- 1 Easy to implement
- 1 Memory efficient
- (3) Token bucket allows a bount of traffic for short periods. A requeron can go through as long as there are tokens left.

Cons

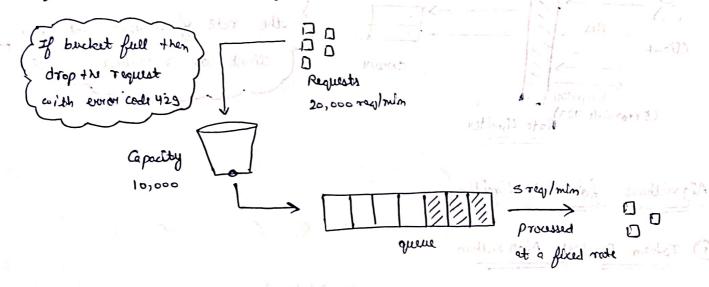
Two parameters: Bucket size & Refall rate

might be challenging to tune them properly.

Pate Ministra

## 2 Leaky Bucket Algorithm

Similar to the token bucket except that requests are processed at a fixed rate. It is usually implemented with a FIFO queue.



Postdore & distinct

#### Pros :-

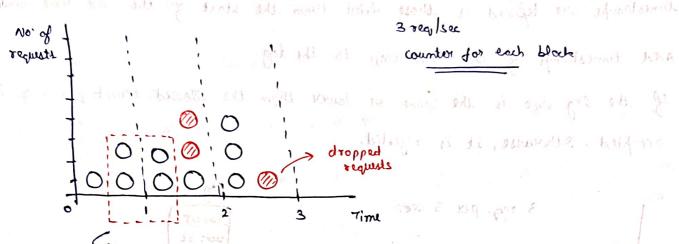
- 1) Memory efficient given the Umbted queue size
- ② Requests are processed at a fixed rate therefore it is suitable for we call that a stable outflow rate is needed

#### Cons

- 1) A burst of traffic fells up the queue with old requand of they are not processed in time, recent requests will be rate limited.
  - 2 Difficult to tune bucket stre 2 veg. processing rate

## 3 Fixed Window Counter Algorithm

- · The algo. divides the timeline into fix-sized time windows & assign a counter for each window.
- · Each req. increments the counter by one
- · Once the counter reacher the pre-defined threshold, new reg, are dropped until a new time window starts.



You are saying that for every second there should only be atmost 3 requests, but if we look at this timeframe (0.5sec-1.5sec)  $\rightarrow$  it is hendling 4 requests in a second.

Problem with this Algorithm

### Pros !-

3

3

- 1 Memory efficient
- 2 Easy to understand
- (3) Resetting available quota at the end of a unit time window fits certain use cashs

Spike in traffic at the edges of a window could cause more requests than the allowed quota to go through.

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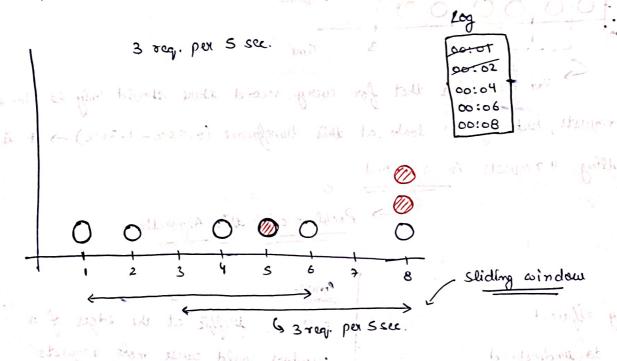
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# 4) Slidling Window Log Algorithm

S fixes the issue of fixed window counter algorithm

- · The algo. Keeps track of req. timestamps. Timestamp data is usually kept in cache such as souled sets of Reolis.
- When a new req. comes in, remove all the outdated timestamps. Outdated timestamps are defined as those older than the start of the cour. time window
- · Add timestamp of the new reg. to the dog.
- . If the log size is the same or lower than the allowed count, a reg. is accepted. Otherwise, it is rejected.



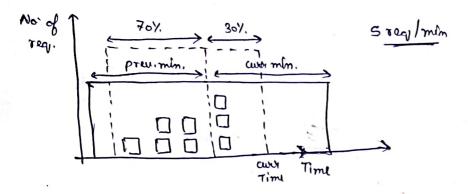
Prosett of the other produce it most.

Rate limiting implemented by this algo. is very accurate. In any rolling window, req. will not exceed the rate limit.

The algo. consumer a lot of memory because even if a req. is rejected, it's timestamp might still be stored in memory.

(5) Slicking Window Counter Algorithm

Shybrid approach of fixed window counter & sliding window log



Assume the rate limiter allows a max. of 7 reg/min & there are 5 reg. in the prev. minute & 3 in the curr. min. For a new reg. that arriver at a 30%. position in the curr. min., the no of reg. in the rolling window is :-

Reg. in the cuts-window + Reg. in the prev. window \* overlap y. of the rolling window & prev. window

=> 3+5\*0.7% = 6.5 request (so this can go through the rate limited as
the limit is 7 reg/min)