

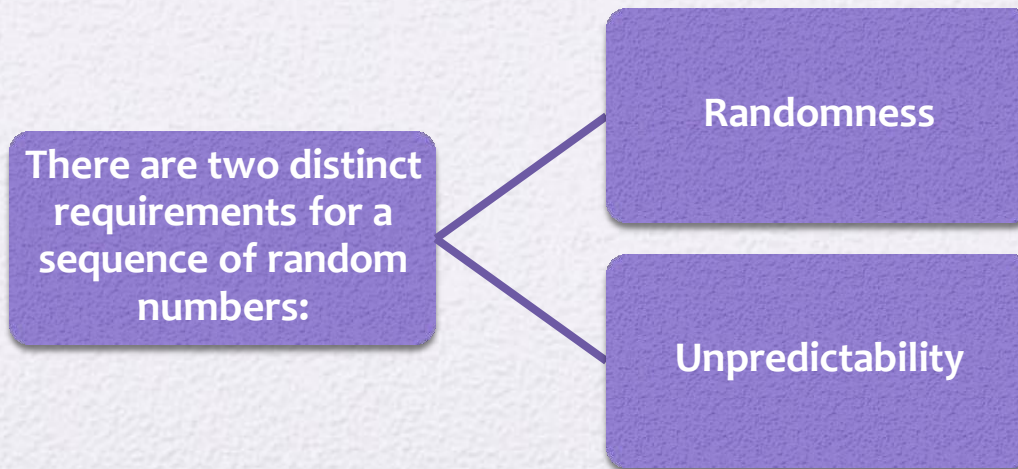


Unit 1

Pseudorandom Number Generators

Random Numbers

- A number of network security algorithms and protocols based on cryptography make use of random binary numbers:
 - Key distribution and reciprocal authentication schemes
 - Session key generation
 - Generation of keys for the RSA public-key encryption algorithm
 - Generation of a bit stream for symmetric stream encryption



Randomness

- The generation of a sequence of allegedly random numbers being random in some well-defined statistical sense has been a concern

Two criteria are used to validate that a sequence of numbers is random:

Uniform distribution

- The frequency of occurrence of ones and zeros should be approximately equal

Independence

- No one subsequence in the sequence can be inferred from the others

Unpredictability

- The requirement is not just that the sequence of numbers be statistically random, but that the successive members of the sequence are unpredictable
- With “true” random sequences each number is statistically independent of other numbers in the sequence and therefore unpredictable
 - True random numbers have their limitations, such as inefficiency, so it is more common to implement algorithms that generate sequences of numbers that appear to be random
 - Care must be taken that an opponent not be able to predict future elements of the sequence on the basis of earlier elements

Pseudorandom Numbers

- Cryptographic applications typically make use of algorithmic techniques for random number generation
- These algorithms are deterministic and therefore produce sequences of numbers that are not statistically random
- If the algorithm is good, the resulting sequences will pass many tests of randomness and are referred to as *pseudorandom numbers*

True Random Number Generator (TRNG)

- Takes as input a source that is effectively random
- The source is referred to as an *entropy source* and is drawn from the physical environment of the computer
 - Includes things such as keystroke timing patterns, disk electrical activity, mouse movements, and instantaneous values of the system clock
 - The source, or combination of sources, serve as input to an algorithm that produces random binary output
- The TRNG may simply involve conversion of an analog source to a binary output
- The TRNG may involve additional processing to overcome any bias in the source

Pseudorandom Number Generator (PRNG)

- Takes as input a fixed value, called the *seed*, and produces a sequence of output bits using a deterministic algorithm
 - Quite often the seed is generated by a TRNG
 - The output bit stream is determined solely by the input value or values, so an adversary who knows the algorithm and the seed can reproduce the entire bit stream
 - Other than the number of bits produced there is no difference between a PRNG and a PRF

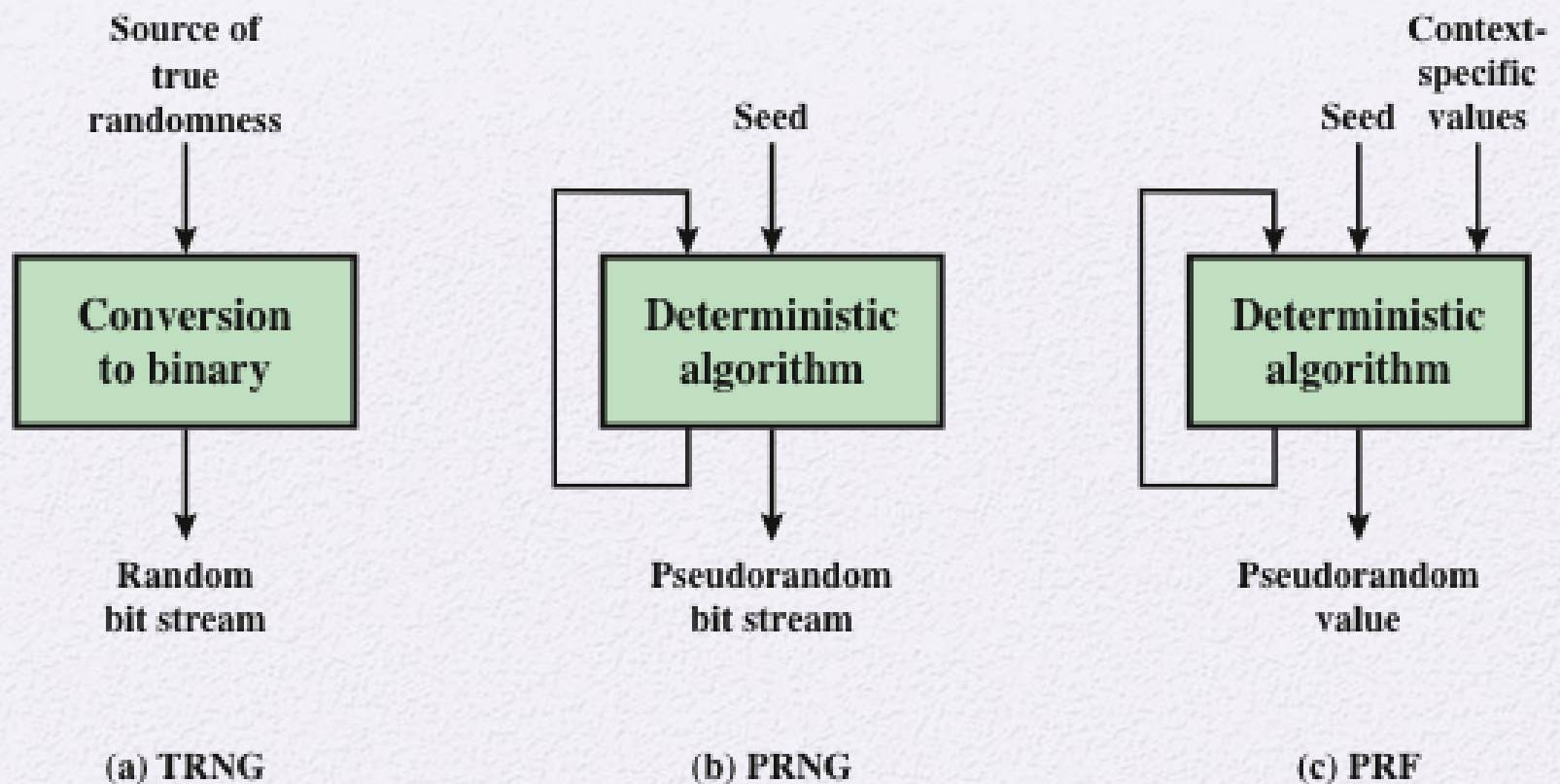
Two different forms of PRNG

Pseudorandom number generator

- An algorithm that is used to produce an open-ended sequence of bits
- Input to a symmetric stream cipher is a common application for an open-ended sequence of bits

Pseudorandom function (PRF)

- Used to produce a pseudorandom string of bits of some fixed length
- Examples are symmetric encryption keys and nonces



TRNG = true random number generator
PRNG = pseudorandom number generator
PRF = pseudorandom function

Figure 7.1 Random and Pseudorandom Number Generators

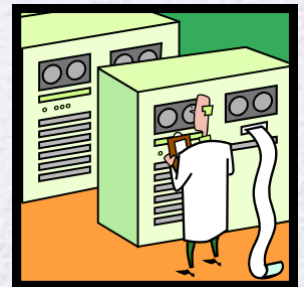
PRNG Requirements

- The basic requirement when a PRNG or PRF is used for a cryptographic application is that an adversary who does not know the seed is unable to determine the pseudorandom string
- The requirement for secrecy of the output of a PRNG or PRF leads to specific requirements in the areas of:
 - Randomness
 - Unpredictability
 - Characteristics of the seed



Randomness

- The generated bit stream needs to appear random even though it is deterministic
- There is no single test that can determine if a PRNG generates numbers that have the characteristic of randomness
 - If the PRNG exhibits randomness on the basis of multiple tests, then it can be assumed to satisfy the randomness requirement
- NIST SP 800-22 specifies that the tests should seek to establish three characteristics:
 - Uniformity
 - Scalability
 - Consistency



Randomness Tests

- SP 800-22 lists 15 separate tests of randomness

Frequency test

- The most basic test and must be included in any test suite
- Purpose is to determine whether the number of ones and zeros in a sequence is approximately the same as would be expected for a truly random sequence

Runs test

- Focus of this test is the total number of runs in the sequence, where a run is an uninterrupted sequence of identical bits bounded before and after with a bit of the opposite value
- Purpose is to determine whether the number of runs of ones and zeros of various lengths is as expected for a random sequence

Maurer's universal statistical test

- Focus is the number of bits between matching patterns
- Purpose is to detect whether or not the sequence can be significantly compressed without loss of information. A significantly compressible sequence is considered to be non-random

The diagram features a central purple circle with the text "Three tests" in white. Three grey arrows point outwards from this circle towards three purple rounded rectangular boxes. The top box is titled "Runs test", the bottom-left box is titled "Frequency test", and the bottom-right box is titled "Maurer's universal statistical test". Each box contains a bulleted list of details about the respective test.

Three
tests

Unpredictability

- A stream of pseudorandom numbers should exhibit two forms of unpredictability:
- Forward unpredictability
 - If the seed is unknown, the next output bit in the sequence should be unpredictable in spite of any knowledge of previous bits in the sequence
- Backward unpredictability
 - It should not be feasible to determine the seed from knowledge of any generated values. No correlation between a seed and any value generated from that seed should be evident; each element of the sequence should appear to be the outcome of an independent random event whose probability is $1/2$
- The same set of tests for randomness also provides a test of unpredictability
 - A random sequence will have no correlation with a fixed value (the seed)

Seed Requirements

- The seed that serves as input to the PRNG must be secure and unpredictable
- The seed itself must be a random or pseudorandom number
- Typically the seed is generated by TRNG



Generation of Seed Input to PRNG

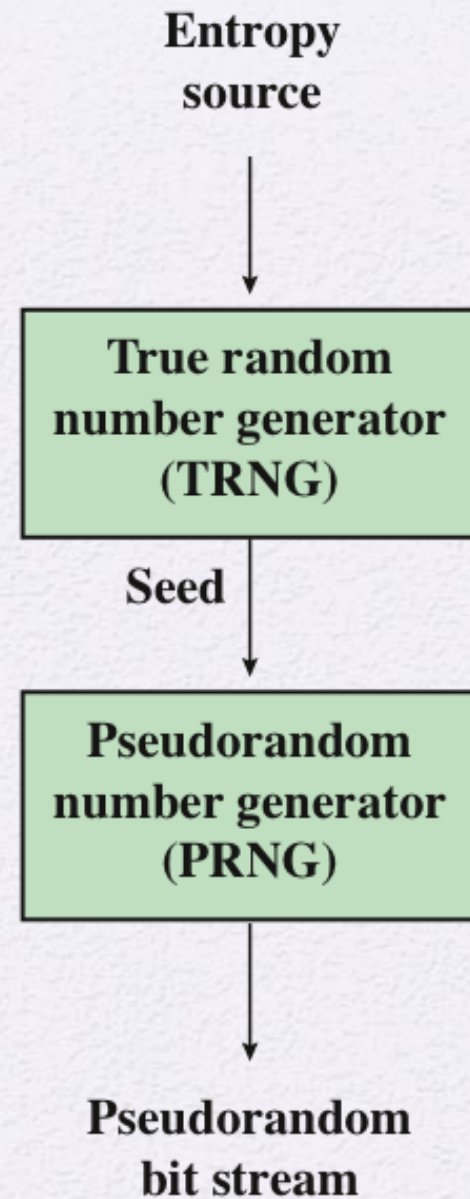


Figure 7.2 Generation of Seed Input to PRNG

Algorithm Design

- Algorithms fall into two categories:
 - Purpose-built algorithms
 - Algorithms designed specifically and solely for the purpose of generating pseudorandom bit streams
 - Algorithms based on existing cryptographic algorithms
 - Have the effect of randomizing input data

Three broad categories of cryptographic algorithms are commonly used to create PRNGs:

- Symmetric block ciphers
- Asymmetric ciphers
- Hash functions and message authentication codes

Linear Congruential Generator

An algorithm first proposed by Lehmer that is parameterized with four numbers:

m the modulus

m > 0

a the multiplier

$0 < a < m$

c the increment

$0 \leq c < m$

X₀ the starting value, or seed

$0 \leq X_0 < m$

- The sequence of random numbers $\{X_n\}$ is obtained via the following iterative equation:

$$X_{n+1} = (aX_n + c) \bmod m$$

- If m , a , c , and X_0 are integers, then this technique will produce a sequence of integers with each integer in the range $0 \leq X_n < m$
- *The selection of values for a , c , and m is critical in developing a good random number generator*

Blum Blum Shub (BBS) Generator

- Has perhaps the strongest public proof of its cryptographic strength of ***any purpose-built algorithm***
- Referred to as a cryptographically secure pseudorandom bit generator (**CSPRBG**)
 - A CSPRBG is defined as one that passes the *next-bit-test* if there is not a polynomial-time algorithm that, on input of the first k bits of an output sequence, can predict the $(k + 1)$ st bit with probability significantly greater than $1/2$
- ***The security of BBS is based on the difficulty of factoring n***

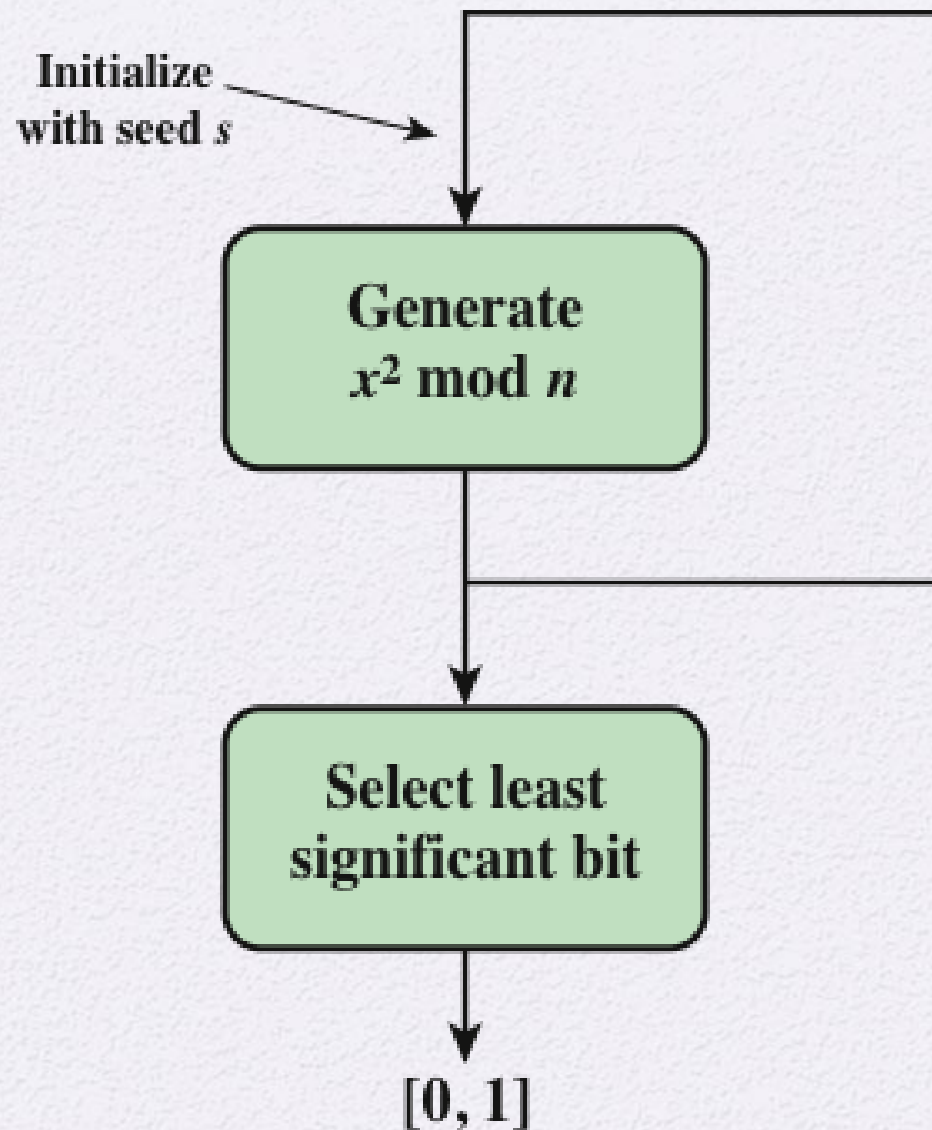


Figure 7.3 Blum Blum Shub Block Diagram

Table 7.1

Example Operation of BBS Generator

i	X_i	B_i
0	20749	
1	143135	1
2	177671	1
3	97048	0
4	89992	0
5	174051	1
6	80649	1
7	45663	1
8	69442	0
9	186894	0
10	177046	0

i	X_i	B_i
11	137922	0
12	123175	1
13	8630	0
14	114386	0
15	14863	1
16	133015	1
17	106065	1
18	45870	0
19	137171	1
20	48060	0

PRNG Using Block Cipher Modes of Operation

- Two approaches that use a block cipher to build a PRNG have gained widespread acceptance:

- **CTR mode (Counter)**

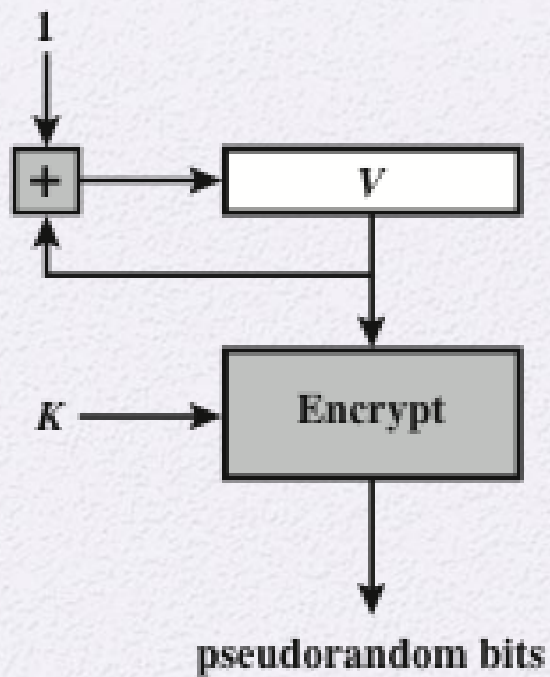
Counter mode turns a [block cipher](#) into a [stream cipher](#). It generates the next [keystream](#) block by encrypting successive values of a "counter". The counter can be any function which produces a sequence which is guaranteed not to repeat for a long time, although an actual increment-by-one counter is the simplest and most popular.

- Recommended in NIST SP 800-90, ANSI standard X.82, and RFC 4086

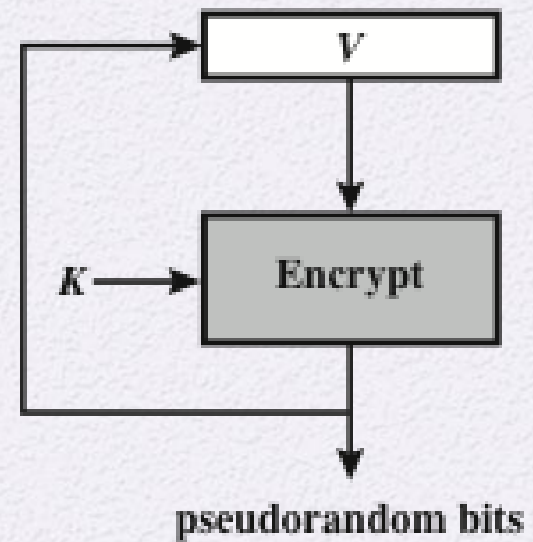
- **OFB mode (Output Feedback)**

The *Output Feedback* (OFB) mode makes a block cipher into a synchronous [stream cipher](#). It generates [keystream](#) blocks, which are then [XORed](#) with the plaintext blocks to get the ciphertext.

- Recommended in X9.82 and RFC 4086



(a) CTR Mode



(b) OFB Mode

Figure 7.4 PRNG Mechanisms Based on Block Ciphers

Table 7.2

Output Block	Fraction of One Bits	Fraction of Bits that Match with Preceding Block
1786f4c7ff6e291dbdfdd90ec3453176	0.57	—
5e17b22b14677a4d66890f87565eae64	0.51	0.52
fd18284ac82251dfb3aa62c326cd46cc	0.47	0.54
c8e545198a758ef5dd86b41946389bd5	0.50	0.44
fe7bae0e23019542962e2c52d215a2e3	0.47	0.48
14fdf5ec99469598ae0379472803accd	0.49	0.52
6aeca972e5a3ef17bd1a1b775fc8b929	0.57	0.48
f7e97badf359d128f00d9b4ae323db64	0.55	0.45

Example Results for PRNG Using OFB

Table 7.3

Output Block	Fraction of One Bits	Fraction of Bits that Match with Preceding Block
1786f4c7ff6e291dbdfdd90ec3453176	0.57	—
60809669a3e092a01b463472fdcae420	0.41	0.41
d4e6e170b46b0573eedf88ee39bff33d	0.59	0.45
5f8fcfc5deca18ea246785d7fadc76f8	0.59	0.52
90e63ed27bb07868c753545bdd57ee28	0.53	0.52
0125856fdf4a17f747c7833695c52235	0.50	0.47
f4be2d179b0f2548fd748c8fc7c81990	0.51	0.48
1151fc48f90eebac658a3911515c3c66	0.47	0.45

Example Results for PRNG Using CTR

ANSI X9.17 PRNG

- One of the strongest PRNGs is specified in ANSI X9.17
 - A number of applications employ this technique including financial security applications and PGP

Input

- Two pseudorandom inputs drive the generator. One is a 64-bit representation of the current date and time. The other is a 64-bit seed value; this is initialized to some arbitrary value and is updated during the generation process.

The algorithm makes use of triple DES for encryption.
Ingredients are:

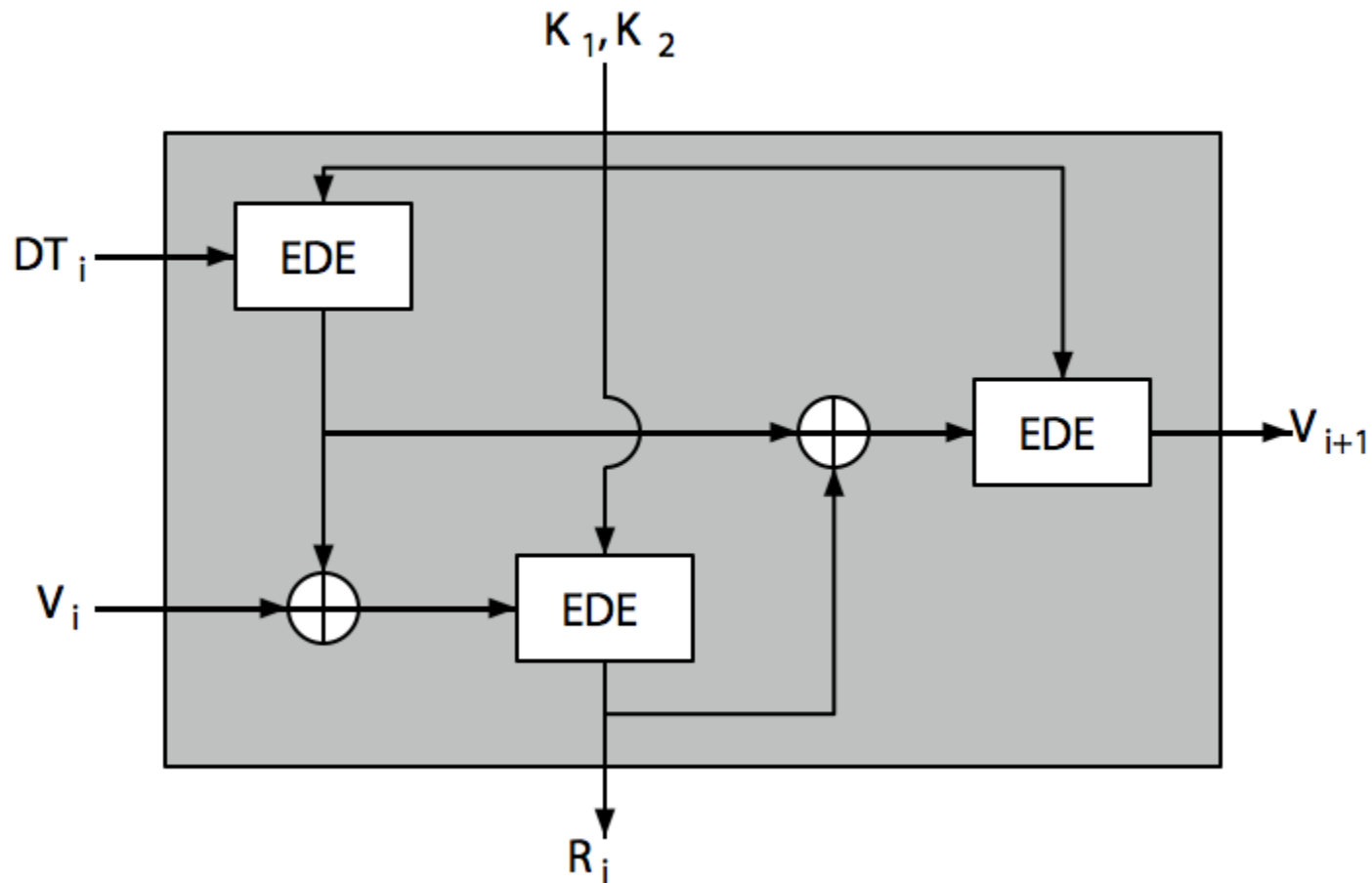
Output

- The output consists of a 64-bit pseudorandom number and a 64-bit seed value.

Keys

- The generator makes use of three triple DES encryption modules. All three make use of the same pair of 56-bit keys, which must be kept secret and are used only for pseudorandom number generation.

ANSI X9.17 PRNG



Dti = date and time