq k \leq (N - 1)$, a sampling point \mathbf{p}_i has a positive reference in the k - t -matrix, if there exists $j \neq i$ such that (a) the j^{th} row of the k - t -matrix is a positive row and (b) the number $+i$ is an element of this j^{th} row.

Definition 6 Given an integer $1 \leq k \leq (N - 1)$, the sample point \mathbf{p}_i is called a local minimiser of f in the k^+ -topograph if the i^{th} row of the k - t -matrix is a positive row.

Definition 7 Given an integer $1 \leq k \leq (N - 1)$, the sample point \mathbf{p}_i is a global minimiser of f in the k^+ -topograph if \mathbf{p}_i is a local minimiser of f in the k^+ -topograph and, in addition, \mathbf{p}_i has no positive references in the k - t -matrix.

The following propositions can be readily demonstrated to show the consistency of the aforementioned definitions [17].

Proposition 1 *Given an integer $1 \leq k \leq (N - 1)$, the sample point \mathbf{p}_i is a global minimiser of f in the k^+ -topograph if and only if the sample point \mathbf{p}_i is the only minimiser of f in the k^+ -topograph which is global.*

Proposition 2 *Given an integer $1 \leq k \leq (N - 1)$, then the i th row of k - t -matrix is the only positive row of this matrix if and only if the sample point \mathbf{p}_i is the only minimiser of f in the k^+ -topograph which is global.*

Corollary 1 *Given an integer $1 \leq k \leq (N - 1)$, if the sample point \mathbf{p}_i is the only local minimiser of f in the k^+ -topograph, then \mathbf{p}_i is a global minimiser of f in this graph.*

In this publication we will use the paradigm that all local minimisers of f in the k^+ -topograph will be used for the local search (Paradigm 2.2 in Henderson et al. [17]). As described in [46] the number of local minimisers of f in the k^+ -topograph is greater than or equal to number of global minimisers in the topograph. We will therefore employ the following definition

Definition 8 Given an integer $1 \leq k \leq (N - 1)$, the minimiser pool \mathcal{M}^k is the set containing all local minimisers \mathbf{p}_i in the k^+ -topograph. The total number of starting points used in the local search step is equal to the cardinality of the minimiser pool $|\mathcal{M}^k|$.

The entire point of using k - t -matrices is because a t -matrix will always have at most one local (and thus global) minimiser. This is undesirable since this sampling point is not necessarily the starting point closest to the true global minimum of the objective function. Henderson et al. [17] developed a semi-empirical formula producing an integer value k_c which is used as an estimate for the optimal value for the integer k .

2.3 Step 3: Local minimisation

Each of the minimisers from the k_c - t -topograph is now used as a starting point in a local minimisation routine. The resulting minima are used to find the global minimum. Conceivably various local optimisation routines can be used to address a broad class of optimisations problems. For problems with non-linear inequality constraints Henderson et al. [17] used the feasible direction interior-point method proposed by Herskovits [19] minimising the objective function f subject to the set of inequality constraint functions \mathbf{g} using the minimiser set as the initial starting points for the algorithm. An algorithm used to solve the feasible direction interior-point method using the set of starting points calculated in step 2 is presented in detail by Henderson et al. [17].

In this publication we will mainly be using the sequential least squares quadratic programming optimisation algorithm (SLSQP) contained in the SciPy library originally developed by Kraft [24, 25]. Our Python implementation of the TGO algorithm published under an open source licence uses this algorithm as implemented in the SciPy library [12, 22].

3 Motivation and a one-dimensional prelude

In this section we will demonstrate how the Euclidean distance criterion in the TGO method disregards useful information about the (approximate) geometry of the objective function and we show how known information can be used effectively both in global optimisation and in mapping the local minima of objective functions as efficiently as possible. We also show how two important hyperparameters used by TGO, namely the number of sampling points N and the choice of k can be iteratively selected by intelligently exploiting information known from the topograph. This draws parallels to other works on iterative versions of TGO (I-TGO) [43] trying to extract information from black-box objective functions. The informal, but intuitive ideas developed here will later be extended more rigorously to higher dimensional surfaces. Note that from Eq. (5) Ω is always a compact space, this fact is important in several proofs used in this section.

Example 1 Consider the following objective function

$$\min_x f(x) = \frac{\sin(x)}{x}, \quad x \in \Omega = [1, 20] \quad (9)$$

In this instance of the bounded optimisation problem there are 3 local minima which we will try to map in as few function evaluations as possible.

Following the TGO procedure we start by generating low-discrepancy sampling points. The first $N = 10$ points in the 1-dimensional Sobol sequence is given by $\mathcal{P} = \{p_1 = 1.0, p_2 = 10.5, p_3 = 15.25, p_4 = 5.75, p_5 = 8.125, p_6 = 17.625, p_7 = 12.875, p_8 = 3.375, p_9 = 4.5625, p_{10} = 14.0625\} \subset \Omega$. After mapping the objective function at the set of sampling points

$$f : \begin{bmatrix} p_1 = 1.0 \\ p_2 = 10.5 \\ p_3 = 15.25 \\ p_4 = 5.75 \\ p_5 = 8.125 \\ p_6 = 17.625 \\ p_7 = 12.875 \\ p_8 = 3.375 \\ p_9 = 4.5625 \\ p_{10} = 14.0625 \end{bmatrix} \rightarrow \begin{bmatrix} f_1 = 0.84147 \\ f_2 = -0.08378 \\ f_3 = 0.02899 \\ f_4 = -0.08840 \\ f_5 = 0.11858 \\ f_6 = -0.05337 \\ f_7 = 0.02359 \\ f_8 = -0.06853 \\ f_9 = -0.21672 \\ f_{10} = 0.07091 \end{bmatrix} \quad (10)$$

the corresponding topograph is constructed

$$\begin{bmatrix} p_1 & -p_8 & -p_9 & -p_4 & -p_5 & -p_2 & -p_7 & -p_{10} & -p_3 & -p_6 \\ p_2 & +p_5 & +p_7 & +p_{10} & +p_3 & -p_4 & -p_9 & +p_6 & +p_8 & +p_1 \\ p_3 & +p_{10} & -p_6 & -p_7 & -p_2 & +p_5 & -p_4 & -p_9 & -p_8 & +p_1 \\ p_4 & -p_9 & +p_5 & +p_8 & +p_1 & +p_2 & +p_7 & +p_{10} & +p_3 & +p_6 \\ p_5 & -p_2 & -p_4 & -p_9 & -p_7 & -p_8 & -p_{10} & +p_1 & -p_3 & -p_6 \\ p_6 & +p_3 & +p_{10} & +p_7 & -p_2 & +p_5 & -p_4 & -p_9 & -p_8 & +p_1 \\ p_7 & +p_{10} & -p_2 & +p_3 & +p_5 & -p_6 & -p_4 & -p_9 & -p_8 & +p_1 \\ p_8 & -p_9 & +p_1 & -p_4 & +p_5 & -p_2 & +p_7 & +p_{10} & +p_3 & +p_6 \\ p_9 & +p_4 & +p_8 & +p_1 & +p_5 & +p_2 & +p_7 & +p_{10} & +p_3 & +p_6 \\ p_{10} & -p_3 & -p_7 & -p_2 & -p_6 & +p_5 & -p_4 & -p_9 & -p_8 & +p_1 \end{bmatrix} \quad (11)$$

The sampling points together with the objective function evaluations are plotted in Fig. 1. Using the empirical relation from Henderson et al. [17] the optimal k_c is calculated at

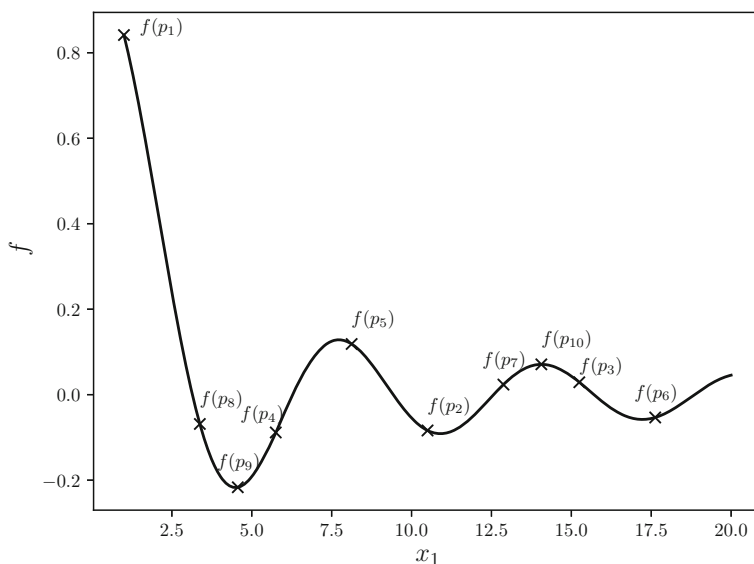


Fig. 1 Test function give by Eq. (9) with 10 Sobol sequenced sampling points

$k_c = 8$. Using Definition 6 we find that the resulting 8- t -matrix has only one minimiser; the global minimiser at $p_9 = 4.5625$. For the local minimisation we use the SLSQP method as implemented in the function `scipy.optimize.minimize` [22] to find the approximate global minimum at $x = 4.4934$.

Observing Fig. 1 it is immediately apparent that the set of 10 sampling points alone provides adequate information to deduce that there are at least 3 local minima. Observe that there are at least two other local minima since $f(p_2) < f(p_7) < f(p_5)$. So at least one local minimum exists in the domain $(p_5, p_7) \subset \mathbb{R}$ since between p_5 and p_2 we must have, by the mean value theorem (MVT), $\frac{df}{dx} < 0$ for some domain $x \in [p_5, p_2] \subset \mathbb{R}$. Similarly for $x \in [p_2, p_7] \subset \mathbb{R}$ we have by MVT $\frac{df}{dx} > 0$. Since f is a smooth, continuous function for $x \in (0, \infty)$ there must exist at least one stationary point $x \in (p_5, p_7) \subset \mathbb{R}$ where $\frac{df}{dx} = 0$. Furthermore we observe $f(p_6) < f(p_3)$ indicating another minimum in the domain $x \in (p_3, 20] \subset \mathbb{R}$ since the minimum must be either on the boundary or in $x \in (p_3, 20] \subset \mathbb{R}$ by the same argument as above.

The empirical relation by Henderson et al. [17] was mainly developed for the purpose of finding the global minimum. Therefore if only 10 sampling points are available, then to find more local minima using the TGO method is required to force a lower k value. Alternatively, since k_c is a function of N , simply sampling more points is sufficient to find all the local minima using Henderson's formula for this test problem. For example at $N = 16$ all 3 local minima are produced by TGO with Henderson's formula. Figure 2 shows the number of minimisers found at different k values for this example. The maximum minimiser set (other than using every sampling point as a starting point) can be trivially extracted by setting $k = 1$ and calculating $|\mathcal{M}^1|$. However, in this Example it leads to more starting points than optimal since at least two minimisers will be in the same convex basin domain and therefore converge to same minimum in the local minimisation step. This results in superfluous function evaluations without extracting more useful information from the objective function.

This idea drives the motivation behind the following definition.

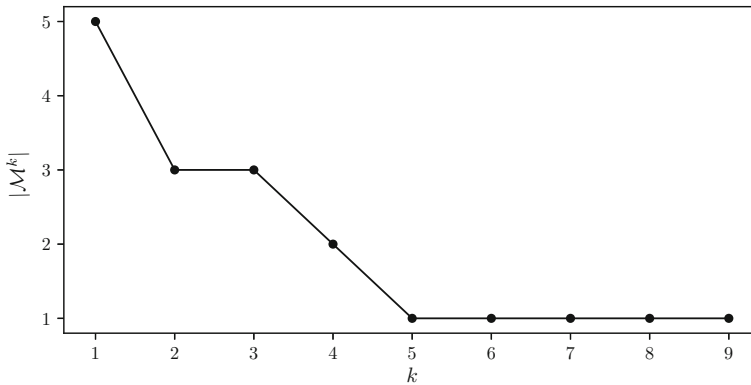


Fig. 2 Number of minimisers $|\mathcal{M}^k|$ found using the TGO method for different k values at $N = 10$

Definition 9 For a given set \mathcal{P} of N sampling points, k_{opt} is any integer $1 \leq k \leq (N - 1)$ that will produce the optimal minimiser set $\mathcal{M}^{k_{opt}}$ containing the maximum set of minimisers such that no two starting points extracted from $\mathcal{M}^{k_{opt}}$ will lead to the same minimum in the local optimisation step for some tolerance ϵ . In other words every element contained in $\mathcal{M}^{k_{opt}}$ should lie in a unique locally convex sub-domain.

Note that for a given N , $\mathcal{M}^{k_{opt}}$ might not produce all the true local minima of an objective function. What's important is that, given the information known from the sampling, the maximum number of local minima are found. In addition, no function evaluations are wasted in the local minimisation step which lead to the same minimum.

In Example 1 for $N = 10$ the optimal k values are $k_{opt} = \{2, 3\}$ which will produce 3 minimisers $|\mathcal{M}^2| = |\mathcal{M}^3| = 3$. We will now show that these lower k values carry unexploited information on the best approximate geometry of the objective function. For example in Fig. 3 we plot the $|\mathcal{M}^k|$ values corresponding to the set $k = \{1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$ for every sampling point range $N \in [2, 50]$.

From Fig. 3 we notice the special property of $k = 3$ for one dimensional objective functions sampled with the Sobol sequence.

Firstly, for a lower number of sampling points N it provides a higher number of starting minimisers than $k > 3$. Note that by inspection of Definition 6 it can be determined that any $k > 3$ value will always produce an equal or lower number of minimisers than $k = 3$. When adding columns to a positive row there are only two possibilities: the next sampling point in the row can either have a positive or a negative sign. All other elements in the row have a positive sign by definition (see Definition 6). If the next sampling point in the row has a positive sign then the row will just remain a positive row and the number of minimisers remain the same. If the point is a negative reference point then the row will no longer be a positive row and thus the point is no longer a minimiser, lowering the total.

Secondly it can be observed that $k = 3$ never calculates a number of starting minimisers higher than optimal unlike $k < 3$. Therefore by using $k = 3$ in Example 1 TGO will always find as many minimisers in as few sampling¹ function evaluations as possible and furthermore all local minima will be found when $N \geq 10$. It should be noted that the total number of function evaluations depends on the particular local minimisation algorithm used. However,

¹ Not necessarily total function evaluations since starting points closer to the local minima may provide better performance for a given local minimisation routines.

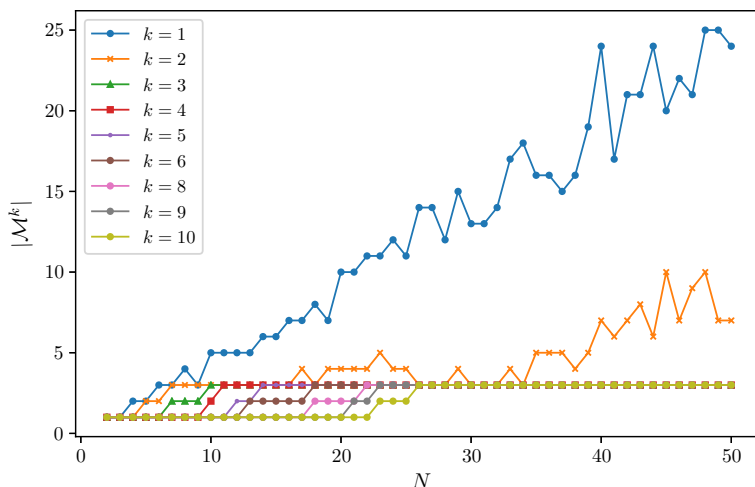


Fig. 3 Number of minimisers $|\mathcal{M}^k|$ found using the TGO method for the given k values at various sampling points N

it is apparent that each minimiser starting point is in a unique locally convex domain. It is tempting for an optimisation practitioner to use the size of the set of minimisers $|\mathcal{M}^3|$ as a stopping criterion for iterative sampling N of one dimensional objective functions. The practical usefulness of this idea can be demonstrated with the following example:

Example 2 The following instance of the optimisation problem has 13 local minima in the given domain

$$\min_x f(x) = -x \sin(x), \quad x \in \Omega = [1, 80] \quad (12)$$

From Fig. 4 we can deduce that the minimum number of sampling points required for $k = 3$ to find all local minima using the Sobol sequence is $N = 40$, this sampling is shown in Fig. 5. If $N < 40$ then there aren't enough sampling points to deduce that there are at least 13 locally convex domains from using the same arguments as in Example 1. Note for example that if we used a sequence that skipped p_1 then $N = 39$ would be adequate since $l = 1 < p_{32} < p_{33}$. Using our Python implementation of TGO [12] with $N = 40$ all 13 local minima of the objective function were found in a total of 285 function evaluations.

An example of a stopping criterion would be to stop sampling if $|\mathcal{M}^3|$ is unchanged after, say, 10 sampling point evaluations. The rate at which the number of elements in $|\mathcal{M}^3|$ grows with increasing N also provides a heuristic for characterising the multimodality and the geometry of the objective function. Objective functions that have a large number of local minima in a small domain (and relatively fewer minima in other larger domains) will have a much smaller growth in $|\mathcal{M}^3|$ for a given low-discrepancy sampling. This idea of continuously classifying and extracting approximate function characteristic information from the sampling points will be formalised and extended to higher dimensions in Sect. 5.

There is a simple reason why the 3- t -matrix has this quality in the first dimension for the optimisation problem given in Eq. (9). However, it is not guaranteed that this property holds for any sampling point distribution. In fact it holds true only under the following conditions:

1. Consider all points in the ordered sampling set from the smallest to greatest x value $\mathcal{P} = \{p_i \mid p_0 < p_1 < p_2 \dots < p_N - 1, p_i \in (x_l, x_u)\}$, excluding the supremum and infimum.

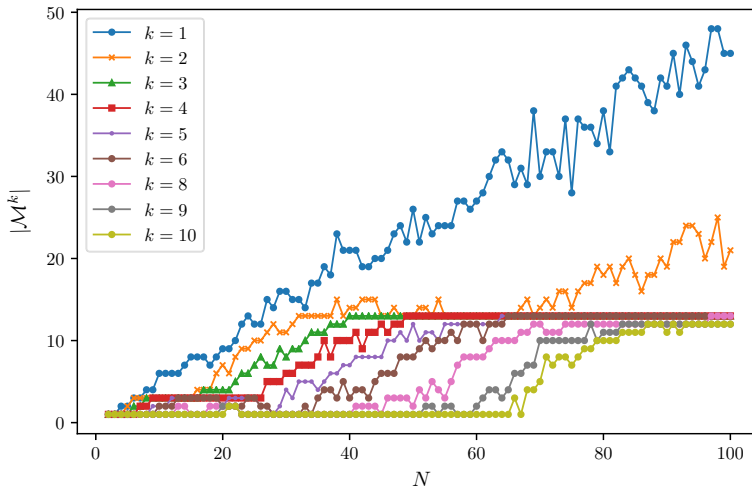


Fig. 4 Number of minimisers $|\mathcal{M}^k|$ found using the TGO method for the given k values at various sampling points N

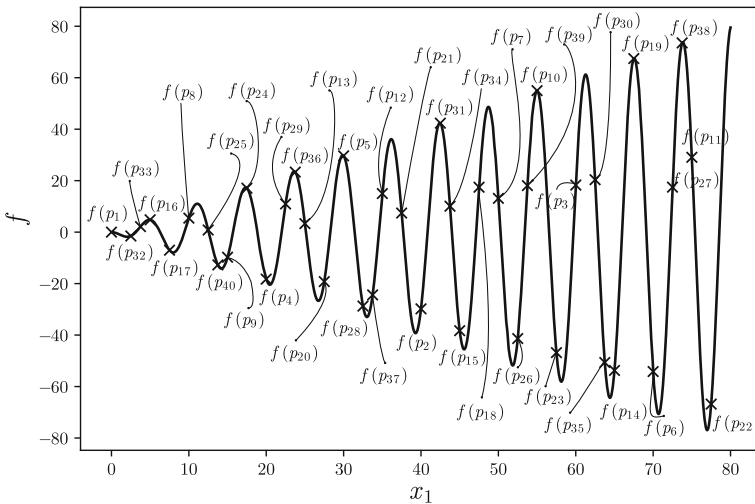


Fig. 5 Plot of the objective function in Example 2 for $N = 40$ sampling points

- For any given point p_i the Euclidean distance between p_i and 2 of its nearest sampling points $p_{i-1} < p_i < p_{i+1}$ should be less than the relative difference between p_i and a fourth point in the sampling sequence $|p_i - p_j|$ where $j \neq i, i-1, i+1$.

In fact it is easy to prove both that for a locally, strictly convex domain of f the 3-topograph construction can produce a larger minimiser pool \mathcal{M}^3 than optimal. It can also be shown that a construction must exist where the optimal number of minimisers will *always* be extracted regardless of the sampling distribution. Furthermore it can be shown that at most 3 sampling points within a locally convex domain $x \in [x_l, x_u]$ is required to produce enough information so that only one minimiser in the domain is produced.

Theorem 1 *There exists a 1-dimensional sampling sequence such that $k = 3$ will produce a minimiser pool larger than optimal as defined by Definition 9.*

Proof Consider a subdomain $x \in [x_l, x_u] \subset \mathbb{R}$ for which f is strictly convex. We define the set of N sampling points \mathcal{P} ordered in such a way that

$$\mathcal{P} = \{p_i \mid p_0 < p_1 < p_2 < \dots < p_{N-1}, p_i \in (x_l, x_u)\}$$

Let $\mathcal{F} = \{f_0, f_1, f_2, \dots, f_{N-1}\}$ be set of one-to-one function values corresponding to the points mapped by $f: \mathcal{P} \rightarrow \mathcal{F}$.

Suppose we have $f_1 < f_0$ and $f_1 < f_2 < f_3, \dots, f_{N-1}$. By construction we have $|p_1 - p_2| < |p_1 - p_3| < |p_1 - p_4| < |p_1 - p_5|$ then by the Definitions 4, 5 and 6 p_1 is a minimiser of the $3 - t$ -topograph. Suppose we have a sampling distribution such that $|p_2 - p_3| < |p_1 - p_2|, |p_2 - p_4| < |p_1 - p_2|$ and $|p_2 - p_5| < |p_1 - p_2|$ then by the Definitions 4, 5 and 6 p_3 is also a minimiser of the $3 - t$ -topograph. Therefore more than two minimisers are produced in the same locally convex sub-domain of $[x_l, x_u]$. We have shown that \mathcal{M}^3 can produce a minimiser pool larger than optimal which concludes the proof. \square

Lemma 1 *A construction exists that will always produce a minimiser pool larger than optimal as defined by Definition 9 for any given 1-dimensional sampling sequence.*

Now suppose that instead of using only the Euclidean distance metric we also invoke knowledge of the nearest point in every cartesian direction. We use the criterion that a minimiser point p_i is a minimiser iff with the ordering constructed in \mathcal{P} and \mathcal{F} we have $f_i < f_{i-1}$ and $f_i < f_{i+1}$. With this definition if the point p_i is a minimiser then no other point meets the criterion since by construction of the sampling in the locally convex domain $f_0 > f_1 > \dots > f_{i-1} > f_i$ and $f_{i+1} < f_{i+2} < f_{i+3} < \dots < f_{N-1}$. This proves Lemma 1.

Finally note that only information from the 3 points in the locally convex sub-domain of $[x_l, x_u]$ and their corresponding function values f_{i-1} , f_i and f_{i+1} are needed to produce a minimiser using this criterion.

An important consequence here is that for low discrepancy sequences in higher dimensions and for less well behaved objective functions the topographs connected with the Euclidean distance metrics will similarly discard available information about the local geometry. This produces larger than optimal minimiser pools leading to very high numbers of function evaluations needed to solve the problem.

In the following section we will develop a more efficient algorithm that will make use of this information. SHGO will always produce equivalent results to this algorithm in the one dimensional case.

4 Axially directed topograph

Based on the observations from Sect. 3 we develop the ATGO (axially directed topographical global optimisation) algorithm that, for a given sampling set, always uses the optimal number of starting minimisers as defined for one dimensional objective functions without requiring *a-priori* specification of the k parameter. Here a new graph structure is proposed and attempts are made to directly extend the idea to higher dimensions by connecting every vertex to the nearest vertex in every cartesian axis direction. In Theorem 2 we show that the one dimensional properties of this algorithm does not extend to higher dimensions which finally

leads us to the built up complexes in Sect. 5. The main conclusion of this section is that simpler graph structures cannot be used to find locally convex sub-domains of a function in the same way that was accomplished in Sect. 3.

The algorithm proceeds in the same way as TGO described in Sect. 2 except for step 2 where a new structure described in Sect. 4.1 replaces the topograph.

4.1 Axially directed topograph

Let \mathcal{F} be the set of scalar outputs mapped by the objective function $f : \mathcal{P} \rightarrow \mathcal{F}$ for a given sampling set $\mathcal{P} \subseteq \Omega \subseteq \mathbb{R}^n$. The scalar elements $f_i \in \mathcal{F}$ have one-to-one correspondence with the vector elements $\mathbf{p}^i \in \mathcal{P}$ where the integer $i \in \{1, 2, 3, \dots, N\}$ indicates the sampling point index. The vector \mathbf{p}^i in turn has components x_j^i where the integer $j \in \{1, 2, 3, \dots, n\}$ indicates the dimension of the scalar value $\forall i (x_1^i, x_2^i, x_3^i, \dots, x_n^i) \in \mathbf{p}^i$.

We wish to construct a graph that is ordered along the coordinate axes, this is done by formally defining the following related partially ordered sets.

Definition 10 Given a finite structured set of N feasible ordered sampling points $\mathcal{P} = (\mathbf{p}^1, \mathbf{p}^2, \dots, \mathbf{p}^N)$ with its corresponding objective function outputs $\mathcal{F} = (f^1, f^2, \dots, f^N)$, the index set of \mathcal{P} is given as the ordered set $\mathcal{I} = (i = \{1, 2, 3, \dots, N\}, \leq)$.

Note that the initial ordering of the index set is arbitrary. What's important is that an ordered index set is defined. This ordering will allow us to keep track of any vertex in the graph to its corresponding sampling point in \mathcal{P} so that the corresponding objective function only needs to be evaluated once. Herein the order is taken as the order that is generated by the Sobol sequence.

Definition 11 Given a set of feasible sampling points $\mathcal{P} \subseteq \Omega \subseteq \mathbb{R}^n$ define X_j for every dimension $j \in \{1, 2, 3, \dots, n\}$ as the partially ordered set $X_j = \{\mathbf{p}^i \mid \forall i (x_j^i < x_j^{i+1})\}$.

The definition is demonstrated with the following numerical example:

Example 3 Given set of the first 5 points in the 2-dimensional Sobol sequence bounded by the 2-cube:

$$\begin{aligned} \mathcal{P} &= ((0, 0), (0.5, 0.5), (0.75, 0.25), (0.25, 0.75), (0.375, 0.375)) \\ &\subseteq [0, 1] \times [0, 1] \subseteq \mathbb{R}^2 \end{aligned}$$

Let $f(x) = x_1^2 + x_2^2$ so that

$$\mathcal{F} = (0, 0.5, 0.625, 0.625, 0.28125)$$

then

$$X_1 = ((0, 0), (0.25, 0.75), (0.375, 0.375), (0.5, 0.5), (0.75, 0.25))$$

and

$$X_2 = ((0, 0), (0.75, 0.25), (0.375, 0.375), (0.5, 0.5), (0.25, 0.75))$$

The corresponding index sets are $\mathcal{I}_1 = (1, 4, 5, 2, 3)$ and $\mathcal{I}_2 = (1, 3, 5, 2, 4)$.

Definition 12 For every dimension j , \mathcal{F}_j is the partially ordered set such that the position of the elements X_j correspond to the original index sampling of \mathcal{P} , $\mathcal{F}_j = \{f_j^{i,k} \mid \forall i (x_j^i < x_j^{i+1}), f_j^{i,k} = f_k \in \mathcal{F}, k \in \mathcal{I}\}$.

That is the first superscript i of the elements $f^{i,k}$ indicate the ordering in \mathcal{F}_j , while the second superscript k indicates the corresponding scalar value of $f^{i,k}$ in \mathcal{F} . Ordering the example we have $\mathcal{F}_1 = (0, 0.625, 0.28125, 0.5, 0.625)$ and $\mathcal{F}_2 = (0, 0.625, 0.28125, 0.5, 0.625)$.

Definition 13 For every dimension j , define the partially ordered sets of cardinality N such that $\mathcal{F}_j^+ = \{f_j^{i,k} - f_j^{i-1,k} \mid \forall i(x_j^i < x_j^{i+1}), f_j^{i,k} = f_k \in \mathcal{F}, i = \{1, 2, \dots, N, k \in \mathcal{I}\}\}$ and $\mathcal{F}_j^- = \{f_j^{i,k} - f_j^{i+1,k} \mid \forall i(x_j^i < x_j^{i+1}), f_j^{i,k} = f_k \in \mathcal{F}, i = \{0, 1, \dots, N-1\}, k \in \mathcal{I}\}$.

These sets are essentially objective function differences between the sampling points along each dimensional Cartesian axis. Continuing from the numerical example we have

$$\begin{aligned}\mathcal{F}_1^+ &= (0.625, -0.34375, 0.21875, 0.125) \\ \mathcal{F}_2^+ &= (-0.625, 0.34375, -0.21875, -0.125) \\ \mathcal{F}_1^- &= (0.625, -0.34375, 0.21875, 0.125) \\ \mathcal{F}_2^- &= (-0.625, 0.34375, -0.21875, -0.125)\end{aligned}$$

We denote the elements as $f_j^{+i,k} \in \mathcal{F}_j^+$ and $f_j^{-i,k} \in \mathcal{F}_j^-$ for every dimension $j \in \{1, 2, 3, \dots, n\}$, cartesian ordering $i \subseteq \mathcal{I}$ and corresponding sampling point $k \in \mathcal{I}$. The usefulness of these abstract constructions is apparent in the following definition.

Definition 14 For a given sampling set \mathcal{P} . The minimiser pool \mathcal{M} is defined as

$$\mathcal{M} = \mathcal{M}_c \cup \mathcal{M}_{lb} \cup \mathcal{M}_{ub}$$

where

$$\begin{aligned}\mathcal{M}_c &= \{\mathbf{p}^i \mid \forall j \left((f_j^{+i} > 0) \wedge (f_j^{-(i+1)} > 0) \right), i = \{1, 2, 3, \dots, N-1\}\} \\ \mathcal{M}_{lb} &= \{\mathbf{p}^i \mid \forall j \left(f_j^{-i} < 0 \right), i = \{0\}\} \\ \mathcal{M}_{ub} &= \{\mathbf{p}^i \mid \forall j \left(f_j^{+i} < 0 \right), i = \{N\}\}\end{aligned}$$

That is, we simply check the finite difference between sampling points in every cartesian direction. In addition we check if the sampling points on the boundaries are minimisers.

Theorem 2 The minimiser pool \mathcal{M} from Definition 14 always produces a set that is either smaller than or equal to the optimum minimiser pool as defined by Definition 9 iff $j = 1$.

Proof The proof for $j = 1$ follows the same argument from Sect. 3. By Definitions 10, 11 and 12 we have the ordering constructed as \mathcal{P} and \mathcal{F}_1 . If a given point \mathbf{p}^i is a minimiser with $f_1^{+i} > 0$ and $f_1^{-i} > 0$, then we have by Definition 13 $f^i < f^{i-1}$ and $f^i < f^{i+1}$, conversely if a given point \mathbf{p}^i is not a minimiser then either $f_1^{+i} < 0$ or $f_1^{-i} > 0$ so that regardless of the sampling method used and the Euclidean distance between points a minimiser will never be generated for any point that has $((f^i > f^{i-1}) \wedge (f^i > f^{i+1})) \vee ((f^i < f^{i-1}) \wedge (f^i < f^{i+1}))$.

If $j > 1$ we have no such guarantee for a higher dimensional locally convex domain. As a counter example consider the set of points

$$\mathcal{P} = ((0, 0), (0.25, 0.25), (0.75, 0.125), (0.125, 0.75))$$

on the same function as above, the minimiser set produced is $\mathcal{M} = \{(0, 0), (0.25, 0.25)\}$ which is clearly larger than optimal and will produce the same local minimum. \square

This unsatisfactory result for higher dimensions could still potentially show good performance for more regular spaced sampling such as grids, however, as we will see in the next section the SHGO algorithm can guarantee that the optimal minimiser set will be produced for any dimension.

4.2 Implementation

Algorithm 1 provides a high-level overview of the ATGO algorithm. A Python implementation of this algorithm can be found in [11].

Algorithm 1 ATGO algorithm

```

1: procedure INITIALISATION
2:   input an objective function  $f$ , constraint functions  $\mathbf{g}$  and variable bounds  $[\mathbf{l}, \mathbf{u}]^n$ .
3:   input  $N$  initial sampling points.
4:   Define a sampling sequence that generates a set  $\mathcal{X}$  of sampling points in the unit hypercube space  $[\mathbf{0}, \mathbf{1}]^n$ 
5: end procedure
6: procedure INITIAL SAMPLING
7:    $\mathcal{P} = \emptyset$ 
8:   while  $|\mathcal{P}| < N$  do
9:     Generate  $N - |\mathcal{P}|$  sequential sampling points  $\mathcal{X} \subset \mathbb{R}^n$ 
10:    Stretch  $\mathcal{X}$  over the lower and upper bounds  $[\mathbf{l}, \mathbf{u}]^n$ 
11:     $\mathcal{P} = \{\mathcal{X}_i \mid \mathbf{g}(\mathcal{X}_i) \geq 0, \forall \mathcal{X}_i \in \mathcal{X}\} \cup \mathcal{P} \triangleright$  (Find  $\mathcal{P}$  in the feasible subset  $\Omega$  by discarding any points mapped outside the linear constraints  $\mathbf{g}$  and adding to the current set of  $\mathcal{P}$ .)
12:    Set  $\mathcal{X} = \emptyset$ 
13:   end while
14:   Find  $\mathcal{F}$  from the objective function  $f : \mathcal{P} \rightarrow \mathcal{F}$ 
15: end procedure
16: procedure CONSTRUCT  $\mathcal{M}$ 
17:   Calculate  $\mathcal{M}$  from the sets  $\mathcal{P}$  and  $\mathcal{F}$  using Definitions 11 through 14.
18: end procedure
19: procedure LOCAL MINIMISATION
20:   Calculate the approximate local minima of  $f$  using a local minimisation routine with the elements of  $\mathcal{M}$  as starting points.  $\triangleright$  These local minimisations can be performed in parallel.
21: end procedure
22: procedure PROCESS RETURN OBJECTS
23:   Order the final outputs of the minima of  $f$  found in the local minimisation step to find the approximate global minimum
24:   output  $f^*, x^*$   $\triangleright$  Here  $f^*, x^*$  refer to the best objective function output found by the algorithm.
25:   output A list of all the minima found in the local minimisation step.
26: end procedure

```

5 Simplicial homology global optimisation

5.1 Overview

The SHGO method strongly relies on constructing a simplicial complex using the sampled points of an objective function f as vertices. From this construction of the complex \mathcal{H} we use the resulting directed subgraph which contains the set of all 1-chains from the elements of $\mathcal{H}^1 \in \mathcal{H}$ to find minimiser pools using definitions similar to the methods demonstrated in the previous sections. This is accomplished by the application of Sperner's lemma [41] allowing us to approximate the domains of stationary points for any objective function in the feasible search space Ω .

We prove that, if provided with an adequate sampling set, the construction of \mathcal{H} will produce the same homology groups. We use this result to show that for the given sampling set of vertices $\mathcal{H}^0 \in \mathcal{H}$ we always extract the optimal minimiser pool similar to the one-dimensional case described in Sect. 3, but extended to higher dimensions.

The algorithm itself consists of four steps which will be described in detail:

1. Uniform sampling point generation of N vertices in the search space within the bounded and constrained subspace of Ω from which the 0-chains of \mathcal{H}^0 are constructed.
2. Construction of the directed simplicial complex \mathcal{H} by triangulation of the vertices.
3. Construction of the minimiser pool $\mathcal{M} \subset \mathcal{H}^0$ by repeated application of Sperner's lemma.
4. Local minimisation using the starting points defined in \mathcal{M} .

We will start by formally defining the construction of \mathcal{H} from a given set of feasible sampling points \mathcal{P} and proving its properties.

5.2 Directed simplicial complex approximation of the objective function

Consider again the general objective function mapping in the continuous domain $f : \mathbb{R}^n \rightarrow \mathbb{R}$. The purpose of this section is to describe a discrete mapping $h : \mathcal{P} \rightarrow \mathcal{H}$ to provide a simplicial approximation for the surface of f . To guide the reader the methods will be demonstrated on the simple 2-dimensional optimisation problem defined in Example 4. The use of a 2-dimensional surface allows for a demonstration of the techniques while the abstractions defined are readily extended to higher dimensions.

We start by formally defining the set of vertices from which 0-chains of the simplicial complex are built and the of edges from which the 1-chains of \mathcal{H} are built.

Definition 15 Let \mathcal{X} be the set of sampling points generated by a sampling sequence in the bounded hyperrectangle $[\mathbf{l}, \mathbf{u}]^n$. The set $\mathcal{P} = \{\mathbf{x} \in \mathcal{X} \mid \mathbf{g}(\mathbf{x}) \geq 0\}$ is a set of points within the feasible set Ω .

Definition 16 For an objective function f , \mathcal{F} is the set of scalar outputs mapped by the objective function $f : \mathcal{P} \rightarrow \mathcal{F}$ for a given sampling set $\mathcal{P} \subseteq \Omega \subseteq \mathbb{R}^n$.

Definition 17 Let \mathcal{H} be a directed simplicial complex. Then $\mathcal{H}^0 := \mathcal{P}$ is the set of all vertices of \mathcal{H} .

Definition 18 For a given set of vertices \mathcal{H}^0 , the simplicial complex \mathcal{H} is constructed by a triangulation connecting every vertex in \mathcal{H}^0 . The triangulation supplies a set of undirected edges E .

Definition 19 The set \mathcal{H}^1 is constructed by directing every edge in E . A vertex $v_i \in \mathcal{H}^0$ is connected to another vertex v_j by an edge contained in E . The edge is directed as $\overline{v_i v_j}$ from v_i to v_j iff $f(v_i) < f(v_j)$ so that $\partial(\overline{v_i v_j}) = v_j - v_i$. Similarly an edge is directed as $\overline{v_j v_i}$ from v_j to v_i iff $f(v_i) > f(v_j)$ so that $\partial(\overline{v_j v_i}) = v_i - v_j$.

For practical computational reasons we must also consider the case where $f(v_i) = f(v_j)$. If neither v_i or v_j is already a minimiser we will make use of rule that the incidence direction of the connecting edge is always directed towards the vertex that was generated earliest by the sampling point sequence. If v_i is not connected to another vertex v_k then we leave the notation $\overline{v_i v_k}$ undefined and let $\partial(\overline{v_i v_k}) = 0$. We let the higher dimensional simplices of \mathcal{H}^k , $k = 2, 3, \dots, n+1$ be directed in any arbitrary direction which completes the construction

of the complex $h : \mathcal{P} \rightarrow \mathcal{H}$. We can now use \mathcal{H} to find the minimiser pool for the local minimisation starting points used by the algorithm:

Definition 20 A vertex v_i is a minimiser iff every edge connected to v_i is directed away from v_i , that is $\partial(\overline{v_i v_j}) = (v_{j \neq i} - v_i) \vee 0 \forall v_{j \neq i} \in \mathcal{H}^0$. The minimiser pool \mathcal{M} is the set of all minimisers.

We will also make extensive use of star notation [15, 18]:

Definition 21 The star of a vertex v_i , written $\text{st}(v_i)$, is the set of points Q such that every simplex containing Q contains v_i .

The k -chain $C(\mathcal{H}^k)$, $k = n + 1$ of simplices in $\text{st}(v_i)$ forms a boundary cycle $\partial(C(\mathcal{H}^{n+1}))$ with $\partial(\partial(C(\mathcal{H}^{n+1}))) = \emptyset$. The faces of $\partial(\mathcal{H}^{n+1})$ are the bounds of the domain defined by $\text{st}(v_i)$.

A visual demonstration of these constructions and notations is provided in the following numerical example:

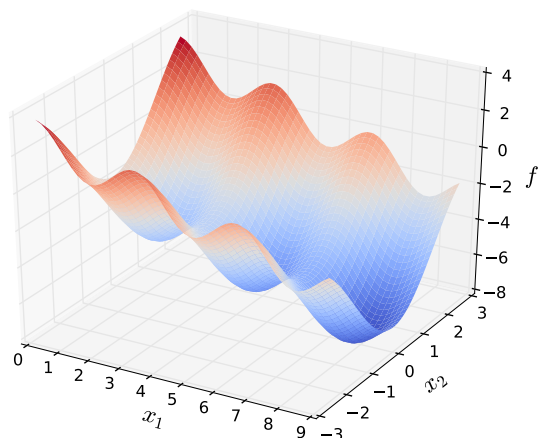
Example 4 The Ursem01 function for two dimensions is defined as follows [14]

$$\min f(\mathbf{x}) = -\sin(2x_1 - 0.5\pi) - 3\cos(x_2) - 0.5x_1, \quad x \in \Omega = [0, 9] \times [-2.5, 2.5]$$

Figure 6 provides a 3 dimensional plot of this function. The function has three local minima within the domain $\mathbf{x} \in [0, 9] \times [-2.5, 2.5]$.

We use a set \mathcal{P} of $N = 15$ sampling points from the 2-dimensional Sobol sequence. First map out the objective function values:

Fig. 6 A 3-dimensional surface plot of the optimisation test function given in Example 4 $f(\mathbf{x}) = -\sin(2x_1 - 0.5\pi) - 3\cos(x_2) - 0.5x_1$ for the domain $\mathbf{x} \in \Omega = [0, 9] \times [-2.5, 2.5]$



$$f : \begin{bmatrix} v_0 = (0.0, -2.5) \\ v_1 = (4.6, 0.0) \\ v_2 = (6.9, -1.25) \\ v_3 = (2.3, 1.25) \\ v_4 = (3.45, -0.625) \\ v_5 = (8.05, 1.875) \\ v_6 = (5.75, -1.875) \\ v_7 = (1.15, 0.625) \\ v_8 = (1.725, -0.9375) \\ v_9 = (6.325, 1.5625) \\ v_{10} = (8.625, -2.1875) \\ v_{11} = (4.025, 0.3125) \\ v_{12} = (2.875, -1.5625) \\ v_{13} = (7.475, 0.9375) \\ v_{14} = (5.175, -0.3125) \end{bmatrix} \rightarrow \begin{bmatrix} f_0 = 3.403 \\ f_1 = -6.275 \\ f_2 = -4.0651 \\ f_3 = -2.208 \\ f_4 = -3.3429 \\ f_5 = -4.051 \\ f_6 = -1.493 \\ f_7 = -3.674 \\ f_8 = -3.591 \\ f_9 = -2.191 \\ f_{10} = -2.606 \\ f_{11} = -5.062 \\ f_{12} = -0.601 \\ f_{13} = -6.239 \\ f_{14} = -6.044 \end{bmatrix} \quad (13)$$

From Definition 17 we find \mathcal{H}^0 from \mathcal{P} . Next we use Delaunay triangulation to find a set of connected edges according to Definition 18. Any triangulation scheme resulting in a simplicial complex can be used. Next the edges are directed from the calculated values of \mathcal{F} using Definition 19. Finally from Definition 20 we find the minimiser set $\mathcal{M} = \{v_1, v_7, v_{13}\}$. The resulting structure is shown in Fig. 7. Also shown in Fig. 7 is the domain of $\text{st}(v_1)$ for a visual description of Definition 21. Next we increase the sampling size to $N = 150$ points and repeat the procedure. The resulting complex is shown in Fig. 8. Notice that while the minimiser vertices have changed (now closer to the true continuous local minima), the cardinality of the minimiser pool $|\mathcal{M}|$ remains unchanged. That is, given an adequate

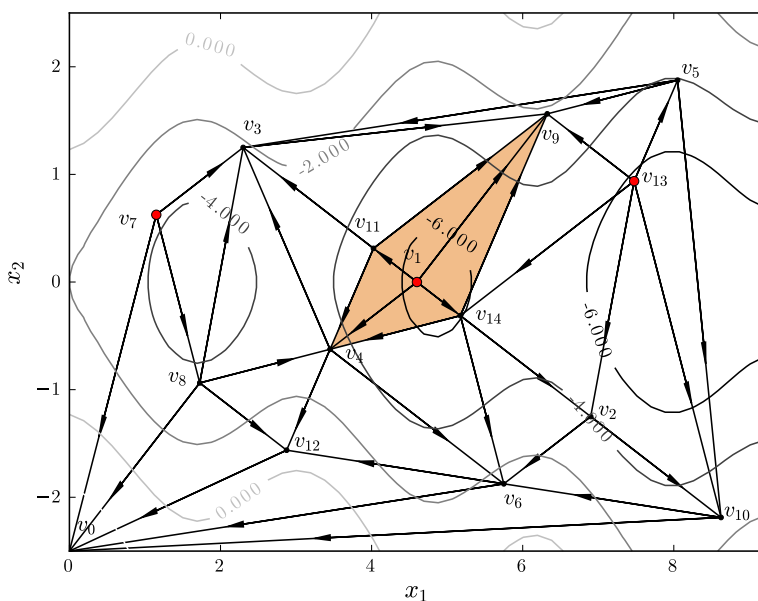


Fig. 7 A directed complex \mathcal{H} with $N = 15$ forming a simplicial approximation for an objective function. There are three minimiser vertices v_1 , v_7 and v_{13} shown by the big red dots. The shaded star domain represents the domain defined by $\text{st}(v_1)$. (Color figure online)

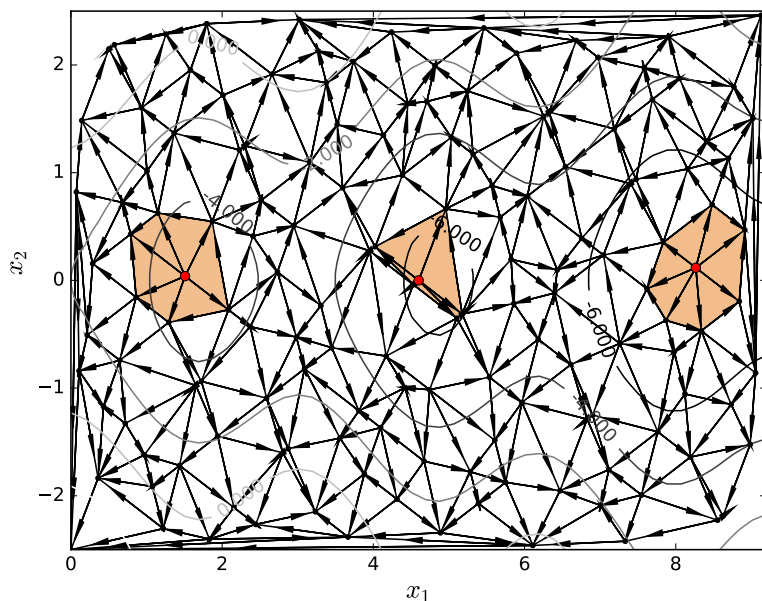


Fig. 8 A directed complex \mathcal{H} forming a simplicial approximation for an objective function with 150 vertices. There are three minimiser vertices given by the big red dots. The refined star domains are shaded in the figure. (Color figure online)

number sampling points $|\mathcal{M}|$ will cease to grow with increasing N , providing a heuristic for the number of sampling points needed to approximately map all minima of an objective function. This useful property of the SHGO algorithm is proven formally in Sect. 5.4.

5.3 Guarantee of stationary points in sub-domains near minimiser points

This section is devoted to proving the following theorem:

Theorem 3 *Given a minimiser $v_i \in \mathcal{M} \subseteq \mathcal{H}^0$ on the surface of a continuous, Lipschitz smooth objective function f with a compact bounded domain in \mathbb{R}^n and range \mathbb{R} , there exists at least one stationary point of f within the domain defined by $st(v_i)$.*

Proof Our strategy relies on finding a simplex with a Sperner labelling where each label represents a different $n + 1$ label in every vector direction of the gradient vector field ∇f of f where of the $n + 1$ Cartesian directions we require only a vector pointing towards a section defined by $n + 1$ hyperplane cuts, the remainder of the proof then proceeds as usual for Brouwer's fixed point theorem [6] found in for example Henle [18, p. 40] utilising Sperner's lemma. \square

Theorem 4 (Sperner's lemma [41]) *Every Sperner labelling of a triangulation of a n -dimensional simplex contains a cell labelled with a complete set of labels: $1, 2, \dots, n+1$.*

Start with the observation that for any minimiser $v_i \in \mathcal{M} \subseteq \mathcal{H}^0$ we have by construction that for any vertex v_j with incidence on a connecting edge $\overline{v_i v_j}$ that $f(v_i) < f(v_j)$, so by the MVT there is at least one point on $\overline{v_i v_j}$ where ∇f points towards a Cartesian direction in a section that can receive a unique Sperner label. If we have $n + 1$ vertices with incidence

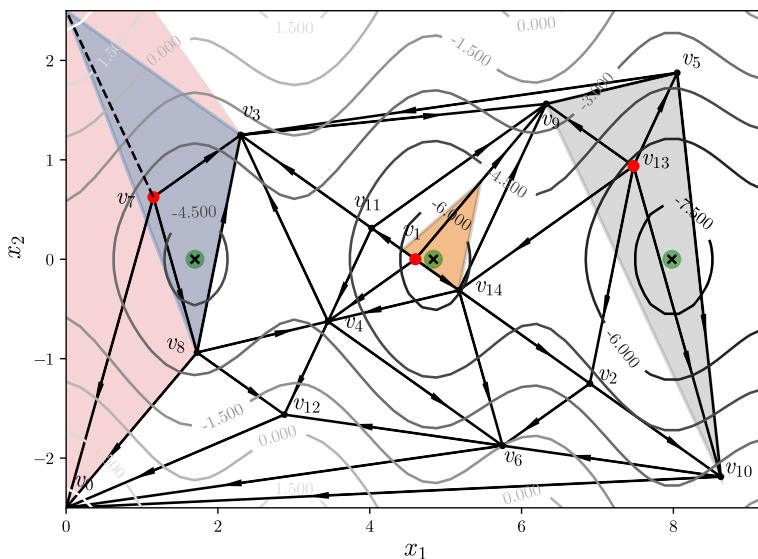


Fig. 9 Visual demonstration of the proof by finding simplices with Sperner labellings. The three circled crosses are the (approximate) minimima of the objective function within the given bounds. The three possible Sperner simplices are contained within the star domains of each minimiser $\text{st}(v_1)$, $\text{st}(v_7)$ and $\text{st}(v_{13})$. v_7 is an example of simplices without complete Sperner labellings, the red shaded area around $\overline{v_7}$ is the bounded domain wherein at least one local minimum exist. (Color figure online)

on an edge $\overline{v_i v_j} \subseteq \mathcal{H}^1$ in every required Cartesian direction then we have a simplex within $\text{st}(v_i)$ with a Sperner labelling.

In the case where we do not have $n + 1$ vertices in every required section then by construction there is no vertex between v_i and the boundary of f defined by Ω in the required section. In the case where the constraint is not active and there exists at least one point v_k boundary where ∇f does not point towards the boundary and by the MVT v_k can receive a unique Sperner label from which we can construct a simplex within $\text{st}(v_i)$ with Sperner labelling.

Following the combinatorial version of Brouwer's fixed point theorem [18] since ∇f is continuous and the domain $\text{st}(v_i)$ is compact we can produce a sequence of complete triangulations with arbitrarily small size in which the size of the simplices decreases toward zero. This sequence produces a sequence of vertices with gradients $\nabla f(V)$ pointing in every $n + 1$ direction. By continuity there is a vector $\nabla f(\mathbf{X})$ near the sequences, since the zero vector is the only vector pointing in all $n + 1$ directions we have a point \mathbf{X} bounded by the domain defined by $\text{st}(v_i)$ where $\nabla f(\mathbf{X}) = \vec{0}$. In the case where the constraint is active a local minimum lies on the constraint which is in the domain defined $\text{st}(v_i)$. This concludes the proof.

Figure 9 provides a visual demonstration of the proof using the complex from Example 4. Here we have divided the plane so that the 3 required directions are $[0, \frac{\pi}{2})$, $[\frac{\pi}{2}, \pi)$ and $[\pi, 2\pi)$. Note that this division is arbitrary and any $n + 1 = 3$ subdivisions can be chosen as long as all possible $n + 1 = 3$ directions can form a simplex in the space are covered. The three possible simplices are contained within the star domains of each minimiser $\text{st}(v_1)$, $\text{st}(v_7)$ and $\text{st}(v_{13})$.

First consider the minimiser v_{13} . There are three possible edges in $[\frac{\pi}{2}, \pi)$ on which a point exists that can be used as a vertex to receive a Sperner labelling for that direction namely

$\overline{v_{13}v_{14}}$, $\overline{v_{13}v_2}$ and $\overline{v_{13}v_{10}}$. The only possible edges in the $[0, \frac{\pi}{2}]$, $[\frac{\pi}{2}, \pi)$ directions are $\overline{v_{13}v_5}$ and $\overline{v_{13}v_9}$ respectively. The simplex $\overline{v_5v_9v_{10}}$ drawn in Fig. 9 is not necessarily the simplex with a Sperner labelling. The three vertices of the Sperner simplex which are proven to exist through the MVT exists on each of the edges $\overline{v_{13}v_{14}}$, $\overline{v_{13}v_2}$ and $\overline{v_{13}v_{10}}$ in a subdomain of this simplex $\overline{v_5v_9v_{10}}$. For example the simplex surrounding the minimiser v_1 is a possible Sperner simplex with vertices on the edges in every required direction.

Note that if the edge $\overline{v_{13}v_{14}}$ was chosen instead of $\overline{v_{13}v_{10}}$ then the local minimum of the function would be outside the domain of the simplex with the Sperner labelling. This is an important observation because it demonstrates that Theorem 3 cannot be used to further refine the location of the local minimum from the domain $\text{st}(v_{13})$ using mechanisms of the proof, it only states that at least one local minimum exists within $\text{st}(v_{13})$.

The boundaries of $\text{st}(v_{13})$ can be found using the 3-chain $C_{13}(\mathcal{H}^3)$ of simplices in $\text{st}(v_{13})$, recall that the directions of simplices higher than dimension 2 are undefined and so the directions can be arbitrarily chosen

$$C_{13}(\mathcal{H}^3) = \overline{v_{13}v_{10}v_5} + \overline{v_{13}v_5v_9} + \overline{v_{13}v_9v_{14}} + \overline{v_{13}v_{14}v_2} + \overline{v_{13}v_2v_{10}}$$

$C_{13}(\mathcal{H}^3)$ clearly forms a cycle, applying the boundary operator we find the faces defining the bounds of the domain of $\text{st}(v_i)$ which in this case is the chain of edges with defined direction

$$\partial(C_{13}(\mathcal{H}^3)) = -\overline{v_{10}v_5} + \overline{v_5v_9} - \overline{v_9v_{14}} + \overline{v_{14}v_2} + \overline{v_2v_{10}}$$

thus $\partial(\partial(C(\mathcal{H}^3))) = \emptyset$.

$v_7 = (1.15, 0.625)$ is an example of a minimiser that does not have all three required directions for a Sperner labelling, the light red shaded area represents the area wherein a local minimum can exist. For example on the lines $x_1 = 0$ for $x_2 \in [0.625, 2.5]$ or $x_2 = 2.5$ for $x_1 \in [0, 1.15]$ there will either exist a point \mathbf{p} where the gradient $\nabla f(\mathbf{p})$ points in any direction pointing towards $[\frac{3}{2}\pi, 0)$ in which case and edge $\overline{v_{13}\mathbf{p}}$ exists that points in the $[\frac{\pi}{2}, \pi)$ direction and we have a simplex with a Sperner labelling. For example the dotted line on Fig. 9 with the Sperner simplex represented by blue shaded around v_7 . If such a point does not exist then all points on those lines points in the $[0, \frac{3}{2}\pi)$ direction and so one or more local minimum lies somewhere on the boundary which is within the defined area.

There have been several developments in the extension of this lemma which could prove useful in applications extending the SHGO algorithm. One of the most interesting is by De Loera et al. [8] where they proved the Atanassov conjecture [3] that for any polytope with N vertices there are $N - n$ simplices that receive a complete set of Sperner labels. Meunier [28] further extended this theorem and more recently Musin [31] extended the theorems to a large class of manifolds with or without boundary. These theorems could prove useful for extending the algorithm to make use of this information. More explicitly, SHGO currently uses knowledge of the objective function evaluations, but only in a Boolean sense (in the form of directed edges). The theorems by Meunier and Musin allow us to extend Sperner's lemma to a simplicial complex built in a $(n + 1)$ -dimensional non-Euclidean space. This would allow the application of ideas from discrete differential geometry. For example the Gauss-Bonnet theorem holds for discrete simplicial surfaces [7]. The Gauss-Bonnet theorem provides a relation between the total Gaussian curvature and the Euler characteristic of a surface. By simple summation of the angle defect around every vertex we can determine the Euler characteristic of a continuous surface. As will be demonstrated in Sect. 5.4 the simplicial complex used by SHGO is homeomorphic to complexes built on other topological hypersurfaces. Therefore when using explicit coordinates of the expected homomorphism

the summation can be used to compare the error with the Euler characteristic which provides a metric for how accurately the objective function surface has been sampled. In global optimisation theory a simplicial complex built in this space can be used for approximating local and global Lipschitz constants for an objective function while still retaining the ability to detect locally convex sub-domains in the search space.

5.4 Invariance of the directed complex within a bounded rectangle

We now have a guarantee of finding stationary points in sub-domains near stationary points. However, we would also like to ensure that SHGO does not generate more than one minimiser starting point per convex sub-domain. This can only be guaranteed when an objective function surface is “adequately sampled”. For black box functions there is no way to know if the number and distribution of sampling points is adequate without more information (for example if the number of local minima are known in the problem). However, it is an important property of the algorithm that $|\mathcal{M}|$ will stop increasing with higher sampling after this point. First we define an adequately sampled surface.

Definition 22 Consider a simplicial complex \mathcal{H} built on an objective function f with a compact feasible set Ω using Definitions 17 through 20. The surface is said to be *adequately sampled* if there is one and only one true stationary point within every domain defined by Theorem 3.

The remainder of this section is devoted to proving the following theorem which holds in the case where $\Omega = [\mathbf{l}, \mathbf{u}]^n$.

Theorem 5 (Invariance of an adequately sampled simplicial complex \mathcal{H}) *For a given continuous objective function f that is adequately sampled by a sampling set of size N . If the cardinality of the minimiser pool extracted from the directed simplex \mathcal{H} is $|\mathcal{M}|$. Then any further increase of the sampling set N will not increase $|\mathcal{M}|$.*

Proof The proof relies on a homomorphism between the simplicial complex \mathcal{H} constructed in the bounded hyperrectangle Ω and the homology (mod 2) groups of a constructed surface S on which we can invoke the invariance theorem.

Define the n -torus S_0 from the compact, bounded hyperrectangle Ω by identification of the opposite faces and all extreme vertices. Now for every strict local minimum point $\mathbf{p} \in \Omega$ puncture a hypersphere and after appropriate identification the resulting n -dimensional manifold S_g is a connected g sum of g tori $S := S_0 \# S_1 \# \dots \# S_{g-1}$ (g times).

Any triangulation \mathcal{K} of the topological space S is homeomorphic to S , $\mathbf{H}_k(\mathcal{K}) \cong \mathbf{H}_k(S) \forall k \in \mathbb{Z}$. Note that this homomorphism is for a mod 2 homology between a triangulation \mathcal{K} and the surface S and is thus undirected. A triangulation corresponding to all vertices and faces of \mathcal{K} can be directed according to Definitions 17, 18 and 19 providing the directed simplicial complex \mathcal{H} . By construction we have, for an adequately sampled simplicial complex \mathcal{H} , an equality which exists between the cardinality of \mathcal{M} and the Betti numbers of S as $|\mathcal{M}| = h_1 = \text{rank}(\mathbf{H}_1(S)) = \text{rank}(\mathbf{H}_1(\mathcal{K}))$. Here we invoke the invariance theorem. \square

Theorem 6 (Invariance theorem[18]) *The homology groups associated with a triangulation \mathcal{K} of the a compact, connected surface S are independent of \mathcal{K} . In other words, the groups $\mathbf{H}_0(\mathcal{K})$, $\mathbf{H}_1(\mathcal{K})$ and $\mathbf{H}_2(\mathcal{K})$ do not depend on the simplices, incidence coefficients, or anything else arising from the choice of the particular triangulation \mathcal{K} ; they depend only on the surface S itself.*

The invariance theorem can be extended to higher dimensional triangulable spaces using singular homology through the Eilenberg–Steenrod Axioms [10, 18]. As a direct consequence any triangulation of \mathcal{S} will produce the same homology groups for \mathcal{K} .

Adding any new sampling point within the corresponding subdomains of $\text{st}(v_i) \forall i (v_i \in \mathcal{M} \subseteq \mathcal{H}^0)$ as defined in Theorem 3 will by Definitions 17 through 20 need to be connected directly to v_i by a new edge or the triangulation is no longer a simplicial complex and thus not increase $|\mathcal{M}|$ since only one vertex will be the new minimiser.

After adding any sampling point outside a domain $\text{st}(v_i)$ then, through the established homomorphism, any construction of \mathcal{H} will produce the same homology groups since $\text{rank}(\mathbf{H}_1(\mathcal{K}))$ remains unchanged and it is thus not possible for a new vertex to be wrongly identified as a minimiser in the triangulation \mathcal{H} .

This concludes the proof that any increase in N will not further increase $|\mathcal{M}|$.

It is important to note that Theorem 5 is only applicable to complexes with adequate sampling as defined, that is to say it is entirely possible that, in complexes with less than adequate sampling, two starting minimiser elements of \mathcal{M} will converge to the same local minimum. This flaw is inherent in the fact that there is insufficient information to completely identify the minima of a surface (and could be overcome if some extra information about f is known).

Theorems 3 and 5 also lead to the following corollary about an optimisation problem:

Corollary 2 *Consider any objective function $f : \Omega \subseteq \mathbb{R}^n \rightarrow \mathbb{R}$. Consider also a local minimisation routine that is guaranteed to converge to a local minimum in the same locally convex domain as the starting point inputted to the algorithm. Alternatively the local minimisation routine is guaranteed to converge to a point within a set of bounds (provided by the boundary of the k -chain around $\text{st}(v_i)$, $\partial(C(\mathcal{H}^k))$, $k = n + 1$). If such a local minimisation routine uses an element $v_i \in \mathcal{M}$ as a starting point and the routine leads to a minimum outside or on $\text{st}(v_i)$ and in addition the minimum is not contained in the set \mathcal{H}^0 . Then it can be concluded that either search space is not adequately sampled or f is not a Lipschitz smooth function.*

Therefore according to Corollary 2 if the number of local minima are known, as in for example phase equilibria problems, then we can extract valuable information about the objective function. In particular it can be determined whether or not the objective function is Lipschitz smooth. Alternatively if the function is known to be Lipschitz smooth then Corollary 2 can be used to prove the sampling is insufficient when the condition is not met. When this happens it is also now known that there are more local minima to be found, one or more of which might possibly be the global minimum. Corollary 2 does not, however, provide any guarantee that the sampling is sufficient when the conditions are met.

5.5 Sampling generation

Using the Sobol sequence sampling point generation proceeds in a similar way as that described in Sect. 2.1. However, rather than only generating an arbitrary number of pre-defined sampling points we will also consider heuristic methods starting with the minimum amount of sampling points required to triangulate an n dimensional space. For example start with the minimum amount of sampling points to construct an n -dimensional simplex and continue sampling while continuously calculating the $\mathbf{H}_1(\mathcal{H})$ homology groups of the com-

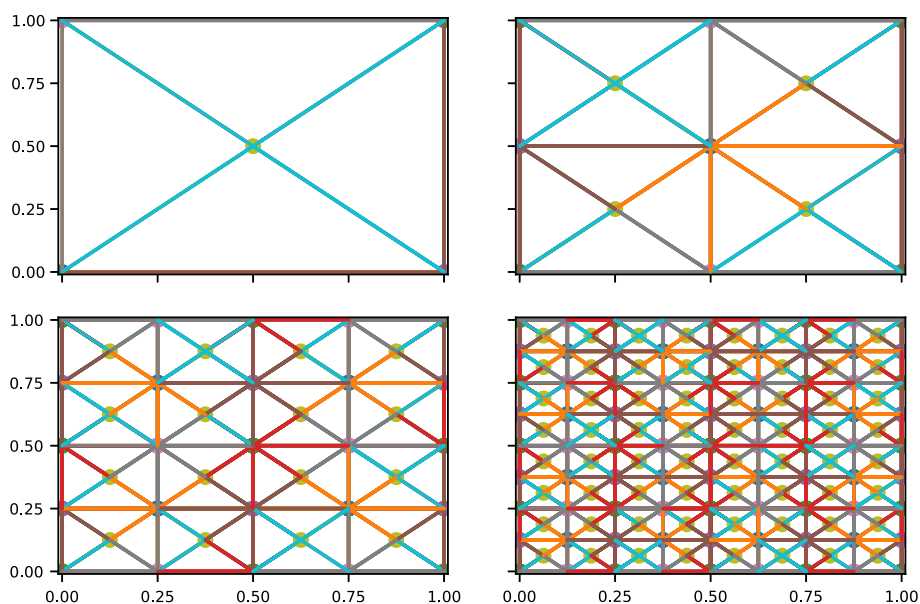


Fig. 10 Triangulation of a unit hypercube shown in 2 dimensions for 4 iterations

plex. Using the definitions described in this section the sampling is continued until the growth rate of the approximated homology groups slows appreciably.

In this publication the Sobol sequenced sampling points are triangulated using Delaunay triangulation as implemented in the SciPy library [22]. A major disadvantage to this triangulation scheme is that it does not scale well to higher dimensions since it relies on solving convex hull using the quickhull method developed by Barber and Dobkin [4]. There are several possibilities for mitigating this problem. Since the Sobol sequence is deterministic the triangulations can be calculated and stored in a database. For SHGO another possibility whereby the convex hull does not need to be solved by using symmetry generated triangulation was developed. Building on the initial n -cube triangulation developed by Paulavičius and Žilinskas [34, 53] and using the symmetry groups S_n , $n = \{1, 2, 3, \dots, n\}$ to generate an initial triangulation. Subsequent uniform sampling that ensures a symmetrical triangulation is generated in the next generation of simplices. This is done by an ordering of edges and using the cycle $(123 \dots n - 1)$ to ensure that we always split every simplex by a hyperplane that goes through a child vertex on the longest edge of simplex and every other vertex in the parent simplex that does not have incidence on the edge. Figure 10 demonstrates the symmetry of this sampling in $n = 2$ where the longest edge in the initial triangulation was sampled. Here an iteration is defined as any generation of sub-triangulations that provides a triangulation symmetrical to the initial triangulation. An implementation of this sampling sequence is available at [11].

In this publication we will use both the Sobol and the hypercube triangulation sampling sequences. Sobol provides a more direct comparison to the TGO algorithm while the second sequence is more similar to the DISIMPL-v algorithm. We will refer to the different uses of sampling sequences as SHGO-Sobol and SHGO-Simpl in the experimental results section in Sect. 6.

5.6 Theoretical comparison to the DISIMPL algorithm

The DISIMPL algorithm developed by Paulavičius and Žilinskas [33,34,36] is based on spatial partitioning of the search space. DISIMPL-v in particular should have a similar initial complex as SHGO-Simpl for box problems since this algorithm samples on the vertices of the simplicial complex (while DISIMPL-c samples at the geometric centre of the simplices which is more appropriate for higher dimensional problems). The graph structure of DISIMPL-v can thus be used to construct the directed complex \mathcal{H} and the homological properties can be calculated and applied. An example of one such application is given in the following paragraph.

At every iteration of the DISIMPL algorithm potentially optimal simplices are selected for refinement by considerations the Lipschitz properties of the optimisation problem. In general a combination of promising simplices with good function evaluations (related to local exploration of the search space) and simplices with larger hypervolumes (related to global exploration of the search space). Gb-DISIMPL [36] is a very promising acceleration technique accomplished by switching between a “global phase” and a “usual phase”. The global phase is focused on exploring simplices with larger hyper volumes and excludes smaller simplices which are potentially optimal in the usual phase. This technique prevents excessive evaluations near local minima as demonstrated in [36]. Local minima can put a “drag” on the progress of refining the minimum because the algorithm selects many neighbouring simplices that are slightly worse on the function values, but also slightly larger in volume. A meta-parameter is used in Gb-DISIMPL to select the simplices to be excluded in the global phase and was shown in [36] to be very efficient. However, using knowledge from the directed complex of \mathcal{H} , the domain containing these simplices near the local minima could also be identified more explicitly through a Sperner labelling if the function is known to be Lipschitz smooth.

5.7 Algorithm implementation

We consider two modes for the SHGO algorithm. In the first a finite number of sampling points N are specified and sampling is continued until an Ω set of cardinality N is produced and no further sampling occurs. This method is demonstrated by Algorithm 2. The main reason for this algorithm is to present a more direct comparison to TGO that can be used in numerical experiments.

For the purposes of global optimisation and local minima exploration Algorithm 3 is more appropriate. By continuously calculating the $\mathbf{H}_1(\mathcal{H})$ homology group several termination criteria can be used to end the sampling. For example if the amount of local minima is known the sampling can be terminated once $|\mathcal{M}|$ is large enough. Another example with many possible heuristics is tracking the historical difference in $|\mathcal{M}|$ over $|\mathcal{P}|$ and terminating sampling if $|\mathcal{M}|$ is unchanged after a certain increase in $|\mathcal{P}|$. In optimisation problems where the global minimum is known we can also use the stopping criteria such as the one defined by [35].

$$pe = 100\% \times \begin{cases} \frac{\min\{\mathcal{F}\} - f^*}{|f^*|}, & f^* \neq 0 \\ \min\{\mathcal{F}\}, & f^* = 0 \end{cases}$$

Here $\min\{\mathcal{F}\}$ is the minimum function evaluation obtained including values obtained in the output of the local minimisation step as shown in the algorithm. Whatever termination criterion is used it requires an input $\mathbf{H}_1(\mathcal{H})$ or $\min\{\mathcal{F}\}$ and should output a Boolean, we will

refer to this function as $\mathbf{TERM}(\mathbf{H}_1(\mathcal{H}), \min\{\mathcal{F}\})$ in Algorithm 3. In the practical implementation of the algorithm the user can also specify a finite number of iterations and/or sampling points. This functionality has been programmed into the $\mathbf{TERM}(\mathbf{H}_1(\mathcal{H}), \min\{\mathcal{F}\})$ function.

Open source Python implementations of both of these algorithms are available and were published under a MIT compatible license [11].

Algorithm 2 SHGO finite sampling algorithm

```

1: procedure INITIALISATION
2:   Input an objective function  $f$ , constraint functions  $\mathbf{g}$  and variable bounds and  $[\mathbf{l}, \mathbf{u}]^n$ .
3:   Input  $N$  initial sampling points.
4:   Define a sampling sequence that generates a set  $\mathcal{X}$  of sampling points in the unit hypercube space  $[\mathbf{0}, \mathbf{1}]^n$ 
5: end procedure
6: procedure INITIAL SAMPLING
7:    $\mathcal{P} = \emptyset$ 
8:   while  $|\mathcal{P}| < N$  do
9:     Generate  $N - |\mathcal{P}|$  sequential sampling points  $\mathcal{X} \subset \mathbb{R}^n$ 
10:    Stretch  $\mathcal{X}$  over the lower and upper bounds  $[\mathbf{l}, \mathbf{u}]^n$ 
11:     $\mathcal{P} = \{\mathcal{X}_i \mid \mathbf{g}(\mathcal{X}_i) \geq 0, \forall \mathcal{X}_i \in \mathcal{X}\} \cup \mathcal{P} \triangleright$  (Find  $\mathcal{P}$  in the feasible subset  $\Omega$  by discarding any points mapped outside the linear constraints  $\mathbf{g}$  and adding to the current set of  $\mathcal{P}$ .)
12:    Set  $\mathcal{X} = \emptyset$ 
13:   end while
14:   Find  $\mathcal{F}$  from the objective function  $f : \mathcal{P} \rightarrow \mathcal{F}$ 
15: end procedure
16: procedure CONSTRUCT DIRECTED COMPLEX  $\mathcal{H}$ 
17:   Calculate  $\mathcal{H}$  from  $h : \mathcal{P} \rightarrow \mathcal{H}$ 
18: end procedure
19: procedure CONSTRUCT  $\mathcal{M}$ 
20:   Find  $\mathcal{M}$  from Definition 20.
21: end procedure
22: procedure LOCAL MINIMISATION
23:   Calculate the approximate local minima of  $f$  using a local minimisation routine with the elements of  $\mathcal{M}$  as starting points.  $\triangleright$  These local minimisations can be performed in parallel.
24: end procedure
25: procedure PROCESS RETURN OBJECTS
26:   Order the final outputs of the minima of  $f$  found in the local minimisation step to find the approximate global minimum
27:   output  $f^*, x^*$   $\triangleright$  Here  $f^*, x^*$  refer to the best objective function output found by the algorithm.
28:   output A list of all the minima found in the local minimisation step.
29: end procedure

```

6 Experimental results

6.1 Comparison to other solvers on problems with linear constraints

In this section we provide experimental comparisons on 22 linearly constrained problems comparing the SHGO, TGO, Lc-DISIMPL [35], LGO [37], PSwarm [47] and DIRECT-L1 [13] algorithms. Note that the data for the Lc-DISIMPL, PSwarm and DIRECT-L1 algorithms was taken from [35]. In [35] Paulavičius and Žilinskas used DIRECT-L1 with the 3 different penalty parameters (p.p.). Due to spatial considerations only the best performing penalty parameters for DIRECT-L1 are shown in the table in this study. The same percentage error of $pe = 0.01\%$ used by [35] was also used in this publication. To provide a fair comparison of TGO to SHGO and the other solvers the TGO algorithm was modified to stop sampling when it produced a minimiser that lead to the global minimum of the problem. Table 1 shows the

Algorithm 3 SHGO homology group growth algorithm

```

1: procedure INITIALISATION
2:   Input an objective function  $f$ , constraint functions  $\mathbf{g}$  and variable bounds and  $[\mathbf{l}, \mathbf{u}]^n$ .
3:   Input  $N$  initial sampling points.
4:   Define a sampling sequence that generates a set  $\mathcal{X}$  of sampling points in the unit hypercube space  $[0, 1]^n$ 
5:   Define the empty set  $\mathcal{M}^E = \emptyset$  of vertices evaluated by a local minimisation.
6: end procedure
7: while  $\text{TERM}(\mathbf{H}_1(\mathcal{H}), \min\{\mathcal{F}\})$  is False do
8:   procedure SAMPLING
9:      $\mathcal{P} = \emptyset$ 
10:    while  $|\mathcal{P}| < N$  do
11:      Generate  $N - |\mathcal{P}|$  sequential sampling points  $\mathcal{X} \subset \mathbb{R}^n$ 
12:      Stretch  $\mathcal{X}$  over the lower and upper bounds  $[\mathbf{l}, \mathbf{u}]^n$ 
13:       $\mathcal{P} = \{\mathcal{X}_i \mid \mathbf{g}(\mathcal{X}_i) \geq 0, \forall \mathcal{X}_i \in \mathcal{X}\} \cup \mathcal{P}$   $\triangleright$  (Find  $\mathcal{P}$  in the feasible subset  $\Omega$  by discarding any
        points mapped outside the linear constraints  $\mathbf{g}$  and adding to the current set of  $\mathcal{P}$ .)
14:      Set  $\mathcal{X} = \emptyset$ 
15:    end while
16:    Find  $\mathcal{F}$  from the objective function  $f : \mathcal{P} \rightarrow \mathcal{F}$  for any new points in  $\mathcal{P}$ 
17:  end procedure
18:  procedure CONSTRUCT/APPEND DIRECTED COMPLEX  $\mathcal{H}$ 
19:    Calculate  $\mathcal{H}$  from  $h : \mathcal{P} \rightarrow \mathcal{H}$   $\triangleright$  (If  $\mathcal{H}$  was already constructed new points in  $\mathcal{P}$  are incorporated
    into the triangulation.)
20:    Calculate  $\mathbf{H}_1(\mathcal{H})$ 
21:  end procedure
22:  procedure CONSTRUCT  $\mathcal{M}$ 
23:    Find  $\mathcal{M}$  from Definition 20.
24:  end procedure
25:  procedure LOCAL MINIMISATION
26:    Calculate the approximate local minima of  $f$  using a local minimisation routine with the elements
    of  $\mathcal{M} \setminus \mathcal{M}^E$  as starting points.  $\triangleright$  Process the most promising points first.
27:     $\mathcal{M}^E = \mathcal{M}^E \cup \mathcal{M}$   $\triangleright$  This excludes the evaluation any element  $v_i \in \mathcal{M}$  that is known to be
    the only point that in the domain  $\partial\text{st}(v_j)$  where  $v_j$  is known to any point already used as a starting point
    in Step 27. If any new  $v_i \in \mathcal{M}$  not in  $\mathcal{M}^E$  is known to be the only point  $\partial\text{st}(v_j)$  it can also be excluded.
28:    Add the function outputs of the local minimisation routine to  $\mathcal{F}$ 
29:  end procedure
30:  Find new value of  $\text{TERM}(\mathbf{H}_1)(\mathcal{H}, \min\{\mathcal{F}\})$ 
31: end while
32: procedure PROCESS RETURN OBJECTS
33:  Order the final outputs of the minima of  $f$  found in the local minimisation step to find the approximate
  global minimum
34:  output  $f^*, x^*$   $\triangleright$  Here  $f^*, x^*$  refer to the best objective function output found by the algorithm.
35:  output A list of all the minima found in the local minimisation step.
36: end procedure

```

results. Here f.e. is the total number of objective function evaluations required to solve the function and p.f.e. is the total number of penalty function evaluations. The PSwarm solver was run 10 times for each test problem.

Note that SHGO-Simpl, SHGO-Sobol, and TGO are hybrid algorithms in the sense that they combine a sampling step with a local minimisation procedure. The total number of sampling evaluations performed in the local minimisation routines are also provided in Table 1. The number of global sampling points on each problem is the total number given in Table 1 minus the local function evaluations given in the brackets in the superscript c . Although the selected local minimisation routine (SLSQP) does not use derivative information, on certain problems another local minimisation routine might be more appropriate, which can change the performance of the algorithm.

Table 1 Function evaluation comparisons for test problems with linear constraints. The results for the Lc-DSIMPL, PSwarm and DIRECT-L1 algorithms were taken from [35]

Problem	shgo-		-sobol	lgo		Lc-DSIMPL ^d		LGO		PSwarm ^d		Average		Maximum		DIRECT-L1 ^d	
	-simpl	f.e.		f.e.	-v	-c	f.e.	f.e.	Minimum		p.f.e.		f.e.	p.f.e.	p.p. = 10	f.e	
			f.e.						p.f.e.	f.e.	p.f.e.						
horst-1	97 ^c (89)	24 ^c (12)	24 ^c (12)	34 ^c (20)	7	249	2457	167	182	1329 ^b (3)	1343 ^b (3)	4100 ^b (3)	4101 ^b (3)	287 ^a			
horst-2	10 ^c (8)	11 ^c (8)	11 ^c (8)	11 ^c (8)	5	171	1645	160	176	424	492	768	867	265 ^a			
horst-3	6 ^c (4)	7 ^c (4)	7 ^c (4)	6 ^c (4)	5	249	3649	42	43	44	45	46	47	5 ^a			
horst-4	10 ^c (5)	25 ^c (15)	25 ^c (15)	24 ^c (20)	8	260	39	90	179	114	194	129	211	58293 ^a			
horst-5	20 ^c (10)	15 ^c (10)	15 ^c (10)	15 ^c (10)	8	259	37	106	150	134	192	214	302	7 ^a			
horst-6	22 ^c (20)	59 ^c (58)	59 ^c (58)	77 ^c (75)	10	28	8476	90	172	110	192	133	227	11 ^a			
horst-7	10 ^c (5)	15 ^c (10)	15 ^c (10)	13 ^c (10)	10	220	5217	188	201	380	403	919	957	7 ^a			
hs021	24 ^c (20)	23 ^c (20)	23 ^c (20)	23 ^c (20)	189	133	13	110	110	189	192	392	405	97			
hs024	24 ^c (20)	15 ^c (14)	15 ^c (14)	36 ^c (34)	3	141	1809	101	153	118	172	138	195	19 ^a			
hs035	37 ^c (36)	41 ^c (36)	41 ^c (36)	35 ^c (31)	630	721	1885	266	311	316	369	327	373	>100000			
hs036	105 ^c (100)	20 ^c (100)	20 ^c (100)	103 ^c (100)	8	314	2756	179	179	396	401	561	574	25 ^a 0			
hs037	72 ^c (31)	63 ^c (62)	258 ^c (236)	186	9129	10516	127	131	160	167	201	574	7 ^a				
hs038	225 ^c (128)	1029 ^c (570)	389 ^c (347)	3379	>100000	221	53662	54445	58576	59821	65677	67660	7401				
hs044	119 ^c (35)	35 ^c (34)	51 ^c (40)	20	440	32464	148 ^b (9)	218 ^b (9)	186 ^b (9)	281 ^b (9)	299 ^b (9)	90283	19135				
hs076	56 ^c (54)	37 ^c (36)	44 ^c (42)	548	4794	221	132	198	203	286	275	341	158	7 ^a			
s224	166 ^c (162)	165 ^c (162)	165 ^c (162)	49	463	24	105	107	121	122	157	4116	1261				
s231	99 ^c (96)	99 ^c (96)	383 ^c (337)	2137	655	1996	542	1011	2366	3020	171	162	236	19 ^a			
s232	24 ^c (20)	15 ^c (14)	22 ^c (20)	3	141	1826	105	144	119	367	375	495	498	25 ^a			
s250	105 ^c (100)	20 ^c (18)	103 ^c (100)	8	314	32	296	296	296	367	375	495	498	25 ^a			

Table 1 continued

Problem	shgo-		tgo	Le-DSIMPL ^d		LGO	PSwarm ^d		Average		Maximum		DIRECT-LI ^d	
	shgo-simpl	-sobol		-v	-c		Minimum		f.e.	p.f.e.	f.e.	p.f.e.	p.p. = 10	f.e
s251	72 ^c (31)	63 ^c (62)	258 ^c (236)	186	9127	10575	83	84	129	137	175	180	7 ^a	
bunnag1	34 ^c (34)	47 ^c (26)	39 ^c (32)	630	721	1884	132	142	214	228	411	438	1529	
bunnag2	46 ^c (18)	36 ^c (18)	35 ^c (12)	16	500	76454	150	153	252	259	410	426	>100000	
Average	63 ^c (46)	85 ^c (56)	97 ^c (86)	366	>5877	6841	2590	2672	3011	3130	3637	3812	>17213	

^a result is outside the feasible region

^b (t) t out of 10 times the global solution was not reached

^c number of function evaluations (out of the given total) that were performed by the local minimisation routines.

^d results produced by Paulavičius and Žilinskas [35]

Table 2 Total and average performance over all 22 test problems

problem	name	f.e.	nlmin	nulmin	runtime (s)
All	shgo-simpl	1463	26	26	0.27294
	shgo-sobol	1864	23	23	0.11225
	tgo	2123	29	25	0.093607
Average	shgo-simplicial	65	1	1	0.012852
	shgo-sobol	88	1	1	0.004144
	tgo	100	1	1	0.004542

The SHGO-Simpl, SHGO-Sobol and TGO (using Henderson’s formula for k_c) algorithms were able to solve all 22 problems. The best results on each problem are bolded. The lowest average number of function evaluations was achieved by SHGO-Simpl followed by SHGO-Sobol and TGO. It can be observed that Lc-DISIMPL-v achieved a better performance than any other algorithm for the horst-1 to horst-6, hs024, hs035, s232, s250 and bunnag2 problems. As noted in [35] the initial triangulation of Lc-DISIMPL-v evaluates the function values at the vertices of the simplices and therefore for some of the tested problems the solutions were found after initial triangulation on one of the vertices of the feasible region. It is also possible to initiate SHGO with such an initial triangulation by definition the first few vertices in \mathcal{X} as the intercepts of the linear constraints in a similar way to [35] and then continuing to add sampling points as normal.

Table 2 provides additional information for SHGO and TGO including the total number of function evaluations required by the algorithm to solve the problem (f.e.), the number of minimisers generated as starting points by the algorithm (nlmin), the number of unique local minima mapped by the algorithm (nulmin) and the total processing time (runtime) in seconds.

It can be seen that neither of the SHGO algorithms produced more starting points leading to the same local minima as predicted by the theory for adequately sampled function surfaces. On the contrary TGO produced more than one starting point in the same locally convex domain on some test problems which lead to extra function evaluations, producing a poorer overall performance. While SHGO-Simpl had the lowest number of average function evaluations, a higher processing run time is observed compared to the other 2 algorithms. This can be explained by the fact the triangulation code for the sampling has not yet been optimised which consumed most of the run time. SHGO-Sobol and TGO use the same sampling generation code and it is observed that SHGO-Sobol has a lower processing run time as expected.

The source code used to produce these results including the scripts that run the test benchmarking suite is publically available at [11]. The specifications of the system used to run the test problems can also be found in [11].

6.2 Function evaluations and comparison to other open source global optimisation algorithms

In this section we present numerical experiments comparing the SHGO and TGO algorithms with the SciPy implementations [22] of basinhopping (BH) [27,48,50,51] and differential evolution (DE) [42]. These algorithms were chosen both because the open source versions are readily available in the SciPy project and because BH is commonly used in energy surface optimisations [49] from which the motivation for developing SHGO grew. DE has

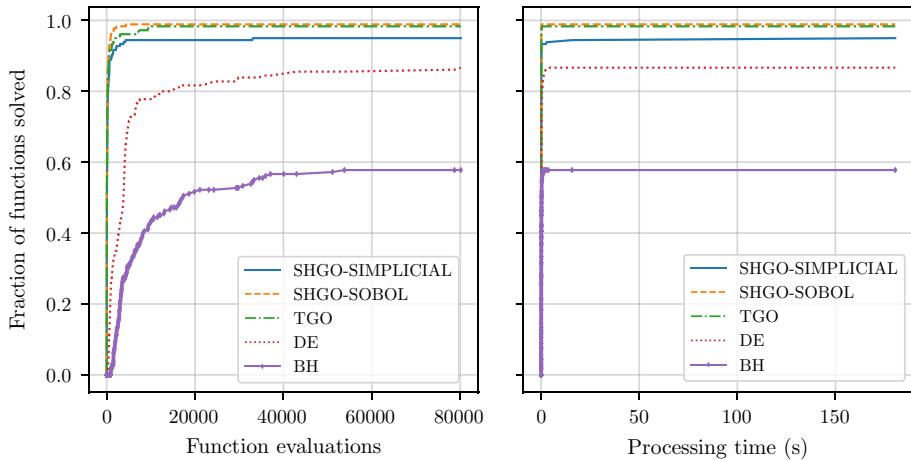


Fig. 11 Performance profiles for SHGO, TGO, DE and BH on SciPy benchmarking test suite

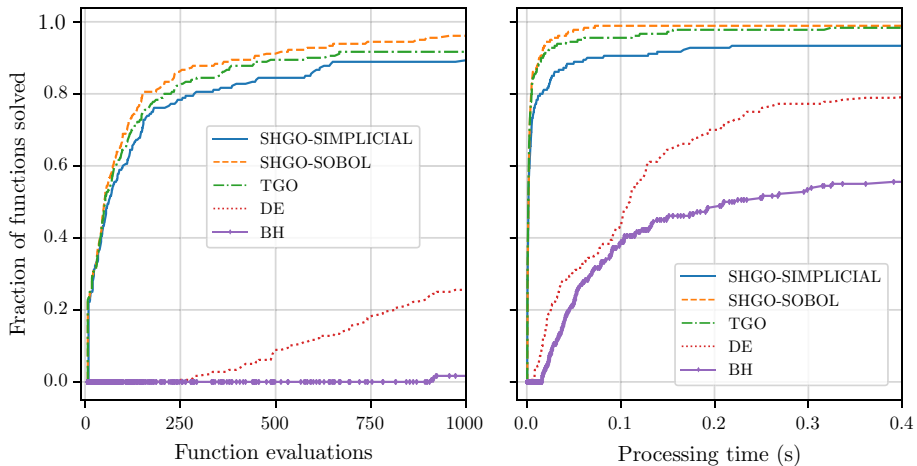


Fig. 12 Performance profiles zoomed in to the range of $f.e. = [0, 1000]$ function evaluations and $[0, 0.4]$ seconds run time

also been applied in optimising Gibbs energy surfaces for phase equilibria calculations [52]. The optimisation problems in “Online Appendix A” were selected from the SciPy global optimisation benchmarking test suite [1, 14, 21, 29, 30, 32]. The test suite contains multi-modal problems with box constraints, they are described in detail in [14]. We again used the stopping criteria $pe = 0.01\%$ for SHGO and TGO. For the stochastic algorithms (BH and DE) the starting points provided by the test suite were used. For every test the algorithm was terminated if the global minimum was not found after 10 minutes of processing time and the test was flagged as a fail. For comparisons we used normalised performance profiles [9] using function evaluations and processing time as performance criteria. In total 180 test problems were used.

From Fig. 11 it can be observed that for this problem set SHGO-Sobol was the best performing algorithm, followed closely by TGO and SHGO-Simpl. Fig. 12 provides a clearer comparison between these three algorithms. While the performance of all 3 algorithms are

comparable, SHGO-Sobol tends to outperform TGO, solving more problems for a given number of function evaluations. This is expected since, for the same sampling point sequence, TGO produced more than one starting point in the same locally convex domain on some test problems which leads to extra function evaluations. In total TGO produced 403 minima of which only 393 minima were unique while all of the 225 minima produced by SHGO-Sobol were unique. SHGO-Simpl produced 238 of which all 238 were unique. It is apparent that SHGO-Simpl performed worse compared to the other sampling methods despite a better performance on the test problem set with linear constraints. There are two reasons for this result. First of all the uniformity properties of the Sobol sequence hold only for hypercubes, therefore it is lost for geometries defined by the search spaces inside linear constraints. Secondly the current code for the triangulation of the simplex cannot add only one sampling point per iteration, but must split all the simplices until the symmetry of the entire complex is restored. This leads to a much higher number of function evaluations during the sampling step of the algorithm.

The Table in [11] shows the raw numerical results. Note that, unlike the data in performance profiles, failed test runs did not get set to the worst case performance criteria by any solver (in order to preserve the raw data). Therefore the total and average function evaluations and processing times are misleading. The Table is mostly useful for comparisons on a particular test problem as well as comparing the total number of minima and unique minima found.

6.3 Invariance and optimum minimiser pool

The following 4 optimisation test problems were used to demonstrate the applications of Theorem 5 and to show the minimiser pool growth compared to TGO over a large number of sampling points. The results plotted in Fig. 13 shows that SHGO performed as expected with the minimiser pool staying at the optimum cardinality to map all the local minima once the sampling is adequate as well as the shortcomings of the TGO especially in the higher dimensional test problems where the the minimiser pool tends to grow rapidly with the number sampling points N .

The Ursem01 function for two dimensions is defined as follows [14]

$$f(\mathbf{x}) = -\sin(2x_1 - 0.5\pi) - 3\cos(x_2) - 0.5x_1, \quad \mathbf{x} \in \Omega = [0, 9] \times [-2, 2] \quad (14)$$

The Paraboloid function for six dimensions is defined as follows

$$f(\mathbf{x}) = \sum_{i=1}^6 x_i^2, \quad \mathbf{x} \in \Omega = [-10, 10]^6 \quad (15)$$

The Bird function for two dimensions is defined as follows [14]

$$f(\mathbf{x}) = (x_1 - x_2)^2 + e^{[1-\sin(x_1)]^2} \cos(x_2) + e^{[1-\cos(x_2)]^2} \sin(x_1), \\ \mathbf{x} \in \Omega = [-2\pi, 2\pi]^2 \quad (16)$$

The Schwefel01 function for six dimensions is defined as follows [14]

$$f(\mathbf{x}) = \left(\sum_{i=1}^n x_i^2 \right)^{\sqrt{\pi}}, \quad \mathbf{x} \in \Omega = [-100, 100]^6 \quad (17)$$

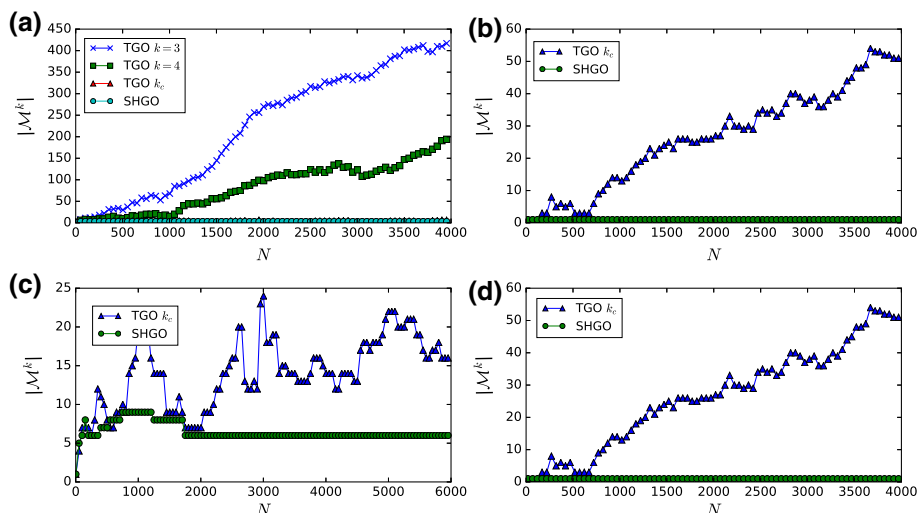


Fig. 13 **a** The minimiser pool growth of the TGO and SHGO algorithms for the smooth objective function described in Example 3 and restated in Eq. (14) for convenience, the SHGO never increases above the optimum of $|\mathcal{M}| = 3$, for TGO 3 different values of the k parameter are shown. **b** The minimiser pool growth for the six dimensional Paraboloid problem defined by Eq. (15), note that even though the problem has only one minimum, the minimiser pool for TGO set at $k = k_c$ tends to increase for increasing sampling points N . In general this problem is exacerbated in higher dimensions while SHGO stays at the optimum $|\mathcal{M}| = 1$. The TGO minimiser pool for $k = 3$ and $k = 4$ are not shown here because the minimiser pool grows too rapidly. **c** The minimiser pool growth for the two dimensional Bird problem defined by Eq. (16), an important observation here is that $|\mathcal{M}|$ is higher than optimum for SHGO before the sampling is adequate as defined by Eq. (5) which happens at the after there are $N = 1722$ Sobol sequenced points after which $|\mathcal{M}|$ stays at the optimum value equal to the number of unique local minima with increasing N . **d** The minimiser pool growth for the six dimensional Schwefel01 problem defined by Eq. (17), here again $|\mathcal{M}|$ for TGO set at k_c grows rapidly with N while $|\mathcal{M}|$ for SHGO stays constant at the optimum

7 Concluding remarks

The SHGO algorithm developed here shows promising properties and performance. On problems with linear constraints it was shown to provide competitive results to the TGO, Lc-DISIMPL, PSwarm and DIRECT-L1 algorithms. The use of a simplicial complex provides access to a wealth of tools from combinatorial topology and the growing field of computational homology. We are hopeful that these will drive further extensions and development of the algorithm. Many challenges remain such as finding the most appropriate sampling sequences for different classes of problems and finding computer resource efficient triangulation schemes. Due to the useful characterisations of objective function hypersurfaces provided by the homology groups of the simplicial complex SHGO allows an optimisation practitioner with a useful visual tool for understanding and efficiently solving higher dimensional black and grey box optimisation problems.

The main initial driving force behind the development of this algorithm grew out of a need for efficient, deterministic and reliable global optimisation methods for applications in phase equilibria modelling and calculations. However, the SHGO algorithm described here is appropriate for solving a wider class of global optimisation problems both those where mapping all the local minima is of interest and where only the global optimum is needed. It is

especially appropriate for computationally expensive black and grey box functions common in science and engineering as described for example by Shan and Wang [39].

Some key features of SHGO are that when the optimisation search space is adequately sampled and enough information is available to determine that all local minima have been mapped it is guaranteed that only one starting point for every locally convex domain will be produced by the algorithm. Note that in optimisation problems where the number of local minima is known, the sampling can stop and the local minimisation step started without superfluous function evaluations while for optimisation problems with an unknown number of local minima is unknown (and thus we can never truly know if all local minima has been found for any finite number of sampling) the guarantee still holds that SHGO will not produce superfluous starting points that lead to the same stationary points. In addition because the homology groups can be calculated as sampling progresses an optimisation practitioner can both visualise the extent of the optimisation problem's multi-modality and use intelligent stopping criteria for the sampling stage.

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Compliance with ethical standards

Conflicts of interest The authors declare that they have no conflict of interest.

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