

# **Sales analysis**

# **MILAVITSA'S SALES ANALYSIS**

**Milavitsa-Clothes** 

Legal Notice:

This document contains privileged and/or confidential information and may not be disclosed, distributed or reproduced without the prior written

permission of EPAM®.

CONFIDENTIAL | Effective Date: 3-Dec-17

# **CONTENTS**

1 BUSINESS DESCRIPTION	4
1.1 BUSINESS BACKGROUND	4
1.2 PROBLEMS BECAUSE OF POOR DATA MANAGEMENT	4
1.3 BENEFITS FROM IMPLEMENTING A DATA WAREHOUSE	
2 DATA DESCRIPTION	6
2.1 DESCRIPTION OF SELECTED SCHEMA	
2.2 DESCRIPTION OF DATA SOURCES	
2.3 DESCRIPTION OF BUSINESS RULES	6
3 DIMENSIONS OF A BUSINESS	7
3.1 DEFINITION OF SELECTED BUSINESS PROCESS	
3.2 DECLARING THE GRAIN	
3.3 IDENTIFYING THE DIMENSIONS	7
3.4 IDENTIFYING THE FACTS	
3.5 DIMENSIONAL MODEL. STAR SCHEMA	8
4 NF-LAYER OF DATA WAREHOUSE	
4.1 3NF-MODEL	. 13
5 OBJECT PARTITIONING	. 17
6 BUSINESS PROCESSES	. 19
6.1 SALES ANALYSIS	. 20
6.2 ANALYSIS OF RECEIPTS	. 21
6.3 LOAD ANALYSIS	. 22
7 DATA MODELLING	
7.1 DETAILING DIAGRAMS FOR 3NF AND STAR/SNOWFLAKE LAYERS	. 23
7.2 TEXTUAL DESCRIPTION OF LAYERS OF DATA WAREHOUSE	. 23
7.2.1 Source layer	. 23
7.2.2 Data staging layer	. 23
7.2.3 Data warehouse layer	. 24
7.2.4 Analysis layer	. 24
7.3 VISUAL DESCRIPTION OF LAYERS OF DATA WAREHOUSE	. 25
8 FACT TABLE PARTITIONING STRATEGY	. 31
9 STRATEGY OF PARALLEL LOAD	. 32

10	REPORT LAYOUTS	37
10.1	3NF-LAYER	37
10.2	DM-LAYER	40

### 1 BUSINESS DESCRIPTION

#### 1.1 BUSINESS BACKGROUND

For many years Milavitsa has been producing lady's lingerie, being one of the biggest lingerie producer in Eastern Europe. The company's products are successfully sold in more than 25 countries around the world and are characterized by high quality, original and fashionable designs.

Company has four main brands:

Milavitsa.

The Milavitsa product portfolio includes bras, knickers, shapewear, knitwear and swimwear. The Milavitsa collection is divided into three categories: classic, fashion and swimwear. The basis of the classic collection is formed by a large variety of everyday styles, combining the basics of classical design, comfortable construction and functional materials. Milavitsa is an expert in creating the styles of large sizes. Fashion and swimwear collections are created for every season, following the fashion trends in design, materials and accessories.

Avelive Collection.

Bras, knickers, shapewear are also manufactured under the Aveline trademark. This is comfortable everyday lingerie at affordable prices. The Aveline products are developed and manufactured by Milavitsa, which is a guarantee of high quality and compliance with the company's standards.

3. Alisee.

Alisee is a French lingerie brand acquired by Milavitsa. Alisee collection is designed and styled by European professionals. Tailored by Milavitsa to fit local market specifics.

4. Hidalgo.

The Hidalgo is men's underwear combines classic shapes and comfortable natural materials.

## 1.2 PROBLEMS BECAUSE OF POOR DATA MANAGEMENT

Company faced next problems:

- 1. absence of business intelligence from several sources;
- 2. absence of sales and inventory information consolidation for the calculation of the optimal order and delivery;
- 3. decreasing query and system performance;
- 4. absence of timely access to data;
- 5. absence of historical intelligence.

#### 1.3 BENEFITS FROM IMPLEMENTING A DATA WAREHOUSE

Better decision-making.

Corporate decision makers will no longer have to make important business decisions based on limited data and hunches. Data warehouse will store credible facts and statistics, and decision makers will be able to retrieve that information from the data warehouse based on their personal needs.

2. Quick and easy access to data.

Speed is an important factor that sets company above its competitors. Business users can quickly access data from multiple sources from a data warehouse, meaning that precious time will not be wasted on retrieving data from multiple sources. This allows company to make quick and accurate decisions, with little or no support from its IT department.

3. Data quality and consistency.

Since data warehouses gather information from different sources and convert it into a single and widely used format, departments will produce results that are in line and consistent with each other. When data is standardized, company can have confidence in its accuracy, and accurate data is what makes for strong business decisions.

## **2 DATA DESCRIPTION**

### 2.1 DESCRIPTION OF SELECTED SCHEMA

The Star schema was chosen for business processes description.

The main reasons:

- simple structure;
- 2. absence of a big number of tables to join;
- 3. denormalized tables are not too large in a specific case of this task;
- 4. widely support by a large number of business intelligence tools.

### 2.2 DESCRIPTION OF DATA SOURCES

Customer and manager information was generated on the site www.mockaroo.com. It allows creating file with 1000 rows.

Additionally, some specific information was modified in the Excel. Information in these columns was created by the specific random formula.

Information for Collection, Line, Product type, Store dimensions is on the official company's site www.milavitsa.com. Information about Sizes can be find via the next links www.milavitsa.com/collections/converter and www.globebrand.com/sizing\_charts

#### 2.3 DESCRIPTION OF BUSINESS RULES

Business rules:

- 1. Every retail sale must be associated with a valid Customer.
- 2. A retail sale is always associated with a Payment Method.
- 3. A retail sale can have one or many Products.
- 4. A retail sale is always associated with a specific Employee.

# **3 DIMENSIONS OF A BUSINESS**

### 3.1 DEFINITION OF SELECTED BUSINESS PROCESS

Business Process for analysis is Milavitsa's sales per different metrics.

### 3.2 DECLARING THE GRAIN

The highest fact granularity is daily sales amount per certain customer and employee in a specific a store.

### 3.3 IDENTIFYING THE DIMENSIONS

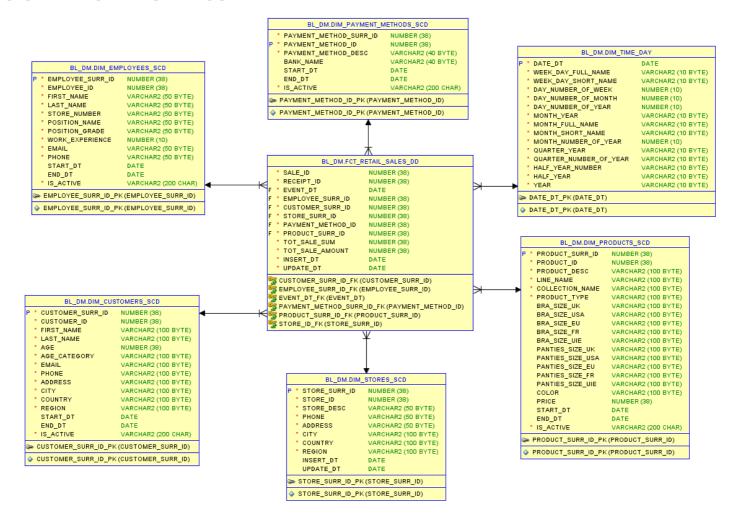
Schema should contains next dimensions:

- Dim\_customers\_scd:
- 2. Dim\_products\_scd:
- 3. Dim\_stores:
- Dim\_payment\_methods\_scd;
- Dim\_employees\_scd;
- 6. Dim\_time\_day.

## 3.4 IDENTIFYING THE FACTS

Schema should contains the fct\_retail\_sales\_dd fact table.

#### 3.5 DIMENSIONAL MODEL. STAR SCHEMA

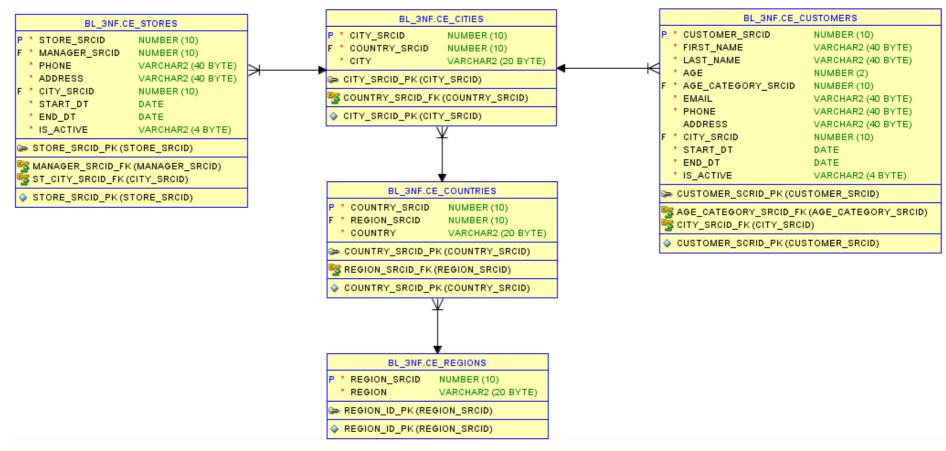


Picture 1 Dimensional model

## **4 NF-LAYER OF DATA WAREHOUSE**

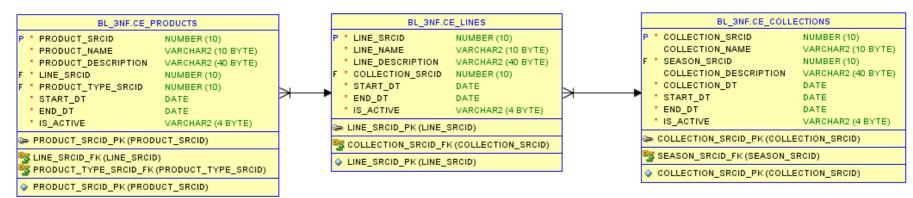
In this model the normalization was made with next steps:

1. Geographical information from the "Stores" and "Customers" was separated on different tables. Tables were created for each geographical object: region -> country -> city. These objects were connected in series: from the lowest in granularity to the highest.



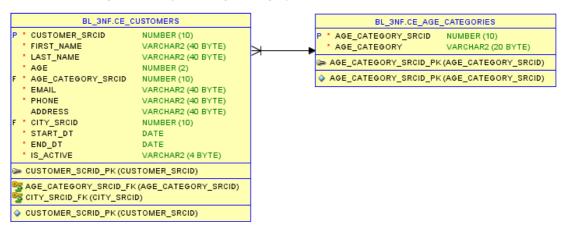
Picture 2 Geography part

2. Information about products was organized in the same way. There were separate tables for each object: collections -> lines -> products. These objects were connected in series: from the lowest in granularity to the highest.



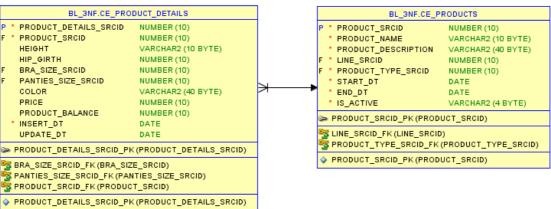
Picture 3 Product hierarchy part

3. A separate "Age\_categories" table was created for Customers. This table describes all possible age ranges that are important for business. Subsequently, each Customer can be assigned to a specific age category.



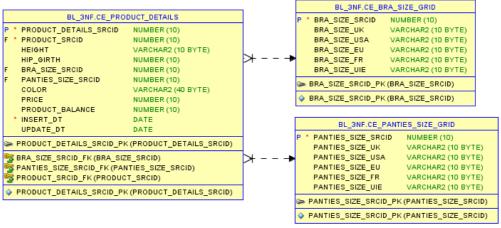
Picture 4 Customer part

4. Information about product characteristics was separated in a specific table "Product\_details". It contains a complete description of each commodity unit of a certain size.



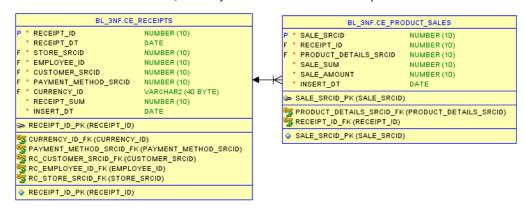
Picture 5 Product part

5. The tables of the size grids "Bra\_size\_grid" and "Panities\_size\_grid", which contain information on the international dimension matching for bra and penities, were separately created.



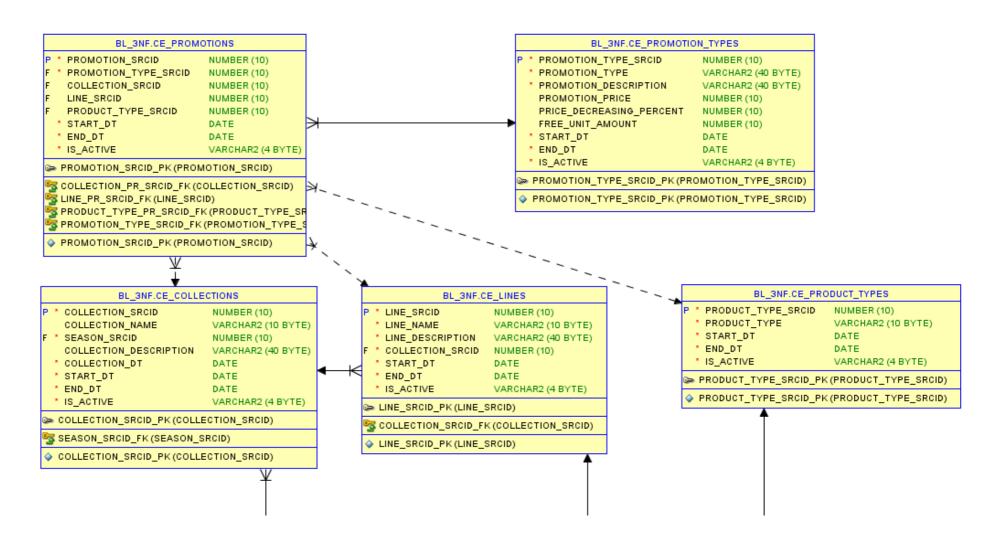
Picture 6 Size grid part

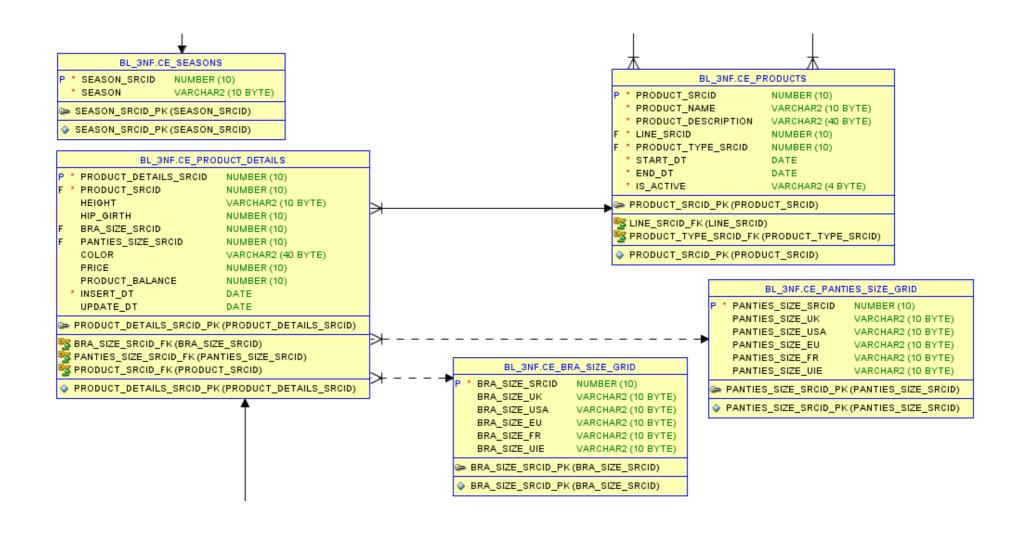
- 6. Additionally receipts information was add. It includes the calculation of their total cost.
- 7. The "Product\_sales" table details the check information, namely it describes a concrete product and amount of units what were purchased.

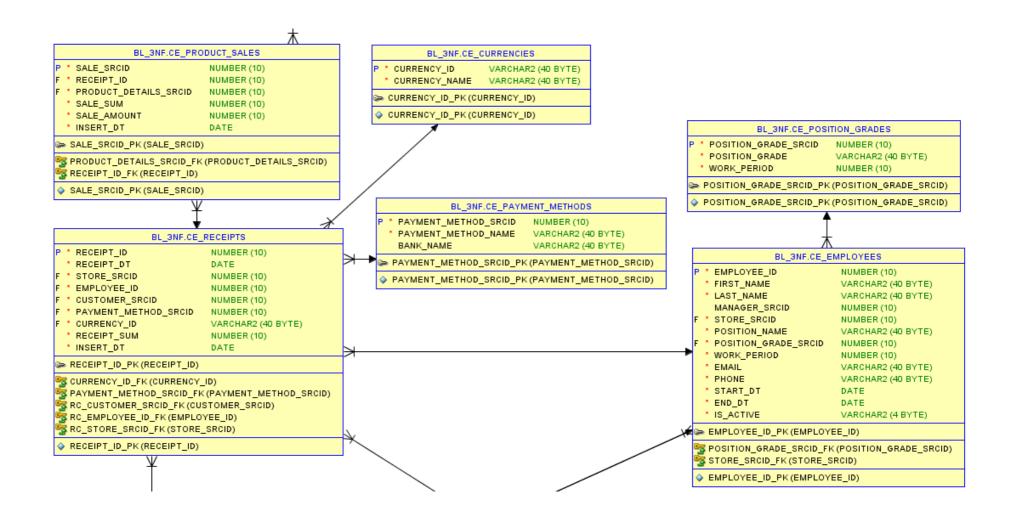


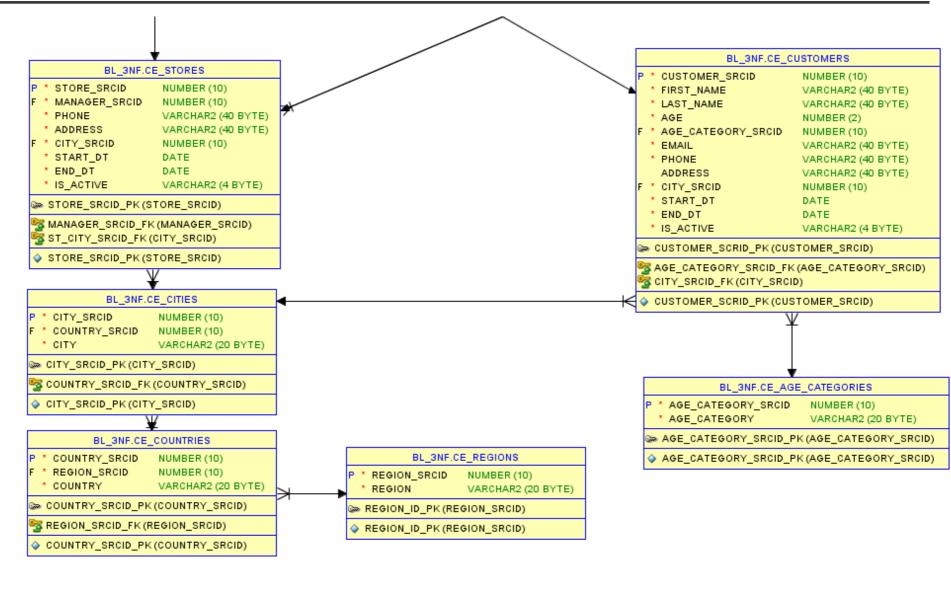
Picture 7 Sales part

#### 4.1 3NF-MODEL









Picture 8 3NF model

## 5 OBJECT PARTITIONING

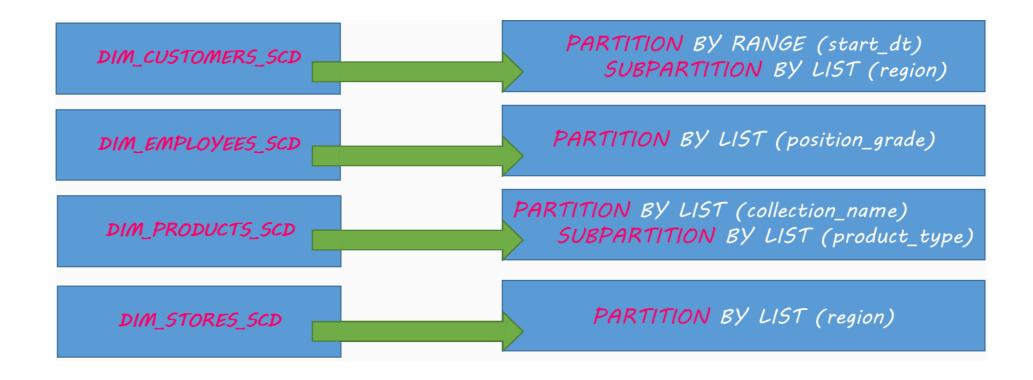
A combination of two methods of data distribution is used to define a table with composite partitioning. Firstly, the table is partitioned by one method of data distribution. Then each section is subdivided into subsections using another method of data distribution. For example, if a range-list is specified for a table, then at the beginning the table is partitioned with keys ranges, and then each section is divided into subsections with the keys lists specified in the second partition. Sections of the table with composite partitioning exist only as metadata and do not provide actual data storage: the subsections of a particular section of a table or an index with composite partitioning are the physical segments of the database where the section data is stored.

In this paper, the partitions of the dimension tables were applied with list and range types of partitions. As indicated in the picture below, the customer dimension was divided into partitions by the year when the client appeared in the database. In turn, each year was divided with sub partitions by the regions. This is very convenient, since it will allow to investigate various cohorts of customers and their activity in time.

Dimension with the employees was divided according to the position grades existing in the company in the timetable. This breakdown will allow to track faster active groups of employees depending on the status.

The product dimension was divided by the collections. Each collection, in turn, is divided by product types, because sales and demand for individual products are usually not tracked, so it's better to take product types into collections as the grain.

The last dimension shops was partitioned by regions. As the sale of a huge holding is usually investigated initially within the region. In a particular region it is easier to find the country of interest than in all dimension.



Picture 9 Partition strategy

## **6 BUSINESS PROCESSES**

In the analysis two business processes will be explored. The first of them is monthly sales of stores within regions, countries and cities, as well as sales by customers, the second - monthly commodity stocks of stores in different geographical sections and per network managers to identify the most valuable personnel.

Three reports are proposed as reports:

1. Sales analysis.

This table will be designed to provide general information about the functioning stores, sales, the number of products in stock, and the number of receipts. The lowest level of granularity will be represented by regions, and then DRILL-DOWN technology will make it possible to deepen the detailed analysis of analytics to countries and cities. The highest level of hierarchy will be presented by shops with a specific address in the city.

The following parameters will be presented in the table: the turnover for the previous period - the number of months of the previous period can be determined by specifying the corresponding value in the "Comparison period" field, the current period can be determined in the "Analysis period" field. The completion percentage will show what percentage of the revenue of the previous period is reached for the current period. The indicator quantity of goods determines the commodity remains, and the number of receipts shows how many sales were made for the analyzed period.

2. Analysis of receipts.

This report will allow user to view value of the main indicators per the metrics selected in the ListBox. In each ListBox, user can select multiple values at the same time. The analysis period can be determined by analogy with the previous report. The choice of month and year is provided using the ListBox.

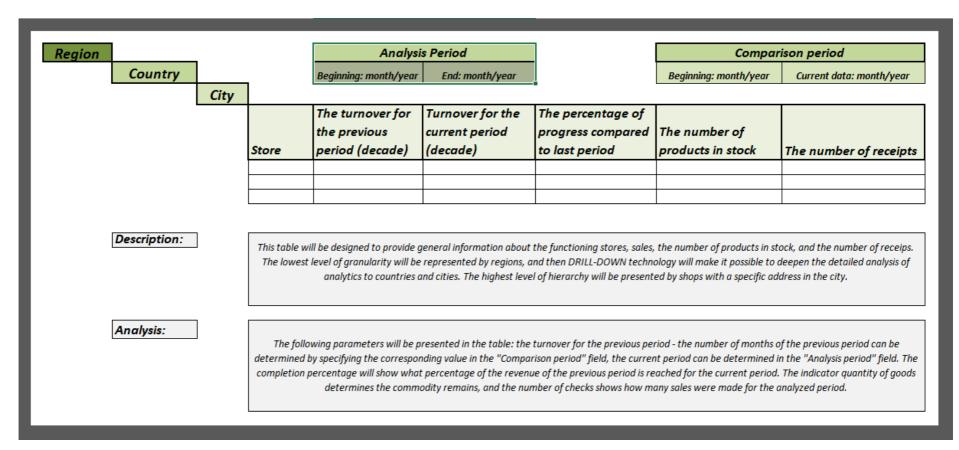
3. Load analysis.

This analysis is used to investigate customers' influx during working days and working hours of the day what will be chosen for analysis. This analysis allows user to see the most favorable working periods for the promotions, as well as the time when it is necessary to have more staff in stores.

4. Data generation.

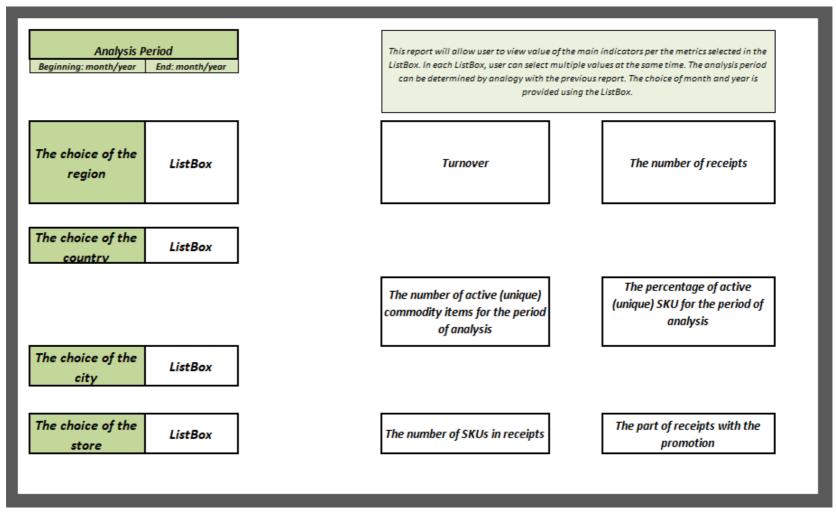
This tab contains tables with generated data to plot graphs in the "Load Analysis" report.

#### **6.1 SALES ANALYSIS**



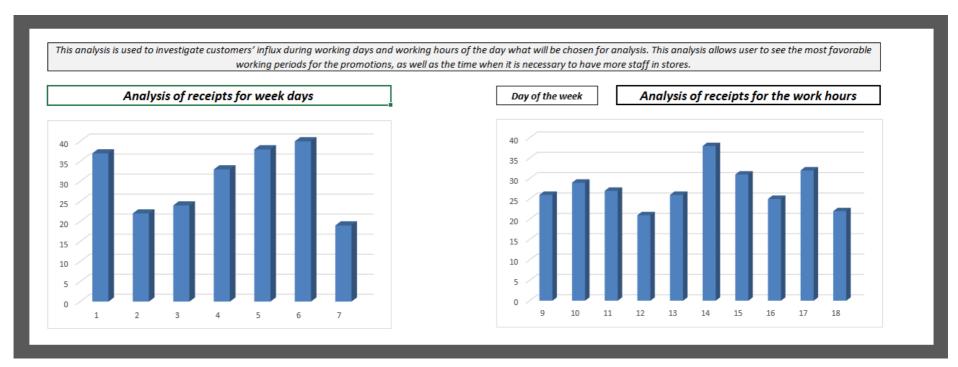
Picture 10 Sales analysis

### **6.2 ANALYSIS OF RECEIPTS**



Picture 11 Analysis of receipts

### **6.3 LOAD ANALYSIS**



Picture 12 Load analysis

## **7 DATA MODELLING**

#### 7.1 DETAILING DIAGRAMS FOR 3NF AND STAR/SNOWFLAKE LAYERS

There were created seven dimensions. They are described in "Business Description #4" document. Several dimensions have a SCD2 type. It was chosen due to It is the next dimensions:

- dim\_customers\_scd;
- dim\_employees\_scd;
- dim\_payment\_methods\_scd;
- dim\_products\_scd;
- dim\_promotions\_scd;

Dim\_time\_day has calendar type. Dim\_stores has a SCD1 type, because it is not necessary to have history about closed stores.

#### 7.2 TEXTUAL DESCRIPTION OF LAYERS OF DATA WAREHOUSE

In this work was used a two-layer architecture.

## 7.2.1 Source layer

There were used flat files and html-code as data sources.

## 7.2.2 Data staging layer

The data stored to sources should be extracted, cleansed to remove inconsistencies and fill gaps, and integrated to merge heterogeneous sources into one common schema.

In this work will be used the following data staging layers step by step.

#### 7.2.2.1 Cleansing layer

This layer was used for data cleansing, filtering wrong data, replace missing values with singletons and performing transformations like code lookups or currency conversions. As the Staging Area, the Cleansing Area contains only data of the last delivery, and data from different sources is not integrated.

## 7.2.2.2 3NF layer

Here reducing data redundancy and improving data integrity were made. Also, the process of simplifying the design of a database, so that it achieves the optimal structure composed of atomic elements, was achieved on this layer.

## 7.2.2.3 Dimensional model layer

The Star schema was chosen as a dimensional model for business processes description.

The main reasons:

- 1. Simple structure;
- 2. Absence of a big number of tables to join;
- 3. Denormalized tables are not too large in a specific case of this task;
- 4. Widely support by a large number of business intelligence tools.

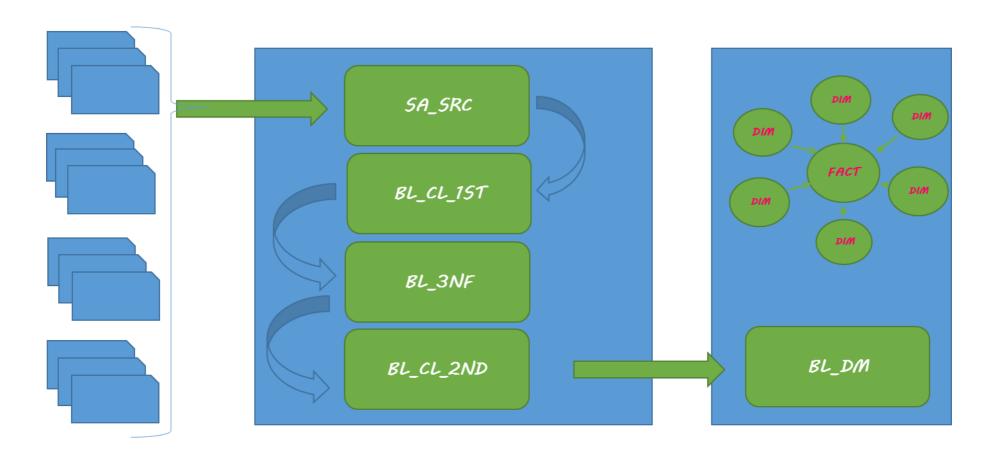
## 7.2.3 Data warehouse layer

Information is stored to one logically centralized single repository: a data warehouse. The data warehouse can be directly accessed, but it can also be used as a source for creating data marts, which partially replicate data warehouse contents and are designed for specific departments. This layer is added after finishing previous steps.

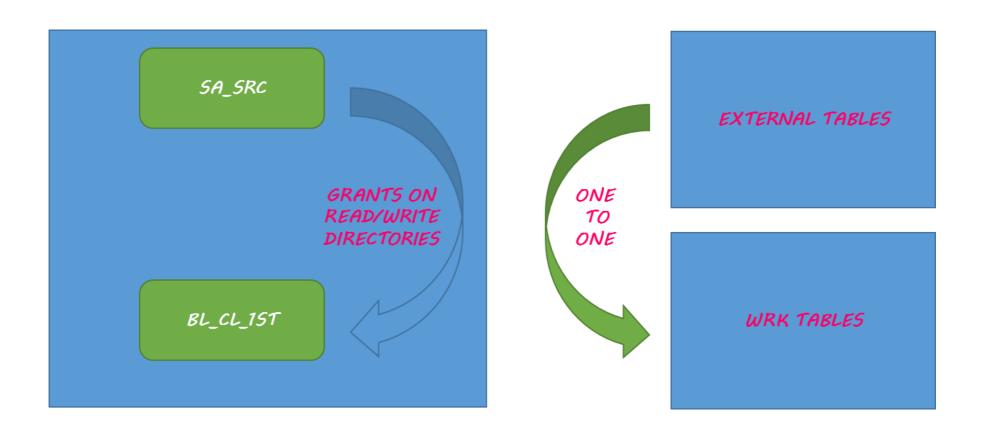
## 7.2.4 Analysis layer

In this layer, integrated data is efficiently and flexibly accessed to issue reports, dynamically analyze information, and simulate hypothetical business scenarios. Template of this layer is represented in "Retail Analysis" document.

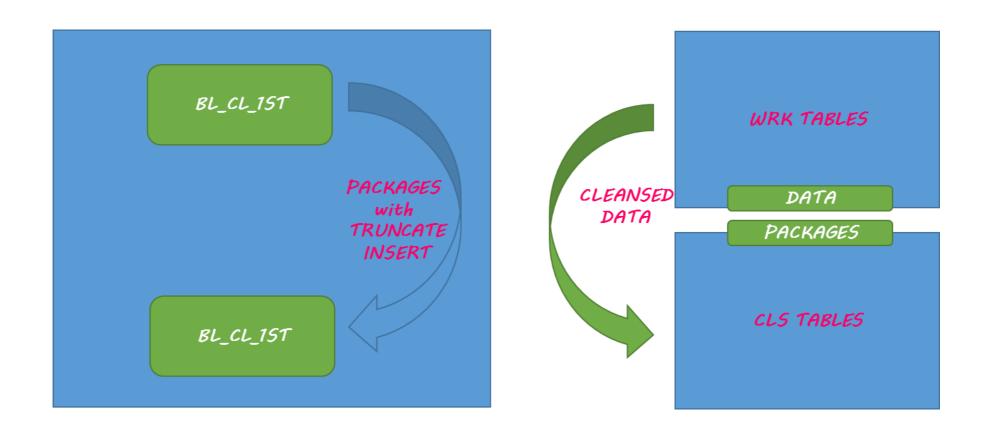
## 7.3 VISUAL DESCRIPTION OF LAYERS OF DATA WAREHOUSE



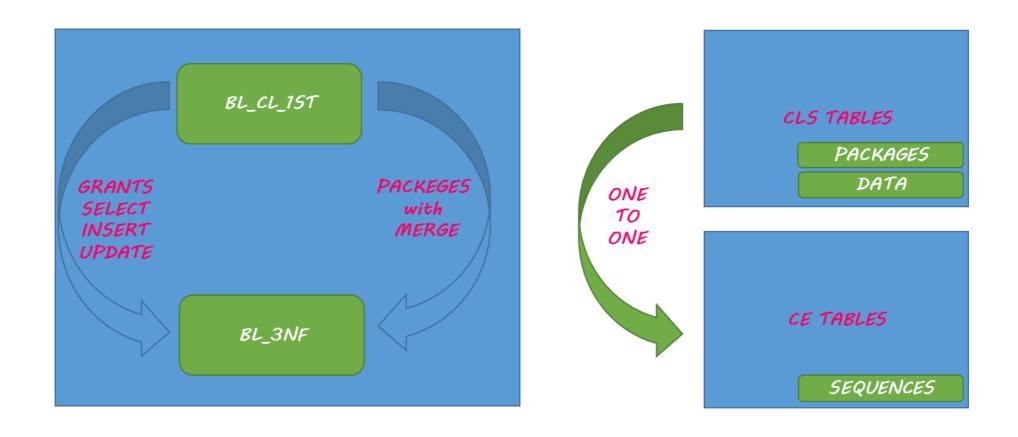
Picture 13 Layers model



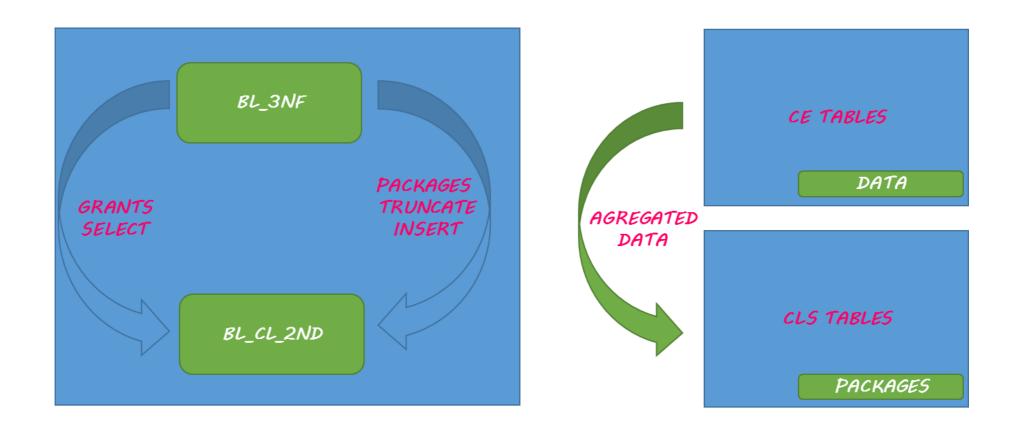
Picture 14 SA SRC and BL CL 1ST



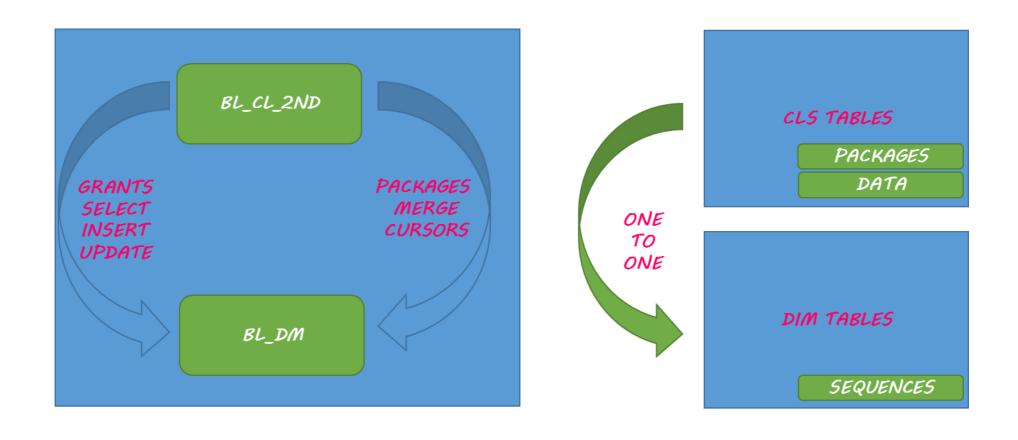
Picture 15 BL CL 1ST



Picture 16 BL CL 1ST and BL 3NF



Picture 17 BL 3NF and BL CL 2<sup>nd</sup>



Picture 18 BL CL 2<sup>nd</sup> and BL DM

## **8 FACT TABLE PARTITIONING STRATEGY**

Among the available partitioning methods, the above-described RANGE was applied in to fact table.

This partitioning method RANGE-is most convenient for the fact table. It will be divided into year's partitions, and every year, in turn, will contain partitions with specific regions where sales are conducted. As the sales network is quite extensive, the granularity in the partitions to the level of the cities is not expedient. For business, it will be much more convenient to track sales within a specific region.



Picture 19 Partition strategy

## 9 STRATEGY OF PARALLEL LOAD

Parallel data loading can improve the performance of DML operators performed radically. With consecutive execution of operations, each subsequent must wait for the end of the preceding one. Parallel execution helps to avoid downtime and allows the next in turn operation to begin processing data that has already been processed by the first operation, even though the first operation was not completed completely.

When loading the generated data, this possibility can be applied to uploading data to the BL\_DM layer. Since the SELECT and INSERT operations for each table are used for this table, PARALLEL LOAD will allow to select data for insertion and load the already selected data simultaneously, which in theory should improve the performance of our queries.

Let's check the time plan for the query execution before applying the hint on the fact table.

1	Pla	an h	nas	h value: 4253285374												
2																
3																
4	1	Id	1	Operation	Name	1	Rows	l E	Bytes	Cost (	€CPU)	Time	Ī	Pstart	Pstop	1
5																
6	I	0	1	INSERT STATEMENT	1	1	1	ı	185 I	13269	(1)	00:00:01	Ī	- 1		1
7	I	1	1	LOAD TABLE CONVENTIONAL	FCT_RETAIL_SALES_DD	1		I	ı		- 1		Ī	- 1		1
8	I	2	1	SEQUENCE	FCT_RETAIL_SALES_DD_SEQ	1		I	- 1		- 1		Ī	- 1		1
9	*	3	1	HASH JOIN		-1	1	I	185	13269	(1)	00:00:01	Ī	- 1		1
10	I	4	1	NESTED LOOPS		-1	1	I	159	12177	(1)	00:00:01	Ī	- 1		1
11	<b> </b> *	5	1	HASH JOIN	1	-1	1	I	146	12177	(1)	00:00:01	Ī	1		1
12	*	6	1	HASH JOIN		-1	1	I	120	10539	(1)	00:00:01	Ī	- 1		1
13	*	7	1	HASH JOIN	1	-1	1	I	94	8901	(1)	00:00:01	I	1		1
14	I	8	1	PARTITION RANGE AL	L	-1	1	I	26	3375	(1)	00:00:01	Ī	1	11	1
15	I	9	1	PARTITION LIST AL	L	-1	1	I	26	3375	(1)	00:00:01	Ī	1	6	1
16	I	10	1	TABLE ACCESS FUL	L  DIM_CUSTOMERS_SCD	1	1	I	26	3375	(1)	00:00:01	Ī	1	66	1
17	I	11	1	TABLE ACCESS FULL	CLS_FCT_RETAIL_SALES_DD	1	2001K	I	129M	5520	(1)	00:00:01	Ī	- 1		1
18	I	12	1	PARTITION LIST ALL	I	-1	1	I	26	1638	(1)	00:00:01	Ī	1	6	1

2,001,998 rows inserted.

Elapsed: 00:06:53.634

Picture 20 Explain plan

The query execution plan using - + PARALLEL (4).

1	Pla	an	has	sh value: 27664	09812									
2														
3														
4	1	Id	ī	Operation			Name	1	Rows	Bytes	Ī	Cost (	%CPU)∣	Time
5														
6	ı	0	1	INSERT STATEME	NT	- 1		1	1	185	Ι	3683	(1)	00:00:01
7	ı	1	1	LOAD TABLE CO	NVENTIONAL	-	FCT_RETAIL_SALES_DD	1	- 1		Τ		I	
8	I	2	-1	SEQUENCE		-	FCT_RETAIL_SALES_DD_SEQ	1	- 1		1		1	
9	I	3	-	PX COORDINA	TOR			1	- 1		1		1	
10	l	4	1	PX SEND QC	(RANDOM)		:TQ10006	1	1	185	1	3683	(1)	00:00:01
11	*	5	1	HASH JOIN		- 1		1	1	185	1	3683	(1)	00:00:01
12	l	6	1	PX RECEI	VE	- 1		1	1	159	1	3380	(1)	00:00:01
13	l	7	-1	PX SEND	BROADCAST	-	:TQ10005	1	1	159	1	3380	(1)	00:00:01
14	I	8	-1	BUFFER	SORT	-		1	902	126F	۲)		1	
15	L	9	1	NESTE	D LOOPS	- 1		1	1	159	1	3380	(1)	00:00:01
16	*	10	1	HASH	JOIN	- 1		1	1	146	1	3380	(1)	00:00:01
17	l	11	-1	PX	RECEIVE	- 1		1	1	120	1	2925	(1)	00:00:01
18	l	12	1	PX	SEND HYBRID HASH	- 1	:TQ10003	1	1	120	1	2925	(1)	00:00:01

2,001,998 rows inserted.

Elapsed: 00:09:02.015

Picture 21 Explain plan

You can see that despite the decrease in the COST parameter, the parallel in this case had the opposite effect - the query execution time increased by two minutes, which is very significant. In this case, the use of a parallel is impractical. Let's try to increase the number of parallel processes to 20.

The query execution plan using - + PARALLEL (10).

1	Pl	lan	ha	ash	value: 2766409812										
2															
3						 									
4	I	Id		1 0	peration	Name		Ro	DWS	1	Bytes	T	Cost	(%CPU)	Time
5						 									
6	Ī		0	l I	NSERT STATEMENT	I		1	1	1	185	Τ	1473	(1)	00:00:01
7	Ī		1	I	LOAD TABLE CONVENTIONAL	FCT_RETAIL	_SALES_DD	1		1		1		- 1	
8	Ī		2	I	SEQUENCE	FCT_RETAIL	_SALES_DD_SEQ	1		1		1		- 1	
9	Ī		3	I	PX COORDINATOR	I		1		1		1		- 1	
10	Ī		4	I	PX SEND QC (RANDOM)	:TQ10006		1	1	1	185	1	1473	(1)	00:00:01
11	19	*	5	I	HASH JOIN	I		1	1	1	185	1	1473	(1)	00:00:01
12	Ī		6	I	PX RECEIVE	I		1	1	1	159	1	1352	(1)	00:00:01
13	I		7	I	PX SEND BROADCAST	:TQ10005		1	1	1	159	1	1352	(1)	00:00:01

2,001,998 rows inserted. Elapsed: 00:20:18.325

Picture 22 Explain plan

This request was carried out for 20 minutes.

Obviously, the time for downloading is only increased using PARALLEL on the fact table. We test the hint on the DIMENSION table DIM\_PRODUCTS\_SCD, because it is the largest of the DIMENSION tables.

The query execution plan without application is + PARALLEL.

1 P	la	n h	ıas	h value: 24526241											
2															
3 -															
4	Ι	d	ī	Operation	ī	Name	1	Rows	ī	Bytes  TempS	pcl	Cost	(%CPU)	Time	ī
5 -															
6		0	1	MERGE STATEMENT	1		Ī	6000	Ī	14M	- 1	182	(2)	00:00:01	- 1
7		1	Ī	MERGE	1	DIM_PRODUCTS_SCD	1		ī	I	- 1		-		1
8		2	T	VIEW	1		1		1	1	-1		- 1		- 1
9		3	Ī	SEQUENCE	1	DIM_PRODUCTS_SEQ	1		T	1	- 1		- 1		- 1
10	*	4	T	HASH JOIN RIGHT OUTER	-1		1	6000	1	14M	-1	182	(2)	00:00:01	. 1
1		5	1	PARTITION LIST ALL	-1		1	1	1	1251	-1	2	(0)	00:00:01	- 1
12		6	T	PARTITION LIST ALL	1		1	1	1	1251	- 1	2	(0)	00:00:01	- 1
13		7	T	TABLE ACCESS FULL	1	DIM_PRODUCTS_SCD	1	1	1	1251	- 1	2	(0)	00:00:01	. 1
14		8	T	VIEW	1		1	6000	1	7183K	-1	180	(2)	00:00:01	. 1
.5 <sub> </sub>		9	Ī	MINUS	1		1		1	1	- 1		- 1		-1
6 1		10		SORT UNIQUE	1		ī	6000	ī	621KI 92	OKI				

6,000 rows merged.

Elapsed: 00:00:01.440

Picture 23 Explain plan

The query execution plan using - + PARALLEL (4).

1	Pla	an l	ha	sh value: 3260348670													
2																	
3																	
4	1	Id	1	Operation	1	Name	1	Rows	1	Bytes  TempSp	cl C	ost (	(%CPU)	Time	L	Pstart	Pstop
5																	
6	I	0	1	MERGE STATEMENT	I		L	6000	I	14M	1	15	(20)	00:00:01	I	- 1	
7	I	1	-1	MERGE	1	DIM_PRODUCTS_SCD	L		I	1	1		-		T	- 1	
8	L	2	-1	VIEW	I		L		I	1	-1		1		L	- 1	
9	I	3	-	SEQUENCE	1	DIM_PRODUCTS_SEQ	L		I	1	1		1		I	- 1	
10	L	4	-1	PX COORDINATOR			L		1	1	-1		1		I	- 1	
11	L	5	-1	PX SEND QC (RANDOM)	L	:TQ10003	L	6000	I	14M	1	15	(20)	00:00:01	T	- 1	
12	*	6	-1	HASH JOIN RIGHT OUTER BUFFERED	1		L	6000	I	14M	-1	15	(20)	00:00:01	1	- 1	
13	L	7	-1	PX PARTITION LIST ALL	I		L	1	I	1251	-1	2	(0)	00:00:01	L	1	3
14	I	8	1	PX PARTITION LIST ALL	I		l	1	I	1251	1	2	(0)	00:00:01	I	1	2
15	L	9	-1	TABLE ACCESS FULL	1	DIM_PRODUCTS_SCD	L	1	I	1251	-1	2	(0)	00:00:01	I	1	6
16	L	10	-1	PX RECEIVE	1		L	6000	1	7183K	1	13	(24)	00:00:01	1	- 1	

6,000 rows merged.

Elapsed: 00:00:01.204

Picture 24 Explain plan

Dimension tables are much smaller than the fact table and are loaded quite quickly, so the use of parallel loading on them is also inappropriate. Acceleration of the request for a millisecond will not make a significant difference to the performance in the projected DWH. Thus, the parallel boot strategy will not be applied in this work.

## **10 REPORT LAYOUTS**

#### **10.1 3NF-LAYER**

Execute the query with the application of analytical functions.

Let's analyze sales for 2018 with grouping by quarter, month and day.

```
SELECT DECODE (GROUPING ID (to char (receipt dt, 'YYYY'),
                           upper(TO CHAR(receipt dt,'YYYY')) || '-' || 'Q' || TO CHAR(receipt dt,'Q'),
                           upper(TO CHAR(receipt dt, 'YYYY') || '-' || TO CHAR(receipt dt, 'Mon') ), receipt dt), 7,
                           'GRAND TOTAL FOR ' || to char(receipt dt, 'YYYY'), ' ') AS year,
       DECODE (GROUPING ID (to char (receipt dt, 'YYYY'),
                           upper(TO CHAR(receipt dt, 'YYYY')) || '-' || 'Q' || TO CHAR(receipt dt, 'Q'),
                          upper (TO CHAR (receipt dt, 'YYYY') || '-' || TO CHAR (receipt dt, 'Mon') ), receipt dt), 3,
                           'GRAND TOTAL FOR ' || upper(TO CHAR(receipt dt,'YYYY')) || '-' || 'Q' ||
TO_CHAR(receipt_dt,'Q'), ' ') AS quarter,
       DECODE (GROUPING ID (to char (receipt dt, 'YYYY'),
                          upper(TO CHAR(receipt dt,'YYYY')) || '-' || 'Q' || TO CHAR(receipt dt,'Q'),
                           upper (TO CHAR (receipt dt, 'YYYY') || '-' || TO CHAR (receipt dt, 'Mon') ), receipt dt), 1,
                           'GRAND TOTAL FOR ' || upper(TO CHAR(receipt dt,'YYYY') || '-' || TO CHAR(receipt dt,'Mon')
), ' ') AS month,
       DECODE (GROUPING (receipt dt), 1, '', receipt dt) AS day,
       TO CHAR(SUM(receipt sum usd), '9,999,999,999') as sales
FROM ce receipts dt
```

			∯ DAY	SALES
363			15-OCT-18	98,632,990
364			16-OCT-18	101,978,967
365			17-OCT-18	99,178,589
366			18-OCT-18	99,693,233
367			19-OCT-18	101,251,898
368			20-OCT-18	99,724,227
369			21-OCT-18	97,660,012
370			22-OCT-18	95,913,192
371			23-OCT-18	100,312,667
372			24-OCT-18	96,909,652
373			25-OCT-18	99,616,695
374			26-OCT-18	101,978,782
375			27-OCT-18	104,302,014
376			28-OCT-18	101,581,387
377			29-OCT-18	103,801,735
378			30-OCT-18	102,354,341
379			31-OCT-18	99,523,147
380		GRAND TOTAL FOR 2018-OCT		3,104,654,916
381	GRAND TOTAL FOR 2018-Q			9,238,938,174
382 GRAND TOTAL FOR 2018				***********
383				***********

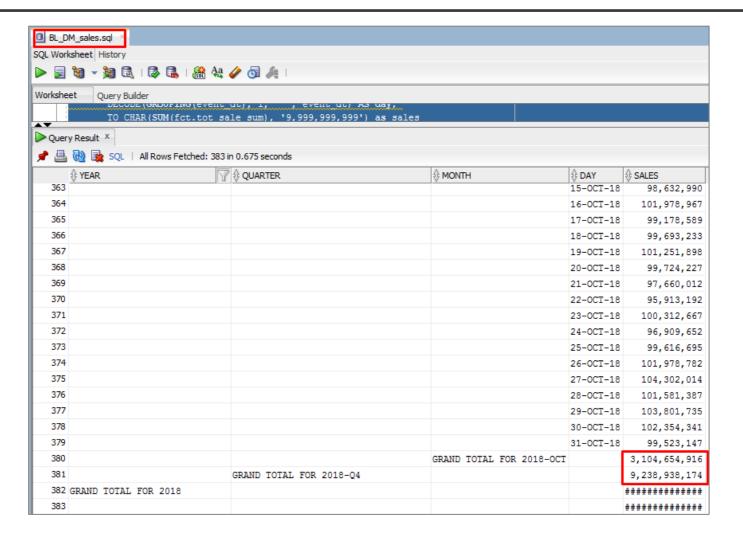
#### All Rows Fetched: 383 in 2.394 seconds

Picture 25 Explain plan

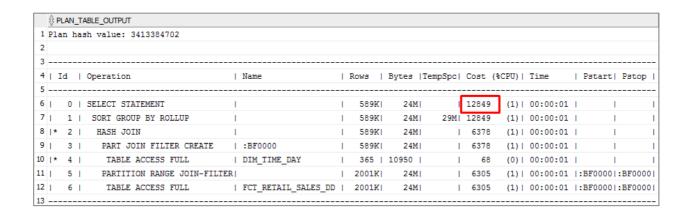
### 10.2 DM-LAYER

Perform a similar query on the DM layer.

```
SELECT DECODE (GROUPING ID (dt.year, dt.quarter year, dt.month year, event dt), 7, 'GRAND TOTAL FOR ' || dt.year, ' ')
AS year,
       DECODE (GROUPING ID (dt.year, dt.quarter year, dt.month year, event dt), 3, 'GRAND TOTAL FOR ' ||
dt.quarter year, ' ') AS quarter,
       DECODE (GROUPING ID (dt.year, dt.quarter year, dt.month year, event dt), 1, 'GRAND TOTAL FOR ' ||
dt.month year, ' ') AS month,
       DECODE (GROUPING (event dt), 1, '', event dt) AS day,
      TO CHAR(SUM(fct.tot sale sum), '9,999,999,999') as sales
FROM fct retail sales dd fct,
       dim time day dt
WHERE dt.date dt = fct.event dt
 AND dt.year = 2018
GROUP BY ROLLUP (
                dt.year,
                dt.quarter year,
                dt.month year,
                event dt
```



All Rows Fetched: 383 in 0.675 seconds



Picture 26 Explain plan

#### Conclusion.

Despite the fact that the COST parameter in 3NF is much lower, the query in 3NF was executed in 2.39 seconds, and in DWH - for 0.68, what demonstrates the efficiency of DWH.