CECS 328: Program #1

For each of the tasks, I have included the algorithm in C# as well as a screenshot of the output. The source code for the project is available here:

 $https://github.com/Vardominator/CSULBProjects/tree/master/CECS328_DataStructures Algorithms$

1. Write a program to calculate S(n) by calculating the values of the Fibonacci sequence recursively.

Source:

```
public static long Sn(long number)
{
    long result = 0;
    for (int i = 1; i <= number; i++)
        {
            result += Fibonacci(i);
        }
    return result;
}

public static long Fibonacci(long number)
{
    if(number <= 2)
        {
        if(number == 0)
            {
                 return 0;
            }
            return 1;
    }

    return Fibonacci(number - 1) + Fibonacci(number - 2);
}</pre>
```

Sample run:

```
#region Problem 1: Calculate S(n) by calculating the values of the Fibonacci sequence recursively 
// S(n) definition: S(n) = f(0) + f(1) + ... + f(n)
int sampleN = 20;
Console.WriteLine($"Problem 1 Result: \{Sn(sampleN)\} \setminus n \setminus n");
#endregion
```

C:\WINDOWS\system32\cmd.exe

Problem 1 Result: 17710

2. Write a non-recursive program to calculate S(n).

Source:

```
public static long NonRecursiveSn(long number)
{
    long previous = 1;
    long previous2 = 1;
    long current = 0;
    long sum = 2;

    for (int i = 3; i <= number; i++)
    {
        current = previous + previous2;
        sum += current;
        previous = previous2;
        previous2 = current;
    }

    return sum;
}</pre>
```

Sample run:

```
#region Problem 2: Write a non-recursive program to calculate S(n)
Console.WriteLine("Problem 2 Result: \n");
Console.WriteLine($"Fibonacci Table for n = {sampleN}:");

for (int i = 0; i <= sampleN; i++)
{
    Console.WriteLine($"f({i}) = {Fibonacci(i)}");
}
Console.WriteLine("-----");
Console.WriteLine($"Sum: {NonRecursiveSn(sampleN)}\n\n");
#endregion</pre>
```

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- 3. Algebraically verify that g(n) is a solution of Equ 1 by substituting g(n) in Equ 1. (This portion is attached to the assignment packet)
- 4. Write a third iterative program by summing g(n) over n.

Source:

```
public static long GrimaldiSum(long number)
{
    long result = 0;
    for (int i = 1; i <= number; i++)
        {
            result += Grimaldi(i);
        }
    return result;
}

public static long Grimaldi(long k)
{
    return (long)((1 / Math.Sqrt(5)) * (Math.Pow(((1 + Math.Sqrt(5)) / 2), k) - Math.Pow(((1 - Math.Sqrt(5)) / 2), k)));
}</pre>
```

Sample run:

```
#region Problem 4: Third possible way to sum Fibonnaci
Console.WriteLine($"Problem 4: {GrimaldiSum(20)}\n\n");
#endregion

C:\WINDOWS\system32\cmd.exe
Problem 4: 17710
```

5. Calculate these values of S for n = 10, 20, 30. Calculates values of f for n = 12, 22, 32.

Sample run:

- 6. Prove the identity S(n) = f(n + 2) 1 using induction.
 - (This portion is attached to the assignment packet)
- 7. Example run using the alternate sum from #6.

Source:

```
public static long SnAlt2(long number)
{
    return Fibonacci(number + 2) - 1;
}
```

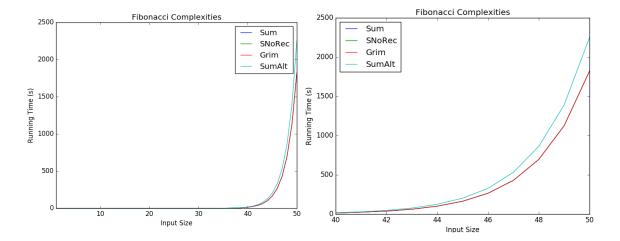
Sample run:

```
#region Problem 7: Fourth way to calcualte S(n)
Console.WriteLine($"Problem 7: {SnAlt2(20)}\n\n");
#endregion

C:\WINDOWS\system32\cmd.exe
Problem 7: 17710
```

- 8. Since I am using Int64 (long in C#) to calculate the sums, the largest n that can be used is such that S(n) does not exceed 9,223,372,036,854,775,807. Another limitation would be the possibility of running out of stack space. Since the Fibonacci calls are recursive, a sufficiently large n would cause the program to crash.
- 9. Compare the running times of the 4 methods.

I computed the sum with each method for values of n starting at 1 and ending at 50. I recorded the running times for each. Finally, I ran a python script to plot the growth of the running times versus the input size.



10. Summary and conclusions.

We see in the plots above that the 4^{th} method grows the fastest by a considerable amount. However, the difference is not much. The drawbacks of the recursive methods are that they have linear growth in space complexity: O(n), as opposed to the non-recursive methods that have a constant space complexity.

As far as time complexity is concerned, the recurrence relation for the Fibonacci sequence is T(n) = T(n-1) + T(n-2) + O(1). Using a recursion tree, we can prove that the time complexity is $O(2^n)$. This means that the recursive Fibonacci sequence grows relatively fast. It is faster than $O(n^3)$, for example.

Finally, it is worth noting that a recursive solution to the sum is slower than the Grimaldi solution. I have shown this in the plots above.