

# **EXPERIMENT NO.10**

## **Sum Of Subset Problem**

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### **Program:-**

```
#include<iostream>
#include<vector>
using namespace std;

bool isSubsetSumDP(vector<int>&set, int sum, vector<int>&subset) {
    int n = set.size();
    int i, j;
    vector<vector<bool>> dp(n + 1, vector<bool>(sum + 1, false));
    for (i = 0; i <= n; ++i)
        dp[i][0] = true;
    for (i = 1; i <= n; ++i) {
        for (j = 1; j <= sum; ++j) {
            if (set[i - 1] > j)
                dp[i][j] = dp[i - 1][j];
            else
                dp[i][j] = dp[i - 1][j] || dp[i - 1][j - set[i - 1]];
        }
    }
    return dp[n][sum];
}
```

```
}

if (!dp[n][sum])

    return false;

for (i=n,j=sum;i>0&&j>0;--i) {

    if (!dp[i- 1][j]) {

        subset.push_back(set[i- 1]);

        j-= set[i- 1];

    }

}

return true;

}

int main() {

    int n,i,sum;

    cout<< "Enter the number of elements in the set: ";

    cin>>n;

    vector<int> set(n);

    cout<< "Enter the elements of the set: ";

    for (i=0;i<n;++i)

        cin>>set[i];

    cout<< "Enter the target sum: ";

    cin>>sum;
```

```
vector<int>subset;

if (isSubsetSumDP(set,sum,subset)){

    cout<< "Found a subset with the given sum!\nSubset: ";

    for (i=0;i<subset.size();++i)

        cout<<subset[i]<< " ";

    cout<<endl;

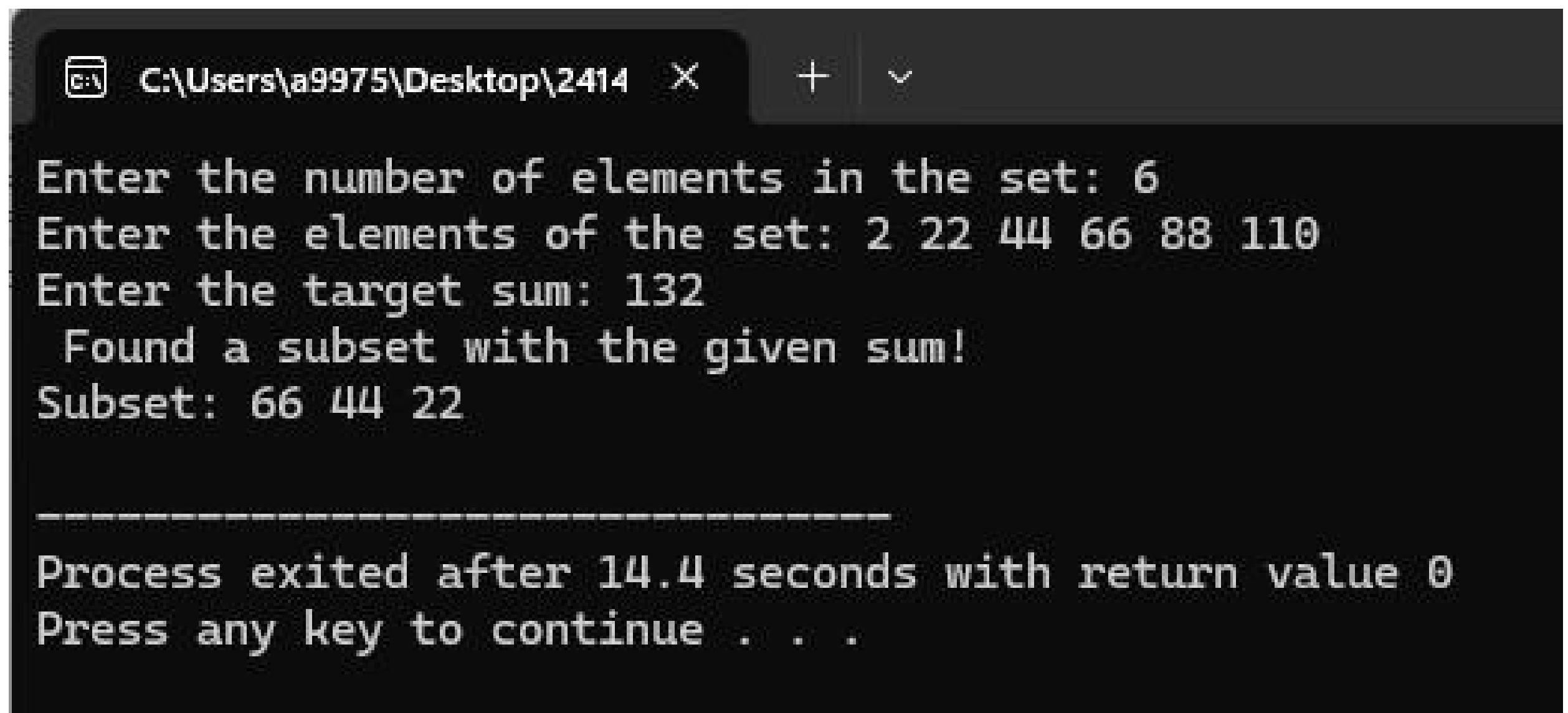
} else{

    cout<< "? No subset found with the given sum." <<endl;

}

return 0;
}
```

## Output:-



```
C:\Users\A9975\Desktop\2414 X + ▾

Enter the number of elements in the set: 6
Enter the elements of the set: 2 22 44 66 88 110
Enter the target sum: 132
Found a subset with the given sum!
Subset: 66 44 22

-----
Process exited after 14.4 seconds with return value 0
Press any key to continue . . .
```

## Algorithm:-

Input: Set of positive integers  $S = \{s_1, s_2, \dots, s_n\}$

Target sum  $d$ .

Output: - All subsets of  $S$  whose elements sum to  $d$   
Step 1: Sort the set  $S$  (optional for optimization).

Step 2: Initialize variables:

$\text{sum} = 0, k = 0$  (starting index).

Step 3: Call  $\text{subset}(k, \text{sum})$  recursively.

Step 4:

$\text{Subset}(k, \text{sum})$ :

1. If  $\text{sum} == d$ :

Print current subset; return.

2. If  $\text{sum} > d$  or  $k >= n$ :

Return (backtrack).

3. Include  $S[k]$  in subset and call  $\text{subset}(k + 1, \text{sum} + S[k])$ .

4. Exclude  $S[k]$  and call  $\text{subset}(k + 1, \text{sum})$ .

Time Complexity:  $O(2^n)$

Space Complexity:  $O(n)$

## List of Applications:-

1. Knapsack problem
2. Resource allocation
3. Budget planning
4. Cryptography
5. Load balancing
6. Data partitioning

7. Combinatorial optimization
8. Decision support systems
9. Scheduling tasks
10. Power set generation