

ARRAYS-2

Question: Given an array with n objects colored 0, 1 or 2, sort them in-place so that objects of the same color are adjacent, with the colors in the order red, white and blue.

YOU HAVE 10 MINUTES

EXAMPLE:

INPUT: [2,0,2,1,1,0]

OUTPUT: [0,0,1,1,2,2]

Link: <https://leetcode.com/problems/sort-colors/>

We can start by maintaining 2 pointers. One pointer for the rightmost 0 in the array and one pointer for the leftmost 2.

Additionally, we can have a pointer called current or iterator, that start from 0th index and moves left -> right.

Whenever $\text{nums}[\text{iterator}] = 0$, we swap $\text{nums}[\text{iterator}]$ with $\text{nums}[\text{rightMost0}]$ and increment the pointer to the right most 0. When $\text{nums}[\text{iterator}] = 2$ we swap $\text{nums}[\text{iterator}]$ with $\text{nums}[\text{leftMost2}]$ and decrement the pointer to the left most 2.

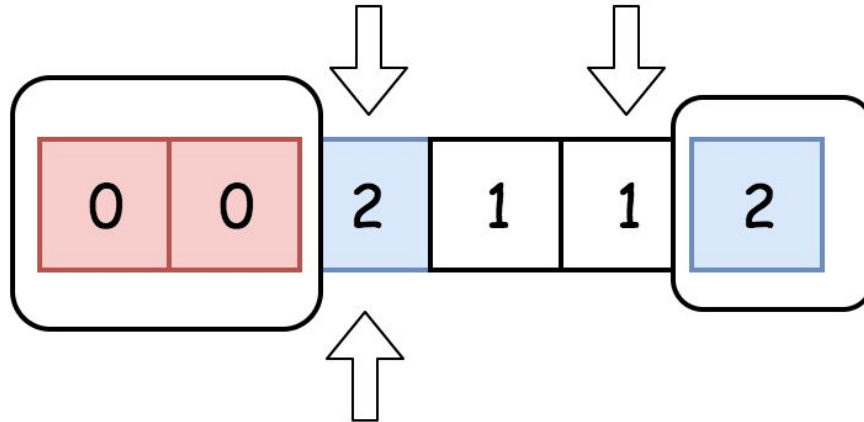
When $\text{nums}[\text{iterator}] = 1$, we do nothing and just move right.

We do this till current/iterator is less than leftMost2.

We will partition the array in 3 groups when we are done.

p_0 = rightmost boundary of 0s

p_2 = leftmost boundary of 2s



$curr$ = index of current element
 $curr \geq p_0$

SOLUTION (SINGLE PASS)

```
1  # Sort colors using Dutch National Flag
2
3  class Solution:
4      def sortColors(self, nums: List[int]) -> None:
5          rightMost0 = 0
6          leftMost2 = len(nums) - 1
7          iteratorPointer = 0
8
9          while iteratorPointer <= leftMost2:
10             if nums[iteratorPointer] == 0:
11                 nums[rightMost0], nums[iteratorPointer] = nums[iteratorPointer], nums[rightMost0]
12                 rightMost0 += 1
13                 iteratorPointer += 1
14             elif nums[iteratorPointer] == 2:
15                 nums[iteratorPointer], nums[leftMost2] = nums[leftMost2], nums[iteratorPointer]
16                 leftMost2 -= 1
17             else:
18                 iteratorPointer += 1
```

TIME COMPLEXITY - $O(N)$ and SPACE COMPLEXITY - $O(1)$. Constant space

Question: Suppose an array sorted in ascending order is rotated at some pivot unknown to you beforehand. You are given a target value to search. If found in the array return its index, otherwise return -1. YOU HAVE 15 MINUTES

EXAMPLE:

INPUT: [4,5,6,7,0,1,2], target = 0

OUTPUT: 4

Link:

<https://leetcode.com/problems/search-in-rotated-sorted-array/>

PLEASE LOOK UP BINARY SEARCH IN THE FOLLOWING
ARTICLE:

<https://www.geeksforgeeks.org/binary-search/>

Essentially a sorted and rotated array can be thought as 2 sorted arrays appended.

We can apply binary search individually on these two sub arrays.

If we can find the pivot, we can simply do binary search between (0, nums[pivot-1]) and (nums[pivot +1], nums[nums.size-1]). Assuming of course nums[pivot] != target. Else we just return pivot.

But we need to find the pivot, in $O(\log N)$! (modify binary search ?)

Essentially the final answer must search using findPivot and regular binarySearch as subroutines/utility functions.

SOLUTION (pivoted-BINARY SEARCH)

```
def search(self, nums: List[int], target: int) -> int:
    if len(nums) == 0:
        return -1
    pivot = self.pivot(nums, 0, len(nums)-1)
    if pivot == -1:
        return self.binarySearch(nums, 0, len(nums) - 1, target)
    if nums[pivot] == target:
        return pivot
    if nums[0] <= target:
        return self.binarySearch(nums, 0, pivot-1, target)
    return self.binarySearch(nums, pivot+1, len(nums) - 1, target)
```

TIME COMPLEXITY - $O(\log(N))$

SOLUTION (find PIVOT)

```
1 class Solution:
2
3     def pivot(self, nums: List[int], low: int, high: int)->int:
4         if high < low:
5             return -1
6         if high == low:
7             return low
8         mid = int((low + high)/2)
9         if mid < high and nums[mid] > nums[mid + 1]:
10             return mid
11         if mid > low and nums[mid] < nums[mid - 1]:
12             return mid - 1
13         if nums[low] >= nums[mid]:
14             return self.pivot(nums, low, mid-1)
15         return self.pivot(nums, mid + 1, high)
```

TIME COMPLEXITY - $O(\log(N))$

Question: Write an efficient algorithm that searches for a value in an $m \times n$ matrix. This matrix has the following properties:

1. Integers in each row are sorted in ascending from left to right.
2. Integers in each column are sorted in ascending from top to bottom.

Constraint -> IN LINEAR TIME

YOU HAVE 15 MINUTES

Link: <https://leetcode.com/problems/search-a-2d-matrix-ii/>

FIND target = 5

| | | | | |
|-----------|-----------|-----------|-----------|-----------|
| 1 | 4 | 7 | 11 | 15 |
| 2 | 5 | 8 | 12 | 19 |
| 3 | 6 | 9 | 16 | 22 |
| 10 | 13 | 14 | 17 | 24 |
| 18 | 21 | 23 | 26 | 30 |

FIND target = 5

| | | | | |
|-----------|-----------|-----------|-----------|-----------|
| 1 | 4 | 7 | 11 | 15 |
| 2 | 5 | 8 | 12 | 19 |
| 3 | 6 | 9 | 16 | 22 |
| 10 | 13 | 14 | 17 | 24 |
| 18 | 21 | 23 | 26 | 30 |

FIND target = 13

| | | | | |
|-----------|-----------|-----------|-----------|-----------|
| 1 | 4 | 7 | 11 | 15 |
| 2 | 5 | 8 | 12 | 19 |
| 3 | 6 | 9 | 16 | 22 |
| 10 | 13 | 14 | 17 | 24 |
| 18 | 21 | 23 | 26 | 30 |

```

1 # Search in 2D Matrix
2 class Solution:
3     def searchMatrix(self, matrix, target):
4         if len(matrix) == 0 or len(matrix[0]) == 0:
5             return False
6
7         noOfRows = len(matrix)
8         noOfColumns = len(matrix[0])
9
10        row = 0
11        col = noOfColumns - 1
12
13        while row < noOfRows and col >= 0:
14            # print(matrix[row][col])
15            if(matrix[row][col] > target):
16                col -= 1
17            elif(matrix[row][col] < target):
18                row += 1
19            else:
20                return True
21        return False
22

```

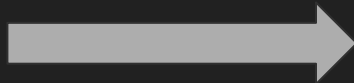
TIME
 COMPLEXITY -
 $O(M + N)$.
 For a $M \times N$
 matrix

Question: You are given an $n \times n$ 2D matrix representing an image.

Rotate the image by 90 degrees (anti - clockwise).

Let's just discuss first.

| | | | |
|----|----|----|----|
| 1 | 2 | 3 | 4 |
| 5 | 6 | 7 | 8 |
| 9 | 10 | 11 | 12 |
| 13 | 14 | 15 | 16 |



| | | | |
|---|---|----|----|
| 4 | 8 | 12 | 16 |
| 3 | 7 | 11 | 15 |
| 2 | 6 | 10 | 14 |
| 1 | 5 | 9 | 13 |

Rotate boundary elements and then move inwards ?

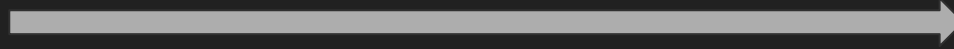
| | | | |
|----|----|----|----|
| 1 | 2 | 3 | 4 |
| 5 | 6 | 7 | 8 |
| 9 | 10 | 11 | 12 |
| 13 | 14 | 15 | 16 |

| | | | |
|----|----|----|----|
| 1 | 2 | 3 | 4 |
| 5 | 6 | 7 | 8 |
| 9 | 10 | 11 | 12 |
| 13 | 14 | 15 | 16 |

| | | | |
|----|----|----|----|
| 1 | 2 | 3 | 4 |
| 5 | 6 | 7 | 8 |
| 9 | 10 | 11 | 12 |
| 13 | 14 | 15 | 16 |



| | | | |
|---|----|----|----|
| 4 | 8 | 12 | 16 |
| 3 | 6 | 7 | 15 |
| 2 | 10 | 11 | 14 |
| 1 | 5 | 9 | 13 |

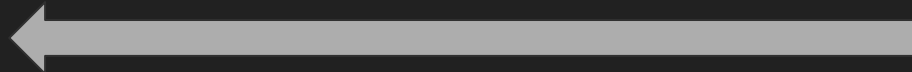


3 SWAPS

| | | | |
|---|----|----|----|
| 4 | 2 | 3 | 16 |
| 5 | 6 | 7 | 8 |
| 9 | 10 | 11 | 12 |
| 1 | 14 | 15 | 13 |

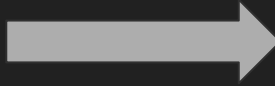


| | | | |
|---|----|----|----|
| 4 | 8 | 3 | 16 |
| 5 | 6 | 7 | 15 |
| 2 | 10 | 11 | 12 |
| 1 | 14 | 9 | 13 |



1 SWAP

| | | | |
|---|----|----|----|
| 4 | 8 | 12 | 16 |
| 3 | 6 | 7 | 15 |
| 2 | 10 | 11 | 14 |
| 1 | 5 | 9 | 13 |



| | | | |
|---|---|----|----|
| 4 | 8 | 12 | 16 |
| 3 | 7 | 11 | 15 |
| 2 | 6 | 10 | 14 |
| 1 | 5 | 9 | 13 |

For $N = 2$, there is only 1 boundary. You will have to swap once.

For $N = 3$, there are 2 boundaries. 2 swaps for outer boundary, 0 swap for inner boundary.

For $N = 4$, there are 2 boundaries. 3 swaps for outer. 1 swap for inner boundary.

For $N = 5$, there 3 boundaries. 4 swaps for outermost. 3 for next boundary. 0 for innermost boundary.

Question: You are given an $n \times n$ 2D matrix representing an image.
Rotate the image by 90 degrees (clockwise).

YOU HAVE 15 MINUTES

Link: <https://leetcode.com/problems/rotate-image/>

```

1  # Rotate a 2 dimensional matrix|
2
3  class Solution:
4      def rotate(self, matrix):
5          N = len(matrix[0])
6          for i in range(0, N // 2): # For boundaries
7              for j in range(i, N - i - 1): # For swaps
8                  tmp = matrix[N - 1 - j][i]
9                  matrix[N - 1 - j][i] = matrix[N - 1 - i][N - j - 1]
10                 matrix[N - 1 - i][N - j - 1] = matrix[j][N - 1 - i]
11                 matrix[j][N - 1 - i] = matrix[i][j]
12                 matrix[i][j] = tmp

```

TIME COMPLEXITY : $O(MN)$, for $M \times N$ matrix