

A Tableau algorithm for \mathcal{ALCSCC}

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1 Introduction

2 Tableau for *ALCSCC*

Tableau Algorithm

Main Idea:

$$x : A \sqcap B \sqcap \exists r.C$$

x must be in A

x must be in B

x must have
an r -successor
in C

ALCSCC: successors

~~$\exists r.C$~~

~~$\forall r.C$~~

~~$\leq m r.C$~~

~~$\geq m r.C$~~

$$\boxed{\text{succ}(c)}$$

c : **set constraint** or **cardinality constraint**

\mathcal{ALCSCC} : constraints

set constraint:

- $r \subseteq s$
- $C \cap r \subseteq D$
- $\text{succ}(C \cap r) \subseteq \text{succ}(D)$

cardinality constraint

- $2 \text{ dvd } |r|$
- $|C \cap r| \leq |D|$
- $|\text{succ}(C \cap r)| \leq |\text{succ}(D)|$

Problem with successors constraints

$$x : succ(|s| > 1) \sqcap succ(|r| = |s|) \sqcap succ(|r| > |s|)$$

s-successors



r-successors



Problem with successors constraints

$$x : \underline{\text{succ}(|s| > 1)} \sqcap \text{succ}(|r| = |s|) \sqcap \text{succ}(|r| > |s|)$$

s-successors



r-successors



Problem with successors constraints

$$x : succ(|s| > 1) \sqcap \underline{succ(|r| = |s|)} \sqcap succ(|r| > |s|)$$

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Problem with successors constraints

$$x : succ(|s| > 1) \sqcap succ(|r| = |s|) \sqcap \underline{succ(|r| > |s|)}$$

s-successors



r-successors



Problem with blocking

$$x : \text{succ}(2 \cdot |r| \leq 5 \cdot |s|) \sqcap \text{succ}(5 \cdot |s| \leq 2 \cdot |r|) \sqcap \text{succ}(|r| > 1)$$

s-successors



r-successors



Problem with blocking

$$x : succ(2 \cdot |r| \leq 5 \cdot |s|) \sqcap succ(5 \cdot |s| \leq 2 \cdot |r|) \sqcap \underline{succ(|r| > 1)}$$

s-successors



r-successors



Problem with blocking

$$x : \underline{\text{succ}(2 \cdot |r| \leq 5 \cdot |s|)} \sqcap \text{succ}(5 \cdot |s| \leq 2 \cdot |r|) \sqcap \text{succ}(|r| > 1)$$

s-successors



r-successors



QFBAPA formula and solver

$$x : \text{succ}(2 \cdot |r| \leq 5 \cdot |s|) \sqcap \text{succ}(5 \cdot |s| \leq 2 \cdot |r|) \sqcap \text{succ}(|r| > 1)$$

QFBAPA formula and solver

$$x : \text{succ}(2 \cdot |r| \leq 5 \cdot |s|) \sqcap \text{succ}(5 \cdot |s| \leq 2 \cdot |r|) \sqcap \text{succ}(|r| > 1)$$



$$2 \cdot |X_r| \leq 5 \cdot |X_s| \wedge 5 \cdot |X_s| \leq 2 \cdot |X_r| \wedge |X_r| > 1$$

QFBAPA formula and solver

$$x : \text{succ}(2 \cdot |r| \leq 5 \cdot |s|) \sqcap \text{succ}(5 \cdot |s| \leq 2 \cdot |r|) \sqcap \text{succ}(|r| > 1)$$



$$2 \cdot |X_r| \leq 5 \cdot |X_s| \wedge 5 \cdot |X_s| \leq 2 \cdot |X_r| \wedge |X_r| > 1$$



$$|X_r| = 5 \text{ and } |X_s| = 2$$

QFBAPA formula and solver

What about

$$2 \cdot |X_r| \leq 5 \cdot |X_s| \wedge 5 \cdot |X_s| \leq 2 \cdot |X_r| \wedge |X_r| > 1$$



$$|X_r| = 10 \text{ and } |X_s| = 4$$

QFBAPA formula and solver

What about

$$2 \cdot |X_r| \leq 5 \cdot |X_s| \wedge 5 \cdot |X_s| \leq 2 \cdot |X_r| \wedge |X_r| > 1$$



$$|X_r| = 100 \text{ and } |X_s| = 40$$

QFBAPA formula and solver

We have infinite possible solutions

QFBAPA formula and solver

We have infinite possible solutions
Do we need all of them?

QFBAPA formula and solver

We have infinite possible solutions

Do we need all of them?

No: Consider formula as ILP and calculate an upper bound

ILP and upper bound

$$2 \cdot |X_r| \leq 5 \cdot |X_s| \wedge 5 \cdot |X_s| \leq 2 \cdot |X_r| \wedge |X_r| > 1$$



$$-5 \cdot |X_s| + 2 \cdot |X_r| \leq 0$$

$$5 \cdot |X_s| - 2 \cdot |X_r| \leq 0$$

$$-|X_r| < -1$$



$$-5 \cdot |X_s| + 2 \cdot |X_r| + l_1 = 0$$

$$5 \cdot |X_s| - 2 \cdot |X_r| + l_2 = 0$$

$$-|X_r| + l_3 = -1$$

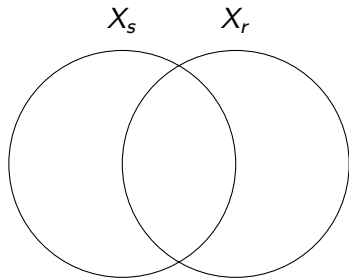
ILP and upper bound

Problem: Are the variables disjoint or not?

ILP and upper bound

Problem: Are the variables disjoint or not?

Solution: Venn region



Venn regions

Two variables X_s and $X_r \rightarrow$ four Venn regions:

$$v_1 = X_s \cap X_r$$

$$v_2 = X_s^\neg \cap X_r$$

$$v_3 = X_s \cap X_r^\neg$$

$$v_4 = X_s^\neg \cap X_r^\neg$$

Venn regions

Two variables X_s and $X_r \rightarrow$ four Venn regions:

$$v_1 = X_s \cap X_r$$

$$v_2 = X_s^\neg \cap X_r$$

$$v_3 = X_s \cap X_r^\neg$$

$$v_4 = X_s^\neg \cap X_r^\neg$$

$$X_s = v_1 \cup v_3 \text{ and } X_r = v_1 \cup v_2$$

ILP

$$\begin{array}{r}
 -5 \cdot |X_s| + 2 \cdot |X_r| + l_1 \\
 5 \cdot |X_s| - 2 \cdot |X_r| + l_2 \\
 -|X_r| + l_3
 \end{array}
 \begin{pmatrix}
 v_1 & v_2 & v_3 & v_4 & l_1 & l_2 & l_3 \\
 -5 & 2 & -5 & 0 & 1 & 0 & 0 \\
 5 & -2 & 5 & 0 & 0 & 1 & 0 \\
 -1 & -1 & 0 & 0 & 0 & 0 & 1
 \end{pmatrix} \cdot \begin{pmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \\ x_5 \\ x_6 \\ x_7 \end{pmatrix}$$

$$= (0 \quad 0 \quad -1)$$

ILP and upper bound

Christos H. Papadimitriou¹:

There exists an upper bound M such that for each ILP $Ax = b$, A a $m \times n$ matrix, $x \in \mathbb{N}^n$ and $b \in \mathbb{R}^m$ there exists a solution $x' \in \{0, \dots, M\}^n$.

¹On the Complexity of Integer Programming. J. ACM, 28(4):765-768, Oct. 1981.

Tableau algorithm

- \Box -rule:
- \sqcup -rule: