

Homework 5

Due 10/10/2024

October 3, 2024

1. Give pseudocode for an algorithm to compute the number of *quadratic nonresidues* modulo n , where a quadratic nonresidue mod n is an integer r in the range of 0 to $n - 1$ such that there is no integer x in this range where $x^2 \bmod n = r$.

For example, the quadratic nonresidues of $n = 6$ are 2 and 5, as shown by the table below:

x	$x^2 \bmod 6$
0	0
1	1
2	4
3	3
4	4
5	1

Thus, the solution for $n = 6$ is 2.

Your algorithm should be as efficient as possible. Be sure to describe which data structure(s) you're using and why.

2. Dr. Snevets is working on an algorithm that uses a Union-Find to maintain a set partition, but her algorithm also needs to efficiently determine the smallest element in each partition. Answer the following about developing an improved Union-Find that efficiently supports a FindMin(x) operation that can return the minimum element in the same partition as x .
 - (a) One way to implement this improved data structure is to store a third array *min* such that *min*[r] is the minimum element in the tree rooted at r , for every root element r . (*min* can equal any value at indexes that are not roots.) Describe how to modify the pseudocode for the Union-Find initialize operation (below) to also initialize *min* appropriately.

Input: n : size of the Union-Find to initialize
Output: a Union-Find of size n where every element points to itself

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1 Algorithm: Union-Find.Initialize
2  $uf = \text{Array}(n)$ 
3 Initialize  $uf$  to  $1..n$ 
4  $size = \text{Array}(n)$ 
5 Initialize  $size$  to 1
6 return ( $uf, size$ )

```

- (b) What is the worst-case time complexity for your modified Initialize algorithm? Justify your answer.
- (c) Describe how to modify the Union operation (below) so that min always stores the minimum element of the tree at every root (min can store anything at an index that is not a root of the Union-Find.)

Input: ($uf, size$): the Union-Find to modify
Input: a : index of one element to union
Input: b : index of another element to union
Output: modified uf that efficiently merges the trees containing a and b

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1 Algorithm:  $uf$ .Union
2  $ra = uf.\text{Find}(a)$ 
3  $rb = uf.\text{Find}(b)$ 
4 if  $size[ra] > size[rb]$  then
5   | Swap  $ra$  and  $rb$ 
6 end
7  $uf[ra] = rb$ 
8  $size[rb] = size[ra] + size[rb]$ 

```

- (d) What is the amortized time complexity of your modified Union algorithm? Show your work. Note that the Find operation does not need to be changed from the optimized Union-Find presented in class.
3. Find the worst-case complexity of the hierarchical clustering algorithm below. You may assume that the distance function takes $\Theta(1)$ time to compute.

Input: $data$: set of data points
Input: n : size of $data$
Input: $distance$: distance function that takes two data points and returns a nonnegative real number
Input: c : desired number of clusters; must be an integer between 1 and n
Output: single-linkage hierarchical clusters for $data$

```

1 Algorithm: SingleHClust
2  $heap = \text{MinHeap}()$ 
3 for  $i = 1$  to  $n - 1$  do
4   | for  $j = i + 1$  to  $n$  do
5   |   | Insert  $distance(data[i], data[j])$  into  $heap$ , along with the
6   |   |   corresponding indexes  $i$  and  $j$ 
7   | end
8 end
9  $uf = \text{UnionFind}(n)$ 
10  $count = n$ 
11 while  $count > c$  do
12   |  $(i, j, dist) = heap.DeleteMin()$ 
13   | if  $uf.Find(i) \neq uf.Find(j)$  then
14   |   |  $uf.Union(i, j)$ 
15   |   |  $count = count - 1$ 
16   | end
17 end
18 return  $uf$ 

```