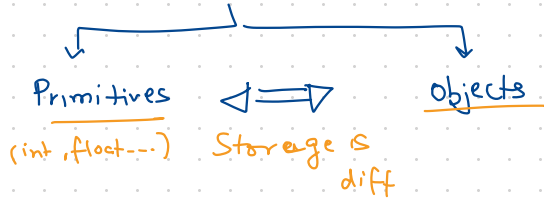


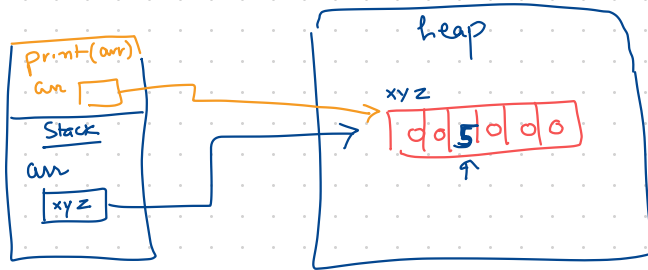
ARRAY



Memory Level

`int [] arr = new int [size];`
Stack heap
`arr[i] = data;` `i=2, data=5`

Pass By Value



Object

Scanner SC; ← Single object

class PaintBrush{

→ color
→ size
→ width

}

brushes[0] = null;
brushes = null;

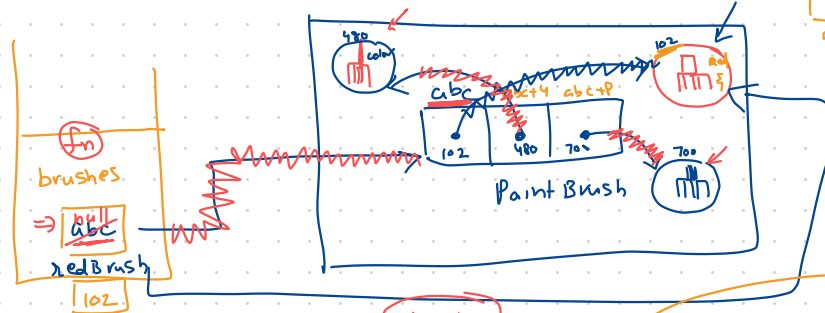
PaintBrush[] brushes = new PaintBrush[3]; ←

Stack

brush[0]

brush[0].color

Stack



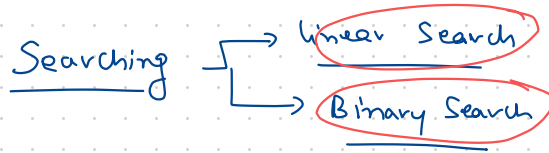
PaintBrush redBrush = new PaintBrush (Red, 5, 2);

Array

brushes[0] = redBrush;

brushes[1] = new PaintBrush(.,.);

return address



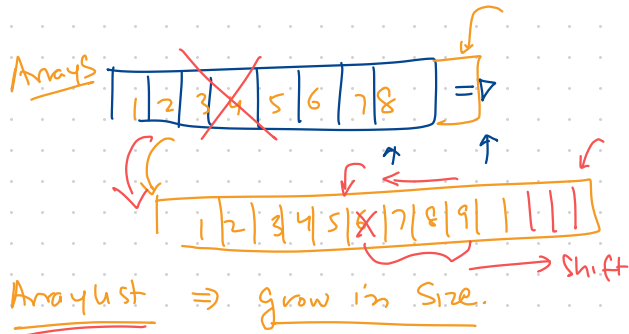
Phone / whatsapp
[] 🔍

Group, Ppl
[] 🔍

Amazon
[] 🔍

Zomato
[] 🔍

Instagram
[] 🔍



✓ ⇒ Insert

✓ ⇒ Search

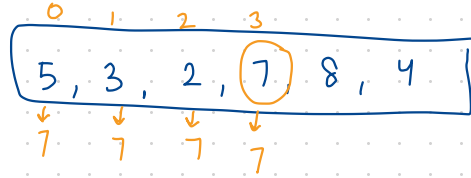
✓ ⇒ Deletion

✓ ⇒ Update



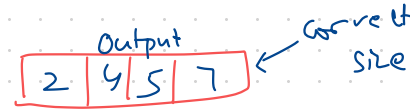
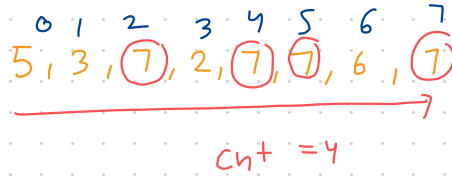
$a[idx] = \text{some-new-val;}$

Linear Search



key
↓
7 where?

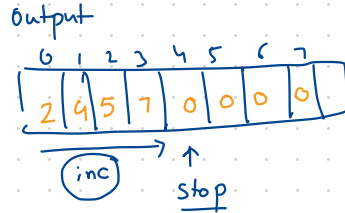
①



$$N + N = 2N \text{ iterations.}$$

→
Note it down the indices

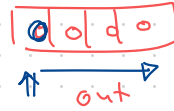
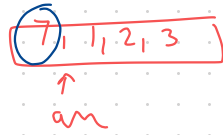
②



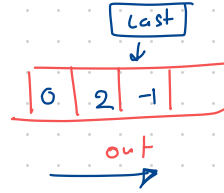
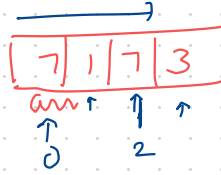
⇒ Interpret this array in matrix

2, 4, 5, 7

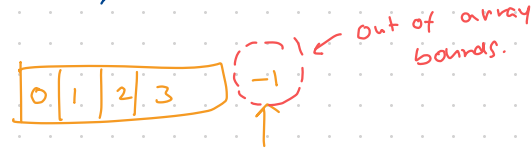
one case it can fail



Confusion?



0, 2



output arr length + 1 ; ⇐

III

ArrayList (grows automatically)

Data

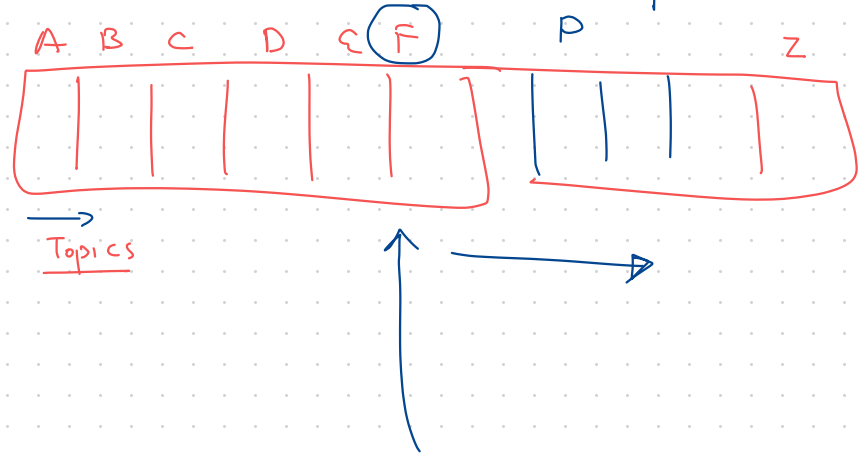
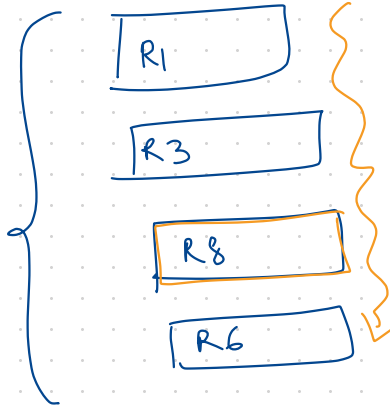
Unsorted
array

"Sorted / ordered"

→ Books in library

Orders

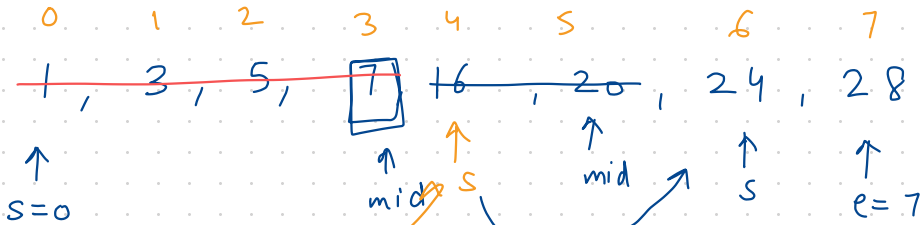
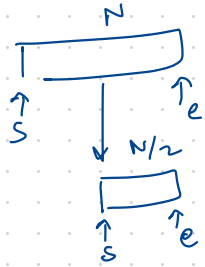
→ Date & Time
→ Most Recent



Binary Search

(only sorted arrays)

inc / dec



Key = 24

✓ 0-7 $mid = \left(\frac{s+e}{2} \right) = \left(\frac{0+7}{2} \right) = 3$

4-7 $mid = \left(\frac{s+e}{2} \right) = \left(\frac{4+7}{2} \right) = 5$

8-7 $mid = \frac{6+7}{2} = 6$

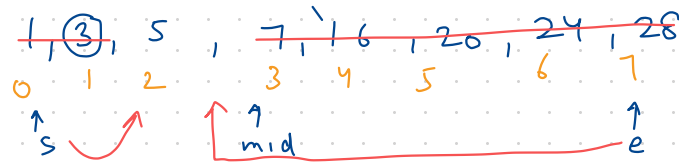
$7 < 24$

$S = mid + 1$ Moved Right

$20 < 24$

$S = mid + 1$

$24 == 24$



0 - 7

$$\text{mid} = \left(\frac{0+7}{2} \right) = 3$$

0 - 2

$$\text{mid} = \left(\frac{0+2}{2} \right) = 1$$

2 - 2

$$\text{mid} = \left(\frac{2+2}{2} \right) = 2$$

key = 5

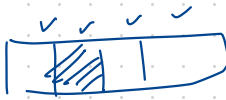
5 (7)

$e = \text{mid} - 1$

3 → 5

$s = \text{mid} + 1$

s == e

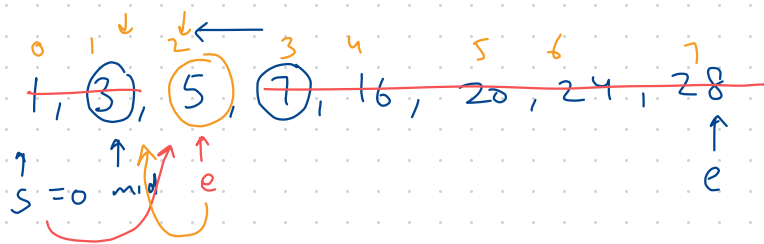


$$\left(\frac{3}{2} \right) \Rightarrow 1$$



e s

Element
Not
Present



key = 4

0 - 7

mid = 3

0 - 2

mid = 1

2 - 2

mid = 2

e = 1
s = 2

conclude?

Not found

4 ← 7

e = mid - 1
3 → s = mid + 1

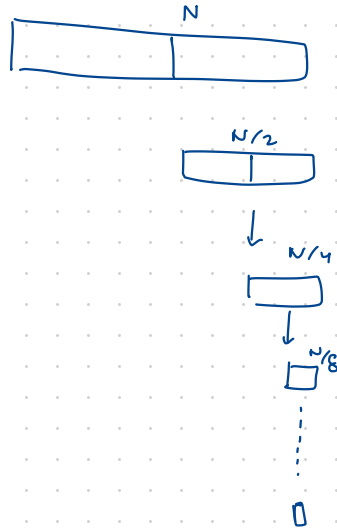
← 5

e = mid - 1

while (s <= e) {

≡

3
Not found



Steps

$$\begin{aligned}
 1 &\rightarrow \frac{N}{2} = \frac{N}{2^1} \\
 2 &\rightarrow \frac{N}{4} = \frac{N}{2^2} \\
 3 &\rightarrow \frac{N}{8} = \frac{N}{2^3} \\
 &\vdots \\
 k \text{ steps} &\rightarrow 1 = \frac{N}{2^k}
 \end{aligned}$$

$$\frac{N}{2^k} = 1$$

$$2^k = N$$

$$\log_2 2^k = \log_2 N$$

$$k \log_2 2 = \log_2 N$$

$$k = \log_2 N$$

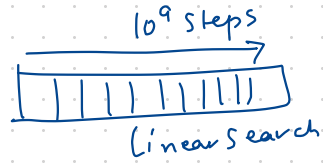
B.S. in worst case takes

$\log_2 N$ steps

$$\log_2 2 = 1$$

(Sorted)

Array Size $\rightarrow 10^9$



1s $\rightarrow 10^8$ ops

$$\boxed{\text{Time}} \quad \frac{10^9}{10^8} = 10 \text{ seconds}$$

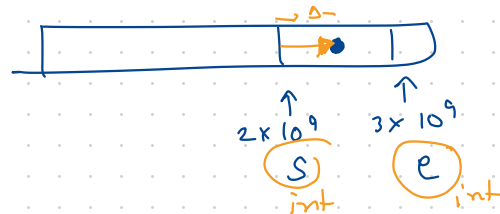
$\log_2 10^9$ steps



$$= \log_2 (10^3)^3$$

=

$$= \frac{3 \log_2 1000}{\text{30 steps}} = 3 \times 10^{-7} \text{ seconds}$$



~~s~~ ~~s~~ $\xrightarrow{\text{too big}} \text{beyond range of int sometime}$

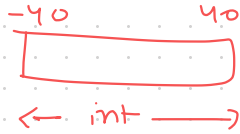
$$\text{mid} = \left(\frac{s + e}{2} \right)$$

$$\boxed{\text{mid} = s + \frac{(e-s)}{2}} \Rightarrow$$

small NO

more error proof

$$\frac{2s + e - s}{2} = \left(\frac{s+e}{2} \right)$$



$$s = 30$$

$$e = 40$$

Interval

$$\text{mid} = \frac{30+40}{2} = \frac{70}{2} = 35$$

$$\text{mid} = 30 + \frac{(40-30)}{2}$$

$$= 30 + \frac{10}{2}$$

$$= 30 + 5 = 35$$

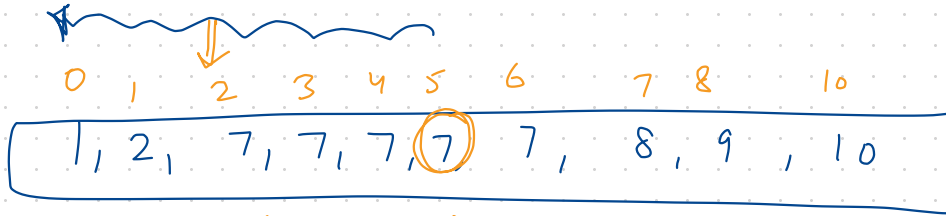
5

↑ ↑ ↑
s mid e

$N/2$



10.40



↑
s

Linear search
after finding
Random idx X

↑
e

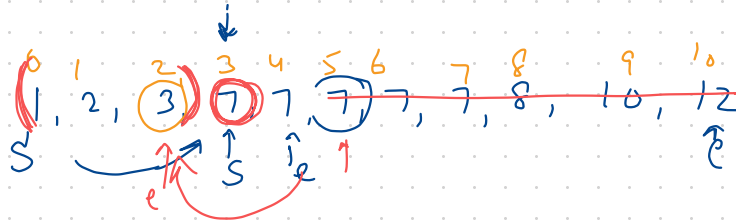
Key = 7

$$mid = \frac{0 + 10}{2} = 5$$

any Random occ.

$\propto N$ iterations.

⇒ First occ



$$0-10 \quad \text{mid} = \frac{0+10}{2} = 5$$

$$0 - 4 \quad \text{mid} = \frac{0 + 4}{2} = 2$$

$$3 - 4 \quad \text{mid} = \frac{3+4}{2} = 3$$

$$\begin{bmatrix} p=2 \\ s=3 \end{bmatrix}$$

→ Stop

stone & go left
↓

ans = ~~5~~ 3

3, 7

update ans

$$e = mid - 1$$

Think:

Given a Number N , find out the square root N
without using pre-built method

20 \rightarrow 4 (int part)

25 \rightarrow 5

49 \rightarrow 7

110 \rightarrow 10

o Linear Search

o Binary Search.

Brute Force / Basic

$i \leftarrow$

$$N = 30$$

↓

Sq's

$$\leq 30$$

Root \boxed{N}

$$i = 0 \quad 0^2 \leq 30$$

$$i = 1 \quad 1^2 \leq 30$$

$$i = 2 \quad 2^2 \leq 30$$

$$i = 3 \quad 3^2 \leq 30$$

$$i = 4 \quad 4^2 \leq 30$$

$$i = 5 \quad 5^2 \leq 30$$

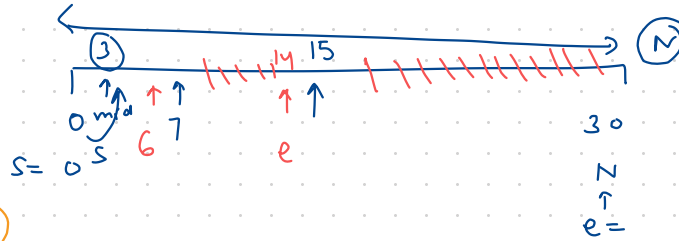
$$i = 6 \quad 6^2 \leq 30 \Rightarrow \text{stop}$$

$i-1$

$$i = 0 \quad \text{---} \quad i^2 = \sqrt{N}$$

$i > \sqrt{N}$ Stop

Binary Search

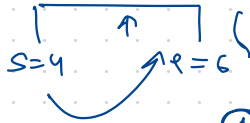


ans = 3 5

(1)

$n/2$

$n/4$



(1)

$e=s, s=e$

mid = 15

mid = 7

mid = 3

mid = 5

mid = 6

Stop

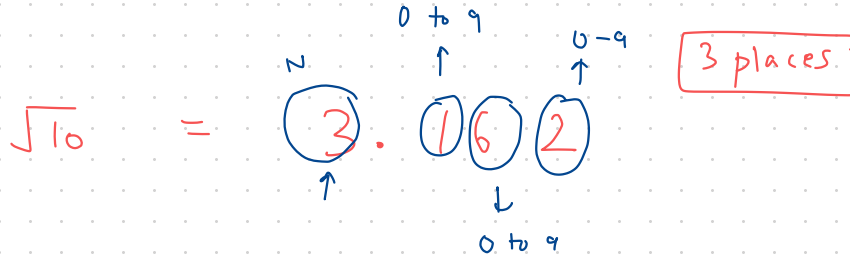
$15^2 \leq 30$ No Left

$7^2 \leq 30$ No Left

$3^2 \leq 30$ Yes Right

$5^2 \leq 30$ Yes Right

$6^2 \leq 30$ No



$3 \cdot 0^2 \leq 10$
 $3 \cdot 1^2 \leq 10$
 $3 \cdot 2^2 \neq 10$

→ Linear Search for each digit
 → Binary Search for each digit

$2^2 \ 2^1 \ 2^0 \ 2^{-1} \ 2^{-2} \ 2^{-3} \ 2^{-4} \ 2^{-5}$
 $1 \ 0 \ 1 \ 0 \ 1 \ 0 \ 1 \dots$
 $0.5 \ 0.25 \ 0.125$

Negative No's

5.25

3.162

101

BS
 $\log N + p \cdot 10$
 $\log N + p \cdot \log 10$

Constants

3.1620000
 $3.1619999 \dots$
 3.162

Monotonic Search Spaces

