Homework 3

Due 09/25/2024

September 19, 2024

- 1. Give pseudocode for an algorithm that accepts a stack stk of size n and an integer k between 0 and n that modifies stk so that its top k elements are reversed, using only a queue for storage. (You may not declare any variables other than a single queue.)
- 2. Give pseudocode for an algorithm that accepts a stack stk of size n and integers i and j such that $0 \le i \le j \le n$ and modifies stk so that all of its entries between indexes i and j (including j but excluding i) are reversed, using only a queue for storage.

Note that if i = 0, the resulting stack will have the top j entries in reverse order. Also, if j = i or j = i + 1, the stack should retain all elements in the same order.

Hint: you may call your solution to problem 1 as a subroutine.

- 3. Give pseudocode for an algorithm that accepts a stack stk of size n and integers i and j such that $1 \le i \le j \le n$ and modifies stk so that indexes i and j (counting from the top) are swapped, using only a queue for storage. Hint: you may call your solution to problem 2 as a subroutine.
- 4. Give pseudocode for a post-order tree traversal for an *m*-ary tree stored in firstChild-nextSibling form.
- 5. Answer the following questions about a modification to a BST so that it can return the min value in $\Theta(1)$ time.
 - (a) What additional data members would you store in the BST? Your answer should indicate what these data members represent and how they can be used to identify the min in $\Theta(1)$.
 - (b) How would you modify insertion so that your data members are updated appropriately? Pseudocode for a basic BST insertion has been given below. Your modification should exhibit the same time complexity.

Your answer may be an English-language description of how you would modify this pseudocode, or it may be the modified pseudocode.

```
Input: ins: new data value to insert
   Input: node: current BST node (originally root)
  Input: BST.Insert
ı if ins \leq node.value then
      if node.left = nil then
          node.left = Node(ins)
 3
          node. {\bf left.parent} = node
 4
      else
5
          BST.Insert(ins, node.left)
 6
      end
 7
8 else
      if node.right = nil then
9
          node.right = Node(ins)
10
          node.right.parent = node
11
12
      else
          BST.Insert(ins, node.right)
13
14
      end
15 end
```

(c) How would you modify deletion? Your modification should exhibit the same time complexity as ordinary deletion.

```
Input: victim: BST node to be deleted
   Input: BST.Delete
 {f 1} Let children be the number of non-nil children of victim
 \mathbf{if} \ children = 0 \ \mathbf{then}
       if victim.parent \neq nil then
          Set victim.parent's matching child pointer to nil
 4
       end
 5
       {\bf delete}\ victim
 7 else if children = 1 then
       Let child be the child of victim
       if victim.parent \neq nil then
 9
          Set victim.parent's matching child pointer to child
10
       end
11
       child. {\tt parent} = victim. {\tt parent}
12
       {\bf delete}\ victim
13
14 else
       lhsMax = victim.left
15
       while lhsMax.right \neq nil do
16
          lhsMax = lhsMax.right
17
       end
18
       Swap victim.data and lhsMax.data
19
       BST.Delete(lhsMax)
20
21 end
```