

$C = \langle x^4 + x^2 + x + 1 \rangle$ , цикличен код с дължина 7, над  $F_2$

$$g(x) = x^4 + x^2 + x + 1, \deg(g(x)) = 4 = n - k, n = 7, k = n - 4 = 7 - 4 = 3,$$

$C$  е  $[7, 3]$  цикличен код над  $F_2$

$$h(x) = \frac{x^7 - 1}{g(x)} = \frac{x^7 - 1}{x^4 + x^2 + x + 1} = x^3 - x - 1 + \frac{2x^2 + 2x}{x^4 + x^2 + x + 1},$$

тъй като  $2x^2 + 2x \equiv 0 \pmod{2}$ , то следва, че  $g(x) \mid x^7 - 1$  без остатък

в  $F_2$ ,  $x^3 - x - 1 \equiv x^3 + x + 1 \pmod{2}$ , което е пораждащ полином на дулания код на  $C$

$$H = \begin{matrix} 1110100 \\ 0111110 \\ 0010001 \end{matrix}$$

a)

$H$  е  $3 \times 7$  матрица, тогава  $n - k = 3, n = 7$ , от тук,  $k = 4$ .  $H_3[7, 4, 3]_2$

b)

$v$  е код, тогава и само тогава, когато  $vH^T = 0$ .

Това са векторите

$$\begin{aligned} &\{(0,0,0,0,0,0,0), (1,1,0,1,0,0,0), (0,1,0,0,1,0,0), (1,0,0,1,1,0,0), \\ &(1,1,0,0,0,1,0), (0,0,0,1,0,1,0), (1,0,0,0,1,1,0), (0,1,0,1,1,1,0), \\ &(0,1,1,0,0,0,1), (1,0,1,1,0,0,1), (0,0,1,0,1,0,1), (1,1,1,1,1,0,1), \\ &(1,0,1,0,0,1,1), (0,1,1,1,0,1,1), (1,1,1,0,1,1,1), (0,0,1,1,1,1,1)\}, \end{aligned}$$

16 на брой.

c)

$$G = \begin{matrix} 1000111 \\ 0100101 \\ 0010110 \\ 0001011 \end{matrix}$$

Намираме  $(0, 1, 0, 1)G = (0, 1, 0, 1, 1, 1, 0)$ ,  $(1, 0, 1, 0)G = (1, 0, 1, 0, 0, 0, 1)$

d)  $n - k = 3, n = 7$ , тогава  $H_3[7, 3, 3]_2$  двоичен код.

e) Умножаваме всички бинарни вектори с дължина 3 с матрицата  $H$  и получаваме

$$\begin{aligned} &\{(0,0,0,0,0,0,0), (0,0,1,0,0,0,1), (0,1,1,1,1,1,0), (0,1,0,1,1,1,1), \\ &(1,1,1,0,1,0,0), (1,1,0,0,1,0,1), (1,0,0,1,0,1,0), (1,0,1,1,0,1,1)\}, \end{aligned}$$

8 на брой.