

Shortest path Algorithm

Let $G=(V,E)$ be a digraph (directed graph).

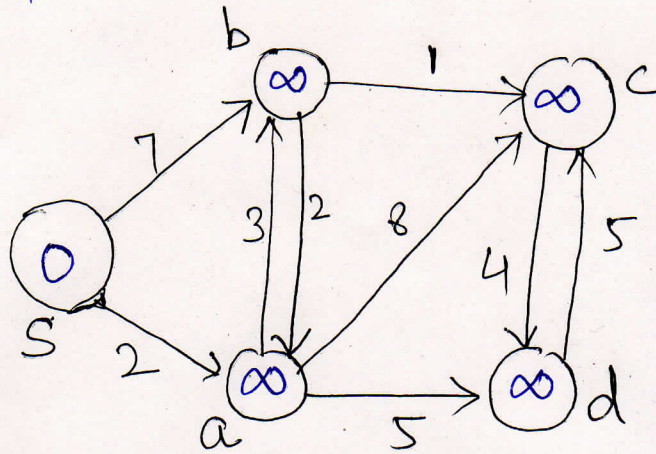
- The Shortest path between two vertices is a path with the shortest length (least number of edges)
- Breadth-first-search is an algorithm for finding shortest (link-distance) paths from a single source vertex to all other vertices.

Dijkstra's Algorithm

- Maintain an estimate $d[v]$ of the length $\delta(s,v)$ of the shortest path for each vertex v .
- Always $d[v] \geq \delta(s,v)$ and $d[v]$ equals the length of a known path
($d[v] = \infty$ if we have no paths so far)
- Initially $d[s] = 0$ and all the other $d[v]$ values are set to ∞ . The algorithm will then process the vertices one by one in some order.
The processed vertex's estimate will be validated as being real shortest distance, i.e. $d[v] = \delta(s,v)$

Here "processing a vertex u " means finding new paths and updating $d[v]$ for all $v \in \text{Adj}[u]$ if necessary
The process by which an estimate is updated is called relaxation.

Example:

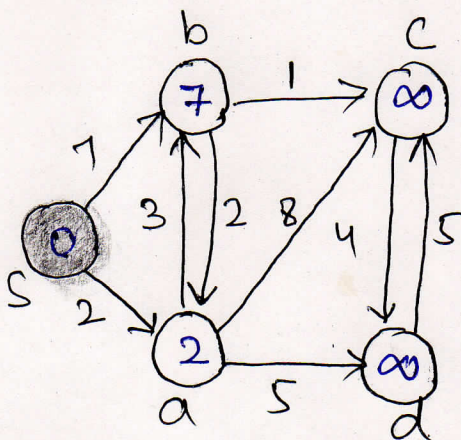


Step 0: Initialization

v	s	a	b	c	d
d[v]	0	∞	∞	∞	∞
pred[v]	nil	nil	nil	nil	nil
color[v]	W	W	W	W	W

Priority Queue

v	s	a	b	c	d
d[v]	0	∞	∞	∞	∞

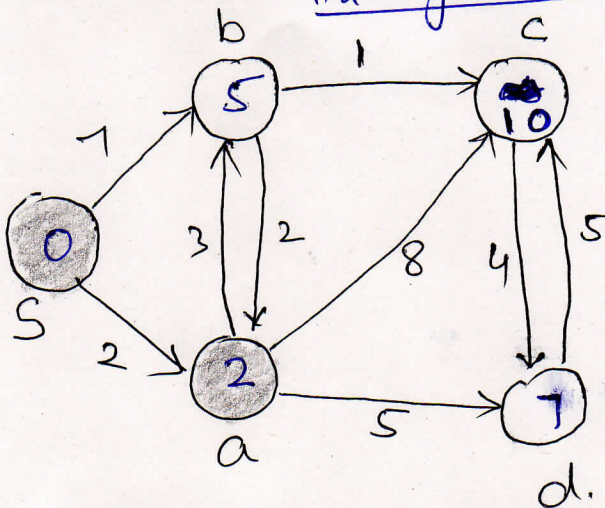


Step 1:

v	s	a	b	c	d
d[v]	0	2	7	∞	∞
pred[v]	nil	s	s	nil	nil
color[v]	B	W	W	W	W

Priority Queue

v	a	b	c	d
d[v]	2	7	∞	∞

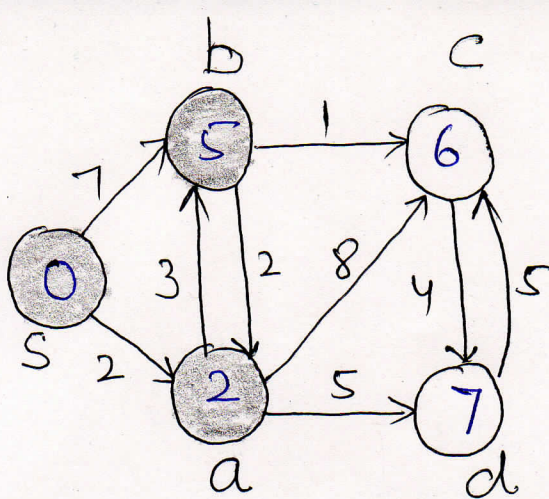


Step 2

v	s	a	b	c	d
d[v]	0	2	5	10	7
pred[v]	nil	s	a	a	a
color[v]	B	B	W	W	W

Priority Queue

v	b	c	d
d[v]	5	10	7

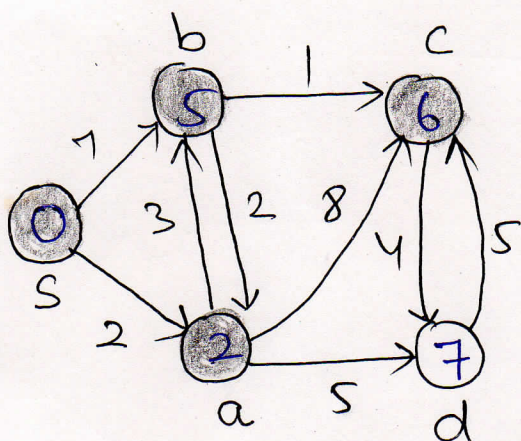


Step 3

v	s	a	b	c	d
d[v]	0	2	5	6	7
pred[v]	nil	s	a	b	a
color[v]	B	B	B	w	w

Priority Queue

v	c	d
d[v]	6	7

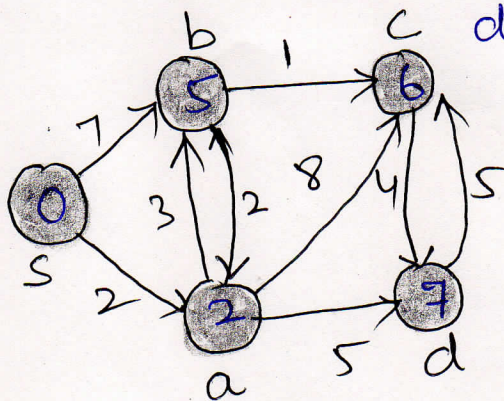


Step 4

v	s	a	b	c	d
d[v]	0	2	5	6	7
pred[v]	nil	s	a	b	a
color[v]	B	B	B	B	w

Priority Queue

v	d
d[v]	7



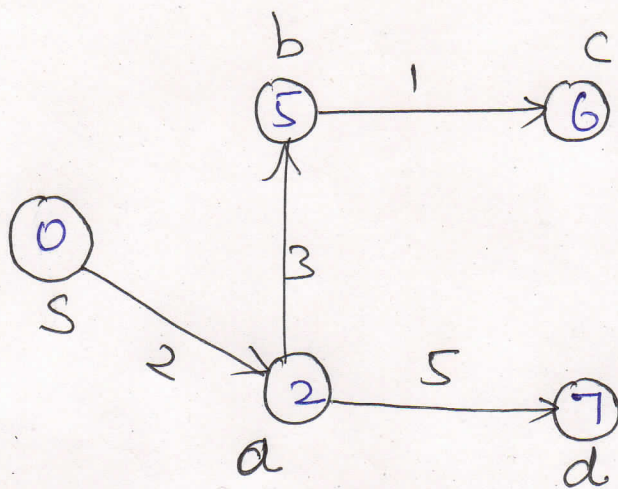
Steps

v	s	a	b	c	d
d[v]	0	2	5	6	7
pred[v]	nil	s	a	b	a
color[v]	B	B	B	B	B

Priority Queue

$Q = \emptyset$

Therefore, Shortest path tree



v	s	a	b	c	d
$d[v]$	0	2	5	6	7
$pred[v]$	nil	s	a	b	a