

TAPL: Homework 1

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Exercise 3.5.16 A different way of formalizing meaningless states of the abstract machine is to introduce a new term called `wrong` and augment the operational semantics with rules that explicitly generate `wrong` in all the situations where the present semantics gets stuck. To do this in detail, we introduce two new syntactic categories.

Show that these two treatments of run-time errors agree by

Question 1 *finding a precise way of stating the intuition that “the two treatments agree”*

Given a full set of states \mathcal{S} and its subset \mathcal{A} containing all the "meaningful states", so the "meaningless states" could be described as \mathcal{B} which is equivalent to $\mathcal{S} - \mathcal{A}$.

The way we define "meaningless states" determines the set \mathcal{B} . And if two treatments agree, intuitionally, we should know that the set \mathcal{B} should be invariant in both treatments. So the key point for proving two treatments are agree is finding a way to prove the two forms of \mathcal{B} are equal.

Theorem 1 *If the two treatments agree, then the set \mathcal{B} is invariant between both solutions.*

Let's prove the two treatments agree by using **Theorem 1**.

Question 2 *proving it*

Let's take \mathcal{B}_1 as the set of meaningless states defined by the first treatment, and \mathcal{B}_2 is defined by using `wrong`.

Hypothesis 1 *There's no one state s in \mathcal{B}_1 which is not in \mathcal{B}_2 .*

$s \in \mathcal{B}_1$
 $\Leftrightarrow s$ is stucked
 $\Leftrightarrow s$ is a normal form but not a value(Definition 3.5.15)
 $s \in \mathcal{B}_2$
 $\Leftrightarrow s$ is not wrong
 $\Leftrightarrow s$ is one of true, false or nv
 $\Leftrightarrow s$ is a normal form and a value
 \Rightarrow *failed*.

Hypothesis 2 *There's no one state s in \mathcal{B}_2 which is not in \mathcal{B}_1 .*

$s \in \mathcal{B}_2$
 $\Leftrightarrow s$ is wrong
 $\Leftrightarrow s$ is a normal form but not a value(Definition 3.5.15)
 $\Leftrightarrow s$ is stucked
 $\Leftrightarrow s \in \mathcal{B}_1$
 \Rightarrow *failed*.

As **Hypothesis 1** and **Hypothesis 2** are both failed.

$s \in \mathcal{B}_1 \Rightarrow s \in \mathcal{B}_2$
 $s \in \mathcal{B}_2 \Rightarrow s \in \mathcal{B}_1$
 $\Rightarrow \mathcal{B}_1 = \mathcal{B}_2$

So the both treatments are agree.