

## 1. Stack using Queue

a. <https://leetcode.com/problems/implement-stack-using-queues/solution/>

b. Push =  $O(1)$  and pop =  $O(n)$

i. Insertion order : 1 2 3 4

Stack : 1 2 3 4 - top

Queue :

Push  
1  
1 2  
1 2 3  
F 1 2 3 4 R

Pop  
F 1 2 3 4 R  
F 4 1 2 3 R Delete and insert 3 times  
F 1 2 3 R Delete one time , returns 4

Pop  
F 1 2 3 R  
F 3 1 2 R Delete and insert 2 times  
F 1 2 R Delete one time , returns 3

ii. Pop : Delete and insert  $Q.size() - 1$  times  
Delete one more time and return it

c. Push =  $O(n)$  and pop =  $O(1)$

i. Insertion order : 1 2 3 4

Stack : 1 2 3 4 - top

Queue :

Push		Pop	
1	1	4 3 2 1	
1 2	2 1	3 2 1	=> 4
2 1 3	3 2 1	2 1	=> 3
3 2 1 4	4 3 2 1	1	=> 2

ii. Push : Delete and insert  $Q.size() - 1$  times  
Insert element

## 2. Queue using Stack

a. <https://leetcode.com/problems/implement-queue-using-stacks/>

## 3. Largest Histogram in Rectangle :

a. [https://leetcode.com/problems/largest-rectangle-in-histogram/discuss/28902/5ms-O\(n\)-Java-solution-explained-\(beats-96\)](https://leetcode.com/problems/largest-rectangle-in-histogram/discuss/28902/5ms-O(n)-Java-solution-explained-(beats-96))

b. Next smaller left side using stack

```
vector<int> ans(A.size());
stack<int> st;
for(int i=0;i<A.size();i++){
    while(!st.empty() && st.top() >= A[i] ) st.pop();
    if(st.empty()) ans[i] = -1;
    else ans[i] = st.top();
    st.push(A[i]);
}
```

c. Using dp

```
int leftsmaller[n], rightsmaller[n];
for(int i=0;i<n;i++){
    int p = i - 1;
    while(p >= 0 && A[p] >= A[i])
        p = leftsmaller[p];
    leftsmaller[i] = p;
}
```

```
for(int i=n-1;i>=0;i--){
    int p = i + 1;
    while(p <= n-1 && A[p] >= A[i])
        p = rightsmaller[p];
    rightsmaller[i] = p;
}
```

#### 4. Postfix Evaluation

- a. <https://leetcode.com/problems/evaluate-reverse-polish-notation/submissions/>
- b. Visit each token of string  
If token is operator
  - i. Pop two element evaluate it and push it backElse
  - ii. Push token

#### 5. Minimum number of bracket reversals needed to make an expression balanced

- a. <https://www.geeksforgeeks.org/minimum-number-of-bracket-reversals-needed-to-make-an-expression-balanced/>
- b. Remove all pairs {} after that -- >}}...{{{
- c.  $\text{Ans} = m/2 + n/2$ , m - no of }, n - no of {

#### 6. Length of the longest valid substring

- a. [geeksforgeeks.org/length-of-the-longest-valid-substring/](https://www.geeksforgeeks.org/length-of-the-longest-valid-substring/)
- b.

```
for (int i=0; i<n; i++)  
{  
    // If opening bracket, push index of it  
    if (str[i] == '(')  
        stk.push(i);  
    else // If closing bracket, i.e., str[i] = ')'   
    {  
        // Pop the previous opening bracket's index  
        stk.pop();  
  
        // Check if this length formed with base of  
        // current valid substring is more than max  
        // so far  
        if (!stk.empty())  
            result = max(result, i - stk.top());  
  
        // If stack is empty. push current index as  
        // base for next valid substring (if any)  
        else stk.push(i);  
    }  
}
```

c.

#### 7. Get min in O(1) using stack

- a. <https://www.interviewbit.com/problems/min-stack/>
- b. The idea is to store the next min below that element in stack so that if we remove any element min value can be updated by next element.  
So when we do push operation and if element x is smaller than current element update min element.  
Push min element and push new element x

#### 8. Sliding window maximum

- a. <https://leetcode.com/problems/sliding-window-maximum/>

- b. Use next greater array and two pointer concept

9. Rain Water Trapped

- a. <https://leetcode.com/problems/trapping-rain-water/>
- b. Each  $\text{height}[i]$  will contribute  $\min(\text{leftMax}[i-1], \text{rightMax}[i+1]) - \text{height}[i]$  amount of water.