

1) Total no of paths in matrix if right and down move allowed

<https://leetcode.com/problems/unique-paths/submissions/>

[https://leetcode.com/problems/unique-paths/discuss/22958/Math-solution-O\(1\)-space](https://leetcode.com/problems/unique-paths/discuss/22958/Math-solution-O(1)-space)

<https://leetcode.com/problems/unique-paths-ii/>(with some obstacles)

$$Dp[i][j] = dp[i-1][j] + dp[i][j-1]$$

Minimum path sum in matrix

a) <https://leetcode.com/problems/minimum-path-sum/>

b) $dp[i][j] = a[i][j] + \min(dp[i-1][j], dp[i][j-1])$

2) Min stair cost

<https://leetcode.com/problems/min-cost-climbing-stairs/>

Each stair has given cost find min cost to reach at top

Input: cost = [1, 100, 1, 1, 1, 100, 1, 1, 100, 1]

Output: 6

$$dp[i] = cost[i] + \min(dp[i-1], dp[i-2])$$

3) Is subsequence

<https://leetcode.com/problems/is-subsequence/>

a) Using dp

i) $dp[i][j] = dp[i-1][j-1]$ if $s[i] = t[j]$
= false else

If one cell is true make entire row true

s\t	h	a	x	b	c
a	F	T	T	T	T
b	F	F	F	T	T
c	F	F	F	F	T

b) Using two pointer

i) Scan s by i and inc j when $s[i] == t[j]$
if $j == t.length()$ found string

4) Robbing house

a) Each house has some value indicated by array element adjacent house can not be robbed find max amount that can be robbed

b) $dp[i] = \max(dp[i-1], \text{value}[i] + dp[i-2])$
Don't rob current house Rob current house

c) Or using inc exc property

d) 2, 1, 1, 2 => (2 + 2 = 4)
1, 2, 3, 4 => (2 + 4 = 6)

e) Circular Robbing house

i) <https://leetcode.com/problems/house-robber-ii/submissions/>

ii) $\text{return max(robhouse(nums, 0, nums.size()-2), robhouse(nums, 1, nums.size()-1))};$

5) Longest valid Parentheses

<https://leetcode.com/problems/longest-valid-parentheses/solution/>

()	(())
0	2	0	0	2	6

```
If s[i] == '('
    Dp[i] = 0
Else
    If s[i] == ')'
        S[i] = s[i-2] + 2
    Else s[i-1] == '('
        if( s[i - dp[i-1] - 1] == '(' )
            dp[i] = dp[i-1] + 2 + dp[i-dp[i-1]-2];
        else
            dp[i] = 0;
```

6) Total no of unique binary search tree

[https://leetcode.com/problems/unique-binary-search-trees/discuss/31666/DP-Solution-in-6-lines-with-explanation.-F\(i-n\)-G\(i-1\)-*-G\(n-i\)](https://leetcode.com/problems/unique-binary-search-trees/discuss/31666/DP-Solution-in-6-lines-with-explanation.-F(i-n)-G(i-1)-*-G(n-i))

Series : 1 2 5 14 42 132

$G(n) = F(1, n) + F(2, n) + \dots + F(n, n)$. // possible roots

$G(n) = G(0) * G(n-1) + G(1) * G(n-2) + \dots + G(n-1) * G(0)$

$G(n) = \sum_{x=1}^n G(x-1) * G(n-x)$ where $1 \leq x \leq n$

Eg : 1 2 3 4 5 6 if 3 is root left side 2 right side 3 element 1,2,3,4,5,6

$G[0] = G[1] = 1$;

```
for(int i=2; i<=n; ++i) {
    for(int j=1; j<=i; ++j) {
        G[i] += G[j-1] * G[i-j];
    }
}
return G[n];
```

7) No of ones in binary representation of number

- $f[i] = f[i >> 1] + (i \& 1)$;
- Explanation.
- Take number X for example, 10011001.
- Divide it in 2 parts:
- <1>the last digit (1 or 0, which is " $i \& 1$ ", equivalent to " $i \% 2$ ")

f) \leq the other digits (the number of 1, which is " f[i] > 1 ", equivalent to " f[i/2] ")

8) Edit distance problem

- Min no of operation to convert word1 to word2
- $dp[i][j] = dp[i-1][j-1]$ if word1[i-1] = word2[j-1]
- $\min(dp[i-1][j], dp[i][j-1], dp[i-1][j-1]) + 1$ else

[https://leetcode.com/problems/edit-distance/discuss/25846/C%2B%2B-O\(n\)-space-DP](https://leetcode.com/problems/edit-distance/discuss/25846/C%2B%2B-O(n)-space-DP)

In above link 3 approaches to solve dp using $O(n^2)$ $O(2n)$ and $O(n)$

9) 0/1 Knapsack Problem

- <https://practice.geeksforgeeks.org/problems/0-1-knapsack-problem/0>
- Here limited supply of item is given so
 if (i==0 || w==0)
 K[i][w] = 0;
 else if (wt[i-1] <= w)
 K[i][w] = max(val[i-1] + K[i-1][w-wt[i-1]], K[i-1][w]);
 else
 K[i][w] = K[i-1][w];

10) Coin change problem

- <https://practice.geeksforgeeks.org/problems/coin-change/0>

	0	1	2	3	4
0 0	1	0	0	0	0
1 1	1	1	1	1	1
2 2	1	1	2	2	3
3 3	1	1	2	3	4

```

if(a[i-1] > j)
    dp[i][j] = dp[i-1][j];
else
    dp[i][j] = dp[i-1][j] + dp[i][j-a[i-1]];
  
```

- <https://leetcode.com/problems/coin-change/submissions/>

	0	1	2	3	4
0 0	0	INF	INF	INF	INF
1 1	0	1	2	3	4

2	2	0	1	1	2	2
3	3	0	1	2	1	2

```

if(a[i-1] > j)
    dp[i][j] = dp[i-1][j];
else
    dp[i][j] = min( dp[i-1][j] , dp[i][j-a[i-1]] + 1); take INT = 99999 to avoid overflow

```

- c) <https://leetcode.com/problems/combination-sum-iv/>
 Find all combinations of numbers having sum = target
 Use backtracking approach and convert it to top down and bottom up approach

11) Longest Inc subsequence nlogn

a) LCS

[https://leetcode.com/problems/longest-increasing-subsequence/discuss/74848/9-lines-C%2B%2B-code-with-O\(NlogN\)-complexity](https://leetcode.com/problems/longest-increasing-subsequence/discuss/74848/9-lines-C%2B%2B-code-with-O(NlogN)-complexity)

```

vector<int> res;
for(int i=0; i<nums.size(); i++) {
    int ind = lower_bound(res.begin(), res.end(), nums[i]) - res.begin();
    if(ind == res.size()) res.push_back(nums[i]);
    else res[ind] = nums[i];
}
return res.size();

```

b) Stack Book

<https://cses.fi/problemset/task/1073/>

```

for(int i=0;i<n;i++){
    int index = upper_bound(dp.begin(), dp.end(), a[i]) - dp.begin();
    if(index == dp.size())
        dp.push_back(a[i]);
    else
        dp[index] = a[i];
}
cout << dp.size() <<en

```

c) Maximum nested doll possible

- <https://leetcode.com/problems/russian-doll-envelopes/>
- Sort by width asc and when equal des by height
- Apply lis on height
- Reason : (6, 1), (6,2), (6, 3) here if we don't sort by height will get ans = 3 so.

d) Maximum no of longest Increasing Subsequence

e. <https://leetcode.com/problems/number-of-longest-increasing-subsequence/>

```
if (nums[i] > nums[j]) {
    if (len[j]+1 > len[i]) {
        len[i] = len[j]+1;
        cnt[i] = cnt[j];
    }
    else if (len[j]+1 == len[i])
        cnt[i] += cnt[j];
}
```

12) Palindrome

a) Longest Palindromic subsequence

i) <https://leetcode.com/problems/longest-palindromic-subsequence/submissions/>

```
if(s[i] == s[j])
    dp[i][j] = dp[i+1][j-1] + 2;
else
    dp[i][j] = max(dp[i][j]-1, dp[i+1][j]);
```

ii)

(1) A X B B A B

A	1	1	1	2	4	4
B	0	1	1	2	2	3
X	0	0	1	2	2	3
B	0	0	0	1	1	3
A	0	0	0	0	1	1
B	0	0	0	0	0	1

b) Longest Palindromic Substring :

i) For each character traverse both side and find odd len palindrome

ii) Similarly traverse both side and find even len palindrome

iii) <https://leetcode.com/problems/longest-palindromic-substring/>

```
if s[i] == s[j]
    dp[i][j] = dp[i+1][j-1]
Else
    dp[i][j] = false
```

A B A A B A

A	1	0	1	0	0	1
B		1	0	0	1	0
A			1	1	0	0
A				1	0	1
B					1	0
A						1

c) Count palindromes

i) <https://practice.geeksforgeeks.org/problems/count-palindrome-sub-strings-of-a-string/0>

```
s = '@' + s + '#';
int n = s.size();
for(int mid = 1; mid <= n-1; mid++){
    i = mid-1, j = mid+1;

    while(s[i] == s[j])
        count++, i--, j++;

    i = mid-1, j = mid;
    while(s[i] == s[j])
        count++, i--, j++;
}
```

```

}
cout << count << endl;

```

d) Min insertion to make string palindrome

- i) <https://www.geeksforgeeks.org/minimum-insertions-to-form-a-palindrome-dp-28/>
- ii) $table[l][h] = (str[l] == str[h]) ? table[l + 1][h - 1] : (\min(table[l][h - 1], table[l + 1][h]) + 1);$
- iii) Find longest palindromic subsequence. Ans = Total length-length of LPS

e) <https://www.geeksforgeeks.org/minimum-number-deletions-make-string-palindrome/>

- i) $table[l][h] = (str[l] == str[h]) ? table[l + 1][h - 1] : (\min(table[l][h - 1], table[l + 1][h]) + 1);$
- ii) Find longest palindromic subsequence. Ans = Total length-length of LPS

f) <https://www.geeksforgeeks.org/find-if-string-is-k-palindrome-or-not/>

Find longest palindromic subsequence. Check length-length of LPS $\leq K$

g) Count Pal substring in range i to j

- i) <https://www.geeksforgeeks.org/count-of-palindromic-substrings-in-an-index-range/>
- ii) <https://www.geeksforgeeks.org/count-palindrome-sub-strings-string/>

13) Longest Repeating Subsequence

a) <https://www.geeksforgeeks.org/longest-repeating-subsequence/>

b) Using dp

```

if (str[i-1] == str[j-1] && i != j)
    dp[i][j] = 1 + dp[i-1][j-1];
else
    dp[i][j] = max(dp[i][j-1], dp[i-1][j]);

```

	a	b	b	a
0	0	0	0	0
a	0	0	1	1
b	0	1	1	1
c	0	1	1	2
d	0	1	1	2

c) Without using dp $O(n)$ for check existence of repeating subsequence

- i) <https://www.interviewbit.com/problems/repeating-subsequence/>
- ii) If any char occurs more than 2 times return true
Now make a string where each char repeats 2 times
If this string is palindrome then then return 0
Else return 1

14) Number of Distinct subsequences t in s

a) <https://leetcode.com/problems/distinct-subsequences/>

- b) if ($t[i-1] == s[j-1]$)
 $dp[i][j] = dp[i-1][j-1] + dp[i][j-1];$
else
 $dp[i][j] = dp[i][j-1];$

c)

	B	A	B	G	B	A	G
1	1	1	1	1	1	1	1
B	0	1	1	2	2	3	3
A	0	0	1	1	1	1	4
G	0	0	0	0	1	1	5

15) Longest inc dec sequence

a) <https://www.interviewbit.com/problems/length-of-longest-subsequence/>

b)

- i) `inc[i]` stores Longest Increasing subsequence ending with `A[i]`
- ii) `dec[i]` stores Longest Decreasing subsequence ending with `A[i]`
- iii) Now we need to find the maximum value of `(inc[i] + dec[i] - 1)`

16) Largest Rectangle with 1's , swapping of columns allowed

a) <https://www.interviewbit.com/problems/largest-area-of-rectangle-with-permutations/>

17) Decode Ways

a) <https://leetcode.com/problems/decode-ways/submissions/>

b)

```
if(x < 27)
    dp[i] = dp[i+1] + dp[i+2];
else
    dp[i] = dp[i+1];
```

18) Interleaving Strings

a) <https://www.interviewbit.com/problems/interleaving-strings/>

b)

```
if(i >= 0 && A[i] == C[k])
    a = isInterleaving(A,B,C,i-1,j,k-1);

if(j >= 0 && B[j] == C[k])
    b = isInterleaving(A,B,C,i,j-1,k-1);

Return a || b
```

19) Wild card character

a) Use map <pair,bool> of dp to avoid memory limit error

b) <https://www.interviewbit.com/problems/regular-expression-match/> (both top down and bottom up)

c)

```
if(text[i-1] == pat[j-1] || pat[j-1] == '?')
    dp[i][j] = dp[i-1][j-1];
else if(pat[j-1] == '*')
    dp[i][j] = dp[i-1][j] || dp[i][j-1];
else
    dp[i][j] = false;
```

20) Regular Expression

a) <https://www.interviewbit.com/problems/regular-expression-ii/> (both top down and bottom up)

b) Top down approach

```
if(text[i-1] == pat[j-1] || pat[j-1] == '.')
    dp[i][j] = dp[i-1][j-1];

else if(pat[j-1] == '*')
{
```

```

        if(pat[j-2] == text[i-1] || pat[j-2] == '.')
            dp[i][j] = dp[i][j-2] || dp[i-1][j-2] || dp[i-1][j];
        Else
            dp[i][j] = dp[i][j-2];
    }

```

21) Scramble String

- a) <https://www.interviewbit.com/problems/scramble-string/>
 b)

```

        if(Scramble(A.substr(0,i),B.substr(0,i)) && Scramble(A.substr(i),B.substr(i)) ||
           Scramble(A.substr(0,i),B.substr(n-i)) && Scramble(A.substr(i), B.substr(0,n-i)))
            return mp[A+B] = true;

```

22) Longest Arithmetic Sequence

- a) <https://leetcode.com/problems/longest-arithmetic-sequence/submissions/>
 b) A,B,C are in AP then $A = 2B - C$
 i) $dp[i][j]$ no of items in arithmetic series ends with $a[i], a[j]$
 ii) $int\ target = 2*a[i] - a[j];$
 iii) $if(index.find(target) == index.end())$
 iv) $dp[i][j] = 2;$
 v) else
 vi) $dp[i][j] = dp[index[target]][i] + 1;$

23) Min palindrome partitioning

- a) <https://www.interviewbit.com/problems/palindrome-partitioning-ii/> (bottom up) $O(n^3)$

```

        i) if(ispal[i][j])
            dp[i][j] = 0;
        else{
            for(int k=i;k<j;k++){
                dp[i][j] = min(dp[i][j], dp[i][k] + dp[k+1][j] + 1);
            }
        }

```

- b) $abcb \Rightarrow a | bcb, ab | cb, abc | b$

- c) <https://leetcode.com/submissions/detail/314502291/> (top down)

```

        i) for(int i=start;i<end;i++){
        ii)     ans = min(ans, 1+findMincut(A,start,i) + findMincut(A,i+1,end));
        iii)    }

```

- d) <https://leetcode.com/problems/palindrome-partitioning-ii/submissions/> $O(n^2)$

24) Word Break Problem

- a) <https://leetcode.com/submissions/detail/314851563/> (top down)

i) `for(int i=1;i<=maxlen;i++)`

`if(dict.find(s.substr(start,i)) != dict.end() && canBreak(s,start+i))`

`return true;`

- b) <https://leetcode.com/submissions/detail/314912714/> (bottom up) $O(n^3)$

i)

`if(dict.find(s.substr(i,len))!=dict.end())`

`dp[i][j] = true;`

`else{`

`for(int k=i;k<j;k++){`

`dp[i][j] = dp[i][k] && dp[k+1][j];`

`if(dp[i][j]) break;`

`}`

`}`

- c) $O(n^2)$ <https://leetcode.com/problems/word-break/discuss/43790/Java-implementation-using-DP-in-two-ways>

25) Word break 2

- a) <https://www.interviewbit.com/problems/word-break-ii/>

b) $Dp[i]$ = string of vector of solution from 0 to i

c) For i in 0 to n-1

i) For j in 0 to i

`Stg = stg(j , i)`

`If (dict contains stg and $dp[j-1] > 0$)`

`For x in $dp[j-1]$`

`$Dp[i] += stg + ' ' + x$`

26) Dungeon Princesses

- a) <https://www.interviewbit.com/problems/dungeon-princess/>

27) Buy and Sell

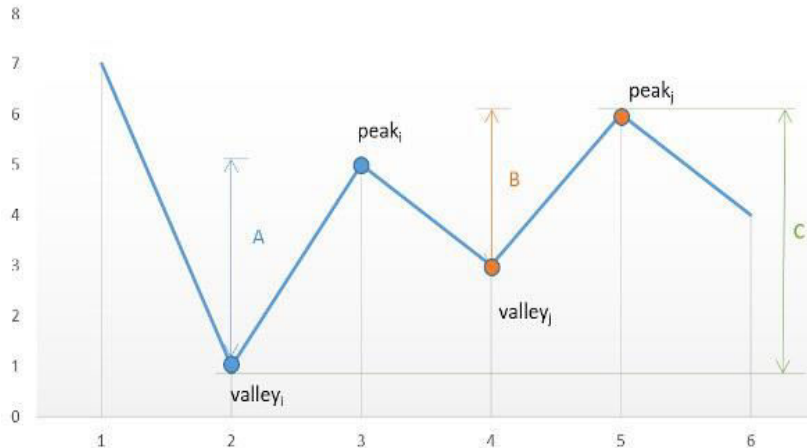
- a) Any no of times

<https://www.interviewbit.com/problems/best-time-to-buy-and-sell-stocks-ii/>

Option 1: Buy stock at valley[i] then sell at peak[i] makes profit A ($peak[i] - valley[i]$). Then buy stock at valley [j] and sell at peak[j] makes profit B ($peak[j] - valley[j]$). So the total profit of this trade option is A + B.

Option 2: Skip the intermediate trades, i.e., we buy stock at valley[i] then sell at peak[j]. In this case, the total profit will be C ($peak[j]-valley[i]$).

Based on the graph shown below, $A + B > C$ (if not, $\text{peak}[i]$ and $\text{valley}[j]$ won't exist). So in order to maximize the profit, we can buy stock at valleys and then sell stock at peaks.



b)

c) At Most two

i) <https://leetcode.com/problems/best-time-to-buy-and-sell-stock-iii/>

d) Exact k transactions

i) <https://leetcode.com/problems/best-time-to-buy-and-sell-stock-iv/>

ii) Either don't do a transaction on jth day or do transaction with maximum profit

iii) If $k \geq n/2$ then maxprofit (a)

Else

$\text{dp}[i][j] = \max(\text{dp}[i][j-1]);$ // i transactions j days

$\max(\text{prices}[j] - \text{prices}[m] + \text{dp}[i-1][m]); 0 \leq m < j$

iv) Optimization using maxvalue refer link

e) <https://leetcode.com/problems/best-time-to-buy-and-sell-stock-with-transaction-fee/discuss/108870/Most-consistent-ways-of-dealing-with-the-series-of-stock-problems>

i) <https://leetcode.com/problems/best-time-to-buy-and-sell-stock-with-cooldown>

ii) <https://leetcode.com/problems/best-time-to-buy-and-sell-stock-with-transaction-fee>

28) Maximum Product subarray

a) www.interviewbit.com/problems/max-product-subarray/

b) $\text{Dpmax}[i] = \text{max subarray ending at index } i$

c) $\text{Dpmin}[i] = \text{min subarray ending at index } i$

d) $\text{Dpmax} = \max(A[i], \text{dpmax} * A[i], \text{dpmin} * A[i])$

e) $\text{Dpmax} = \min(A[i], \text{dpmax} * A[i], \text{dpmin} * A[i])$

29) Smallest k no whose prime factors are only Primes p1 p2 and p3

a) <http://www.interviewbit.com/problems/smallest-sequence-with-given-primes/>

```
dp[0] = 1;
for(int i=1; i<D+1; i++){
    int mine = min(dp[p1]*A, min(dp[p2]*B, dp[p3]*C));
    if(mine == dp[p1]*A) p1++;
    if(mine == dp[p2]*B) p2++;
    if(mine == dp[p3]*C) p3++;
    dp[i] = mine;
}
```

30) Flip min no so resultant array is min no negative

a) <https://www.interviewbit.com/problems/flip-array/>

b) $S1 + s2 = \text{sum}$

$$S1 - s2 = 0$$

$$\text{So } s2 = \text{sum}/2$$

Using knapsack find no of items with max target possible