

CONNECTIVITY

Vertex Connectivity


The vertex connectivity of a graph G denoted by $k(G)$ is the minimum vertices whose removal from G disconnects G or a trivial graph.

- A graph G is k -connected if its connectivity is at least k .

(i) For any connected graph, $k(G) \geq 1$.

(ii) For any trivial graph, $k(G) = 0$.

(iii) For any disconnected graph, $k(G) = 0$.

(iv) If G is a complete graph of order ' n ', then $k(G) = n-1$. For K_4 , $k(G) = 3$.

(v) If G is not a complete graph of order ' n ', then $k(G) \leq n-2$.

(vi) If G is a complete bipartite graph, then $k(G = K_{m,n}) = \min\{m, n\}$. The connectivity of $K_{3,3}$ is 3, means the graph is 1-connected, 2-connected, and 3-connected, but not 4-connected.

Definition

A graph G is said to be k -connected if $k(G) = k$.

Example

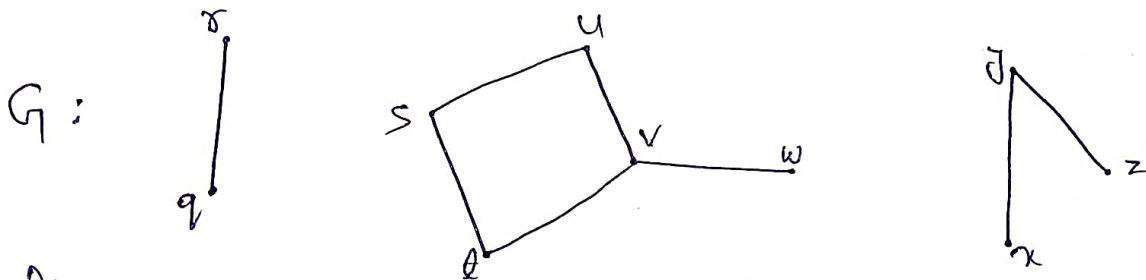
(vii) The vertex connectivity of a graph of order

at least ~~three~~ there is one if and only if it has a cut vertex.

(vi) Vertex connectivity of a path is one and that of cycle C_n ($n \geq 4$) is two.

Problem

Find the (i) vertex sets of components
(ii) Cut-vertices
(iii) Cut-edges of the graph given below.



Solⁿ: The graph has three components. The vertex set of the components are $\{q, r\}$, $\{s, t, u, v, w\}$ and $\{x, y, z\}$.

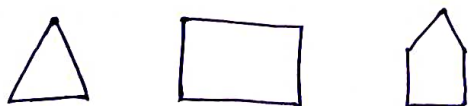
The cut vertices are v and y .

The cut edges are qr , vw , xy and yz .

Definition

A graph G is said to be k -connected if $k(G) = k$.

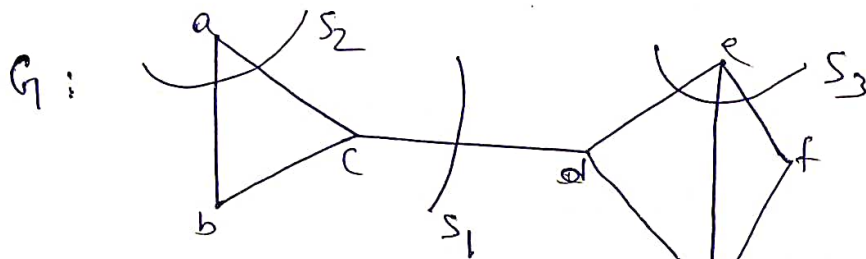
Example: All cycle graph C_n , $n \geq 3$ are 2-Connected



If we remove any two vertices, then graph is disconnected. So it is 2-connected. (19)

Edge Connectivity $k'(G)$

The number of edges in the smallest subset of a graph G is defined as the edge connectivity of G and it is denoted by $k'(G)$.



Here $k'(G) = 1$, because if we remove $\{c, d\}$ edge then G is disconnected.

- If G is a trivial graph, then $k'(G) = 0$.
- If G is a disconnected graph, then $k'(G) = 0$.
- In complete graph K_n , $k'(G) = n - 1$.

Ex:



K_5

Here $k'(K_5) = 4$.

Definition

A graph G is said to be k -edge connected if $k'(G) = k$.

- All nontrivial tree are 1-edge connected.

- All cycle's are 2-edge connected.

Notation

- ① $\delta(G)$ is known as minimum degree of a graph G .
- ② $\Delta(G)$ is known as maximum degree of a graph G .

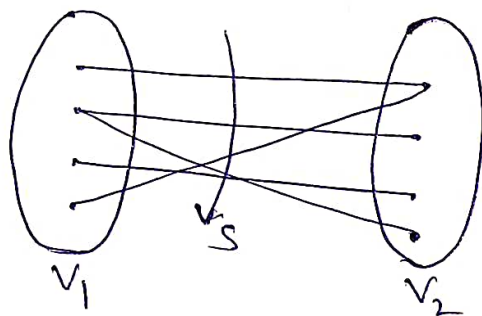
Whitney Inequality

For any simple graph G ,

$$k(G) \leq k'(G) \leq \delta(G). \text{ Equality holds for complete graph.}$$

Proof:

Let $k'(G) = \alpha$. Then there exists a cutset S having α -edges such that $G-S$ is disconnected.



Let V_1 and V_2 be the partition of $V(G)$ with respect to S . By removing of most α -vertices from V_1 (or V_2) on which the edges in S are incident, we have

$$k(G) \leq \alpha = |S| = k'(G)$$

$$\Rightarrow k(G) \leq k'(G) \quad \text{--- (1)}$$

Next, we have to show that $k'(G) \leq \delta(G)$.

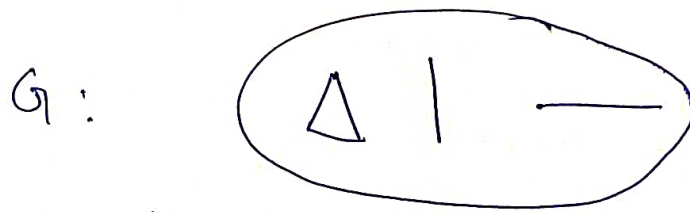
Case-1

If G has no edges, then $k'(G) = 0$ and $\delta(G) = 0$.

$$\therefore k'(G) = \delta(G)$$

Case-II

If G is a disconnected graph having some edges



Here $k'(G) = 0$ and $\delta(G) \geq 0$.

So $k'(G) \leq \delta(G)$

Case-III

If G is a connected graph with a vertex v having minimum degree $\delta(G)$, then a disconnected graph results when all the incident on v are removed.

$$\therefore k'(G) \leq \delta(G) \quad \text{--- (2)}$$

Combining (1), & (2), we get

$$k(G) \leq k'(G) \leq \delta(G) \quad \text{(proved)}$$