

## Problems for Lecture 1

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### problem 1.

(a) Statement `char *s[4] = {"this", "that", "we", "!"};`

declares a pointer array each of which is pointing to the beginning of a string. In this case, `s` is also a pointer pointing to the beginning of array `s[4]`. Then `s[2][1]` means `*(*(s+2)+1)`, which gives the value `'e'`. similarly, `s[0][0]` is `'t'` and `s[0][4]` is `'\0'`.

(b) As said in (a), `s` is a pointer to pointer, so we can declare a function like `void myfunc(char **s)`, and use `myfunc(s)` to call the function.

If the array is declared as `char t[4][5]`, it can be passed to a function like `void myfunc(char t[4][5])` by calling `myfunc(s)`.

### problem 2.

(a) **0.0:** `0 00000000 000000000000000000000000` (the exponent part is coded by eight zeros because of stipulation of IEEE)

**1.0:** `0 01111111 000000000000000000000000`

**1/3:** The science notation of 1/3 in binary system is:

**0.01010...**

.

If we code this as

`0 01111111 01010101010101010101010`,

because now  $0 < e < 255$ , the computer will retrieve the float by  $(1 + f) \times 2^{e-127}$ . It's wrong because the " $1 + f$ ".

So we need to rewrite 1/3 as

**$1.01010... \times 2^{-2}$**

,

then the exponent is now  $-2 + 127 = 125$ .

So finally, the answer is

`0 01111101 01010101010101010101010`

(b)

machine epsilon  $2^{-23}$ .

smallest positive number:  $1 \times 2^{-23} \times 2^{-126} = 2^{-149}$  (form 1) or  $(1 + 0) \times 2^{-126}$  (form 2).

Because  $e = 255$  is for special cases, so the  $e$  for largest positive number is  $e = 254$ , then the real exponent is  $254 - 127 = 127$ .

So the largest positive number can be calculated by  
 $(1 + 2^{-1} + \dots + 2^{-23}) \times 2^{127} = (2 - 2^{-23}) \times 2^{127}$ .

### problem 3.

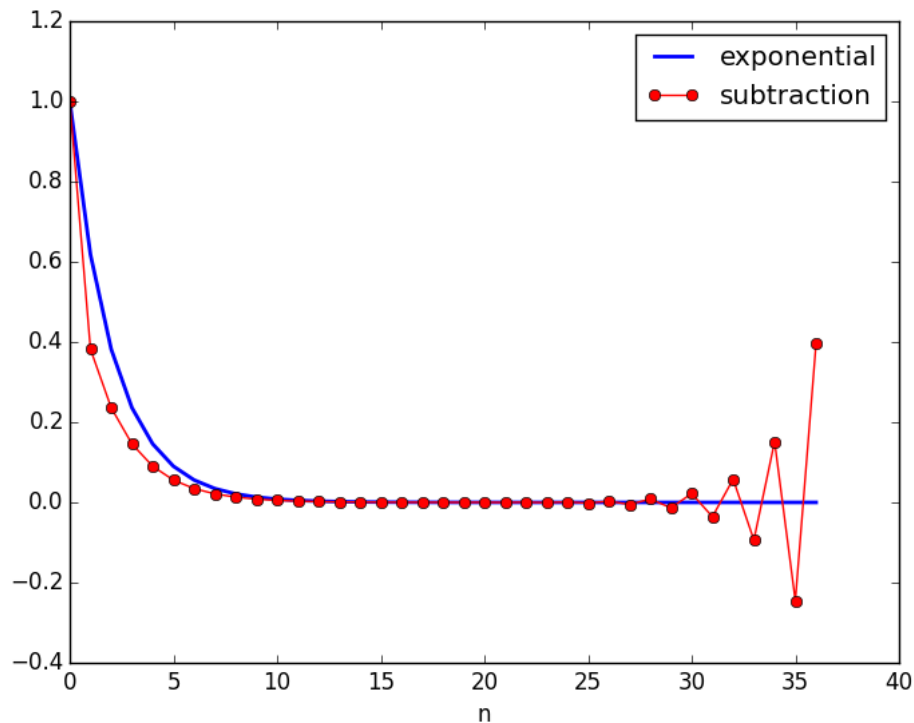
Substitute  $F_n = Ar^n$  to the recursion formula, then it's easy to find

$$r^2 + r - 1 = 0$$

The solution is  $r = \frac{-1 \pm \sqrt{5}}{2}$ . So the general solution is

$$F_n = A \left( \frac{-1 \pm \sqrt{5}}{2} \right)^n$$

For the golden-mean case,  $r = \frac{-1 + \sqrt{5}}{2}$  and  $A = 1$ . Because  $r < 1$ , the  $\Phi_n$  will go to 0 when increasing  $n$  and two adjacent number will be closer and closer. This just leads the instability of computation, which is clearly shown in the following figure.



(32-bit-float numbers are used in this plot.)

