

02	M	T	W	T	F	S	S
05						1	2
06	3	4	5	6	7	8	9
07	10	11	12	13	14	15	16
08	17	18	19	20	21	22	23
09	24	25	26	27	28	29	

## Operating System

## Assignment - 2

Date of Submission :- 02/04/2020

### Question / Answer

Q1] Implement the Dekker's solution in order to solve the critical section problem in Producer - Consumer with finite buffer

Ans :- Dekker's solution implementation on producer process  $\Rightarrow$

```
while (true) {  
    while (counter == BUFFER_SIZE)  
        ; // do nothing  
    buffer[in] = nextProduced;  
    in = (in + 1) % BUFFER_SIZE;  
    flag[i] = true;  
    while (flag[j] == true)  
    {  
        if (turn == j)  
        {  
            flag[i] = false;  
            while (turn == j)  
            {  
            }  
            flag[i] = true;  
        }  
    }  
}
```

NOTES

```

00      }
00      }
00      counter ++ ;      // Critical Section
00      turn = j ;
11 00      flag[i] = false
12 00      //remainder section
12 00  }

```

\* Dekker's solution implementation on consumer process  $\Rightarrow$

```

2 00      while (true) {
3 00          while (counter == 0)
4 00              ; // do nothing
4 00          next consumed = buffer[out];
5 00          out = (out + 1) % BUFFER_SIZE;
5 00          flag[i] = true;
6 00          while (flag[j] == true)
6 00              {
7 00                  if (turn == j)
7 00                      {
8 00                          flag[i] = false;
8 00                          while (turn == j)
8 00                              {
9 00                                  }
9 00                          flag[i] = true;
9 00                      }
9 00              }
9 00          counter -- ; // consumer critical section

```

NOTES



02	M	T	W	T	F	S	S
03						1	2
04	3	4	5	6	7	8	9
05	10	11	12	13	14	15	16
06	17	18	19	20	21	22	23
07	24	25	26	27	28	29	

2020  
Saturday  
January

11

011-355 • WK 02

```
turn = j;
flag[i] = false;
// Remainder section
```

}

Q2] Implement the Peterson's solution in order to solve the critical section problem in Producer Consumer with finite buffer.

Ans:— Peterson's solution implementation on Producer process  $\Rightarrow$

```
while (true) {
    while (counter == BUFFER_SIZE)
        ; // do nothing
    buffer[in] = nextProduced;
    in = (in + 1) % BUFFER_SIZE;
    flag[i] = true;
    turn = j;
    while (flag[j] && turn == j)
        ;
}
```

```
counter++; // Critical section
flag[i] = false;
// Remainder section
```

}

NOTES

Peterson's solution implementation on  
Consumer process :-

```

while (true) {
    while (counter == 0)
        ; // do nothing
    nextConsumed = buffer[out];
    out = (out + 1) % BUFFER_SIZE;
    flag[i] = true;
    turn = i;
    while (flag[j] && turn == j)
        ;
    counter--; // critical section
    flag[i] = false;
    // Remainder section
}

```

Q3] What is Dining Philosophers Problem?

Ans → Consider five philosophers who spend their lives thinking and eating. The philosophers share a circular table surrounded by five chairs, each belonging to one philosopher. In the centre of the table is a bowl of rice, and the table is laid with five single chopsticks.

NOTES



Vaani Verma

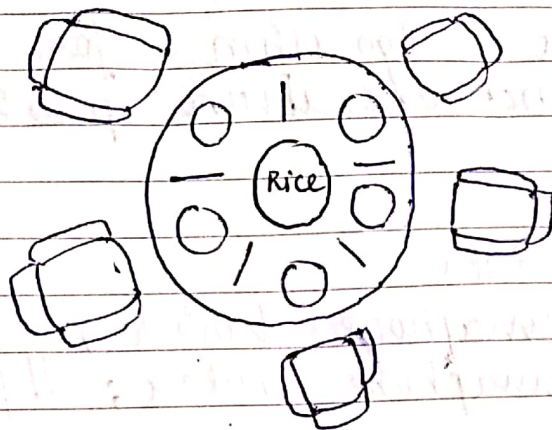
February 2020

02	M	T	W	T	F	S	S
03						1	2
04	3	4	5	6	7	8	9
05	10	11	12	13	14	15	16
06	17	18	19	20	21	22	23
07	24	25	26	27	28	29	

2020  
Tuesday  
January

14

014-352 • WK 03



Situation of the  
dinning  
philosophers.

When a philosopher thinks, she does not interact with her colleagues. From time to time, a philosopher gets hungry and tries to pick up the two chopsticks that are closest to her (the chopsticks that are b/w her and her left & right neighbours). A philosopher may pick up only one chopstick at a time. Obviously, she cannot pick up a chopstick that is already in the hand of a neighbour. When a hungry philosopher has both her chopsticks at the same time, she eats without releasing her chopsticks. When she is finished eating, she puts down both of her chopsticks and starts thinking again.

NOTES

15

015-351 • WK 03

2020

Wednesday

January

Varun Verma

01

January 2020

wk	M	T	W	T	F	S	S
01			1	2	3	4	5
02	6	7	8	9	10	11	12
03	13	14	15	16	17	18	19
04	20	21	22	23	24	25	26
05	27	28	29	30	31		

9.00 Q4] write an algorithm for the sleeping  
Barber synchronization problem using  
10.00 semaphores

11.00 Ans → shared data

12.00 semaphore barber;  
semaphore mutex; // for mutual  
exclusion

1.00 semaphore customer; // no. of customer waiting  
2.00 int Numberofseats; // no. of seats vacant

Initially

3.00 barber = 0;  
4.00 mutex = 1;  
customer = 0;  
5.00 Numberofseats = n;

The structure of Barber Process

do {

7.00 wait (customer); // if none avail. go to sleep.  
wait (mutex);

Numberofseats ++;

signal (barber);

signal (mutex);

//haircut

NOTES } while (true);



02	M	T	W	T	F	S	S
03						1	2
04	3	4	5	6	7	8	9
05	10	11	12	13	14	15	16
06	17	18	19	20	21	22	23
07	24	25	26	27	28	29	

2020  
Thursday  
January

16

016-350 • WK 03

## The structure of Customer Process

do {

wait(mutex); // enter critical section  
if (Number of Seats > 0)

{

Number of Seats --;

signal (customer);

signal (mutex);

wait (barber);

// haircut

}

else {

signal (mutex);

// leave without a haircut

}

} while (true);

Q5) Give the solution of second Reader-Writer problem by using the concept of semaphore?

Ans → shared Data

int readcount;

int writecount;

semaphore rmutex;

semaphore wmutex;

semaphore readtry;

semaphore resource;

wk	M	T	W	T	F	S	S
01			1	2	3	4	5
02	6	7	8	9	10	11	12
03	13	14	15	16	17	18	19
04	20	21	22	23	24	25	26
05	27	28	29	30	31		

Initially

```
readcount = 0;
writecount = 0;
rmutex = 1;
wmutex = 1;
readtry = 1;
resource = 1;
```

The Structure of a Reader Process

do {

```
wait(readtry);
wait(rmutex);
readcount++;
if (readcount == 1)
    wait(resource);
signal(rmutex);
signal(readtry);
```

// Reading is performed

```
wait(wmutex);
readcount--;
if (readcount == 0)
    signal(resource);
signal(wmutex);
} while (true);
```



02	M	T	W	T	F	S	S
wk						1	2
05						8	9
06	3	4	5	6	7	8	9
07	10	11	12	13	14	15	16
08	17	18	19	20	21	22	23
09	24	25	26	27	28	29	

## The structure of a writer Process

do {

wait (wmutex);

wcount ++;

if (wcount == 1)

wait (readtry);

signal (wmutex);

wait (resource);

// writing is performed

signal (resource);

wait (wmutex);

wcount --;

if (wcount == 0)

signal (readtry);

signal (wmutex);

} while (true);