

Rank Annotated Trees

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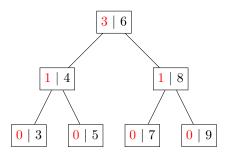
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- Proofs

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Inorder traversal and getting rank Definition of the function Proof of correctness

Definition



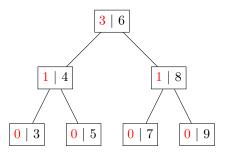
(Type definition)

datatype 'a rtree = Leaf | Node "'a rtree" nat 'a "'a rtree"

Example 1

 $\langle \langle \langle \langle \rangle, 0, 3, \langle \rangle \rangle, 1, 4, \langle \rangle \rangle, 3, 6 :: nat, \langle \langle \langle \rangle, 0, 7, \langle \rangle \rangle, 1, 8, \langle \langle \rangle, 0, 9, \langle \rangle \rangle \rangle \rangle$

Definition



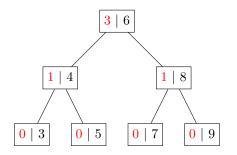
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Example :

 $\langle\langle\langle\langle\rangle,0,3,\langle\rangle\rangle,1,4,\langle\rangle\rangle,3,6::$ nat, $\langle\langle\langle\rangle,0,7,\langle\rangle\rangle,1,8,\langle\langle\rangle,0,9,\langle\rangle\rangle\rangle\rangle$

Definition



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Example 1

 $\langle\langle\langle\langle\rangle,0,3,\langle\rangle\rangle,1,4,\langle\rangle\rangle,3,6::\mathtt{nat},\langle\langle\langle\rangle,0,7,\langle\rangle\rangle,1,8,\langle\langle\rangle,0,9,\langle\rangle\rangle\rangle\rangle$

First functions

```
fun num_nodes :: "'a rtree \Rightarrow nat" where "num_nodes \langle \rangle = 0" | "num_nodes \langle 1, _, _, r\rangle = 1 + num_nodes 1 + num_nodes r"
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```
fun rbst :: "('a::linorder) rtree ⇒ bool" where
   "rbst ⟨⟩ = True" |
   "rbst ⟨ l, n, x, x⟩ = ((∀a ∈ set_rtree l. a < x) ∧
        (∀a ∈ set_rtree r. x < a) ∧
        rbst l ∧
        rbst r ∧
        n = num_nodes l)"</pre>
```

First lemmas

Some useful lemmas

```
lemma set rtree rbst:
```

"rbst $\langle 1, n, x, r \rangle \Rightarrow a \in \text{set_rtree } \langle 1, n, x, r \rangle \Rightarrow a < x \Rightarrow a \in \text{set rtree } 1$ "

lemma rins_set: "set_rtree (rins x t) = insert x (set_rtree t)"

lemma num nodes rins notin:

"x \notin set_rtree t \Rightarrow rbst t \Rightarrow num_nodes (rins x t) = 1 + num_nodes t"

lemma rins_invar: "x \notin set_rtree t \Rightarrow rbst t \Rightarrow rbst (rins x t)"

```
lemma set_rtree_rbst:

"rbst \langle 1, n, x, r \rangle \Rightarrow a \in set_rtree \langle 1, n, x, r \rangle \Rightarrow a < x \Rightarrow a \in set_rtree 1"
```

```
lemma rins_set: "set_rtree (rins x t) = insert x (set_rtree t)"
```

```
lemma num_nodes_rins_notin:

"x \notin \text{set\_rtree } t \Rightarrow \text{rbst } t \Rightarrow \text{num\_nodes } (\text{rins } x \ t) = 1 + \text{num\_nodes } t"
```

```
lemma rins_invar: "x \notin set_rtree t \Rightarrow rbst t \Rightarrow rbst (rins x t)"
```

```
lemma set_rtree_rbst:  
"rbst \langle1, n, x, r\rangle \Rightarrow a \in set_rtree \langle 1, n, x, r\rangle \Rightarrow a < x \Rightarrow a \in set_rtree 1"
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lemma rins_set: "set_rtree (rins x t) = insert x (set_rtree t)"
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lemma rins_invar: "x \notin set_rtree t \Rightarrow rbst t \Rightarrow rbst (rins x t)"
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lemma set_rtree_rbst:

"rbst \langle 1, n, x, r \rangle \Rightarrow a \in set_rtree \langle 1, n, x, r \rangle \Rightarrow a < x \Rightarrow a \in set_rtree l"
```

```
lemma rins_set: "set_rtree (rins x t) = insert x (set_rtree t)"
```

```
lemma num_nodes_rins_notin:
```

```
"x \notin set_rtree t \Rightarrow rbst t \Rightarrow num_nodes (rins x t) = 1 + num_nodes t"
```

```
\texttt{lemma rins\_invar:} \quad \texttt{"x} \notin \texttt{set\_rtree t} \Rightarrow \texttt{rbst t} \Rightarrow \texttt{rbst (rins x t)"}
```

Inorder traversal and getting rank

- Tree traversal: inorder function (in-order DFS)
- Getting rank: rank function w.r.t. the structure of the rank annotated tree

```
fun rank:: "'a::linorder \Rightarrow 'a rtree \Rightarrow nat" where

"rank a \langle \rangle = 0" |

"rank a \langle  1, n, x, r\rangle =

(if a = x then n

else if a > x then 1 + n + rank a r

else rank a 1)"
```

• Selection : select function

- Tree traversal: inorder function (in-order DFS)
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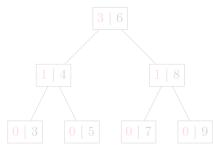
(if a = x then n
else if a > x then 1 + n + rank a r
else rank a 1)"
```

Selection: select function

(Selection)

```
fun sel:: "nat ⇒ 'a::linorder rtree ⇒ 'a" where
   "sel _{-}\langle\rangle = undefined" |
   "sel i \langle 1, n, x, r \rangle =
       (if i = n then x
       else if i < n then sel i l
       else sel (i - n - 1) r)"
```

inorder
$$t = [3, 4, 5, 6, \frac{7, 8, 9}{r}]$$



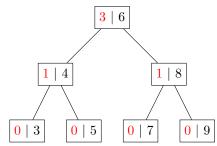
Secure Coding

(Selection)

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sel 4 t

inorder t =
$$\underbrace{[3, 4, 5, 6, \frac{7, 8, 9]}_{\text{offset: n + 1}}}_{\text{offset: n + 1}}$$



lemma select correct:

```
"rbst t \implies i \le length (inorder t) \implies sel i t = inorder t!i"
```

Idea of the proof:

- By induction on t with i arbitrary
- Cases are split w.r.t. the body of the function sel
- In the third case, we show:
 - o sel i $\langle 1, n, x, r \rangle$ = sel (i n 1) r
 - o sel (i n 1) r = inorder r!(i n 1)

☐ Proof of correctness

```
lemma rank_sel_id:  
"rbst t \implies i < length (inorder t) \implies rank (sel i t) t = i"
```

Idea of the proof:

- By induction on t with i arbitrary
- Trivialized with lemma select_correct