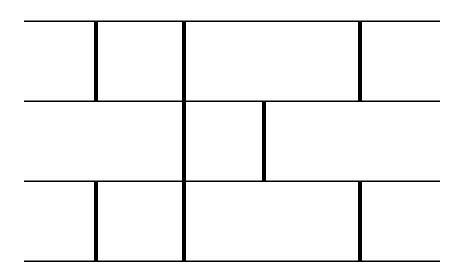
Mreže za sortiranje

Sortiranje

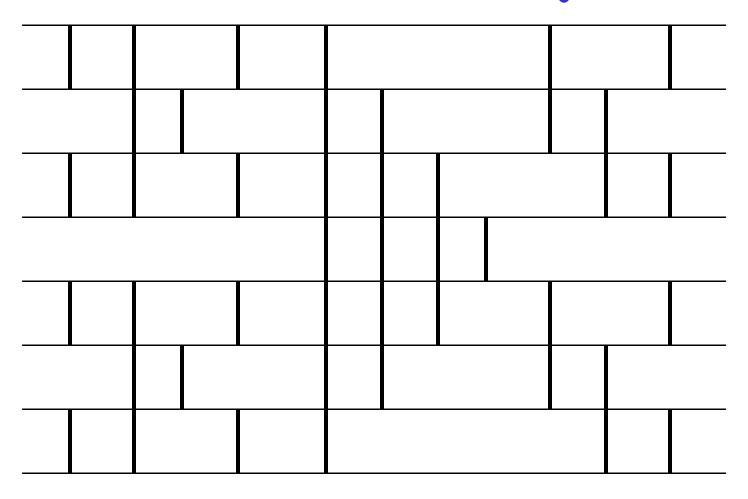
4		10
9		9
3		8
1		7
2		4
10		3
8		2
7		1

Mreža za sortiranje



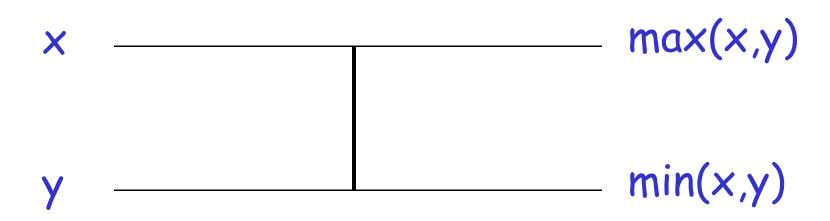
(Mreža za bitonik sortiranje)

Mreža za sortiranje

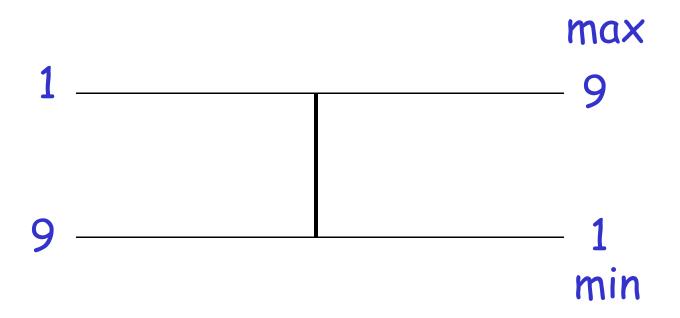


(Mreža za bitonik sortiranje)

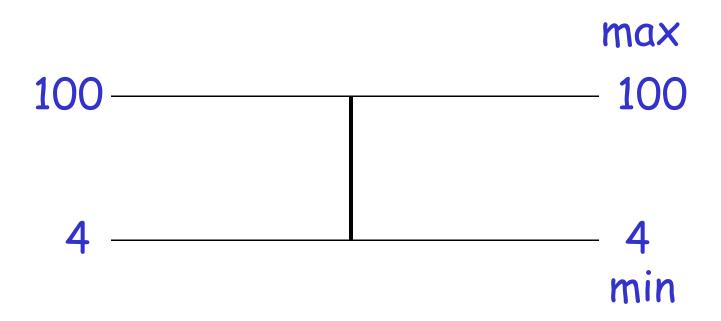
Komparator

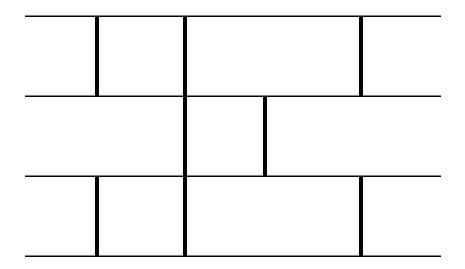


Komparator

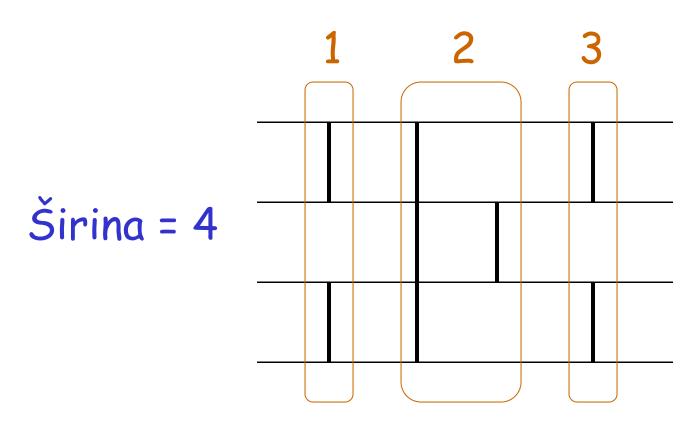


Komparator

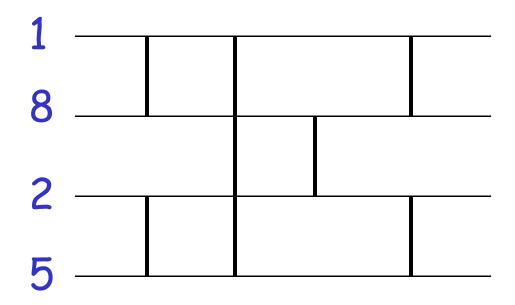


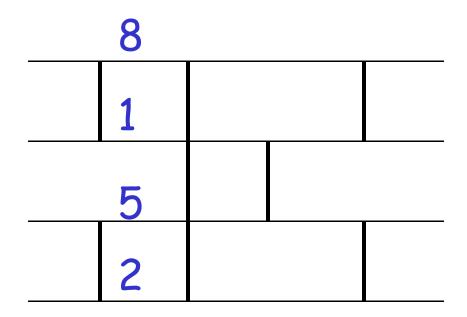


Nivoi

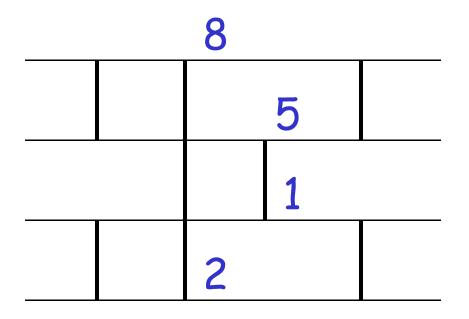


Dubina = 3

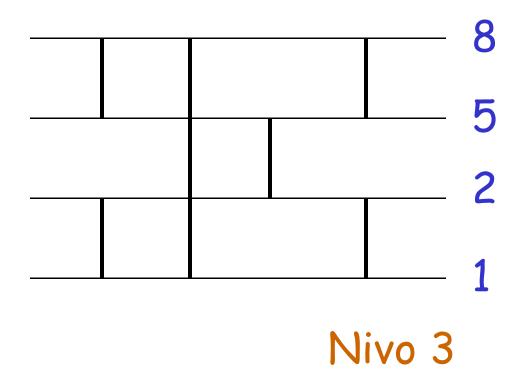




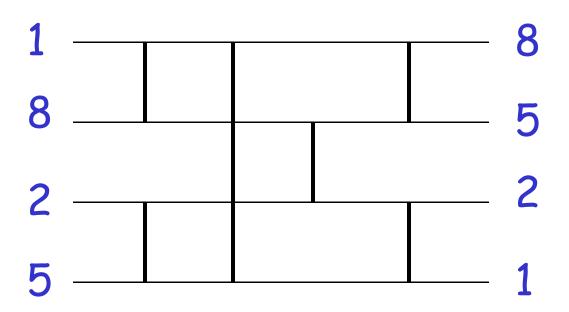
Nivo 1



Nivo 2

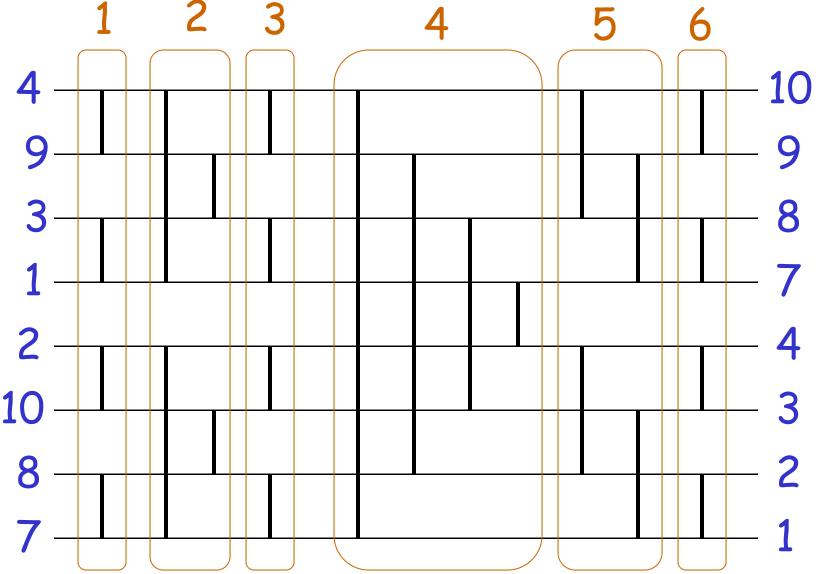


Ulaz Izlaz

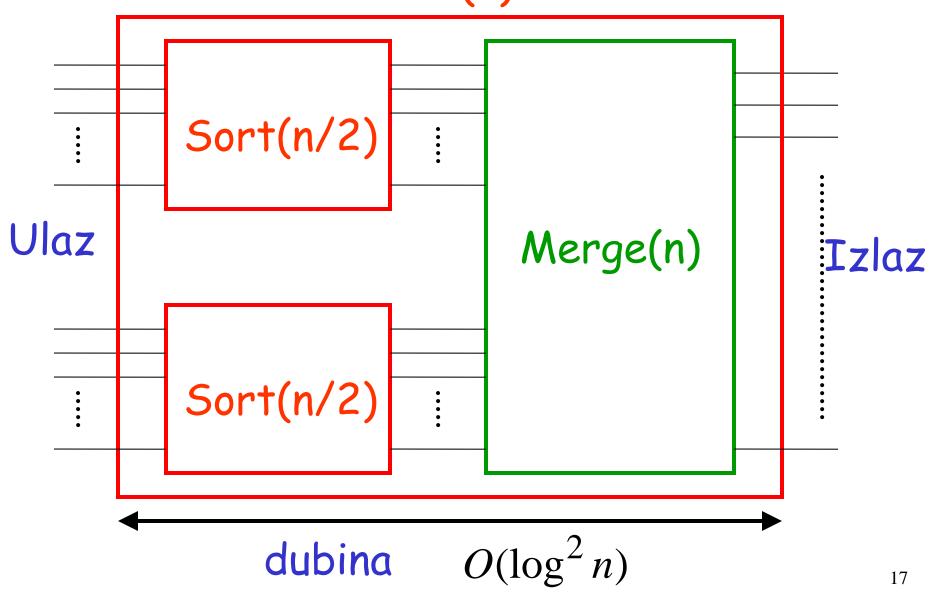


Ulaz Izlaz 10

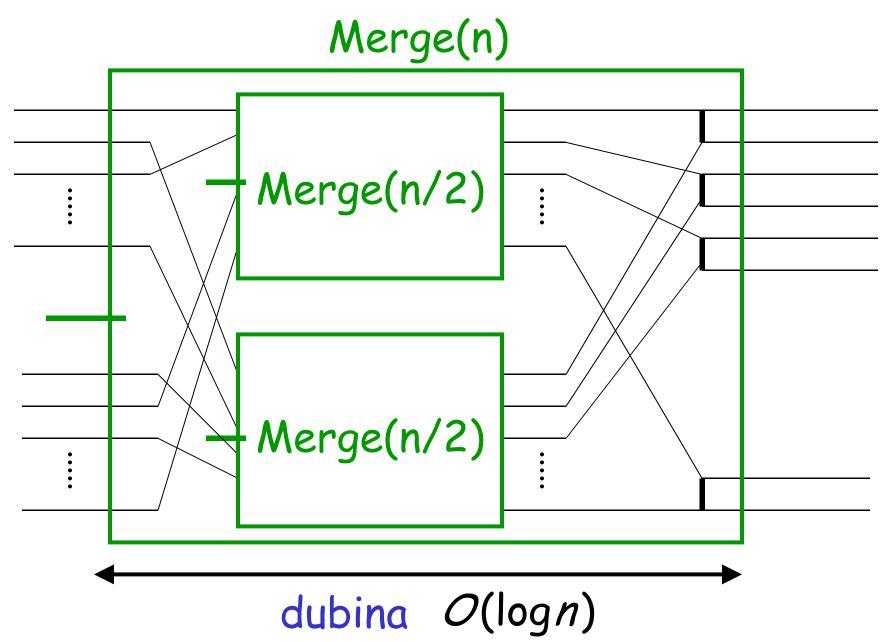
nivoi



Rekurzivna konstrukcija Mreža za bitonik sortiranje Sort(n)



Rekurzivna konstrukcija Mreže za spajanje



Osnova za indukciju

Širina n

Dubina mreže za sortiranje: $O(\log^2 n)$

Širina spajača
$$n \frac{n}{2} \frac{n}{4}$$

Ukupna
$$\log n + \log \frac{n}{2} + \log \frac{n}{4} + ... + \log 2 = O(\log^2 n)$$
 dubina

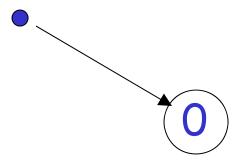
Mreža za brojanje

Problem brojanja

0

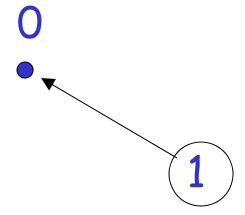
Deljena promenljiva

Token = Zahtev za inkrement



Deljena promenljiva

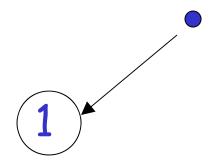
Zahtev za inkrement



Deljena promenljiva

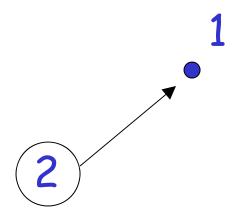
Deljena promenljiva

Zahtev za inkrement



Deljena promenljiva

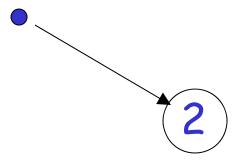
Zahtev za inkrement



Deljena promenljiva

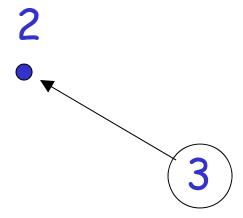
Deljena promenljiva

Zahtev za inkrement



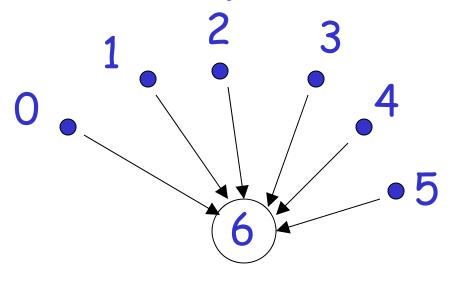
Deljena promenljiva

Zahtev za inkrement



Deljena promenljiva

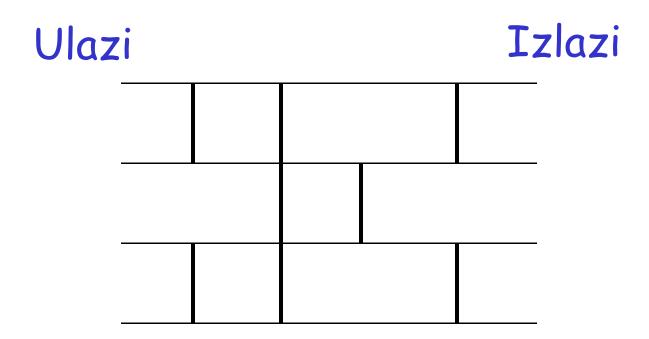
Sekvencijalno usko-grlo



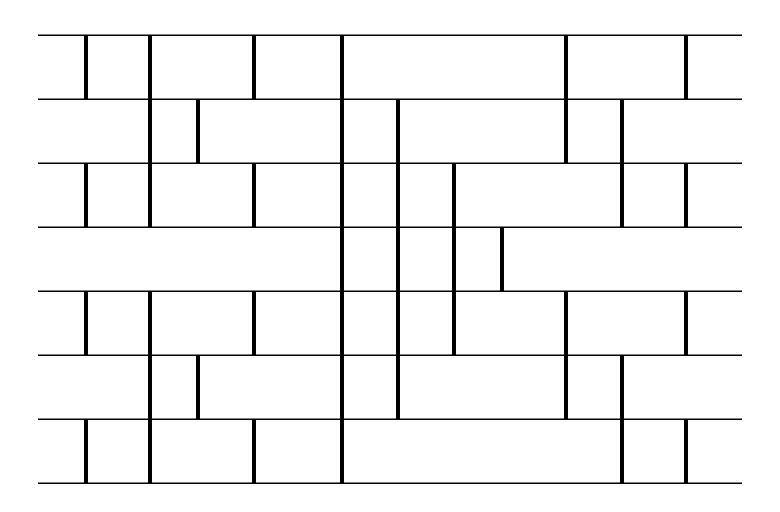
Deljena promenljiva

Zahtevi moraju biti serializovani

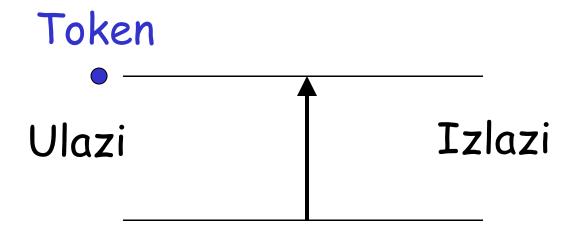
Mreža za brojanje



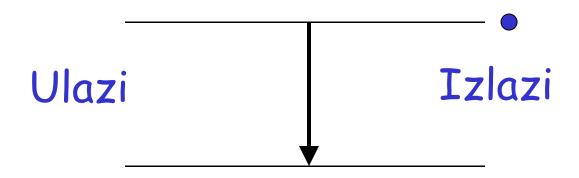
Mreža za brojanje



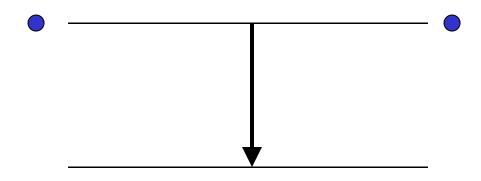
Balanser

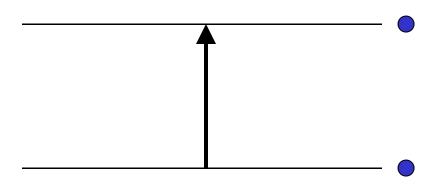


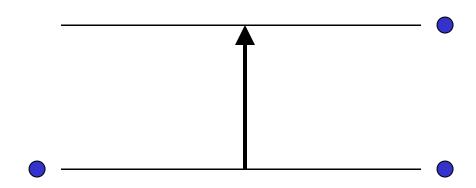
Balanser

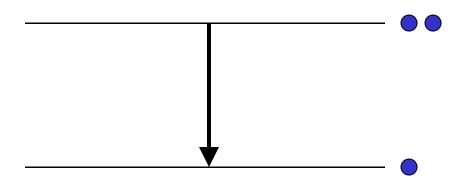


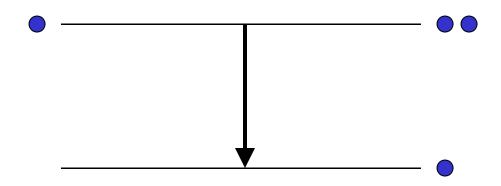
Balanser

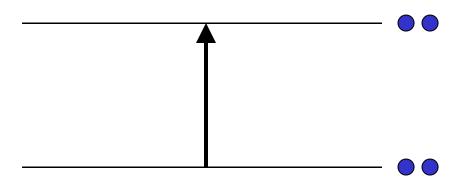


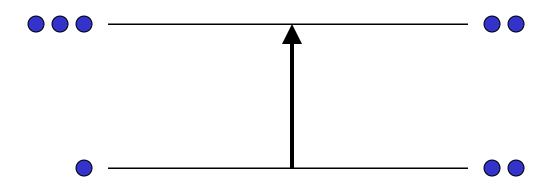




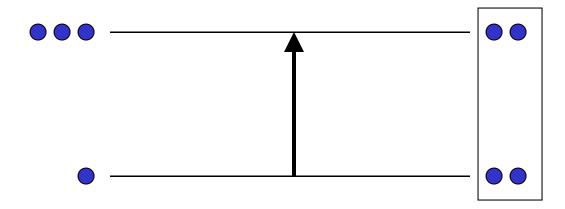


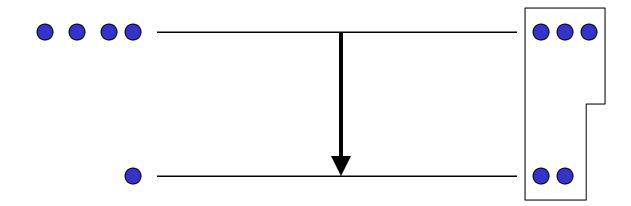






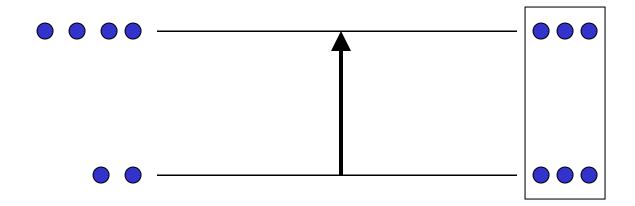
Svi tokeni zajedno





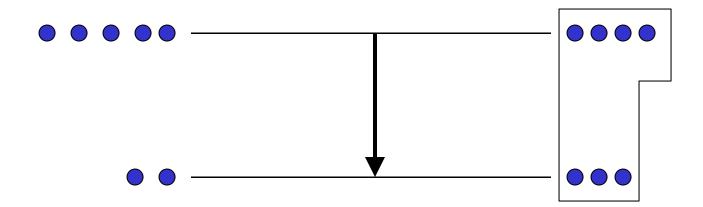
Osobina koraka

Sledeći primer



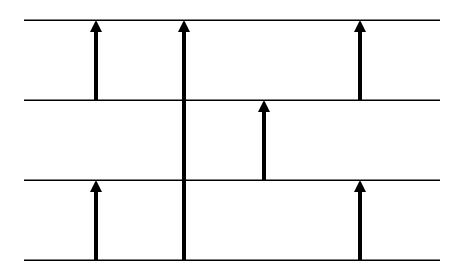
Osobina koraka

Sledeći primer

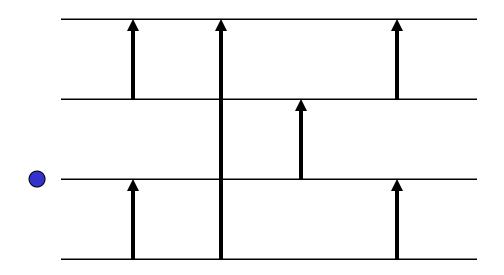


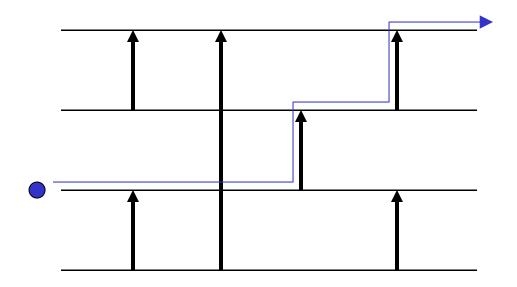
Osobina koraka

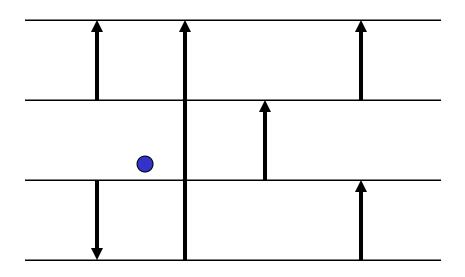
Sledeći primer

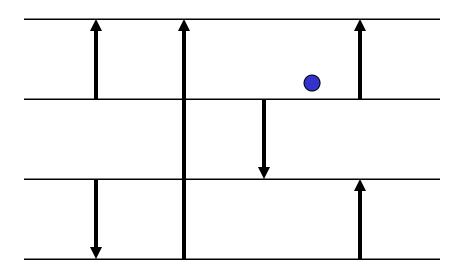


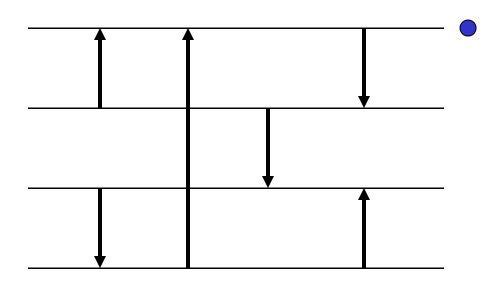
(Mreža za bitonik brojanje)

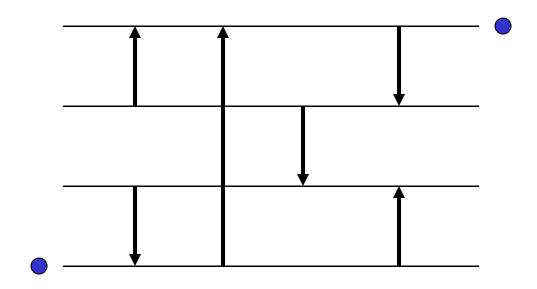


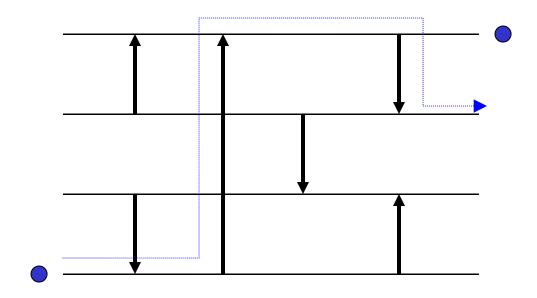


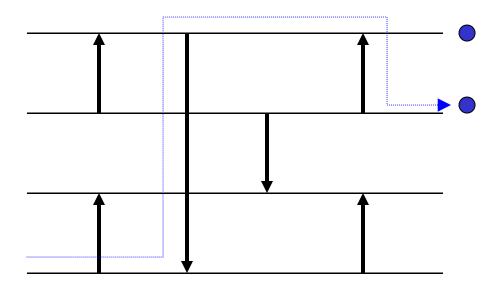


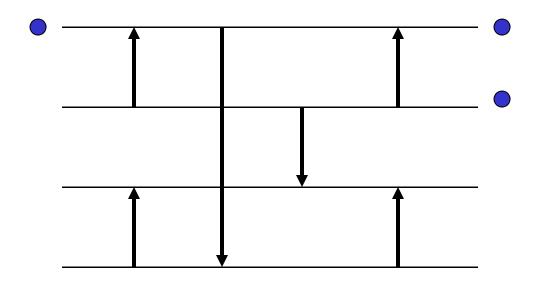


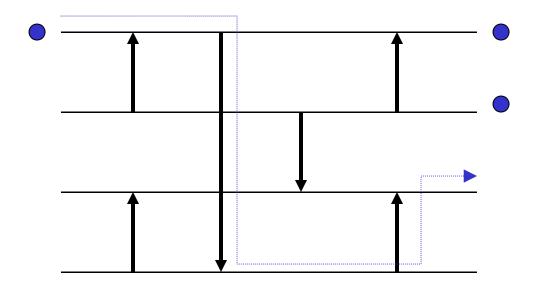


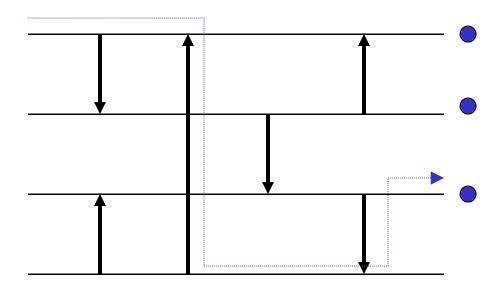


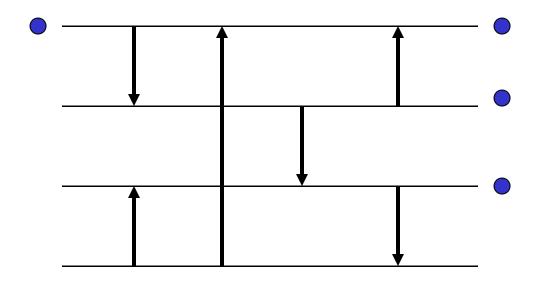


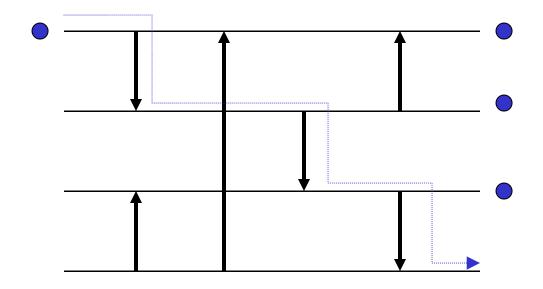


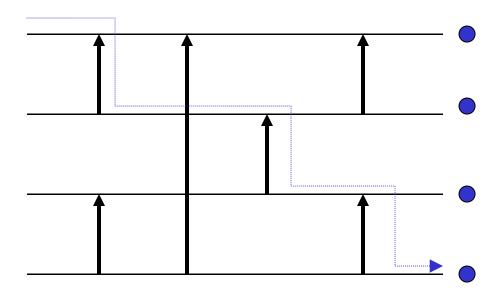


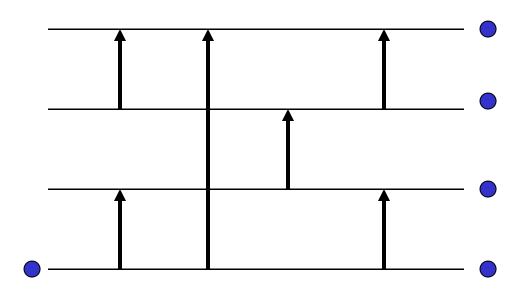


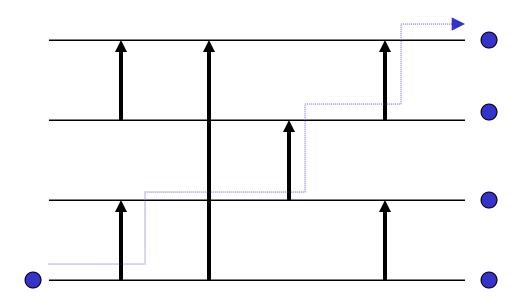


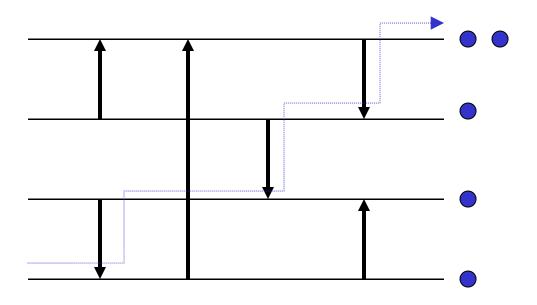


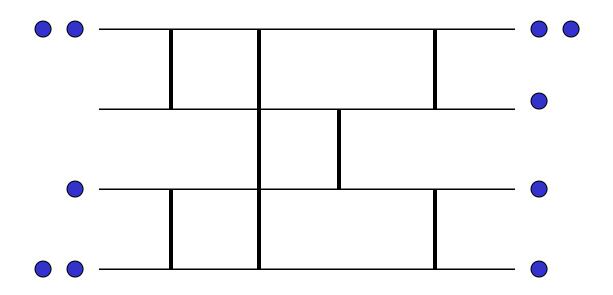




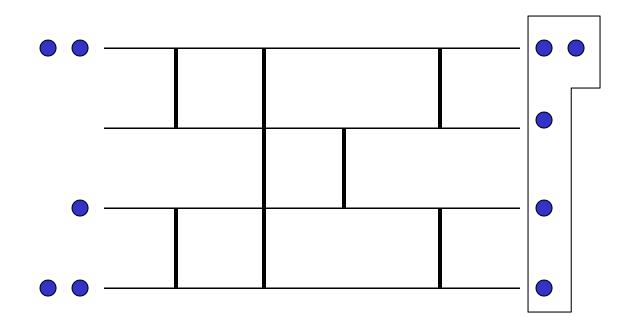




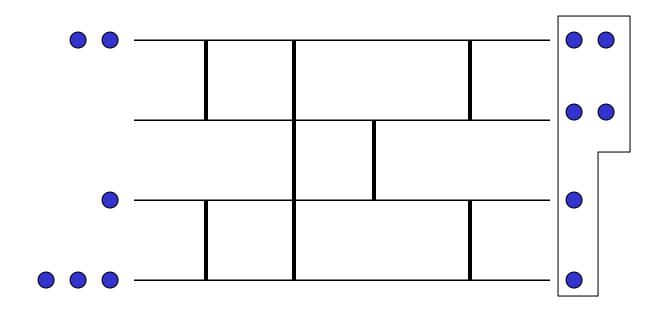




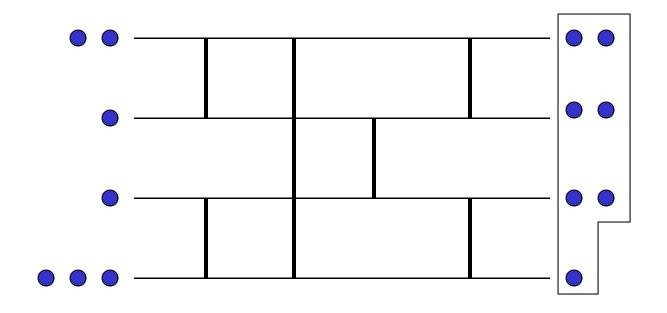
Svi tokeni



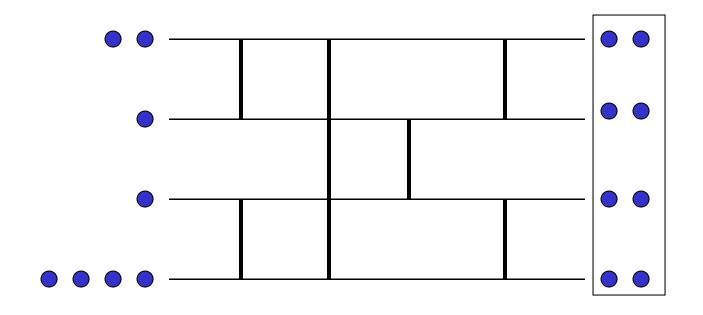
Svi tokeni



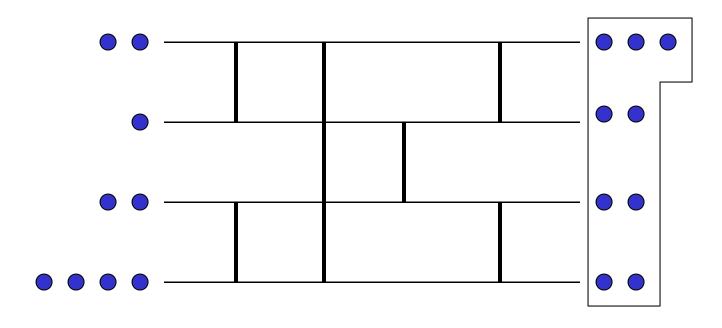
Sledeći primer



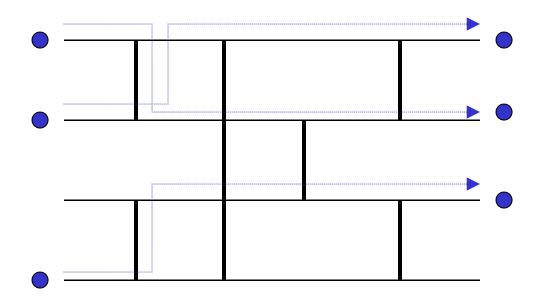
Sledeći primer



Sledeći primer



Sledeći primer



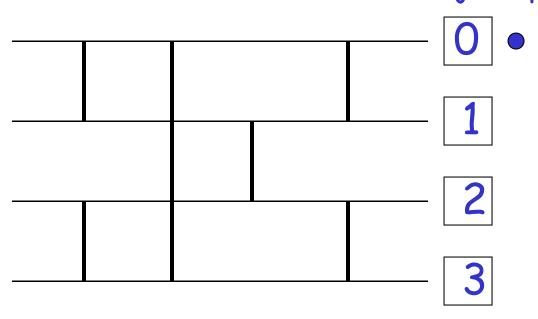
Paralelizam

Mnogi zahtevi za inkrement se obrađuju simultano

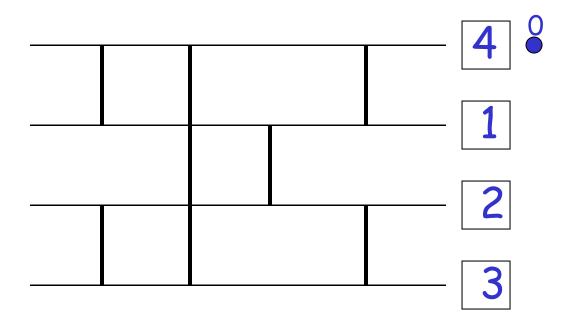
Brojanje Izlaz Deljene promenljive

Token = Zahtev za inkrement

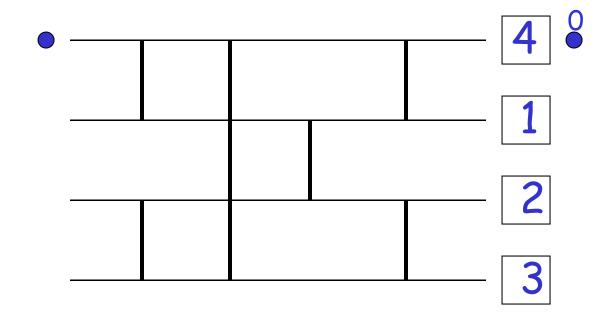
Izlaz Deljena promenljive



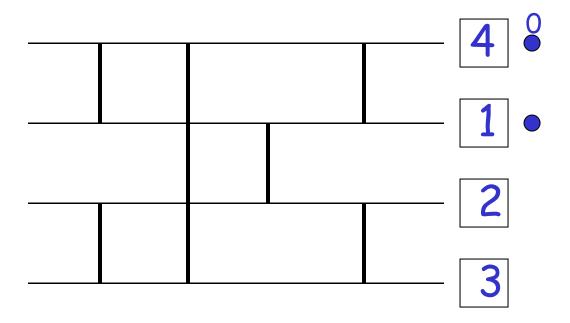
Povratna vrednost

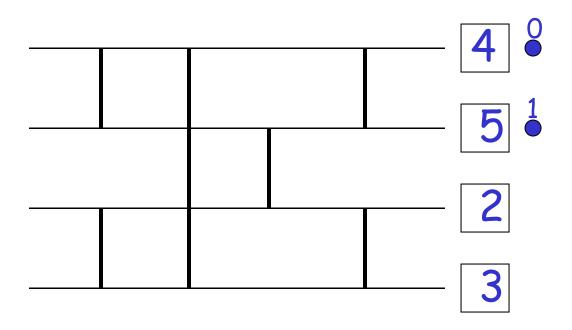


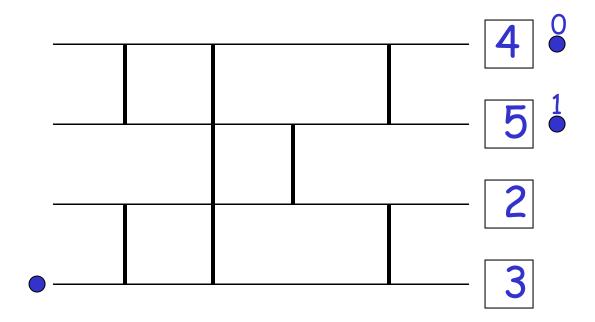
Vrednost deljene promenljive se poveća za 4 (širina izlaza)



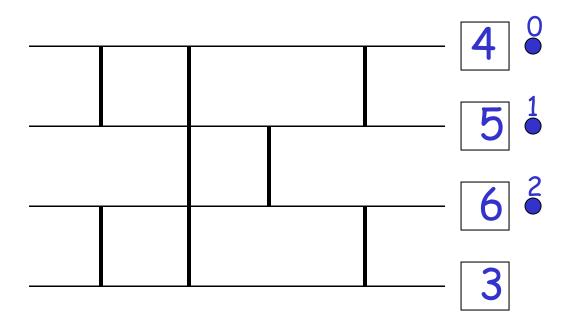
Zahtev za inkrement

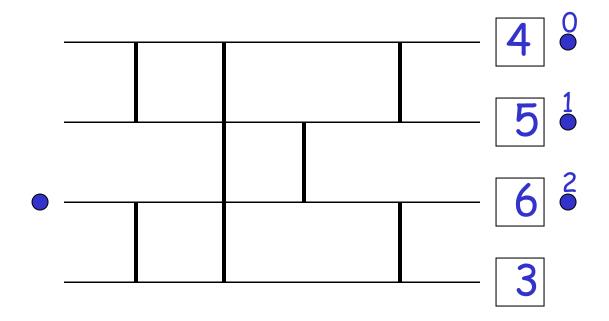


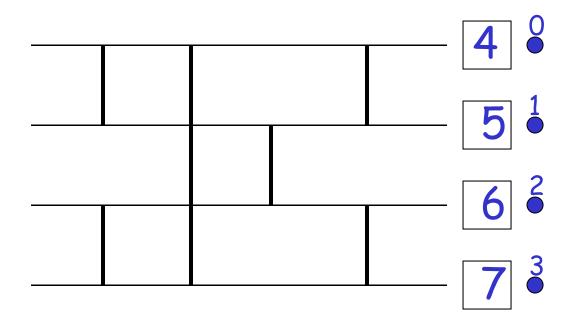


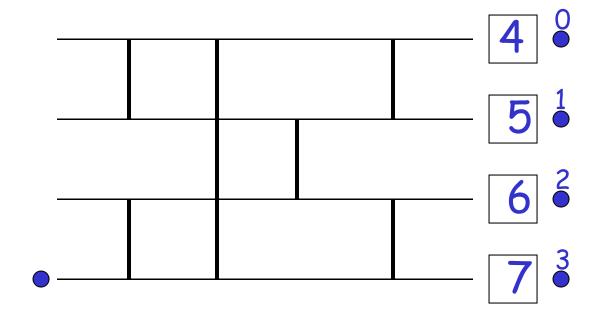


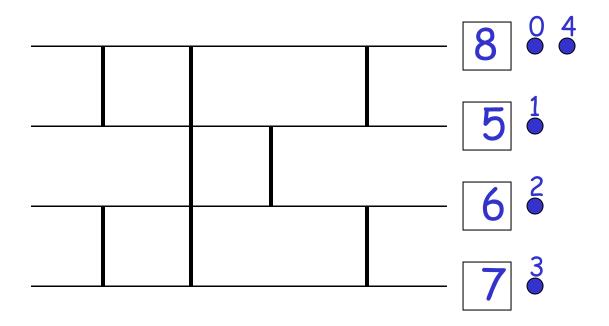
Zahtev za inkrement

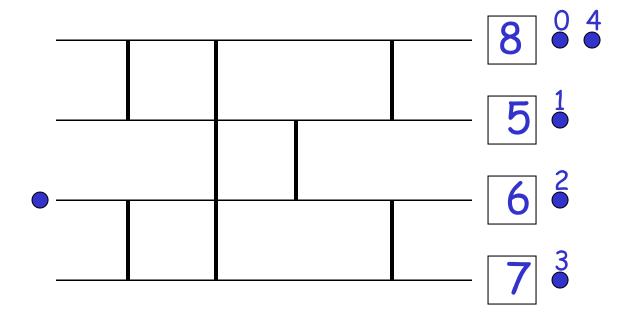


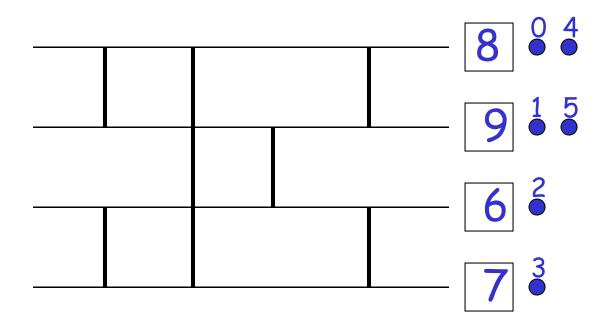


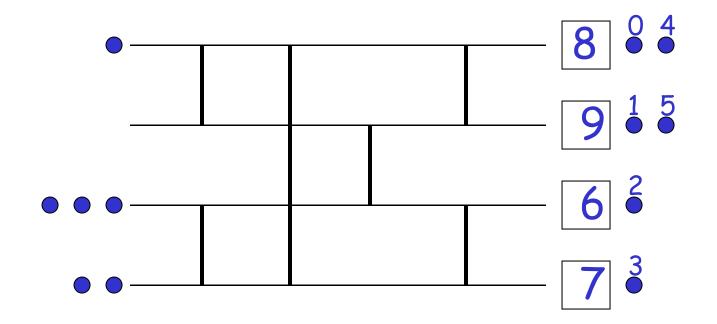










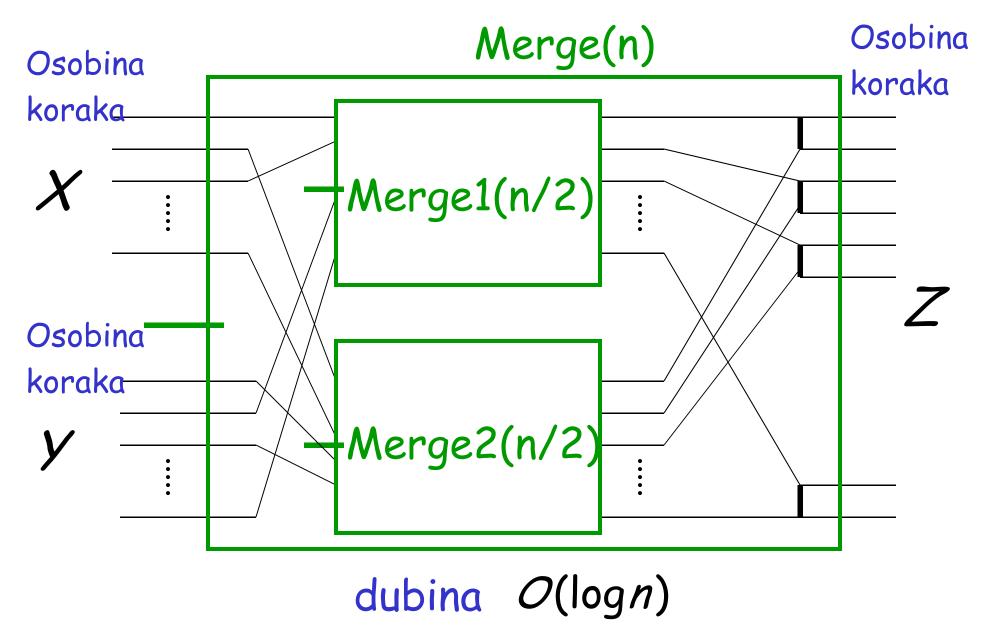


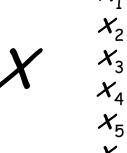
Svi tokeni

Mreža za bitonik brojanje

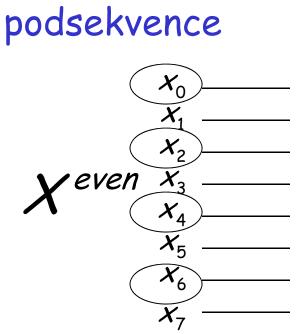
Izomorfna sa Mrežom za bitonik sortiranje

Count(n) Count(n/2) Ulaz Merge(n) Izlaz Count(n/2) $O(\log^2 n)$ dubina

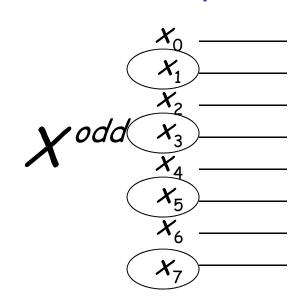


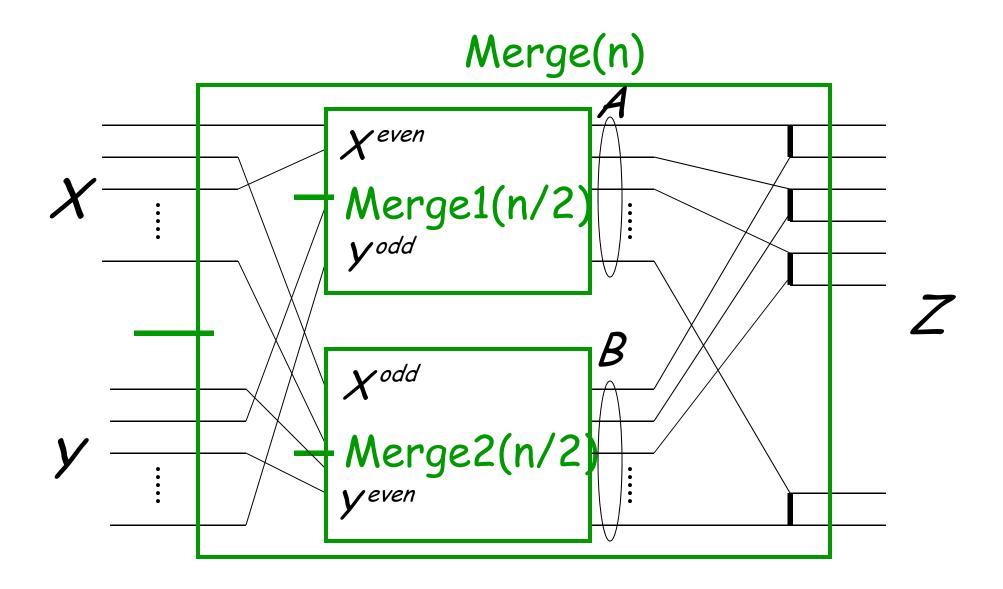


Neparne podsekvence



Parne





Teorema:

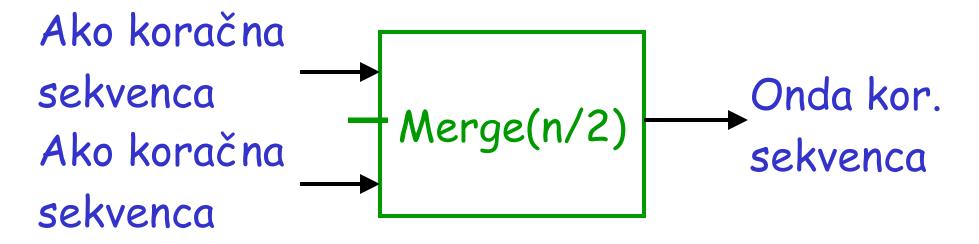
merger(n) proizvodi izlaz Z sa osobinom koraka ako oba ulaza X i Y imaju osobnu koraka

Dokaz:

Dokaz je pomoću indukcije na nOsnova indukcije:

Za n=2 spajač je balanser i tvrdnja je trivijalno zadovoljena

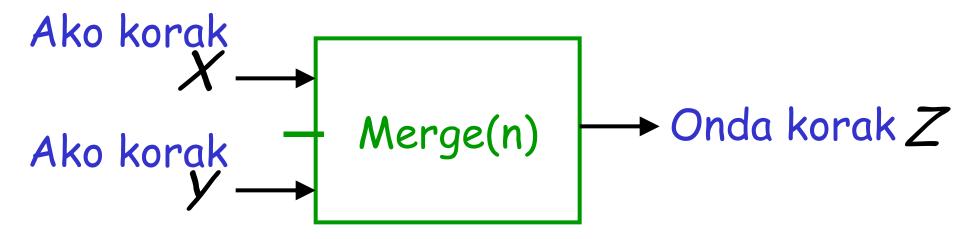
Indukciona hipoteza:



Predpost. da svaki spajač veličine n/2 i manje ispravno obavlja spajanje

Indukcioni korak:

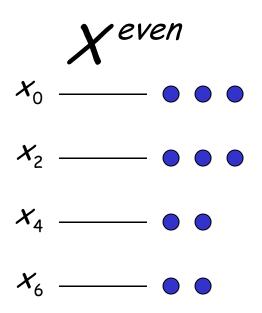
Želimo da pokažemo da je:

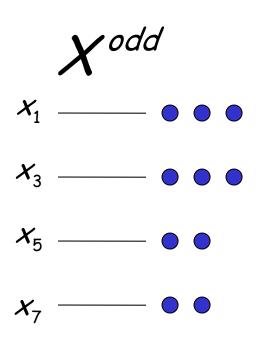


Predpost. da X i Y imaju osobinu koraka. Onda ćemo pokazati da Z ima osobinu koraka

X_0 ... X_0 ... X_1 ... X_2 ... X_2 ... X_3 ... X_4 ... X_5 ... X_6 ... X_6 ... X_7 ...

Ako X ima osobinu koraka onda X^{even} i X^{odd} imaju osobinu koraka





Zato, iz indukcione hipoteze:





Prvo, pokažimo: $-1 \le |A| - |B| \le 1$

Gde | A | označava ukupan broj tokena u sekevnci A

$$|A| = |X^{even}| + |Y^{odd}|$$

Spajač 2
$$|B| = |Y^{even}| + |X^{odd}|$$

Pošto X ima osobinu koraka:

$$|X^{even}| = |X^{odd}|$$

$$|X^{even}| = |X^{odd}| + 1$$

$$X_0$$
 X_1
 X_2
 X_3
 X_4
 X_5
 X_6
 X_7

$$X_0$$
 X_1
 X_2
 X_3
 X_4
 X_5
 X_6
 X_7

Zato je:
$$0 \le |X^{even}| - |X^{odd}| \le 1$$

Slično:
$$0 \le |y^{even}| - |y^{odd}| \le 1$$

$$0 \leq |X^{even}| - |X^{odd}| \leq 1$$

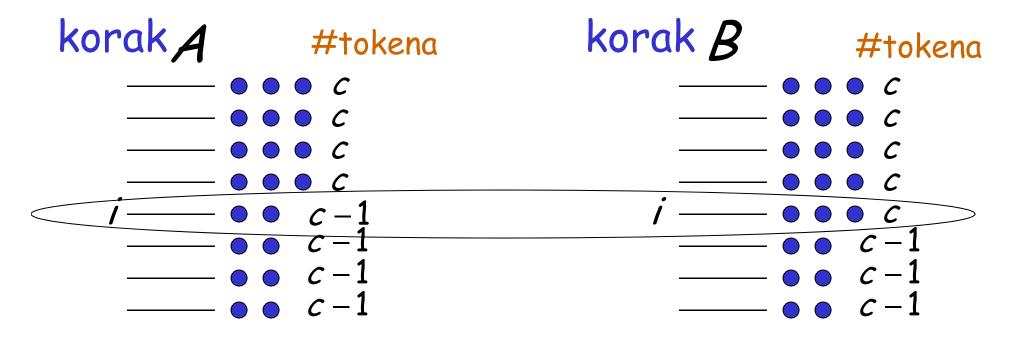
$$0 \leq |Y^{even}| - |Y^{odd}| \leq 1$$

$$|A| - |B| = (|X^{even}| - |X^{odd}|) + (|Y^{odd}| - |Y^{even}|)$$

$$-1 \le |A| - |B| \le 1$$

Sada, pokažimo da Z ima osobinu koraka

$$-1 \le |A| - |B| \le 1$$



Postoji najviše jedna žica i gde se dve sekvence razlikuju

