Lab 3

Da) R(A, B, C, D) where AB > C but A > D. The latter does not fulfill 2NF.

b) R(A,B,C,D) where AB > C and C > D. Again, the latter does not fulfill 3NF, as C is determinant for D.

c) R(A, B, L, D) where  $ABC \rightarrow D$ but  $D \rightarrow C$ .

2) ABCPEF

a)  $A \rightarrow D$ ,  $D \rightarrow AE \mid = A \rightarrow AE$  transitive  $A \rightarrow AE \mid = A \rightarrow E$  reflexive

 $A \rightarrow E, E \rightarrow F \mid = A \rightarrow F$  transitive

A > E, A > F, A > D |= A > DEF union

A = AB = AB -> A reflexive

AB -> A, A -> DEF |= AB -> DEF transitive

AB >DEF, AB > C = AB > CDEF union



3) BOOK (Tille#, Title, Author#, BookType,
Price, AuthorName, Publisher) {{Title# > Title, BookType, Publisher}, {Author# > AuthorName}, {BookType > Price}} Title# and Author# are condidate keys.

b) No attribute depends on both key attributes, meaning we split the table, and make Author# a foreign key.

2NF: Book (Title#, Title, Author#, BookType, Price, Publisher)

AUTHOR (Author # Author Name)

3NF: BOOK (Title#, Title, Author#, BookType, Publisher)

AUTHOR (Author#, AuthorName)

BOOKTYPE (BookType, Price)

BCNF: No changes for BCNF

In the previous hand-in, we only had AB. In addition to this, there should also be BD:  $D \to AE = D \to A$  (decomposition)

 $D \rightarrow AE \mid = D \rightarrow A$  (decomposition)  $D \rightarrow AE \mid = D \rightarrow E$  (decomposition)  $D \rightarrow E$ ,  $E \rightarrow F \mid = D \rightarrow F$  (transitive)  $D \rightarrow A$ ,  $AB \rightarrow C \mid = BD \rightarrow C$  (psuedotransitive)  $D \rightarrow A$ ,  $D \rightarrow E$ ,  $D \rightarrow F \mid = D \rightarrow AEF$  (union)  $D \Rightarrow BD \mid = BD \Rightarrow D$  (reflexive)  $BD \rightarrow D$ ,  $D \Rightarrow AEF \mid = BD \Rightarrow AEF$  (transitive)  $BD \rightarrow AEF$ ,  $BD \rightarrow C \mid = BD \Rightarrow AEF$  (union)

This gives A, B and D as prime attributes, and the remaining L, E and F as non-prime.

2) b) As we now know D to be prime, it can remain in R, even though it is only dependent on part of the bey.

R=(A, B, C, D, E, F), where candidate beys are AB, BD. As E depends on D, it is dependent on part of a CK and violates 2NF.

 $R_1 = (A, B, C, D)$   $R_2 = (D, E, F)$  (K: AB, BD) (K: D)  $\{\{AB \rightarrow C\}, \{A \rightarrow D\}, \{D \rightarrow A\}\}$   $\{\{D \rightarrow E\}, \{E \rightarrow F\}\}$ Now in 2NF

C) Ry is already in 3NF.

Rz must be broken up.

$$\ell_2 = (D, E)$$
  $\ell_3 = (E, F)$ 

(K: D LK: E

$$\{\{D \rightarrow F\}\}$$

Now in 3HF

d) Rz and Rz alrealy in BCNF Rz must be broken up, as D is determined only by A.

$$R_1 = (A,B,C)$$
  $R_4 = (A,D)$   
 $(K:AB)$   $(K:A)$ 

$$\{\{AB \rightarrow C\}\}$$

3) B) As authorst does not depend on titlett, it was incorrect to make authorst a toreign key of the BOOK entity. More correct would have been to split the entity into two completely disjunct entities, meaning we mitially have three tables for 2NF:

BOOK = (Title#, Title, BookType, Publisher)
AUTHOR = (Author#, AuthorName)
PUBLICATION = (Title#, Author#)

For 3NF, more out BookType and Prize.

BOOK = (Title#, Title, BookType)

BOOKTYPE = (BookType, Price)

Other two tables as in 2NF

No changes for BLNF.