

Quadratically constrained quadratic problem

$$\begin{array}{ll}\min_{x \in \mathbb{R}^N} & x^T C x \\ \text{s.t.} & x^T F : x \geq g_i, i = 1, \dots, p \\ & x^T B : x \leq b_i, i = 1, \dots, l \\ & x^T H : x = h, i = 1, \dots, k\end{array}$$

The matrices are real and symmetric.

Remember, that $f(x) = x^T C x$ is convex if $c \succeq 0$.

What makes the feasible set convex?

$$\{x \in \mathbb{R}^N : x^T F x \leq 1\}$$

If $F \succeq 0$ the set is convex. Is the following set convex?

$$\{x \in \mathbb{R}^N : x^T F x \geq 2\}$$

No, as this set is all the points outside of the circle with radius 2. If the inequality is flipped the set will be convex.

Let C and F be PSD

$$\begin{array}{ll}\min & x^T C x \\ \text{s.t.} & x^T F x \geq 1\end{array}$$

This problem is not convex.

$$\begin{array}{ll}\min_{x \in \mathbb{R}^N} & x^T C x \\ \text{s.t.} & x_i^2 = 1, i = 1, \dots, p\end{array}$$

The set is discrete as only few points fulfill the properties. The problem is not convex.

Trace of a matrix

The trace tr of a matrix is the sum of the elements on the main diagonal.

$$x^T C x = \text{tr}(x^T C x)$$

The tr -operator is invariant to cyclic permutation:

$$\text{tr}(ABCD) = \text{tr}(DABC)$$

This means that

$$x^T C = \text{tr}(x^T C x) = \text{tr}(x x^T C)$$

Notice that $x x^T \in \mathbb{R}^{N \times N}$ and $x^T x \in \mathbb{R}$. Let $w = x x^T$, which gives $\text{rank}(w) = 1$.

Making a non-convex problem convex

$$\begin{array}{ll} \min_{x \in \mathbb{R}^N} & x^T C x \\ \text{s.t.} & x^T F x \geq g \\ \Downarrow & \\ \min_{w \in \mathbb{R}^{N \times N}} & \text{tr}(C w) \\ \text{s.t.} & \text{tr}(F w) \geq g \\ & w \succeq 0 \\ & \text{rank}(w) = 1 \end{array}$$

The second problem is linear in w but the problem is just as hard as the first problem as $\text{rank}(w) = 1$. If $\text{rank}(w) = 1$ is skipped, it is called semidefinite relaxation:

$$\begin{array}{ll} \min_{w \in \mathbb{R}^{N \times N}} & \text{tr}(C w) \\ \text{s.t.} & \text{tr}(F w) \geq g \\ & w \succeq 0 \end{array}$$

The problem is that $\text{rank}(w)$ is usually large and the problem therefore does not give the same solutions as the original problem.

Let $\text{rank}(w) = r$.

$$\begin{aligned} w &= \sum_{i=1}^r \lambda_i v_i v_i^T \\ \lambda_1 &\geq \lambda_2 \geq \dots \geq \lambda_r \\ v_i, & i = 1, \dots, r \end{aligned}$$

Then $w \approx \lambda_1 \lambda_1 v_1^T$. This means that $x \approx \sqrt{\lambda_1} v_1 \sqrt{\lambda_1} v_1^T$.

Let $S^{N \times N}$ be the set of all symmetric $N \times N$ -matrices.
Let $\xi \in \mathbb{R}^N \sim \mathcal{N}(0, Q)$.

$$\begin{array}{ll} \min & E [\xi^T C \xi] \\ \text{s.t.} & E[\xi^T F \xi] \geq g \\ & Q \succeq 0 \end{array}$$

Remember that $E [\xi \xi^T]$ and $\xi^T C \xi = \text{tr}(C \xi \xi^T)$. As both are linear:

$$\begin{aligned} \text{tr} (E [\xi^T C \xi]) &= E [\text{tr} (\xi^T C \xi)] \\ &= E [\text{tr} (C \xi \xi^T)] \\ &= \text{tr} (C E [\xi \xi^T]) \\ &= \text{tr}(CQ) \end{aligned}$$

The problem is now

$$\begin{array}{ll} \min_Q & \text{tr}(CQ) \\ \text{s.t.} & \text{tr}(FQ) \geq g \\ & Q \succeq 0 \end{array}$$

And $Q = w$.