

# A Tableau algorithm for $\mathcal{ALCSCC}$

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- 2 Difficulties with  $\mathcal{ALCSCC}$
- 3 Dealing with successor constraints
- 4 Tableau for  $\mathcal{ALCSCC}$
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# Motivation

- Reasoning in data base

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- Satisfiability check for reasoning
- Tableau algorithm for satisfiability check

# Tableau Algorithm

Main Idea (for  $\mathcal{ALCQ}$ ):

$$x : A \sqcap B \sqcap \exists r.C$$

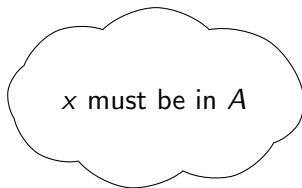


$x$  must be in  $A$

# Tableau Algorithm

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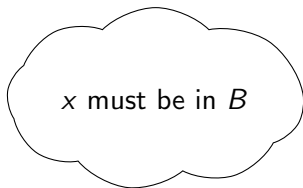
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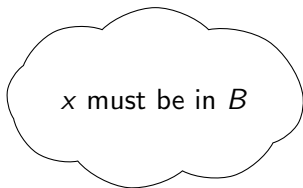


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$$\textcircled{x} \quad A, B$$

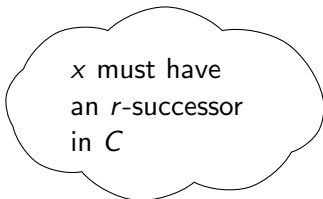


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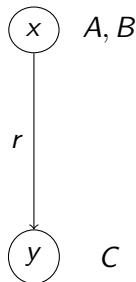
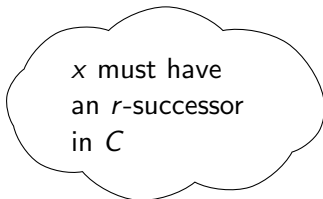
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# Tableau Algorithm

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# Goal

Tableau algorithm for  $\mathcal{ALCSCC}$  concepts

*ALCSCC*: successors

~~$\exists r.C$~~

~~$\forall r.C$~~

~~$\leq mr.C$~~

~~$\geq mr.C$~~

$succ(c)$

$c$ : **set constraint** or **cardinality constraint**

# *ALCSCC*: constraints

## set constraint:

- $hasChild \cap Female \subseteq Mother$
- $succ(2\ dvd\ |Edges|) \subseteq succ(|Edges| = |Nodes|)$

## cardinality constraint

- $2\ dvd\ |hasLegs|$
- $|Edges| \leq |Nodes|$
- $succ(2\ dvd\ |Edges|) \subseteq succ(|Edges| = |Nodes|)$

## Problem with successors constraints

$$x : succ(|s| > 1) \sqcap succ(|r| \leq |s|) \sqcap succ(|r| > |s|)$$

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- endless loop of adding  $r$ - and  $s$ -successors



## Problem with successors constraints

$$x : succ(|s| > 1) \sqcap succ(|r| \leq |s|) \sqcap succ(|r| > |s|)$$

- endless loop of adding  $r$ - and  $s$ -successors
- blocking?

## Problem with blocking

$$x : succ(2 \cdot |r| \leq 5 \cdot |s|) \sqcap succ(5 \cdot |s| \leq 2 \cdot |r|) \sqcap succ(|r| > 1)$$

*s*-successors



*r*-successors



## Problem with blocking

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## QFBAPA formula and solver

$$x : \text{succ}(2 \cdot |r| \leq 5 \cdot |s|) \sqcap \text{succ}(5 \cdot |s| \leq 2 \cdot |r|) \sqcap \text{succ}(|r| > 1)$$

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$$2 \cdot |X_r| \leq 5 \cdot |X_s| \wedge 5 \cdot |X_s| \leq 2 \cdot |X_r| \wedge |X_r| > 1$$



$$X_r = \{r_1, r_2, \dots, r_5\} \text{ and } X_s = \{s_1, s_2\}$$

## QFBAPA formula and solver

What about

$$2 \cdot |X_r| \leq 5 \cdot |X_s| \wedge 5 \cdot |X_s| \leq 2 \cdot |X_r| \wedge |X_r| > 1$$



$$|X_r| = 10 \text{ and } |X_s| = 4$$

## QFBAPA formula and solver

What about

$$2 \cdot |X_r| \leq 5 \cdot |X_s| \wedge 5 \cdot |X_s| \leq 2 \cdot |X_r| \wedge |X_r| > 1$$



$$|X_r| = 100 \text{ and } |X_s| = 40$$

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→ upper bound

→ ILP

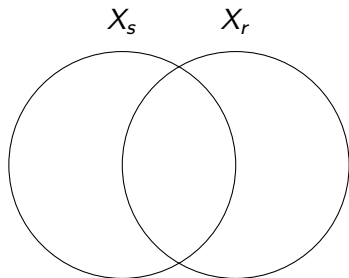
## Venn regions

**Problem:** Are the variables disjoint or not?

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**Solution:** Venn region



## Venn regions

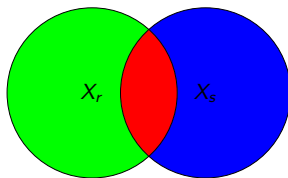
Two variables  $X_s$  and  $X_r \rightarrow$  four Venn regions:

$$v_1 = X_s \cap X_r$$

$$v_2 = X_s^{\neg} \cap X_r$$

$$v_3 = X_s \cap X_r^{\neg}$$

$$v_4 = X_s^{\neg} \cap X_r^{\neg}$$





## Venn regions

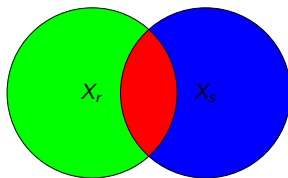
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$$X_s = v_1 \cup v_3 \text{ and } X_r = v_1 \cup v_2$$

# Venn regions

Franz Baader <sup>1</sup>:

For every QFBAPA formula  $\phi$  there is a number  $N$ , which is polynomial in the size of  $\phi$  and can be computed in polynomial time such, that for every solution  $\sigma$  of  $\phi$  there exists a solution  $\sigma'$  of  $\phi$  such that:

- $|\{v \mid v \text{ is a Venn region and } \sigma'(v) \neq \emptyset\}| \leq N$
- $\{v \mid v \text{ is a Venn region and } \sigma'(v) \neq \emptyset\} \subseteq \{v \mid v \text{ is a Venn region and } \sigma(v) \neq \emptyset\}$

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<sup>1</sup>A New Description Logic with Set Constraints and Cardinality Constraints on Role Successors. Springer International Publishing: 43-59, 2017

# Venn regions

In short:

$2^n$  Venn regions,

$n$  : amount of set variables



$N$  Venn regions,

$N$  : polynomial in the size of formula



## Transforming formula into ILP

$$2 \cdot |X_r| \leq 5 \cdot |X_s| \wedge 5 \cdot |X_s| \leq 2 \cdot |X_r| \wedge |X_r| > 1$$



$$-5 \cdot |X_s| + 2 \cdot |X_r| \leq 0$$

$$5 \cdot |X_s| - 2 \cdot |X_r| \leq 0$$

$$|X_r| > 1$$

## Transforming formula into ILP



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$$-5 \cdot |X_s| + 2 \cdot |X_r| + l_1 = 0$$

$$5 \cdot |X_s| - 2 \cdot |X_r| + l_2 = 0$$

$$|X_r| - l_3 = 2$$

## Transforming formula into ILP



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$$-5 \cdot |v_1 \cup v_3| + 2 \cdot |v_1 \cup v_2| + l_1 = 0$$

$$5 \cdot |v_1 \cup v_3| - 2 \cdot |v_1 \cup v_2| + l_2 = 0$$

$$|v_1 \cup v_2| - l_3 = 2$$

# ILP

$$\begin{array}{l}
 -5 \cdot |X_s| + 2 \cdot |X_r| + l_1 \\
 5 \cdot |X_s| - 2 \cdot |X_r| + l_2 \\
 |X_r| - l_3
 \end{array}
 \begin{pmatrix}
 & v_1 & v_2 & v_3 & v_4 & l_1 & l_2 & l_3 \\
 -5 & +2 & 2 & -5 & 0 & 1 & 0 & 0 \\
 5 & -2 & -2 & 5 & 0 & 0 & 1 & 0 \\
 1 & 1 & 0 & 0 & 0 & 0 & 0 & -1
 \end{pmatrix} \cdot \begin{pmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \\ x_5 \\ x_6 \\ x_7 \end{pmatrix}$$

$$= \begin{pmatrix} 0 \\ 0 \\ 2 \end{pmatrix}$$

## Upper bound

Christos H. Papadimitriou<sup>2</sup>:

For each ILP  $Ax = b$ ,  $A$   $m \times n$  matrix,  $x \in \mathbb{N}^n$  and  $b \in \mathbb{R}^m$  there exists an upper bound  $M$  and a solution  $x' \in \{0, \dots, M\}^n$  such that

$$Ax' = b$$

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<sup>2</sup>On the Complexity of Integer Programming. J. ACM, 28(4):765-768, Oct. 1981.



## Tableau for $\mathcal{ALCSCC}$

Now we consider the rules for the tableau algorithm.

## Decomposing rules

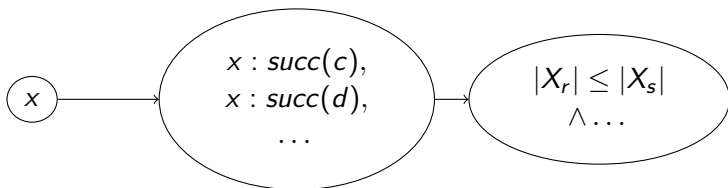
- $\sqcap$ -rule:  
If  $x : A \sqcap B$  is in ABox then  $x : A$  and  $x : B$  must be in ABox
- $\sqcup$ -rule:  
If  $x : A \sqcup B$  is in ABox then  $x : A$  or  $x : B$  must be in ABox

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Always apply when possible (higher priority)

## successor-rule

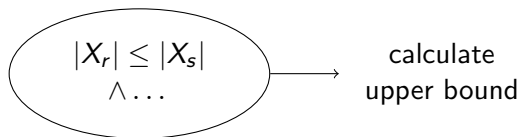


rule not  
applied yet

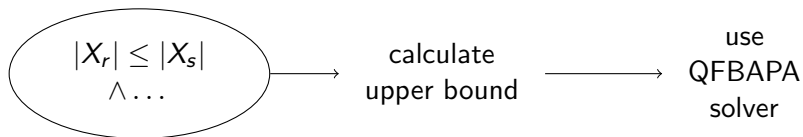
set of  
successor constraints

QFBAPA formula

## *successor-rule*



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## *successor-rule*

use QFBAPA solver

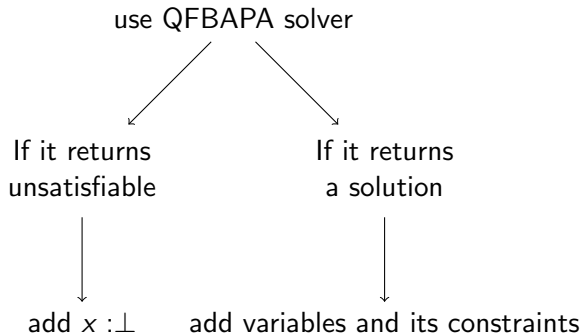


If it returns  
unsatisfiable



add  $x : \perp$

## *successor-rule*





## Example: Concept

*pancake* :  $\text{succ}(|\text{contains} \cap \text{Milk}| \geq |\text{contains} \cap \text{Flavour}|) \sqcap$

$\text{succ}(|\text{contains} \cap \text{Milk}| \leq 2 \cdot |\text{contains} \cap \text{Flavour}|) \sqcap$

$\text{succ}(|\text{contains} \cap \text{Egg}| \geq 1)$

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$$pancake : succ(|contains \cap Milk| \geq |contains \cap Flavour|),$$

$$pancake : succ(|contains \cap Milk| \leq 2 \cdot |contains \cap Flavour|),$$

$$pancake : succ(|contains \cap Egg| \geq 1)$$

## Example: *successor*-rule

$$\begin{aligned} & succ(|contains \cap Milk| \geq |contains \cap Flavour|) \\ & succ(|contains \cap Milk| \leq 2 \cdot |contains \cap Flavour|) \\ & succ(|contains \cap Egg| \geq 1) \end{aligned}$$

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$$\begin{aligned} & |X_{\text{contains}} \cap X_{\text{Milk}}| \geq |X_{\text{contains}} \cap X_{\text{Flavour}}| \wedge \\ & |X_{\text{contains}} \cap X_{\text{Milk}}| \leq 2 \cdot |X_{\text{contains}} \cap X_{\text{Flavour}}| \wedge \\ & |X_{\text{contains}} \cap X_{\text{Egg}}| \geq 1) \end{aligned}$$

## Example: *successor*-rule

Assume the solver returns

$$X_{Milk} = \{cup_m\},$$

$$X_{Flavour} = \{cup_f\}$$

$$X_{egg} = \{egg1, egg2\} \text{ and}$$

$$X_{contains} = X_{Milk} \cup X_{Flavour} \cup X_{Egg}$$

## Example: *successor*-rule

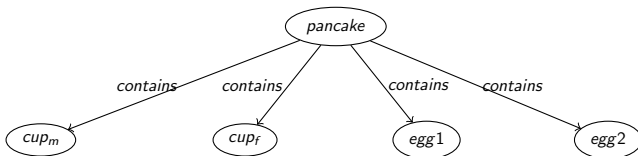
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- Tableau for *ALCSCC*
- 2ExpSpace because of upper bound

For the future:

- use/find smaller upper bound
- extend tableau for whole knowledge base
- tableau without QFBAPA solver