

(16)

with no  $(r+1)$ -clique,  $T_{n,r}$  has the maximum numbers of edges.

### Edge - Coloring

An Edge - Coloring of a graph is an assignment of ~~the~~ colors to the edges of the graph so that no two incident edges have the same color.

OR, An edge coloring of a graph  $G$  is a function  $f: E(G) \rightarrow C$ , where  $C$  is a set of distinct colors. For any positive integer  $k$ , a  $k$ -edge coloring is an edge coloring that uses exactly  $k$  different colors.

### Proper Edge Coloring

A proper edge coloring of a graph is an edge coloring such that no two adjacent edges are assigned the same color. Thus a proper edge coloring of  $G$  is a function

$$f: E(G) \rightarrow C$$

such that

$$f(e) \neq f(e')$$

whenever edges  $e$  and  $e'$  are adjacent in  $G$ .

## Chromatic Index

The chromatic index of a graph  $G$ , denoted by  $\chi'(G)$ , is the minimum number of different colors required for a proper edge colouring of  $G$ .  $G$  is  $k$ -edge-chromatic if  $\chi'(G) = k$ .

→ Chromatic index of Bipartite graph is always  $\Delta(G)$ .

### Theorem

If  $G$  is any simple graph, then  $\chi'(G) \geq \Delta$ .

Proof: Given that  $G$  is any simple graph. We know  $\Delta$  is the maximum degree of the graph  $G$ . Therefore, there exist a vertex  $v$  such that

$$d(v) = \Delta.$$

So there are  $\Delta$  edges incident on this vertex  $v$ , i.e., all the  $\Delta$  edges are adjacent. So we require at least  $\Delta$  colours to proper colouring of the ~~edge~~ of  $G$ .

Thus  $\chi'(G) \geq \Delta$ .

Lemma: Let  $c = (E_1, E_2, \dots, E_k)$  be the optimal  $k$ -edge colouring of  $G$ . If there is a vertex  $v$  in  $G$  and colours  $i$  and  $j$  such that  $i \neq j$  and  $i$  is not represented at  $v$  and  $j$  is represented twice at  $v$ . Then the component  $H$  of  $G$  ( $E_i \cup E_j$ ) that

that contain  $u \Rightarrow$  an odd cycle.

### Theorem

If  $G$  is bipartite, then  $\chi'(G) = \Delta$

Proof: Suppose  $G$  is bipartite graph with  $\chi'(G) > \Delta$ .  
 Let  $C = (E_1, E_2, \dots, E_\Delta)$  be an optimal  $\Delta$ -edge colouring of  $G$ .

Since  $\chi'(G) > \Delta$ , this colouring can not be proper.  
 Let  $u$  be a vertex such that  $c(u) < d(u)$ . Now,  
 consider  $\Delta$ -edge colouring out  $u$ . Some colour is not  
 used and  $j \in$  represented twice.

By Lemma,  $G$  contains an odd cycle, which  
 $\Rightarrow$  a ~~contradiction~~ contradiction to the fact that  
 $G$  has no odd cycle. We know that for  
 any simple graph  $\chi'(G) \geq \Delta$ .

Therefore  $\chi'(G) = \Delta$ .

□

### Vizing Theorem

If  $G$  is simple, then either  $\chi'(G) \geq \Delta$  or  
 $\chi'(G) \leq \Delta + 1$ .

Proof: Let  $G$  be a simple graph. If  $G$  is  
 bipartite, then  $\chi'(G) = \Delta$ .

We need only to show that  $\chi'(G) \leq \Delta + 1$ .

Suppose  $\chi'(G) > \Delta + 1$ . Let  $C = (E_1, E_2, \dots, E_{\Delta+1})$  be an optimal  $\Delta+1$ -edge colouring of  $G$  and  $u$  be the vertex such that  $C(u) < d(u)$ . Then there exist colours  $i_0$  and  $i_1$  such that  $i_0$  is not represented at  $u$  and  $i_1$  is represented at least twice at  $u$ .

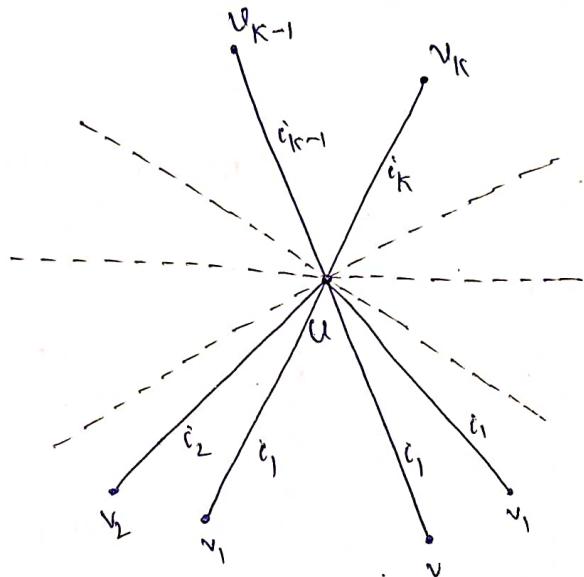


Fig: a

Let  $uv_1$  have colour  $i_1$  as shown in fig(a). Since  $d(v_1) < \Delta + 1$ , some colour  $i_2$  is not represented at  $v_1$ . Now  $i_2$  must be represented at  $u$ , since otherwise by recolouring  $uv_1$  with  $i_2$  we would obtain an improvement on  $C$ . Thus some edge  $uv_2$  has colour  $i_2$ .

Again  $d(v_2) < \Delta + 1$ , some colour  $i_3$  is not represented at  $v_2$  and  $i_3$  must be represented at  $u$ . Since, otherwise by recolouring  $uv_1$  with  $i_2$  and  $uv_2$  with  $i_3$  we would obtain an improved  $\Delta+1$  edge colouring.

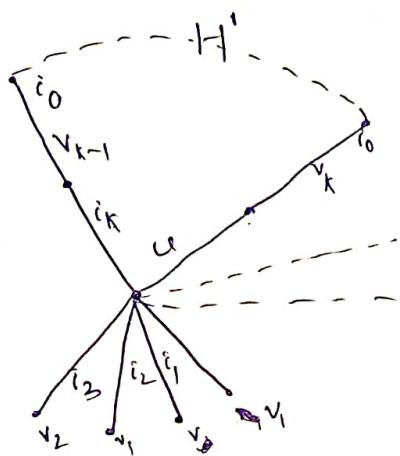
Thus some edge  $uv_3$  has colour  $i_3$ , continuing this procedure we construct a sequence  $(v_1, v_2, \dots)$  of vertices and  $(i_1, i_2, \dots)$  of colours such that

- (i)  $uv_j$  has colour  $i_j$

(ii)  $i_{j+1}$  is not represented at  $v_i$  (since  $d(v)$  is finite, there exist a smallest integer  $\ell$  such that for some  $k < \ell$ ).

$$(iii) i_k = i_{\ell+1}$$

Now we recolour  $G$  as follows for  $1 \leq i \leq k-1$ ,  $uv_i$  with colour  $i_{j+1}$  yielding a new  $A+1$ -edge colouring  $c' = (E'_1, E'_2, \dots, E'_{A+1})$  of  $G$ .



Clearly,  $c'(v) \geq c(v)$  for every  $v \in V$  and  $c'$  is also an optimal  $A+1$  edge colouring of  $G$ .

The component  $H'$  of  $G [E'_{i_k} \cup E'_{i_0}]$  that contains  $u$  is an odd cycle.

Now, in addition  $uv_j$  with the colour  $i_{j+1}$ ,  $k \leq j \leq \ell-1$  and  $uv_1$  with the colour  $i_k$  to obtain a  $A+1$  edge colouring  $c'' = (E''_1, E''_2, \dots, E''_{A+1})$  as the following fig : c.

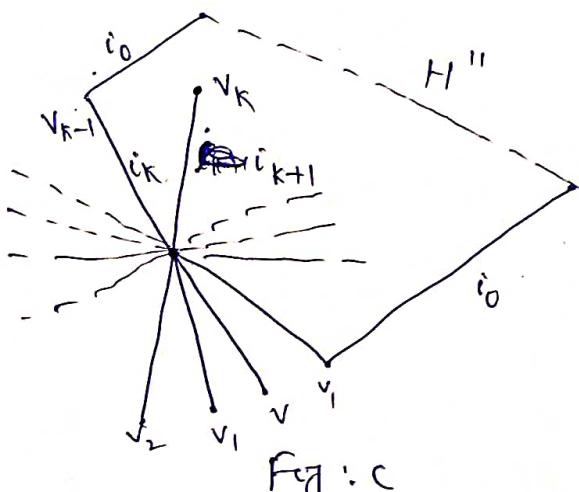


Fig : c

$C''(v) \geq c(v)$  for every  $v \in V$  and the component  
 $H''$  of  $G[B_{i_k}'' \cup E_{i_0}'']$  that contains  $u$  is an odd cycle.  
But since  $v_k$  has degree 2 in  $H'$ , clearly  $v_k$   
has degree 1 in  $H''$ . But  $H''$  is an odd cycle  
contains  $u$ , which is a contradiction.

This prove that our assumption is wrong.

Thus  $\chi'(G) \geq A$  and  $\chi'(H) \leq A+1$ .

(proved)